
Artificial Intelligence: Adversarial Search *part 4* Minimax

- Russell & Norvig: Chapter 5

Today

■ Adversarial Search

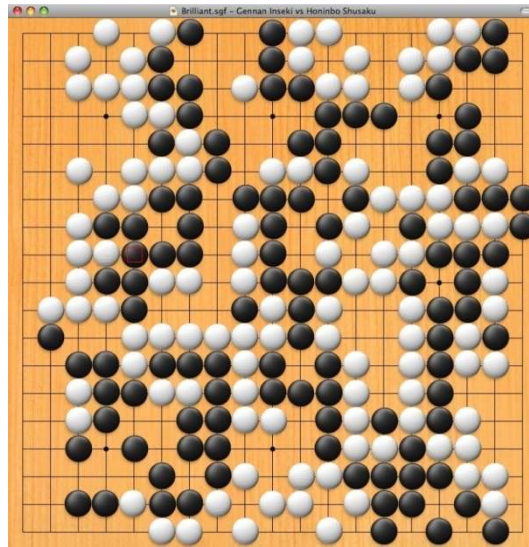
YOU ARE HERE!

1. Minimax) 4.1
2. Alpha-beta pruning) 4.2
3. Other Adversarial Search
 1. Multiplayer Games
 2. Stochastic Games
 3. Monte Carlo Tree Search

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4.3

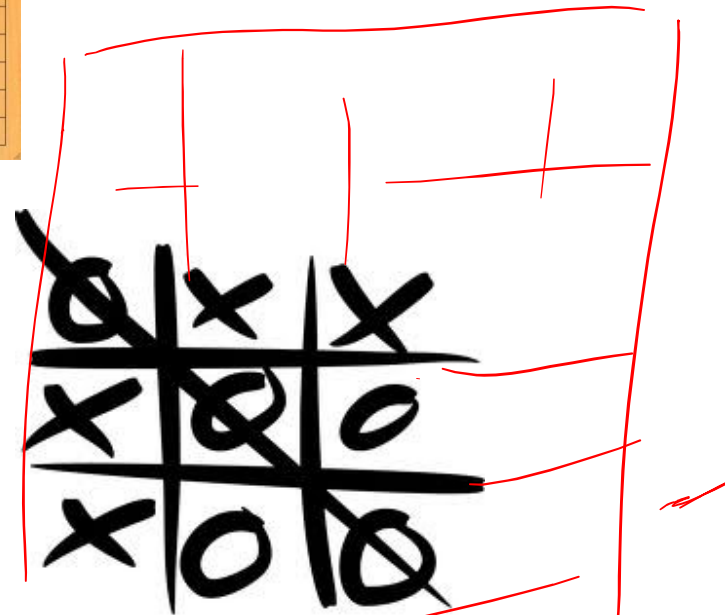
Motivation



GO



chess



tic-tac-toe

Games

1. Assumptions

1. Zero-sum game:

- no player has a game advantage. If the total gains of one player are added up, and the total losses are subtracted, they will sum to 0.

2. Players play rationally (i.e. to win)

2. Characteristics of Games

□ Number of players? 1, 2, 3+

□ Deterministic:

- the outcome of the game is only dependent on the moves of the players

□ Stochastic:

- chance involved

□ Perfect information:

- all players know the state of the game and all possible moves a.k.a. fully observable

□ Imperfect information:

- eg. a player is hiding their game aka. partially observable

	Deterministic	Stochastic
Perfect information	chess, checkers, go	backgammon, monopoly
Imperfect information	battleship	scrabble, poker, bridge

Techniques

Nb of players	Deterministic / Stochastic ?	Perfect / Imperfect information?	Example	Technique
1	deterministic	perfect info	8-puzzle	heuristic search
2	deterministic	perfect info	chess, checkers, go	minimax & alpha-beta
3+	deterministic	perfect info	Chinese Checkers	max ⁿ
1	stochastic	yes	2048	expectimax
2	stochastic	yes	backgammon	expectiminimax
3+	stochastic	yes		modified expectimax

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Minimax Search

- Game between two opponents, MIN and MAX
 - MAX tries to win, and
 - MIN tries to minimize MAX's score
- Existing heuristic search methods do not work
 - would require a helpful opponent
 - need to incorporate "hostile" moves into search strategy

- 2 flavors:

1. exhaustive Minimax



2. n-ply Minimax with Heuristic



Exhaustive Minimax Search

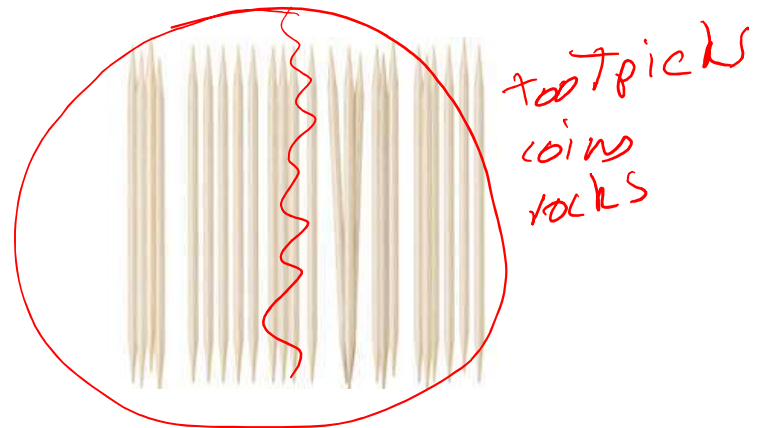


- For small games where exhaustive search is feasible
- Procedure:
 1. ✓ build complete game tree
 2. ✓ label each level according to player's turn (MAX or MIN)
 3. ✓ label leaves with a utility function to determine the outcome of the game
 - e.g., ^{MAX takes} (0, 1) or ^{MAX wins} (-1, 0, 1) ^{MAX wins}
 4. ✓ propagate this value up:
 - if parent=MAX, give it max value of children
 - if parent=MIN, give it min value of children
 5. ✓ Select best next move for player at the root as the move leading to the child with the highest value (for MAX) or lowest values (for MIN)

