

EXERCISES

CHAPTER 7

SEAN LI ¹

1. Redacted

Reference - Propositional Logic in λC

$$\begin{array}{c} \frac{A \quad \neg A}{\perp} \text{ } \perp\text{I or } \neg\text{E} \qquad \frac{\perp}{A} \perp\text{E} \qquad \frac{A \quad B}{A \wedge B} \wedge\text{I} \qquad \frac{A \wedge B}{A} \wedge\text{EL} \qquad \frac{A \wedge B}{B} \wedge\text{ER} \\[10pt] \frac{a}{a \vee b} \vee\text{IL} \qquad \frac{b}{a \vee b} \vee\text{IR} \qquad \frac{A \Rightarrow B \quad A}{B} \Rightarrow\text{E} \\[10pt] \frac{A \vee B \quad A \rightarrow C \quad B \rightarrow C}{C} \vee\text{E} \qquad \frac{a \in S \quad P(a)}{\exists a \in S, P(a)} \exists\text{I} \qquad \frac{\begin{array}{l} 1. \quad A \\ 2. \quad \left| \begin{array}{l} \dots \\ \perp \end{array} \right. \\ 3. \end{array}}{\neg A} \neg\text{I} \\[10pt] \frac{\begin{array}{l} 1. \quad A \\ 2. \quad \left| \begin{array}{l} \dots \\ B \end{array} \right. \\ 3. \end{array}}{A \Rightarrow B} \Rightarrow\text{I} \qquad \frac{\begin{array}{l} 1. \quad a \in S \\ 2. \quad \left| \begin{array}{l} \dots \\ P(a) \end{array} \right. \\ 3. \end{array}}{\forall a \in S, P(a)} \forall\text{I} \\[10pt] \frac{\exists x \in S, P(x) \quad \forall x \in S, (P(x) \Rightarrow A)}{A} \exists\text{E} \qquad \frac{a \in S \quad \forall x \in S, P(x)}{P(a)} \forall\text{E} \\[10pt] \frac{}{\neg\neg A \Rightarrow A} \text{DN (Classical)} \qquad \frac{}{A \vee \neg A} \text{ET (Classical)} \end{array}$$

Reference - 2nd Encoding for Propositional Logic

| Proposition | Minimal Propositional Logic |
|-------------------------|---|
| \perp | $\forall A, A$ |
| $A \Rightarrow B$ | $A \Rightarrow B$ |
| $\neg A$ | $A \Rightarrow \perp$ |
| $A \wedge B$ | $\forall C, (A \Rightarrow B \Rightarrow C) \Rightarrow C$ |
| $A \vee B$ | $\forall C, (A \Rightarrow C) \Rightarrow (B \Rightarrow C) \Rightarrow C$ |
| $\forall a \in S, P(a)$ | $\forall a \in S. P(a)$ |
| $\exists a \in S, P(a)$ | $\forall \alpha, (\forall a \in S, (P(a) \Rightarrow \alpha)) \Rightarrow \alpha$ |

Problem

(7.1 a) Prove in natural deduction and λC the tautology

$$B \Rightarrow (A \Rightarrow B)$$

Solution.

Natural Deduction.

1. B
2. $\left| \begin{array}{l} A \\ \left| \begin{array}{l} B \\ \hline A \Rightarrow B \end{array} \right. \end{array} \right. \Rightarrow \mathbf{I}$
3. $\left| \begin{array}{l} A \\ \left| \begin{array}{l} B \\ \hline A \Rightarrow B \end{array} \right. \end{array} \right. \Rightarrow \mathbf{I}$
4. $\left| \begin{array}{l} A \\ \left| \begin{array}{l} B \\ \hline A \Rightarrow B \end{array} \right. \end{array} \right. \Rightarrow \mathbf{I}$
5. $B \Rightarrow (A \Rightarrow B) \Rightarrow \mathbf{I}$

■

λC . Assuming context $\Gamma \equiv A : *, B : *$. By the PAT paradigm the proof is equivalent to an inhabitant of $B \rightarrow A \rightarrow B$.

1. $A : *, B : *$
2. $\left| \begin{array}{l} x : B \\ \left| \begin{array}{l} y : A \\ \left| \begin{array}{l} x : B \\ \hline \lambda y : A. x : A \rightarrow B \end{array} \right. \end{array} \right. \end{array} \right. \mathbf{Weak}$
3. $\left| \begin{array}{l} x : B \\ \left| \begin{array}{l} y : A \\ \left| \begin{array}{l} x : B \\ \hline \lambda y : A. x : A \rightarrow B \end{array} \right. \end{array} \right. \mathbf{4 Abst}$
4. $\left| \begin{array}{l} x : B \\ \left| \begin{array}{l} y : A \\ \left| \begin{array}{l} x : B \\ \hline \lambda y : A. x : A \rightarrow B \end{array} \right. \end{array} \right. \mathbf{5 Abst}$
5. $\left| \begin{array}{l} x : B \\ \left| \begin{array}{l} y : A \\ \left| \begin{array}{l} x : B \\ \hline \lambda y : A. x : A \rightarrow B \end{array} \right. \end{array} \right. \mathbf{5 Abst}$
6. $\left| \begin{array}{l} x : B \\ \left| \begin{array}{l} y : A \\ \left| \begin{array}{l} x : B \\ \hline \lambda y : A. x : A \rightarrow B \end{array} \right. \end{array} \right. \mathbf{5 Abst}$

■

Problem

(7.1 b) Prove in natural deduction and λC the tautology

$$\neg A \Rightarrow (A \Rightarrow B)$$

Solution.

