The Z specification language

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Introduction

- Z (pronounced "zed") is a *specification language* based on predicate logic and set theory.
- A formal specification in Z describes *what* the system does, without saying *how* it does it.
- The formal specification is implementation independent.

Example: Birthday book

- In this example we are defining the specification of a birthday book.
- We introduce the set of all names, and the set of all dates, as basic types in the specification:

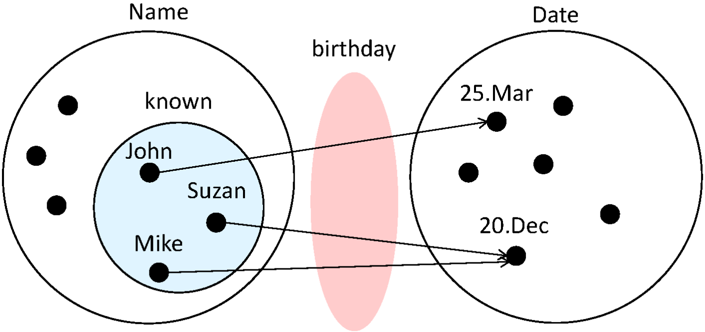
$$[Name, Date]$$

- One possible state of the system has three people in the set *known*, with their birthdays recorded by the function *birthday*:

$$known = \{John, Mike, Suzan\}$$

$$\begin{array}{l} birthday = \\ \quad \{ \\ \quad\quad John \mapsto 25.Mar, \\ \quad\quad Mike \mapsto 20.Dec, \\ \quad\quad Suzan \mapsto 20.Dec \\ \quad \} \end{array}$$

Birthday book /cont. - Function *birthday*



Birthday book: State schema

- A specification is decomposed into *schemata* (plural for *schema*).
- A Z *schema* can provide a high-level description of a model.

$$
\begin{array}{|l}
\multicolumn{1}{l}{\textit{BirthdayBook}} \\
\hline
\textit{known} : \mathbb{P}\,\textit{Name} \\
\textit{birthday} : \textit{Name} \nrightarrow \textit{Date} \\
\hline
\textit{known} = \mathrm{dom}\ \textit{birthday} \\
\hline
\end{array}
$$

- Variable *known* is the set of names with birthdates recorded and *birthday* is a function which, when applied to certain names, gives the birthdays associated with them.
- Note that *birthday* is a *partial function*: It does not relate every element of *Name* to *Date*.

Schemata /cont.

- We define a schema for both the data of the model as well as for each of the system's operations.

$$\begin{array}{|l}
\hline \textit{Name} \\
\textit{Declaration part (state)} \\
\hline
\textit{Predicate part} \\
\hline
\end{array}$$

- A schema consists of
- **Name**
- **Declaration part (state)**: specifies a set of variable assignments and their types (called the *state* of the system).
- **Predicate part**
  - Properties of the model expressed in terms of the variables defined in the declaration part.
  - These properties constitute *state invariants* (must be true at all times).

Schemata /cont.

- Alternatively a schema can be expressed in a linear notation as

$$\mathit{name} \mathrel{\hat{=}} [\mathit{declaration\ part} \mid \mathit{predicate\ part}]$$

Decorations: Before and after state

- *Decorations* give identifiers conventional meanings.
- A variable decorated with the symbol ' denotes its value after having being manipulated by an operation.
- Thus, *known* and *birthday* denote the states of their corresponding sets before an operation is invoked, whereas $known'$ and $birthday'$ denote the states of their corresponding sets after an operation has been successfully terminated.

A birthday book /cont.

- The predicate $known = \mathrm{dom}\ birthday$ states that the set *known* is the same as the domain of the function *birthday*, i.e. the set of names that the function can apply.
- This relationship is an *invariant* of the system.

$$
\begin{array}{l}
\textit{BirthdayBook} \\
\hline
known : \mathbb{P}\ Name \\
birthday : Name \nrightarrow Date \\
\hline
known = \mathrm{dom}\ birthday
\end{array}
$$

A birthday book: State schema

- In our schema we have managed to capture the information that

1. Each person can have only one birthday (because *birthday* is a function), and
2. Two (or more) people can share the same birthday (as in the example).

- The invariant $\textit{known} = \operatorname{dom} \textit{birthday}$ is satisfied, because *birthday* records a date for exactly the three names in set *known*.