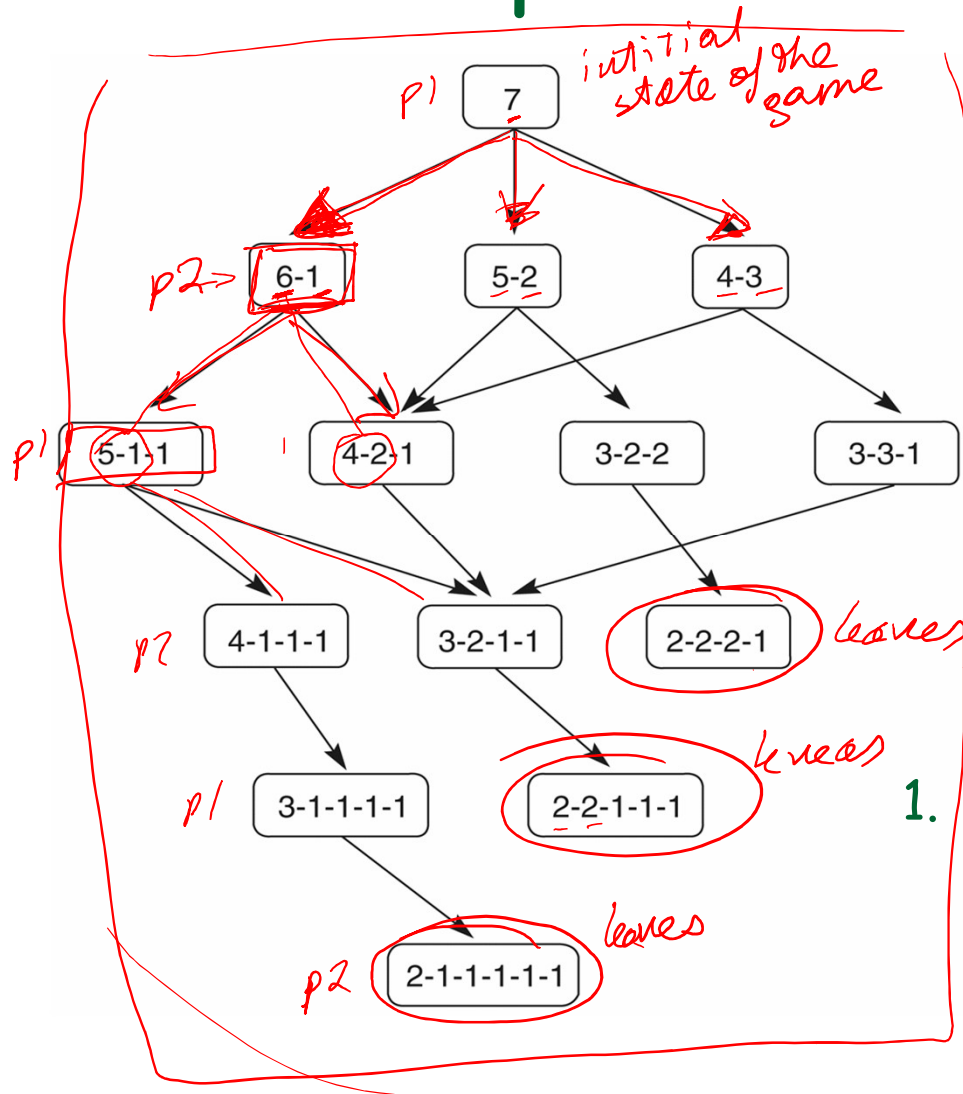
Example: Game of Nim

■ Rules

- ❑ 2 players start with a pile of tokens
- ❑ move: split (any) existing pile into two non-empty differently-sized piles
- ❑ game ends when no pile can be unevenly split
- ❑ player who cannot make their move loses