Natural Deduction.

1. $\neg A$
2. $\begin{array}{|l} A \end{array}$
3. $\begin{array}{|l} \neg A \end{array}$
4. $\begin{array}{|l} A \end{array}$
5. $\begin{array}{|l} \perp \end{array} \quad \perp \text{I}$
6. $\begin{array}{|l} B \end{array} \quad \perp \text{E}$
7. $\begin{array}{|l} A \Rightarrow B \end{array} \quad \Rightarrow \text{I}$
8. $\neg A \Rightarrow (A \Rightarrow B) \quad \Rightarrow \text{I}$

■

λC . Assuming context $\Gamma \equiv A : *, B : *$. By the PAT paradigm the proof is equivalent to an inhabitant of $(A \rightarrow \perp) \rightarrow A \rightarrow B$.

1. $A : *, B : *$
2. $\begin{array}{|l} x : \neg A \end{array}$
3. $\begin{array}{|l} y : A \end{array}$
4. $\begin{array}{|l} x y : \Pi \alpha : *. \alpha \end{array} \quad \mathbf{2,3 \text{ App (Neg Elim)}}$
5. $\begin{array}{|l} x y B : B \end{array} \quad \mathbf{4,1 \text{ App (Ex Falso)}}$
6. $\begin{array}{|l} \lambda y : A. x y B : A \rightarrow B \end{array} \quad \mathbf{5 \text{ Abst}}$
7. $\begin{array}{|l} \lambda x : \neg A. \lambda y : A. x y B : \neg A \rightarrow A \rightarrow B \end{array} \quad \mathbf{6 \text{ Abst}}$

■

Problem

(7.1 c) Prove in natural deduction and λC the tautology

$$(A \Rightarrow \neg B) \Rightarrow ((A \Rightarrow B) \Rightarrow \neg A)$$

Solution.

Natural Deduction.

1. $A \Rightarrow \neg B$
2. $\begin{array}{|l} A \Rightarrow B \end{array}$
3. $\begin{array}{|l} A \end{array}$
4. $\begin{array}{|l} \neg B \end{array}$ **1,3 \Rightarrow E**
5. $\begin{array}{|l} B \end{array}$ **2,3 \Rightarrow E**
6. $\begin{array}{|l} \perp \end{array}$ **5,4 \perp I**
7. $\begin{array}{|l} \neg A \end{array}$ **3,6 \neg I**
8. $\begin{array}{|l} (A \Rightarrow B) \Rightarrow \neg A \end{array}$ **2,7 \Rightarrow I**
9. $(A \Rightarrow \neg B) \Rightarrow ((A \Rightarrow B) \Rightarrow \neg A)$ **1,8 \Rightarrow I**

■

λC . Assuming context $\Gamma \equiv A : *, B : *$. The proof should be equivalent to an inhabitant of $(A \rightarrow B \rightarrow \perp) \rightarrow (A \rightarrow B) \rightarrow A \rightarrow \perp$.

1. $A : *, B : *$
2. $\begin{array}{|l} h : A \rightarrow \neg B \end{array}$
3. $\begin{array}{|l} q : A \rightarrow B \end{array}$
4. $\begin{array}{|l} a : A \end{array}$
5. $\begin{array}{|l} q a : B \end{array}$ **3,4 App**
6. $\begin{array}{|l} h a : B \rightarrow \perp \end{array}$ **2,4 App**
7. $\begin{array}{|l} h a (q a) : \perp \end{array}$ **6,5 App (Neg Elim)**
8. $\begin{array}{|l} \lambda a : A . h a (q a) : \neg A \end{array}$ **7 Abst (Neg Intro)**
9. $\begin{array}{|l} \lambda q : A \rightarrow B . \lambda a : A . h a (q a) : A \rightarrow B \rightarrow \neg A \end{array}$ **8 Abst**
10. $\begin{array}{|l} \lambda h : A \rightarrow \neg B . \lambda q : A \rightarrow B . \lambda a : A . h a (q a) : (A \rightarrow \neg B) \rightarrow A \rightarrow B \rightarrow \neg A \end{array}$ **9 Abst**

■

Problem

(7.1 d) Prove in natural deduction and λC the tautology

$$\neg(A \Rightarrow B) \Rightarrow \neg B$$

Solution.