A birthday book: Operation *AddBirthday*

- Once the state space is defined, we must provide a schema for each of the operations of the system.
- Operation *AddBirthday* adds a new birthday into the system:

$$
\begin{array}{|l}
\hline
\textit{AddBirthday} \\
\Delta \textit{BirthdayBook} \\
\textit{name}? : \textit{Name} \\
\textit{date}? : \textit{Date} \\
\hline
\textit{name}? \notin \textit{known} \\
\textit{known}' = \textit{known} \cup \{\textit{name}?\} \\
\underline{\textit{birthday}' = \textit{birthday} \cup \{\textit{name}? \mapsto \textit{date}?\}} \\
\hline
\end{array}
$$

- The declaration Δ*BirthdayBook* captures the fact that the schema describes a *state change*: it introduces variables *known*, *birthday*, and *known*′, *birthday*′.
- The first is an observation of the state before the change, and the last two are observations of the state after the change.

A birthday book: Operation *AddBirthday*

- Next come the declarations of the two inputs to the operation and by convention, the names of inputs end in a question mark.

$$name? : Name$$
$$date? : Date$$

A birthday book: Operation *AddBirthday*

- The first line under the dividing line defines the precondition to the operation:

$$\underline{\mathit{name}? \notin \mathit{known}}$$

- The name to be added must not already be one of those known to the system, as each person can have only one birthday.
- The specification can be extended to include error conditions in the case of a precondition failure.

A birthday book: Operation *AddBirthday*

- If the precondition is satisfied, the next two lines specify the postcondition to the operation:

$$\boxed{\begin{array}{l} known' = known \cup \{name?\} \\ birthday' = birthday \cup \{name? \mapsto date?\} \end{array}}$$

1. $name?$ is now a member of the set $known$.
2. The birthday function is extended to map the new name to the given date.

A birthday book: Operation *FindBirthday*

- Operation *FindBirthday* finds the birthday of a person known to the system:

$$
\begin{array}{|l}
\hline
\textit{FindBirthday} \\
\Xi \textit{BirthdayBook} \\
\textit{name}? : \textit{Name} \\
\textit{date}! : \textit{Date} \\
\hline
\textit{name}? \in \textit{known} \\
\textit{date}! = \textit{birthday}(\textit{name}?) \\
\hline
\end{array}
$$

- The declaration $\Xi BirthdayBook$ indicates that this is an operation in which the state of the system does not change.
- It makes the following statements redundant:

$$
\begin{aligned}
&\textit{known}' = \textit{known} \\
&\textit{birthday}' = \textit{birthday}
\end{aligned}
$$

Operation *FindBirthday*: Declaring an output

- The use of a name ending in an exclamation mark defines an output:

$$date! : Date$$

Operation *FindBirthday*: Precondition and output

- The precondition of the operation is that $\mathit{name}?$ is one of the names known to the system:

  $$\mathit{name}? \in \mathit{known}$$

- If this is so, then the output $\mathit{date}!$ is the value of the birthday function given argument $\mathit{name}?$.

  $$\mathit{date}! = \mathit{birthday}(\mathit{name}?)$$

A birthday book: Operation *Remind*

- Operation *Remind* finds which people have a birthday on a given date:

$$
\begin{array}{|l}
\hline
\textit{Remind} \\
\Xi \textit{BirthdayBook} \\
\textit{today}? : \textit{Date} \\
\textit{cards}! : \mathbb{P}\, \textit{Name} \\
\hline
\textit{cards}! = \{ n : \textit{known} \mid \textit{birthday}(n) = \textit{today}? \} \\
\hline
\end{array}
$$

- This time there is no precondition.
- The output *cards*! is defined as the set of all values drawn from the set *known* such that the value of the birthday function at *n* is *today*?.

Initializing the system

- To complete the specification, we must state the initial state of the system:

InitBirthdayBook
BirthdayBook
$known = \emptyset$

Handling errors

- What will happen if the user tries to add a birthday for someone already known to the system? Or if they try to find the birthday of someone not known?
- Our current specification makes no claim about what should happen: The system may ignore incorrect input, or it may break down (among other things).
- In addition to the basic specification, we can describe, separately the errors which might be detected and the desired responses to them.
- We can then use operations to combine the two descriptions into a stronger (robust) specification.