State Space of Game Nim



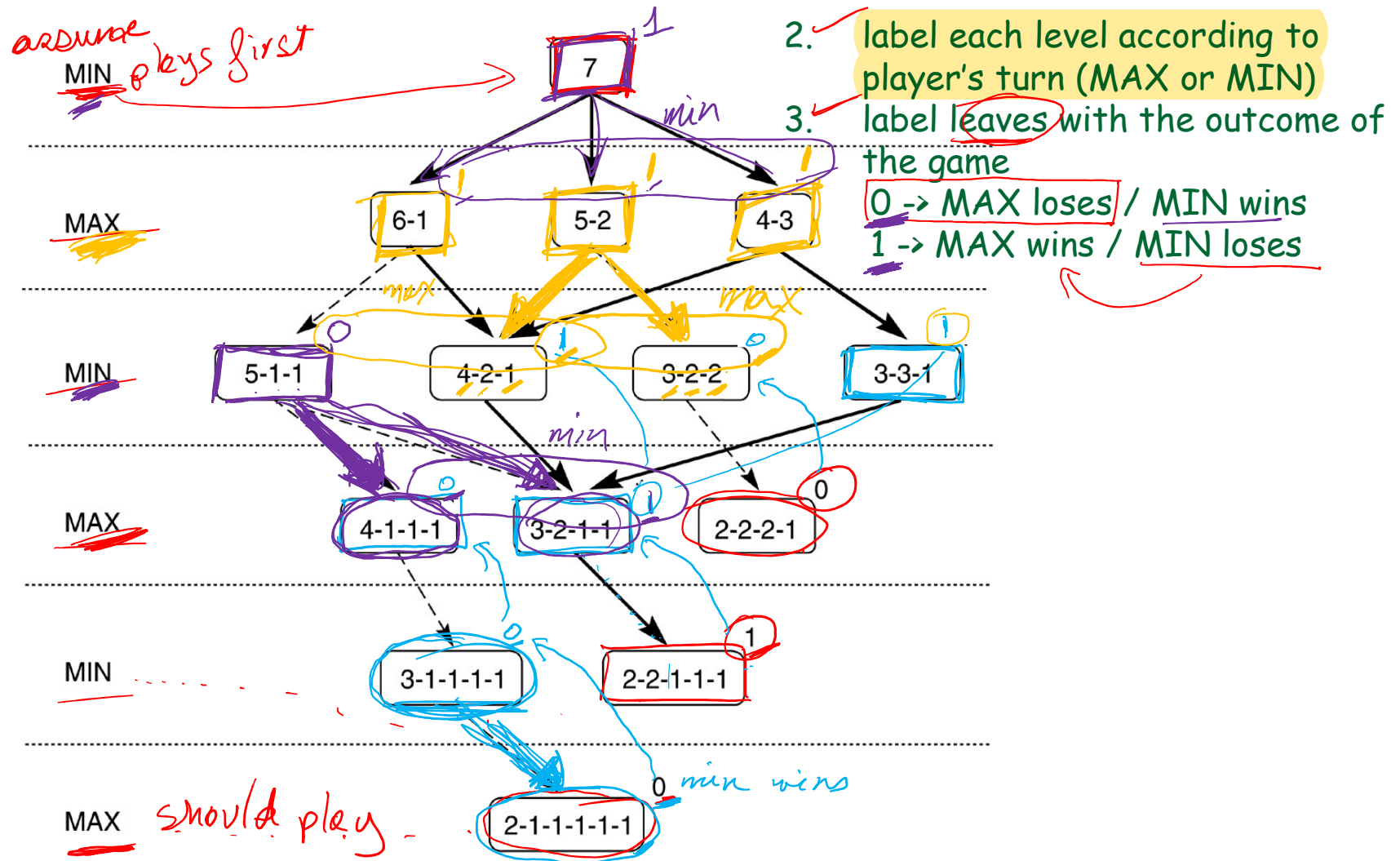
eg. start with one pile of 7 tokens

each step has to divide one pile into 2 non-empty piles of different sizes

player without a move left loses game

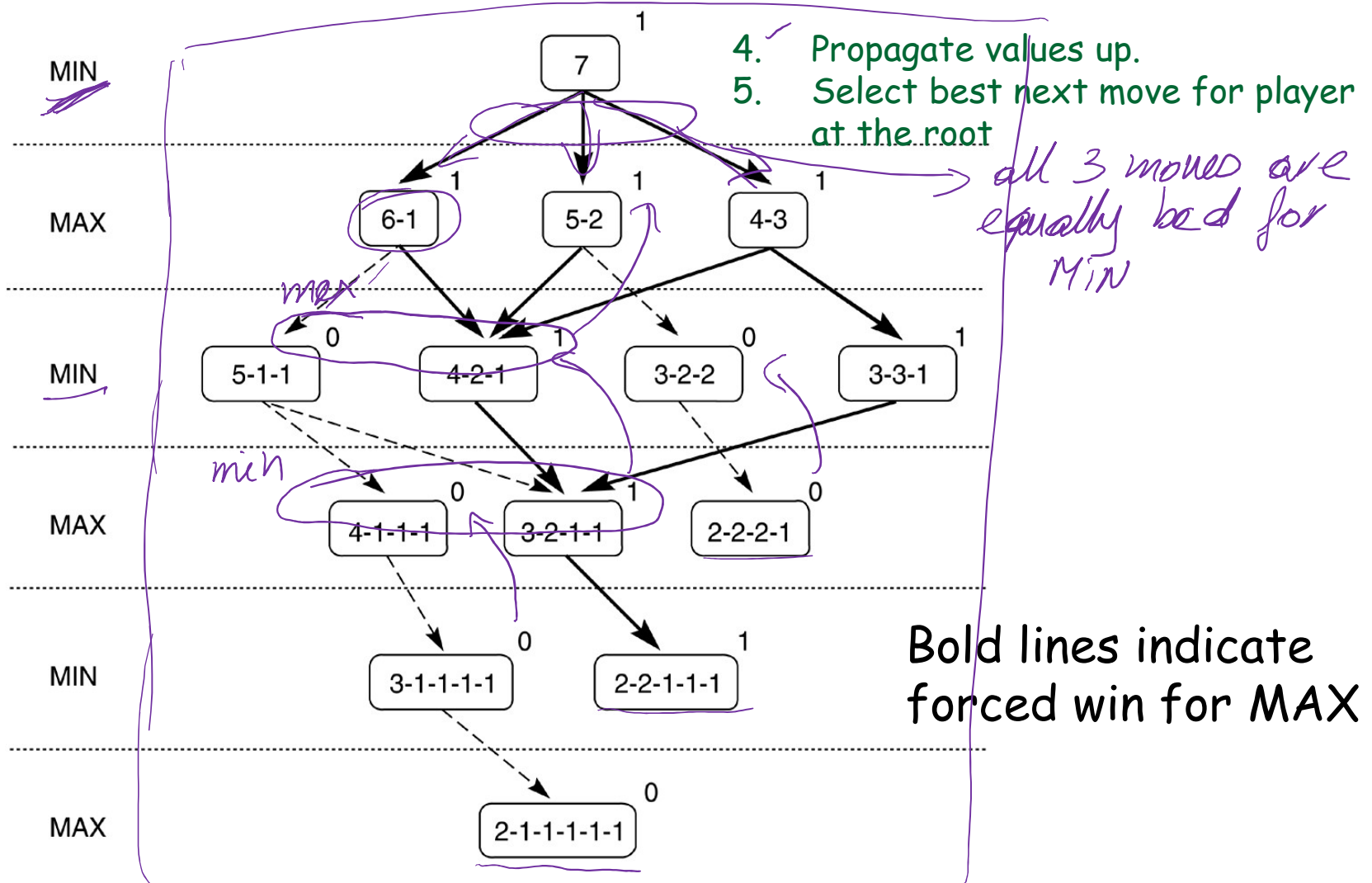
1. build complete game tree ✓

Exhaustive Minimax for Nim



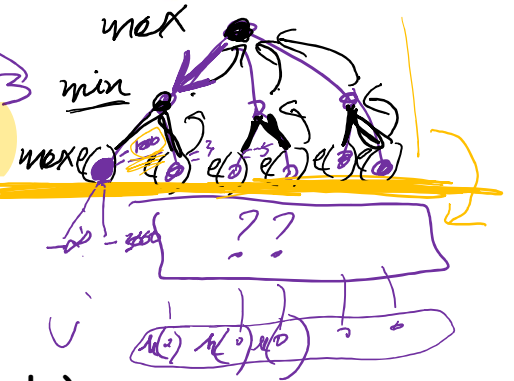
source: G. Luger (2005)

Exhaustive Minimax for Nim



source: G. Luger (2005)

- ## horizon

exact value of the leafs

nodes are evaluated with heuristics function $e(n)$

$e(n)$ indicates how good a state seems to be for MAX compared to MIN

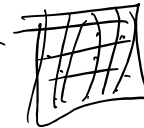
- suffers from the horizon effect

$$e\left(\begin{pmatrix} \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot \end{pmatrix}\right) = 30$$



Heuristic Function for 2-player games

- simple strategy:
 - try to maximize difference between MAX's game and MIN's game
- typically called $e(n)$
- $e(n)$ is a **heuristic** that estimates how favorable a node n is for MAX with respect to MIN
 - $e(n) > 0$ --> n is favorable to MAX
 - $e(n) < 0$ --> n is favorable to MIN
 - $e(n) = 0$ --> n is neutral



Choosing a Heuristic Function $e(n)$

- Usually $e(n)$ is a weighted sum of various features:

$$e(n) = \sum w_i f_i(n)$$

✓ dependent on
actual game

- E.g. of features:

- f_1 = number of pieces left on the game for MAX
- f_2 = number of possible moves left for MAX
- f_3 = \ominus (number of pieces left on the game for MIN)
- f_4 = \ominus (number of possible moves left for MIN)



- E.g. of weights:

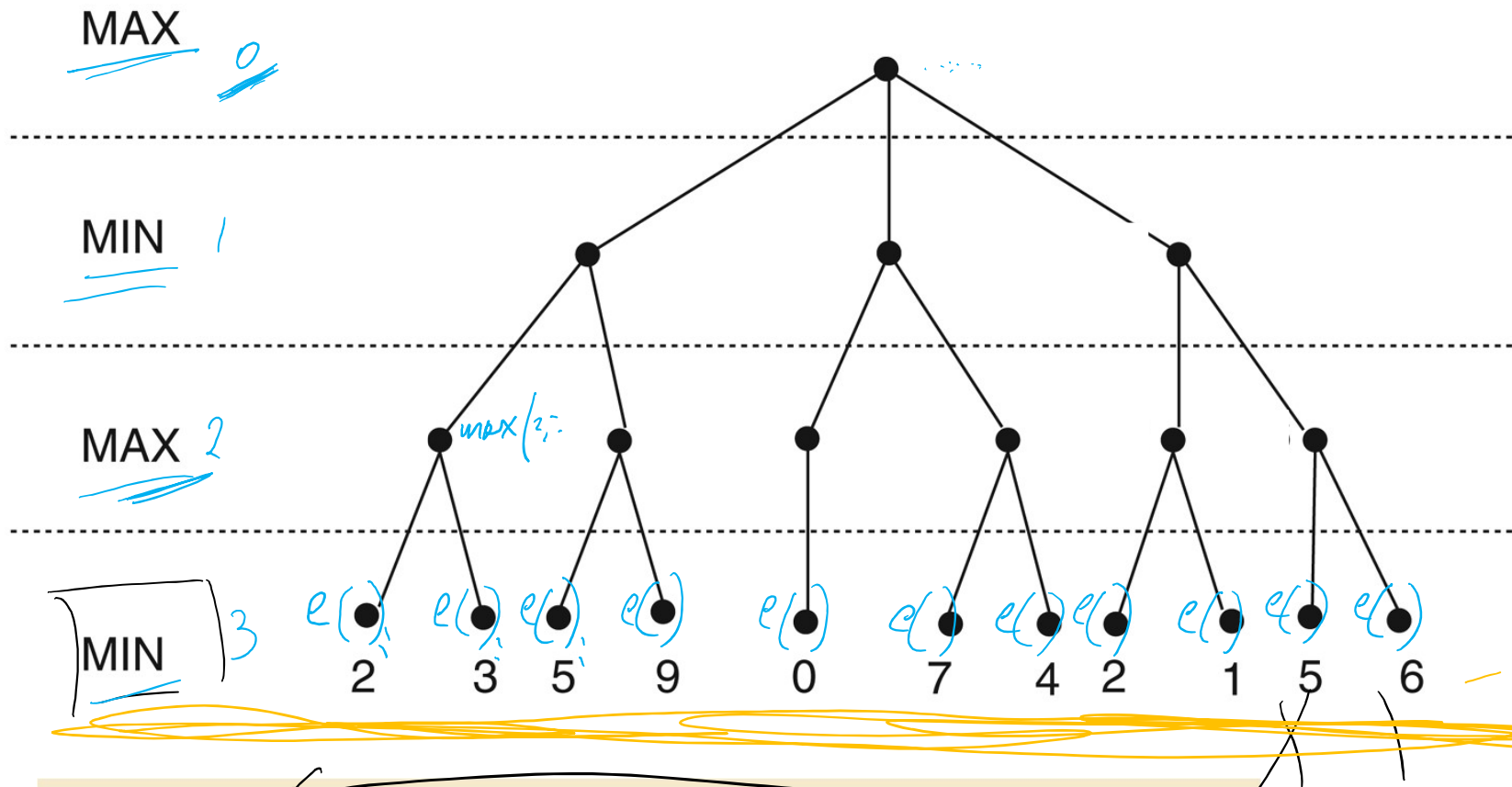
- $w_1 = 0.5$ // f_1 is a very important feature
- $w_2 = 0.2$ // f_2 is not very important
- $w_3 = 0.2$ // f_3 is not very important
- $w_4 = 0.1$ // f_4 is really not important

f_5 : queen still alive (1/0)

f_6 : # of bishops left

$w_5 = 0.3$

Minimax with 3-ply look-ahead

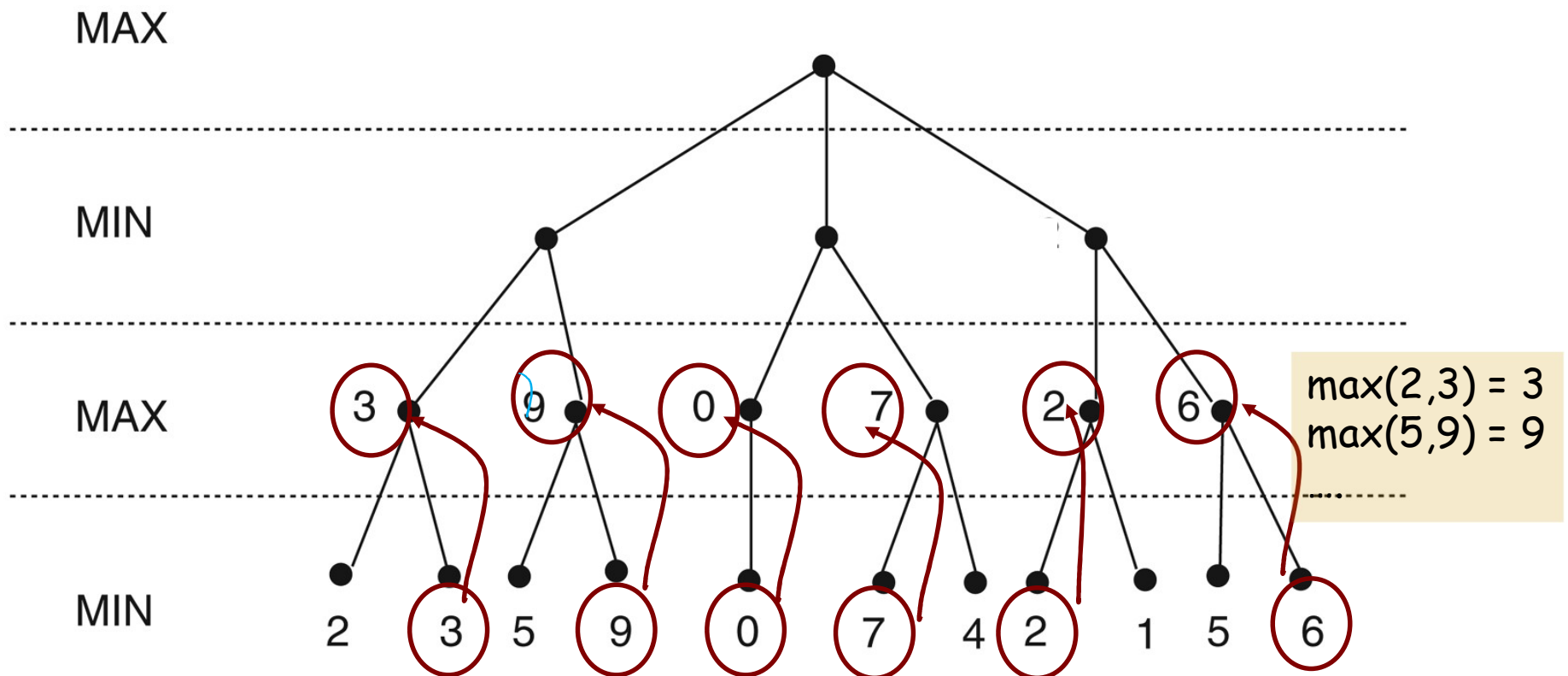


Leaf nodes show the actual heuristic value $e(n)$

source: G. Luger (2005)