Natural Deduction.

$$\begin{array}{ll}
1. & \neg(A \Rightarrow B) \\
2. & \begin{array}{l} | B \\ | \\ 3. & | A \\ | & | \\ 4. & | \quad | B \\ | & | \\ 5. & | A \Rightarrow B \quad \quad 3,4 \Rightarrow I \\ | & | \\ 6. & | \perp \quad \quad 5,1 \perp I \\ | & | \\ 7. & | \neg B \quad \quad 6 \neg I \end{array} \\
8. & \neg(A \Rightarrow B) \Rightarrow \neg B \quad 1,7 \Rightarrow I
\end{array}$$

■

λC . Assuming context $\Gamma \equiv A : *, B : *$. The proof should be equivalent to an inhabitant of $((A \rightarrow B) \rightarrow \perp) \rightarrow B \rightarrow \perp$.

$$\begin{array}{ll}
1. & n : \neg(A \rightarrow B) \\
2. & \begin{array}{l} | b : B \\ | \\ 3. & | a : A \\ | & | \\ 4. & | \quad | b : B \quad \quad \text{Weak} \\ | & | \\ 5. & | \lambda a : A . b : A \rightarrow B \quad \quad 4 \text{ Abst} \\ | & | \\ 6. & | n (\lambda a : A . b) : \perp \quad \quad 1,5 \text{ App (Neg Elim)} \\ | & | \\ 7. & | \lambda b : B . n (\lambda a : A . b) : \neg B \quad \quad 6 \text{ Abst (Neg Intro)} \end{array} \\
8. &
\end{array}$$

$$\begin{array}{l}
\lambda n : \neg(A \rightarrow B) . \lambda b : B . n (\lambda a : A . b) \\
: \neg(A \rightarrow B) \rightarrow \neg B \quad \quad 7 \text{ Abst}
\end{array}$$

■

Problem

(7.2) Formulate the double negation law as an axiom in λC , and prove the following tautology in λC with DN.

$$(\neg A \Rightarrow A) \Rightarrow A$$

Solution. The rule

$$\frac{}{\neg\neg A \Rightarrow A} \text{DN-E}$$

Could be translated into lambda calculus as

$$\Pi A : * . ((A \rightarrow \perp) \rightarrow \perp) \rightarrow A$$

Proof. Assume context $\Gamma \equiv A : *$.

| | | |
|----|---|--------------------------------|
| 1. | $A : *$ | |
| 2. | $h : \neg A \rightarrow A$ | |
| 3. | $x : \neg A$ | |
| 4. | $h x : A$ | 2,3 App |
| 5. | $x (h x) : \perp$ | 3,4 App (Contradiction) |
| 6. | $\lambda x : \neg A . x (h x) : \neg \neg A$ | 5 Abst (Neg Intro) |
| 7. | $\text{DN } A : \neg \neg A \rightarrow A$ | 1,1 App |
| 8. | $\text{DN } A (\lambda x : \neg A . x (h x)) : A$ | App (Axiom DN) |
| 9. | $\lambda h : \neg A \rightarrow A . \text{DN } A (\lambda x : \neg A . x (h x)) : (\neg A \rightarrow A) \rightarrow A$ | 8 Abst |

■

Problem

(7.3 a) Prove the following tautology in classical logic using λC

$$(A \Rightarrow B) \Rightarrow (\neg B \Rightarrow \neg A)$$

Solution. Assume context $\Gamma \equiv A : *, B : *$. It suffices the proof to find an inhabitant of type

$$(A \rightarrow B) \rightarrow (\neg B \rightarrow \neg A)$$

Proof.

| | | |
|-----|--|--------------------------------|
| 1. | $A : *, B : *$ | |
| 2. | $h : A \rightarrow B$ | |
| 3. | $b : \neg B$ | |
| 4. | $a : A$ | |
| 5. | $h a : B$ | 5,2 App |
| 6. | $b (h a) : \perp$ | 6,3 App (Contradiction) |
| 7. | $b (h a)(\neg A) : \neg A$ | 7,5 App (Ex Falso) |
| 8. | $\lambda a : A . b (h a)(\neg A) : A \rightarrow \neg A$ | 7 Abst |
| 9. | $a : \neg A$ | |
| 10. | $a : \neg A$ | Var |
| 11. | $\lambda a : \neg A . a : (\neg A \rightarrow \neg A)$ | 10 Abst |

| | | |
|-----|--|-----------------------|
| 12. | ET $A : A \vee \neg A$ | App (Axiom ET) |
| | ET $A (\neg A) : (A \rightarrow \neg A) \rightarrow$ | |
| 13. | $(\neg A \rightarrow \neg A) \rightarrow \neg A$ | 12 App |
| | ET $A (\neg A)(\lambda a : A . b (h a)(\neg A)) :$ | |
| 14. | $(\neg A \rightarrow \neg A) \rightarrow \neg A$ | 13,8 App |
| | ET $A (\neg A)$ | |
| | $(\lambda a : A . b (h a)(\neg A))$ | |
| 15. | $(\lambda a : \neg A . a) : \neg A$ | 14,11 App |
| | $\lambda b : \neg B . \text{ET } A (\neg A)$ | |
| | $(\lambda a : A . b (h a)(\neg A))$ | |
| 16. | $(\lambda a : \neg A . a) : \neg B \rightarrow \neg A$ | 15 Abst |
| | $\lambda h : A \rightarrow B . \lambda b : \neg B . \text{ET } A (\neg A)$ | |
| | $(\lambda a : A . b (h a)(\neg A))$ | |
| | $(\lambda a : \neg A . a) :$ | |
| 17. | $(A \rightarrow B) \rightarrow \neg B \rightarrow \neg A$ | 16 Abst |

■

Problem

(7.3 b) Prove the following tautology in classical logic using λC

$$(\neg B \Rightarrow \neg A) \Rightarrow (A \Rightarrow B)$$

Solution. Assume context $\Gamma \equiv A : *, B : *$. It suffices the proof to find an inhabitant of type

$$(\neg B \rightarrow \neg A) \rightarrow A \rightarrow B$$

Proof.