Handling errors /cont.

- We define Report as an enumerated set of values that an output variable $result!$ may assume:

$$Report ::= ok \mid already\_known \mid not\_known.$$

- We can now define a schema that specifies that the result should be ok.

$$\begin{array}{|l}
\multicolumn{1}{l}{\textit{Success}} \\
\hline
\Xi BirthdayBook \\
result! : Report \\
\hline
result! = ok \\
\hline
\end{array}$$

- We can now combine this description with $AddBirthday$ as

$$AddBirthday \land Success$$

Handling errors /cont.

- For each precondition error we must define a schema which describes the conditions under which the error occurs and specifies the appropriate report to be produced.
- Schema *AlreadyKnown* specifies that the report *already_known* should be produced when the input *name*? is already a member of *known*:

$$
\begin{array}{|l}
\hline
\textit{AlreadyKnown} \\
\Xi \textit{BirthdayBook} \\
\textit{name}? : \textit{Name} \\
\textit{result}! : \textit{Report} \\
\hline
\textit{name}? \in \textit{known} \\
\textit{result}! = \textit{already\_known} \\
\hline
\end{array}
$$

Strengthening the specification: *RAddBirthday*

- We can now provide a robust version of *AddBirthday*:

$$RAddBirthday \mathrel{\hat{=}} (AddBirthday \land Success) \oplus AlreadyKnown$$

- This definition introduces a new schema obtained by combining the schemas on the right-hand side.
- Alternatively, operation *RAddBirthday* could be specified directly by writing a single schema which combines the predicate parts of the three schemas *AddBirthday*, *Success* and *AlreadyKnown*.

Strengthening the specification: *RFindBirthday*

- A robust version of operation *FindBirthday* must be able to report if the input name is not known:

$$
\begin{array}{|l}
\textit{NotKnown} \\
\hline
\Xi \textit{BirthdayBook} \\
\textit{name}? : \textit{Name} \\
\textit{result}! : \textit{Report} \\
\hline
\textit{name}? \notin \textit{known} \\
\textit{result}! = \textit{not\_known} \\
\hline
\end{array}
$$

- The robust operation either behaves as described by *FindBirthday* and reports success, or reports that the name was not known:

$$\textit{RFindBirthday} \mathrel{\widehat{=}} (\textit{FindBirthday} \land \textit{Success}) \oplus \textit{NotKnown}$$

Strengthening the specification: *RRemind*

- Operation *Remind* never results in an error, so its robust version need only add the reporting of success:

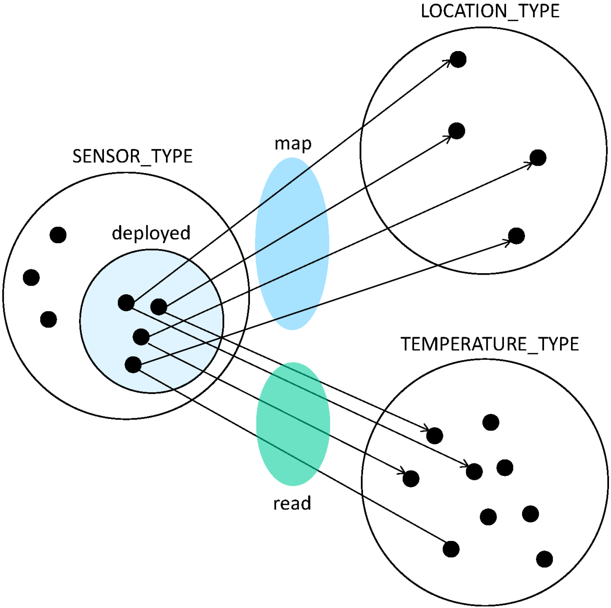
$$RRemind \mathrel{\hat{=}} Remind \land Success$$