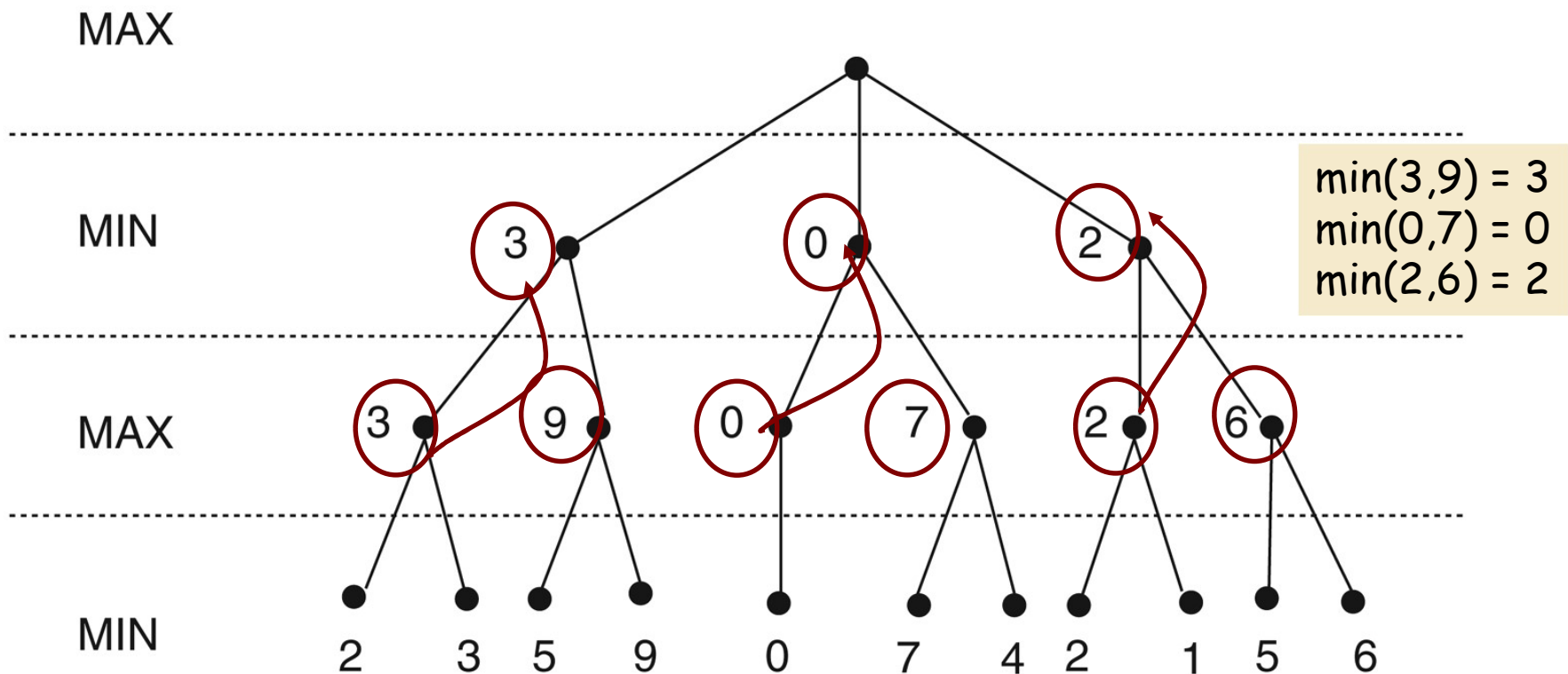
leaves = end of the game

Minimax with 3-ply look-ahead

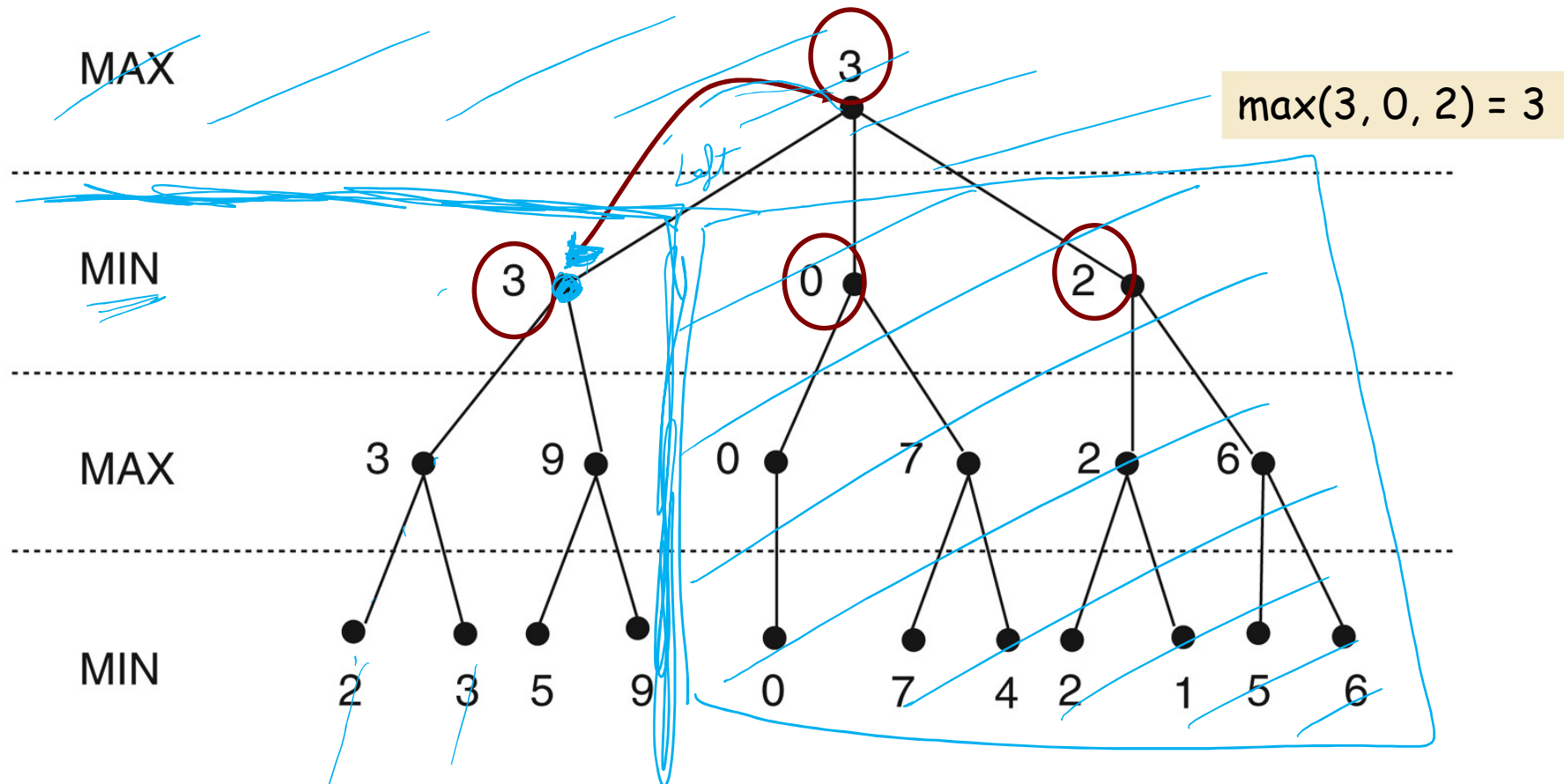


Leaf nodes show the actual heuristic value $e(n)$
Internal nodes show back-up heuristic value