| | | |
|----|---------------------------------------|-------------|
| 1. | $A : *, B : *$ | |
| 2. | $h : \neg B \rightarrow \neg A$ | |
| 3. | $a : A$ | |
| 4. | $b : B$ | |
| 5. | $b : B$ | Weak |
| 6. | $\lambda b : B . b : B \rightarrow B$ | |
| 7. | $b : \neg B$ | |

| | | | | | |
|-----|--|--|--|--|--------------------|
| 8. | | | | $h\ b : \neg A$ | 2,7 App |
| 9. | | | | $h\ b\ a : \perp$ | 8,2 App (Neg Elim) |
| 10. | | | | $h\ b\ a\ B : B$ | 9 App (Ex Falso) |
| 11. | | | | $\lambda b : \neg B . h\ b\ a\ B : \neg B \rightarrow B$ | 10 Abst |
| 12. | | | | ET $B : B \vee \neg B$ | 1 App (Axiom ET) |
| 13. | | | | ET $B\ B : (B \rightarrow B) \rightarrow (\neg B \rightarrow B) \rightarrow B$ | 12,1 App |
| 14. | | | | ET $B\ B\ (\lambda b : B . b) : (\neg B \rightarrow B) \rightarrow B$ | 13,6 App |
| 15. | | | | ET $B\ B\ (\lambda b : B . b)(\lambda b : \neg B . h\ b\ a\ B) : B$ | 14,11 App |
| | | | | $\lambda a : A . \text{ET } B\ B\ (\lambda b : B . b)$ | |
| 16. | | | | $(\lambda b : \neg B . h\ b\ a\ B) : A \rightarrow B$ | 15 Abst |
| | | | | $\lambda h : \neg B \rightarrow \neg A . \lambda a : A .$ | |
| | | | | ET $B\ B\ (\lambda b : B . b)(\lambda b : \neg B . h\ b\ a\ B)$ | |
| 17. | | | | $: (\neg B \rightarrow \neg A) \rightarrow A \rightarrow B$ | 16 Abst |

■

Problem

(7.4) Derive \wedge -EL and \wedge -ER in λC .

Solution. The derivation is the same as proving

$$M : \Pi C . (A \rightarrow B \rightarrow C) \rightarrow C \vdash N_1 : A$$

$$M : \Pi C . (A \rightarrow B \rightarrow C) \rightarrow C \vdash N_2 : B$$

Left Projection.

| | | |
|----|---|--|
| 1. | $M : \Pi C . (A \rightarrow B \rightarrow C) \rightarrow C$ | |
| 2. | | $M\ A : (A \rightarrow B \rightarrow A) \rightarrow A$ |
| 3. | | $x : A$ |
| 4. | | |
| 5. | | |
| 6. | | |
| 7. | | |
| 8. | | |

Weak

5 Abst

6 Abst

2,7 App

■

Right Projection.

| | | |
|----|---|----------------|
| 1. | $M : \Pi C . (A \rightarrow B \rightarrow C) \rightarrow C$ | |
| 2. | $M B : (A \rightarrow B \rightarrow B) \rightarrow B$ | |
| 3. | $x : A$ | |
| 4. | $b : B$ | |
| 5. | $b : B$ | Weak |
| 6. | $\lambda b : B . b : B \rightarrow B$ | 5 Abst |
| 7. | $\lambda x : A . \lambda b : B . b : A \rightarrow B \rightarrow B$ | 6 Abst |
| 8. | $N_2 \equiv M B (\lambda x : A . \lambda b : B . b) : B$ | 2,7 App |

■

Problem

(7.5 a) Prove tautology under classical logic

$$\neg(A \Rightarrow B) \Rightarrow A$$

Solution. Assume context $\Gamma \equiv A : *, B : *$. Therefore the proof suffices to find an inhabitant of

$$((A \rightarrow B) \rightarrow \perp) \rightarrow A$$

Proof.

| | | |
|-----|--|--------------------------------|
| 1. | $A : *, B : *$ | |
| 2. | $h : \neg(A \rightarrow B)$ | |
| 3. | $a : A$ | |
| 4. | $a : A$ | Weak |
| 5. | $\lambda a : A . a : A \rightarrow A$ | 4 Abst |
| 6. | $a : \neg A$ | |
| 7. | $n : A$ | |
| 8. | $a n : \perp$ | 6,7 App (Contradiction) |
| 9. | $a n B : B$ | 8 App (Ex Falso) |
| 10. | $\lambda n : A . a n B : A \rightarrow B$ | 9 Abst |
| 11. | $h (\lambda n : A . a n B) : \perp$ | 10,2 App (Contra.) |
| 12. | $h (\lambda n : A . a n B) A : A$ | 11 App (Ex Falso) |
| | $\lambda a : \neg A . h (\lambda n : A . a n B) A$ | |
| 13. | $: \neg A \rightarrow A$ | 12 Abst |
| 14. | $ET : \Pi S : * . S \vee \neg S$ | Axiom ET |

| | | |
|-----|---|------------------|
| 15. | ET $A : A \vee \neg A$ | 14,1 App |
| 16. | ET $A A : (A \rightarrow A) \rightarrow (\neg A \rightarrow A) \rightarrow A$ | 15,1 App |
| | ET $A A (\lambda a : A . a)$ | |
| 17. | $: (\neg A \rightarrow A) \rightarrow A$ | 16,5 App |
| | ET $A A (\lambda a : A . a)$ | |
| 18. | $(\lambda a : \neg A . h (\lambda n : A . a \ n B) A) : A$ | 17,13 App |
| | $\lambda h : \neg(A \rightarrow B) . \text{ET } A A (\lambda a : A . a)$ | |
| | $(\lambda a : \neg A . h (\lambda n : A . a \ n B) A)$ | |
| 19. | $: \neg(A \rightarrow B) \rightarrow A$ | 18 Abst |