Example: Temperature monitoring

- A system maintains a number of sensors, where each is deployed in a separate location in order to read the location's temperature.
- Before the system is deployed, all locations are marked on a map, and each location will be addressed by a sensor.
- The formal specification of the system introduces the following three types:

$$SENSOR\_TYPE,\\ LOCATION\_TYPE,\\ TEMPERATURE\_TYPE$$

Visualization of the system



State schema

$$
\begin{array}{|l}
\hline
\textit{TempMonitor} \\
\hline
\textit{deployed} : \mathbb{P}\,\textit{SENSOR\_TYPE} \\
\textit{map} : \textit{SENSOR\_TYPE} \nrightarrow \textit{LOCATION\_TYPE} \\
\textit{read} : \textit{SENSOR\_TYPE} \nrightarrow \textit{TEMPERATURE\_TYPE} \\
\hline
\textit{deployed} = \operatorname{dom} \textit{map} \\
\textit{deployed} = \operatorname{dom} \textit{read} \\
\hline
\end{array}
$$

Operation DeploySensorOK

- Operation `DeploySensorOK` places a new sensor to a unique location.
- We may assume that some (default) temperature is also passed as an argument.

$$
\begin{array}{|l}
\hline
\textit{DeploySensorOK} \\
\Delta \textit{TempMonitor} \\
\textit{sensor?} : \textit{SENSOR\_TYPE} \\
\textit{location?} : \textit{LOCATION\_TYPE} \\
\textit{temperature?} : \textit{TEMPERATURE\_TYPE} \\
\hline
\textit{sensor?} \notin \textit{deployed} \\
\textit{location?} \notin \operatorname{ran} \textit{map} \\
\textit{deployed}' = \textit{deployed} \cup \{\textit{sensor?}\} \\
\textit{map}' = \textit{map} \cup \{\textit{sensor?} \mapsto \textit{location?}\} \\
\textit{read}' = \textit{read} \cup \{\textit{sensor?} \mapsto \textit{temperature?}\} \\
\hline
\end{array}
$$

Operation ReadTemperatureOK

- Operation `ReadTemperatureOK` obtains the temperature reading from a sensor, given the sensor's location.

$$
\begin{array}{|l}
\hline
\textit{ReadTemperatureOK} \\
\Xi \textit{TempMonitor} \\
\textit{location}? : \textit{LOCATION\_TYPE} \\
\textit{temperature}! : \textit{TEMPERATURE\_TYPE} \\
\hline
\textit{location}? \in \operatorname{ran} \textit{map} \\
\textit{temperature}! = \textit{read}(\textit{map}^{-1}(\textit{location}?)) \\
\hline
\end{array}
$$

Success and error schemata

- We introduce an enumerated type *MESSAGE* which will assume values that correspond to success and error messages.

$$
\begin{array}{|l}
\hline
\textit{Success} \\
\hline
\Xi \textit{TempMonitor} \\
\textit{response}! : \textit{MESSAGE} \\
\hline
\textit{response}! = \textit{'ok'} \\
\hline
\end{array}
$$

Error schema SensorAlreadyDeployed for DeploySensorOK

- Operation `DeploySensorOK` places the following precondition: $sensor? \notin deployed$
- We provide an error schema for the case where this precondition fails:

$$
\begin{array}{|l}
\hline
\textit{SensorAlreadyDeployed} \\
\Xi \textit{TempMonitor} \\
sensor? : SENSOR\_TYPE \\
response! : MESSAGE \\
\hline
sensor? \in deployed \\
response! = \text{'Sensor deployed'} \\
\hline
\end{array}
$$

Error schema LocationAlreadyCovered for DeploySensorOK

- Operation `DeploySensorOK` places the following precondition: $location? \notin \mathrm{ran}\ map$
- We provide an error schema for the case where this precondition fails:

$$
\begin{array}{|l}
\hline
\textit{LocationAlreadyCovered} \\
\Xi \textit{TempMonitor} \\
location? : LOCATION\_TYPE \\
response! : MESSAGE \\
\hline
location? \in \mathrm{ran}\ map \\
response! = \textit{'Location already covered'} \\
\hline
\end{array}
$$

Error schema LocationUnknown for ReadTemperatureOK

- Operation `ReadTemperatureOK` places the following precondition: $\mathit{location}? \in \operatorname{ran} \mathit{map}$
- We provide an error schema for the case where this precondition fails:

$$
\begin{array}{|l}
\hline
\multicolumn{1}{l}{\mathit{LocationUnknown}} \\
\Xi \mathit{TempMonitor} \\
\mathit{location}? : \mathit{LOCATION\_TYPE} \\
\mathit{response}! : \mathit{MESSAGE} \\
\hline
\mathit{location}? \notin \operatorname{ran} \mathit{map} \\
\mathit{response}! = \mathit{'Location\ not\ covered'} \\
\hline
\end{array}
$$