Minimax with 3-ply look-ahead

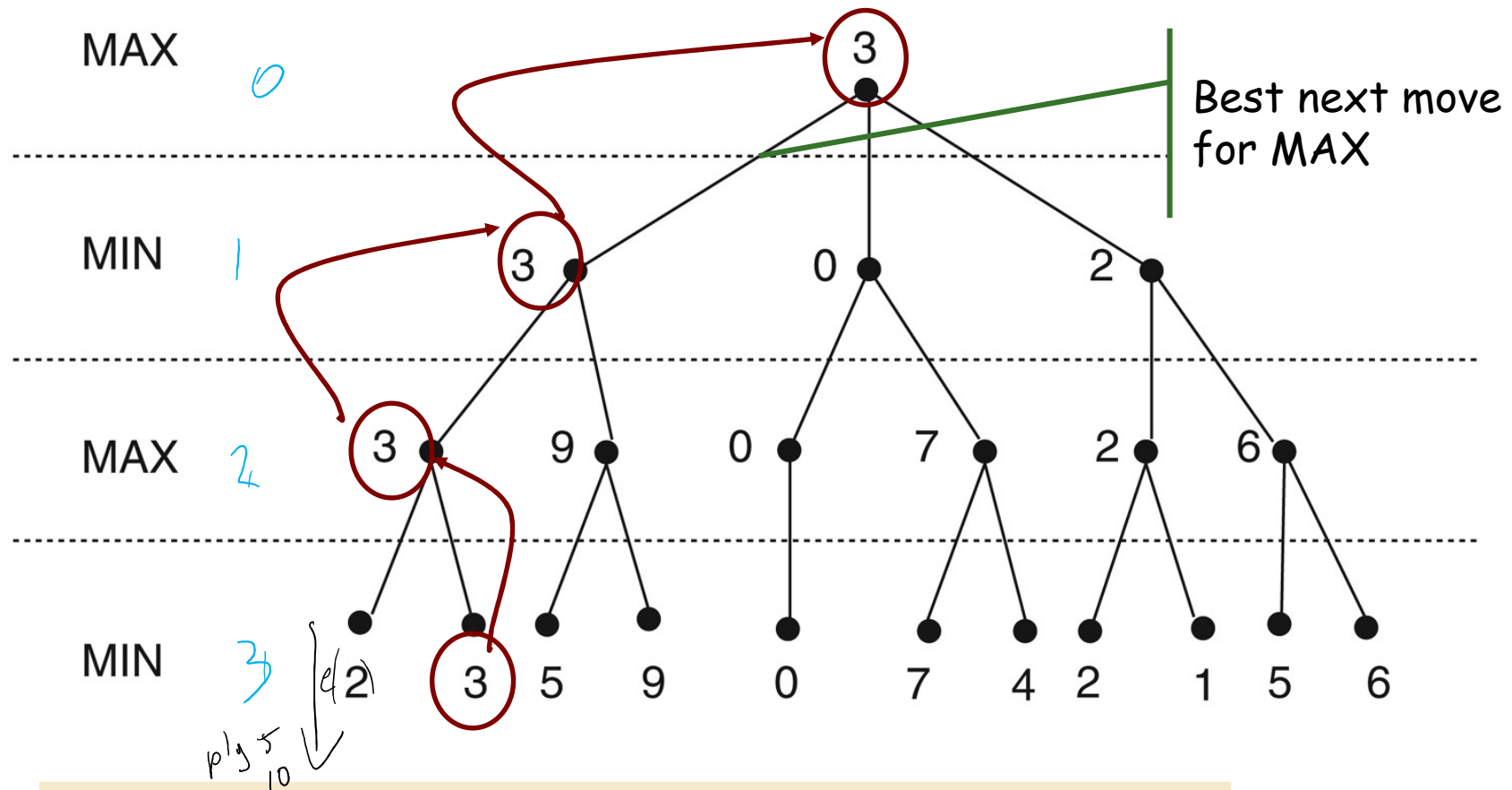


Leaf nodes show the actual heuristic value $e(n)$
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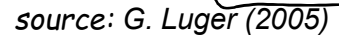
Leaf nodes show the actual heuristic value $e(n)$
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Minimax with 3-ply look-ahead



Leaf nodes show the actual heuristic value $e(n)$
Internal nodes show back-up heuristic value

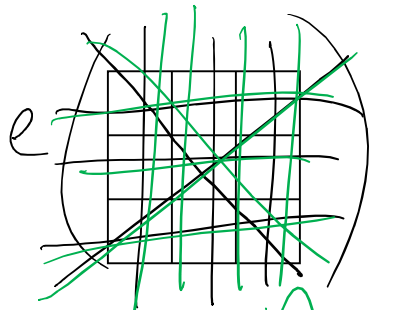
4 ply ahead



Example: $e(n)$ for Tic-Tac-Toe

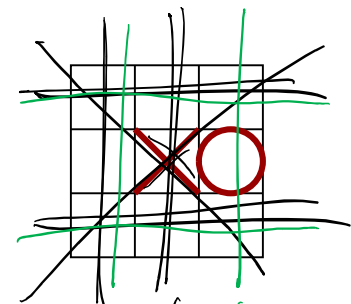
- assume MAX plays X
- possible $e(n)$

$$e(n) = \begin{cases} \text{number of rows, columns, and diagonals open for MAX} \\ - \text{number of rows, columns, and diagonals open for MIN} \\ +\infty \text{ if } n \text{ is a forced win for MAX} \\ -\infty \text{ if } n \text{ is a forced win for MIN} \end{cases}$$



