

# INF2080

## Oblig 1

**Deadline:** February 20, 2017

### Hand-in and deadline

Hand in a single PDF file in [Devilry](#). Deadline is **February 20, at 23:59**.

We recommend L<sup>A</sup>T<sub>E</sub>X, but all major text editors allows exporting to PDF. You can get help with L<sup>A</sup>T<sub>E</sub>X at the group sessions. You can also download the L<sup>A</sup>T<sub>E</sub>X source (`.tex`) for this assignment at the assignments page.

### Problem 1: Regular languages

Let  $A$  and  $B$  be regular languages defined by DFAs  $\mathcal{A}$  and  $\mathcal{B}$ . Let  $n_{\mathcal{A}}$  and  $n_{\mathcal{B}}$  be the number of states in  $\mathcal{A}$  and  $\mathcal{B}$ .

#### Problem 1a

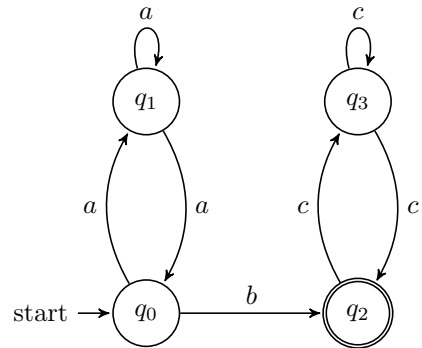
What are the highest number of states you would need in **DFAs** for the languages  $A \cap B$  and  $A^*$ ?

#### Problem 1b

What are the highest number of states you would need in **NFAs** for the languages  $A \cap B$ ,  $AB$  and  $A^*$ ?

#### Problem 1c

Create a regular expression defining the same language as the NFA



### Problem 1d

Create a DFA for the language

$$\{w \mid w \text{ contains equally many occurrences of the substrings } 01 \text{ and } 10\}.$$

### Problem 2: all-NFAs

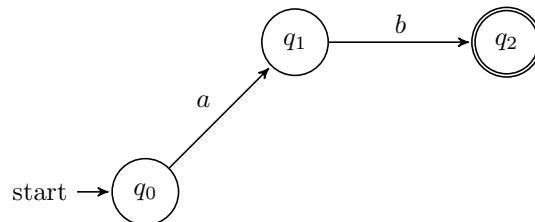
An all-NFA is defined in Sipser, problem 1.43 as a 5-tuple  $(Q, \Sigma, \delta, q_0, F)$  that accepts  $x \in \Sigma^*$  if *every* possible state that  $M$  could reach after reading input  $x$  is in  $F$  (as opposed to *at least one*).

If any branch in an all-NFA computation reaches an implicit or explicit sink state, the input is not accepted.

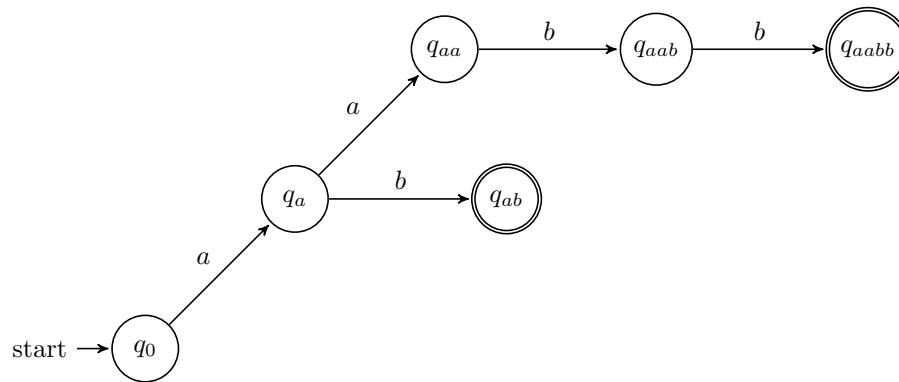
Show how any all-NFA can be converted to a DFA.

### Problem 3: Non-regular languages

The DFA



defines the language  $\{ab\}$ . The DFA



defines the language  $\{ab, aabb\}$ .

### Problem 3a

Create a DFA that defines the language  $\{ab, aabb, aaabbb\}$ .

### Problem 3b

Create a deterministic *infinite* automaton that defines the language

$$\{a^n b^n \mid n \in \mathbb{N}\}.$$

### Problem 3c

Using the pumping lemma, give a detailed proof that  $\{a^n b^n \mid n \in \mathbb{N}\}$  is not a regular language; that is, no deterministic *finite* automaton may define it.

### Problem 3d

Show that  $\{a^n b^n \mid n \in \mathbb{N}\}$  is a context free language.