■

Problem

(7.5 b) Prove tautology under classical logic

$$\neg(A \Rightarrow B) \Rightarrow (A \wedge \neg B)$$

Solution. Assume context $\Gamma \equiv A : *, B : *$. Therefore the proof suffices to find an inhabitant of

$$((A \rightarrow B) \rightarrow \perp) \rightarrow \Pi C : * . ((A \rightarrow (B \rightarrow \perp) \rightarrow C) \rightarrow C)$$

Proof.

| | | |
|-----|---|----------------------------|
| 1. | $A : *, B : *$ | |
| 2. | $p : \neg(A \rightarrow B)$ | |
| 3. | $C : *$ | |
| 4. | $h : A \rightarrow \neg B \rightarrow C$ | |
| 5. | $hb : B$ | |
| 6. | $ha : A$ | |
| 7. | $hb : B$ | Weak |
| 8. | $\lambda ha : A . hb : A \rightarrow B$ | 7 Abst |
| 9. | $p (\lambda ha : A . hb) : \perp$ | 2,8 App (Contra.) |
| 10. | $\lambda hb : B . p (\lambda ha : A . hb) : \neg B$ | 9 Abst (Neg Intro) |
| 11. | $hna : \neg A$ | |
| 12. | $ha' : A$ | |
| 13. | $hna ha' : \perp$ | 11,12 App (Contra.) |

| | | |
|-----|---|----------------------------|
| 14. | $\boxed{hna\ ha' B : B}$ | 13,1 App (Ex Falso) |
| 15. | $\lambda ha' : A . hna\ ha' B : A \rightarrow B$ | 14 Abst |
| 16. | $p(\lambda ha' : A . hna\ ha' B) : \perp$ | 2,15 App (Contra.) |
| 17. | $\boxed{p(\lambda ha' : A . hna\ ha' B) A : A}$ | 2,15 App (Contra.) |
| | $\lambda hna : \neg A . p$ | |
| | $(\lambda ha' : A . hna\ ha' B) A$ | |
| 18. | $: \neg A \rightarrow A$ | 17 Abst |
| 19. | $ha'' : A$ | |
| 20. | $\boxed{ha'' : A}$ | Var |
| 21. | $\lambda ha'' : A . ha'' : A \rightarrow A$ | 20 Abst |
| 22. | $ET : \Pi S : * . S \vee \neg S$ | Axiom ET |
| 23. | $ET\ A : A \vee \neg A$ | 22,1 App |
| 24. | $ET\ A\ A : (A \rightarrow A) \rightarrow (\neg A \rightarrow A) \rightarrow A$ | 23,1 App |
| | $ET\ A\ A$ | |
| | $(\lambda ha'' : A . ha'')$ | |
| 25. | $: (\neg A \rightarrow A) \rightarrow A$ | 24,21 App |
| | $ET\ A\ A$ | |
| | $(\lambda ha'' : A . ha'')$ | |
| | $(\lambda hna : \neg A . p$ | |
| 26. | $(\lambda ha' : A . hna\ ha' B) A) : A$ | 25,18 App |
| | $h\ (ET\ A\ A$ | |
| | $(\lambda ha'' : A . ha'')$ | |
| | $(\lambda hna : \neg A . p$ | |
| | $(\lambda ha' : A . hna\ ha' B) A))$ | |
| 27. | $: \neg B \rightarrow C$ | 4,26 App |
| | $h\ (ET\ A\ A$ | |
| | $(\lambda ha'' : A . ha'')$ | |
| | $(\lambda hna : \neg A . p$ | |
| | $(\lambda ha' : A . hna\ ha' B) A))$ | |
| 28. | $(\lambda hb : B . p(\lambda ha : A . hb)) : C$ | 27,10 App |

| | | |
|-----|---|----------------|
| 29. | $ \begin{array}{l} \lambda h : A \rightarrow \neg B \rightarrow C . h \text{ (ET } A \ A \\ (\lambda ha'' : A . ha'') \\ (\lambda hna : \neg A . p (\lambda ha' : A . hna \ ha' \ B) \ A)) \\ (\lambda hbb : B . p (\lambda ha : A . hb)) \\ : (A \rightarrow \neg B \rightarrow C) \rightarrow C \end{array} $ | 28 Abst |
| 30. | $ \begin{array}{l} \lambda C : * . \lambda h : A \rightarrow \neg B \rightarrow C . \\ h \text{ (ET } A \ A \\ (\lambda ha'' : A . ha'') \\ (\lambda hna : \neg A . p (\lambda ha' : A . hna \ ha' \ B) \ A)) \\ (\lambda hbb : B . p (\lambda ha : A . hb)) : A \wedge \neg B \end{array} $ | 29 Abst |
| 31. | $ \begin{array}{l} \lambda p : \neg(A \rightarrow B). \\ \lambda C : * . \lambda h : A \rightarrow \neg B \rightarrow C . \\ h \text{ (ET } A \ A \\ (\lambda ha'' : A . ha'') \\ (\lambda hna : \neg A . p (\lambda ha' : A . hna \ ha' \ B) \ A)) \\ (\lambda hbb : B . p (\lambda ha : A . hb)) \\ : \neg(A \rightarrow B) \rightarrow (A \wedge \neg B) \end{array} $ | 30 Abst |