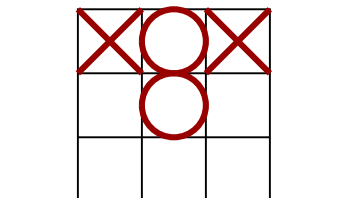
$$e(n) = 8 - 8 = 0$$

max min



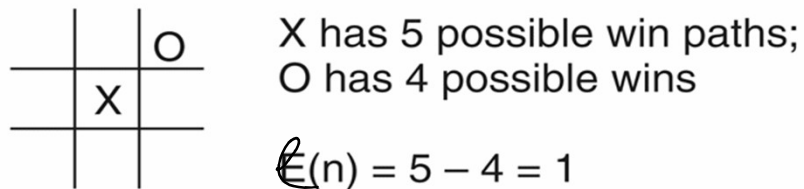
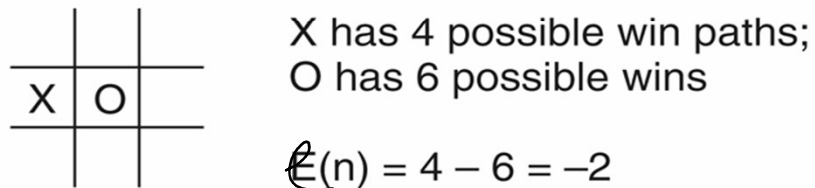
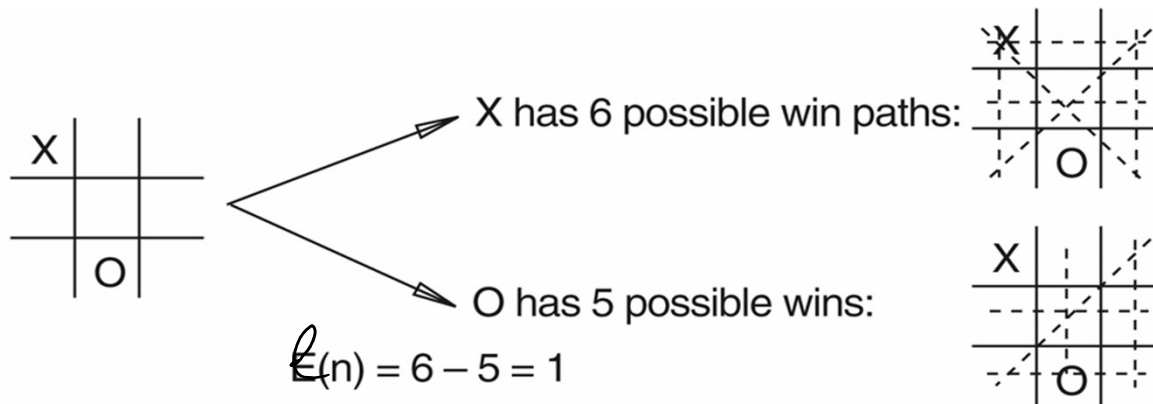
$$e(n) = 6 - 4 = 2$$