Defining reliable operations

$$\begin{aligned}
&\mathit{DeploySensor} \mathrel{\hat{=}} \\
&\quad (\mathit{DeploySensorOK} \wedge \mathit{Success}) \oplus \\
&\quad (\mathit{SensorAlreadyDeployed} \vee \mathit{LocationAlreadyCovered})
\end{aligned}$$

$$\begin{aligned}
&\mathit{ReadTemperature} \mathrel{\hat{=}} \\
&\quad (\mathit{ReadTemperatureOK} \wedge \mathit{Success}) \oplus \mathit{LocationUnknown}
\end{aligned}$$

Bibliography

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Appendix A: Operations on sets

Operator	Synopsis	Meaning
$\in$	$x \in S$	set membership
$\cup$	$S_1 \cup S_2$	set union
$\cap$	$S_1 \cap S_2$	set intersection
$\setminus$	$S_1 \setminus S_2$	set difference
#	#S	cardinality of a set
$\subseteq$	$S_1 \subseteq S_2$	subset
$\subset$	$S_1 \subset S_2$	proper subset
$=$	$S_1 = S_2$	set equality
$\bigcup$	$\bigcup SS$	generalized union of sets SS
$\bigcap$	$\bigcap SS$	generalized intersection of sets SS
$\mathbb{P}$	$\mathbb{P}\, S$	power set of the set S
$\mathbb{F}$	$\mathbb{F}\, S$	finite subsets of the set S

Appendix B: Notations for functions

Symbol	Meaning
→	Total function
⇸	Partial function
↣	Total injective function
⤔	Partial injective function
↠	Total surjective function
⤀	Partial surjective function
⤗	Partial bijective function
⤖	Total bijective function
⇻	Finite partial function
⤕	Finite partial injective function

Appendix C: Operations on relations and functions

Operator	Synopsis	Meaning
$\leftrightarrow$	$X \leftrightarrow Y$	declaration of a binary relation between X and Y
$\mapsto$	$x \mapsto y$	maplet
dom	dom R	domain of the relation R
ran	ran R	range of the relation R
id	id X	identity relation
$\fatsemi$	$R_1 \fatsemi R_2$	relational composition
$\circ$	$R_1 \circ R_2$	backward relational composition
$\lhd$	$S \lhd R$	domain restriction
$\rhd$	$R \rhd S$	range restriction
$\ndres$	$S \ndres R$	domain subtraction (domain anti-restriction)
$\nrres$	$R \nrres S$	range subtraction (range anti-restriction)
$\sim$	$R^\sim$	relational inverse
$\_(\!\|\_\|\!)$	$R (\!\| S \|\!)$	relational image
$\oplus$	$R_1 \oplus R_2$	relational overriding
	R^k	relational iteration
	R^+	transitive closure of the relation R
	R^*	reflexive transitive closure of the relation R

Appendix D: Operations on sequences

Operator	Synopsis	Meaning
#	# S	length of the sequence S
$\frown$	$S_1 \frown S_2$	concatenation of sequence S_1 with S_2
rev	*rev* S	reverse of the sequence S
head	*head* S	first element of the sequence S
last	*last* S	last element of the sequence S
tail	*tail* S	sequence S with its first element removed
front	*front* S	sequence S with its last element removed
$\frown/$	$\frown/$ SS	distributed concatenation of the sequence of sequences SS
$\subseteq$	$S \subseteq T$	S is a sequence forming the prefix of the sequence T
suffix	S suffix T	S is a sequence forming the suffix of the sequence T
in	S in T	S is a segment inside the sequence T
$\upharpoonleft$	$U \upharpoonleft S$	extract the elements from the sequence S corresponding to the index set U; the result is also a sequence, maintaining the same order as in S
$\upharpoonright$	$S \upharpoonright V$	extract the elements of the set V from the sequence S; the result is also a sequence, maintaining the same order as in S
disjoint	disjoint SeqSet	SeqSet is an indexed family of mutually distinct sets
partitions	SeqSet partitions T	the indexed family of mutually disjoint sets whose distributed union is T