■

Problem

(7.6 a) Verify that the following expressions is a tautology in constructive logic

$$\neg A \Rightarrow \neg(A \wedge B)$$

Solution. Suppose context $A : *, B : *$. The proof suffices to give a inhabitant of type

$$\neg A \rightarrow \neg(A \wedge B) \equiv (A \rightarrow \perp) \rightarrow (\Pi C : * . (A \rightarrow B \rightarrow C) \rightarrow C) \rightarrow \perp$$

Proof.

| | | |
|----|---|----------------|
| 1. | $A : *, B : *$ | |
| 2. | $na : \neg A$ | |
| 3. | $h : \Pi C : * . (A \rightarrow B \rightarrow C) \rightarrow C$ | |
| 4. | $h \ A : (A \rightarrow B \rightarrow A) \rightarrow A$ | 3,1 App |
| 5. | $ha : A$ | |
| 6. | $hb : B$ | |

| | | | | | |
|-----|--|--|--|--|---------------------------------|
| 7. | | | | $ha : A$ | |
| 8. | | | | $\lambda hb : B . ha : B \rightarrow A$ | 7 Abst |
| 9. | | | | $\lambda ha : A . \lambda hb : B . ha : A \rightarrow B \rightarrow A$ | 8 Abst |
| 10. | | | | $h A (\lambda ha : A . \lambda hb : B . ha) : A$ | 4,9 App |
| 11. | | | | $na (h A (\lambda ha : A . \lambda hb : B . ha)) : \perp$ | 2,10 App (Contradiction) |
| | | | | $\lambda h : A \wedge B .$ | |
| | | | | $na (h A (\lambda ha : A . \lambda hb : B . ha))$ | |
| 12. | | | | $: A \wedge B \rightarrow \perp$ | 11 Abst |
| | | | | $\lambda na : \neg A . \lambda h : A \wedge B .$ | |
| | | | | $na (h A (\lambda ha : A . \lambda hb : B . ha))$ | |
| 13. | | | | $: \neg A \rightarrow \neg(A \wedge B)$ | 12 Abst |

■

Problem

(7.6 b) Verify that the following expressions is a tautology in constructive logic

$$\neg(A \wedge \neg A)$$

Solution. Suppose context $A : *, B : *$. The proof suffices to give a inhabitant of type

$$(\Pi S : *. (A \rightarrow (A \rightarrow \perp) \rightarrow S) \rightarrow S) \rightarrow \perp$$

Proof.

| | | |
|-----|---|----------------|
| 1. | $A : *, B : *, \perp : \square$ | |
| 2. | $h : \Pi S : *. (A \rightarrow \neg A \rightarrow S) \rightarrow S$ | |
| 3. | $h \perp : (A \rightarrow \neg A \rightarrow \perp) \rightarrow \perp$ | 2,1 App |
| 4. | $a : A$ | |
| 5. | $n : \neg A$ | |
| 6. | $na : \perp$ | 5,4 App |
| 7. | $\lambda n : \neg A . na : \neg A \rightarrow \perp$ | 6 Abst |
| 8. | $\lambda a : A . \lambda n : \neg A . na : A \rightarrow \neg A \rightarrow \perp$ | 7 Abst |
| 9. | $h \perp (\lambda a : A . \lambda n : \neg A . na) : \perp$ | 3,8 App |
| 10. | $\lambda h : A \wedge \neg A . h \perp (\lambda a : A . \lambda n : \neg A . na) : A \wedge \neg A \rightarrow \perp$ | 9 Abst |

■

Problem

(7.7) Derive \forall -**IL** and \forall -**IR**.

Solution. The derivation is the same as proving

$$m : A, B : * \vdash N : \Pi C . (A \rightarrow C) \rightarrow (B \rightarrow C) \rightarrow C$$

$$m : B, A : * \vdash N : \Pi C . (A \rightarrow C) \rightarrow (B \rightarrow C) \rightarrow C$$

Left Introduction.

| | | |
|----|--|----------------|
| 1. | $m : A, B : *$ | |
| 2. | $C : *$ | |
| 3. | $hl : A \rightarrow C$ | |
| 4. | $hr : B \rightarrow C$ | |
| 5. | $hl\ m : C$ | 3,1 App |
| 6. | $\lambda hr : B \rightarrow C . hl\ m : (B \rightarrow C) \rightarrow C$ | 5 Abst |
| 7. | $\lambda hl : A \rightarrow C . \lambda hr : B \rightarrow C . hl\ m$ $: (A \rightarrow C) \rightarrow (B \rightarrow C) \rightarrow C$ | 6 Abst |
| 8. | $\lambda C : * . \lambda hl : A \rightarrow C . \lambda hr : B \rightarrow C . hl\ m$ $: \Pi C : * . (A \rightarrow C) \rightarrow (B \rightarrow C) \rightarrow C$ | 7 Abst |