max min

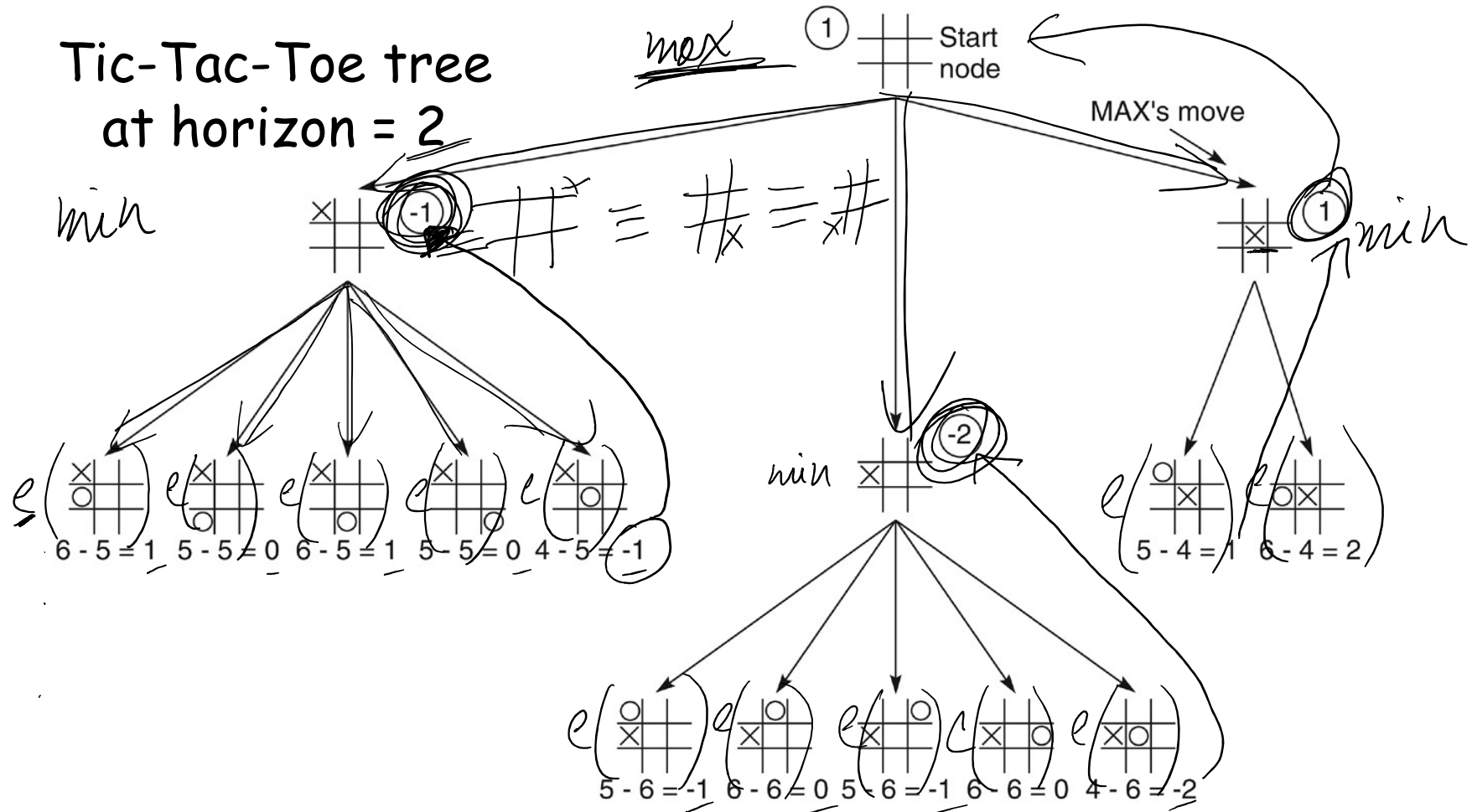


$$e(n) = 3 - 3 = 0$$

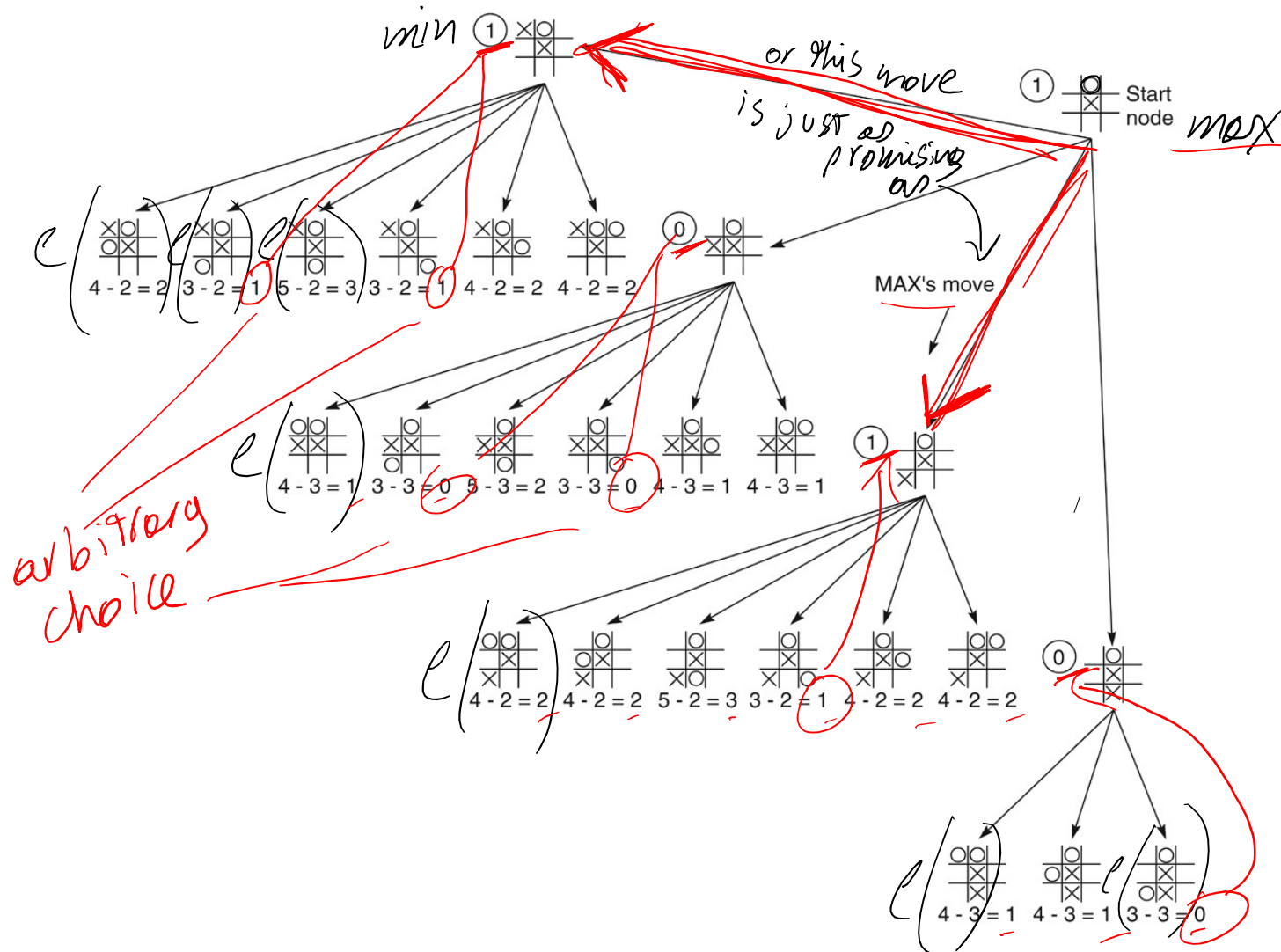
More examples...



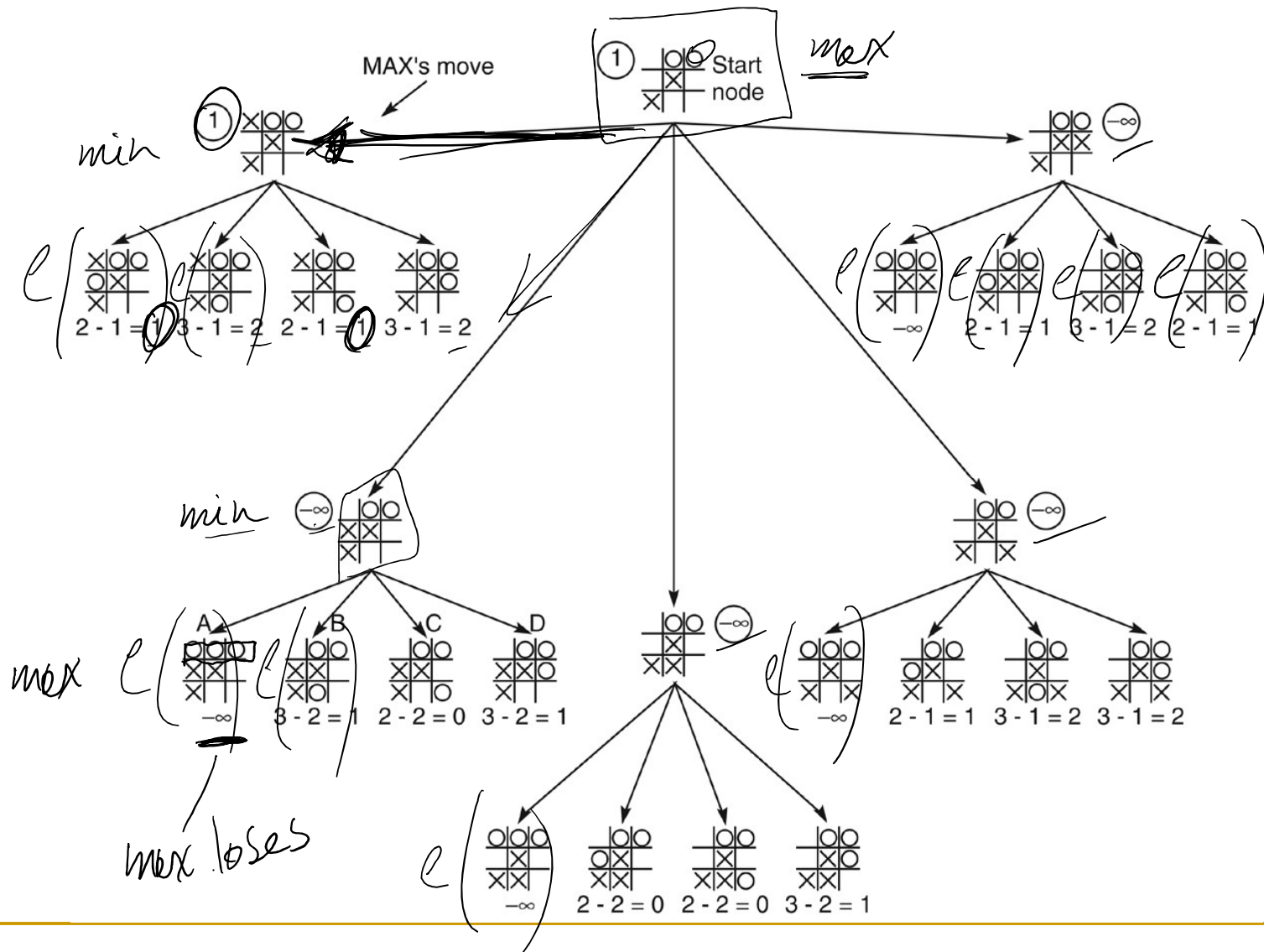
2-ply Minimax for Opening Move



2-ply Minimax: MAX's possible 2nd moves



2-ply Minimax: MAX's move at end



source: G. Luger (2005)

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Up Next

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