■

Right Introduction.

| | | |
|----|--|----------------|
| 1. | $m : B, A : *$ | |
| 2. | $C : *$ | |
| 3. | $hl : A \rightarrow C$ | |
| 4. | $hr : B \rightarrow C$ | |
| 5. | $hr\ m : C$ | 4,1 App |
| 6. | $\lambda hr : B \rightarrow C . hr\ m : (B \rightarrow C) \rightarrow C$ | 5 Abst |
| 7. | $\lambda hl : A \rightarrow C . \lambda hr : B \rightarrow C . hr\ m$ $: (A \rightarrow C) \rightarrow (B \rightarrow C) \rightarrow C$ | 6 Abst |
| 8. | $\lambda C : * . \lambda hl : A \rightarrow C . \lambda hr : B \rightarrow C . hr\ m$ $: \Pi C : * . (A \rightarrow C) \rightarrow (B \rightarrow C) \rightarrow C$ | 7 Abst |

■

Problem

(7.8 a) Give λC derivations verifying the intuitionistic tautology below

$$(A \vee B) \Rightarrow (B \vee A)$$

Solution.

Proof.

| | | |
|-----|---|---------|
| 1. | $A : *, B : *$ | |
| 2. | $h : A \vee B$ | |
| 3. | $C : *$ | |
| 4. | $hb : B \rightarrow C$ | |
| 5. | $ha : A \rightarrow C$ | |
| 6. | $h C : (A \rightarrow C) \rightarrow (B \rightarrow C) \rightarrow C$ | 2,3 App |
| 7. | $h C ha : (B \rightarrow C) \rightarrow C$ | 6,5 App |
| 8. | $h C ha hb : C$ | 7,4 App |
| 9. | $\lambda ha : A \rightarrow C . h C ha hb : (A \rightarrow C) \rightarrow C$ | 8 Abst |
| 10. | $\lambda hb : B \rightarrow C . \lambda ha : A \rightarrow C . h C ha hb : (B \rightarrow C) \rightarrow (A \rightarrow C) \rightarrow C$ | 9 Abst |
| 11. | $\lambda C : * . \lambda hb : B \rightarrow C . \lambda ha : A \rightarrow C . h C ha hb : \Pi C : * . (B \rightarrow C) \rightarrow (A \rightarrow C) \rightarrow C$ | 10 Abst |
| 12. | $\lambda h : A \vee B . \lambda C : * . \lambda hb : B \rightarrow C . \lambda ha : A \rightarrow C . h C ha hb : A \vee B \rightarrow B \vee A$ | 11 Abst |

■

Problem

(7.8 b) Give λC derivations verifying the intuitionistic tautology below

$$\neg(A \vee B) \Rightarrow (\neg A \wedge \neg B)$$

Solution.

Proof.

| | |
|----|---|
| 1. | $A : *, B : *, \perp : \square$ |
| 2. | $h : \neg(\Pi S : * . (A \rightarrow S) \rightarrow (B \rightarrow S) \rightarrow S)$ |

| | | |
|-----|---|-------------------|
| 3. | $C : *$ | |
| 4. | $p : (\neg A \rightarrow \neg B \rightarrow C)$ | |
| 5. | $a : A$ | |
| 6. | $C : *$ | |
| 7. | $ha : A \rightarrow C$ | |
| 8. | $hb : B \rightarrow C$ | |
| 9. | $ha a : C$ | 7,5 App |
| 10. | $\lambda hb : B \rightarrow C . ha a : (B \rightarrow C) \rightarrow C$ | 9 Abst |
| 11. | $\lambda ha : A \rightarrow C . \lambda hb : B \rightarrow C . ha a$ $: (A \rightarrow C) \rightarrow (B \rightarrow C) \rightarrow C$ | 10 Abst |
| 12. | $\lambda C : * . \lambda ha : A \rightarrow C . \lambda hb : B \rightarrow C . ha a$ $: \Pi C : * . (A \rightarrow C) \rightarrow (B \rightarrow C) \rightarrow C$ | 11 Abst |
| 13. | $h (\lambda C : * . \lambda ha : A \rightarrow C . \lambda hb : B \rightarrow C . ha a) : \perp$ | 2,12 App (Contr.) |
| 14. | $\lambda a : A . h$ $(\lambda C : * . \lambda ha : A \rightarrow C . \lambda hb : B \rightarrow C . ha a) : \neg A$ | 13 Abst (Neg I) |
| 15. | $b : B$ | |
| 16. | $C : *$ | |
| 17. | $ha : A \rightarrow C$ | |
| 18. | $hb : B \rightarrow C$ | |
| 19. | $hb b : C$ | 18,15 App |
| 20. | $\lambda hb : B \rightarrow C . hb b : (B \rightarrow C) \rightarrow C$ | 19 Abst |
| 21. | $\lambda ha : A \rightarrow C . \lambda hb : B \rightarrow C . hb b$ $: (A \rightarrow C) \rightarrow (B \rightarrow C) \rightarrow C$ | 20 Abst |
| 22. | $\lambda C : * . \lambda ha : A \rightarrow C . \lambda hb : B \rightarrow C . hb b$ $: \Pi C : * . (A \rightarrow C) \rightarrow (B \rightarrow C) \rightarrow C$ | 21 Abst |
| 23. | $h (\lambda C : * . \lambda ha : A \rightarrow C . \lambda hb : B \rightarrow C . hb b) : \perp$ | 2,22 App (Contr.) |
| 24. | $\lambda b : B . h$ $(\lambda C : * . \lambda ha : A \rightarrow C . \lambda hb : B \rightarrow C . hb b) : \neg B$ | 23 Abst (Neg I) |
| 25. | $p (\lambda a : A . h$ $(\lambda C : * . \lambda ha : A \rightarrow C . \lambda hb : B \rightarrow C . ha a))$ $: \neg B \rightarrow C$ | 4,14 App |
| 26. | $p (\lambda a : A . h$ $(\lambda C : * . \lambda ha : A \rightarrow C . \lambda hb : B \rightarrow C . ha a))$ $(\lambda b : B . h$ $(\lambda C : * . \lambda ha : A \rightarrow C . \lambda hb : B \rightarrow C . hb b))) : C$ | 25,24 App |

| | | |
|-----|--|---------|
| 27. | $ \begin{array}{l} \lambda p : \neg A \rightarrow \neg B \rightarrow C . p \\ (\lambda a : A . h \\ (\lambda C : * . \lambda ha : A \rightarrow C . \lambda hb : B \rightarrow C . ha a)) \\ (\lambda b : B . h \\ (\lambda C : * . \lambda ha : A \rightarrow C . \lambda hb : B \rightarrow C . hb b)) \\ : (\neg A \rightarrow \neg B \rightarrow C) \rightarrow C \end{array} $ | 26 Abst |
| 28. | $ \begin{array}{l} \lambda C : * . \lambda p : \neg A \rightarrow \neg B \rightarrow C . p \\ (\lambda a : A . h \\ (\lambda C : * . \lambda ha : A \rightarrow C . \lambda hb : B \rightarrow C . ha a)) \\ (\lambda b : B . h \\ (\lambda C : * . \lambda ha : A \rightarrow C . \lambda hb : B \rightarrow C . hb b)) \\ : \Pi C : * . (\neg A \rightarrow \neg B \rightarrow C) \rightarrow C \end{array} $ | 27 Abst |
| 29. | $ \begin{array}{l} \lambda h : \neg(A \vee B) . \lambda C : * . \lambda p : \neg A \rightarrow \neg B \rightarrow C . p \\ (\lambda a : A . h \\ (\lambda C : * . \lambda ha : A \rightarrow C . \lambda hb : B \rightarrow C . ha a)) \\ (\lambda b : B . h \\ (\lambda C : * . \lambda ha : A \rightarrow C . \lambda hb : B \rightarrow C . hb b)) \\ : \neg(A \vee B) \rightarrow (\neg A \wedge \neg B) \end{array} $ | 28 Abst |

■

Problem

(7.8 c) Give λC derivations verifying the intuitionistic tautology below

$$(\neg A \wedge \neg B) \Rightarrow \neg(A \vee B)$$

Solution.

Proof.

1. $A : *, B : *, \perp : \square$
2. $h : \Pi C : * . (\neg A \rightarrow \neg B \rightarrow C) \rightarrow C$
3. $p : \Pi S : * . (A \rightarrow S) \rightarrow (B \rightarrow S) \rightarrow S$
4. $h \perp : \neg(\neg A \rightarrow \neg B \rightarrow \perp)$ **2,1 App**
5. $p \perp : (A \rightarrow \perp) \rightarrow (B \rightarrow \perp) \rightarrow \perp$ **3,1 App**
6. $h \perp(p \perp) : \perp$ **4,5 App (Cont.)**

| | | |
|----|---|---------------|
| 7. | $\frac{\lambda p : A \vee B . h \perp (p \perp) : \neg(A \vee B)}{\lambda h : \neg A \wedge \neg B . \lambda p : A \vee B . h \perp (p \perp)}$ | 6 Abst |
| 8. | $: \neg A \wedge \neg B \rightarrow \neg(A \vee B)$ | 7 Abst |

■

Problem

(7.9 a) Verify that the following expression is a tautology in constructive logic by giving a proof in first-order natural deduction and a flag styled derivation.

$$\forall x \in S, (\neg P(x) \Rightarrow (P(x) \Rightarrow Q(x) \wedge R(x)))$$

Solution.

Natural Deduction.

| | | |
|----|--|------------------------------------|
| 1. | $x \in S$ | |
| 2. | $\neg P(x)$ | |
| 3. | $P(x)$ | |
| 4. | \perp | 3,2 \perpI |
| 5. | $Q(x) \wedge R(x)$ | 4 \perpE |
| 6. | $P(x) \Rightarrow Q(x) \wedge R(x)$ | 5 \RightarrowI |
| 7. | $\neg P(x) \Rightarrow (P(x) \Rightarrow Q(x) \wedge R(x))$ | 6 \RightarrowI |
| 8. | $\forall x \in S, (\neg P(x) \Rightarrow (P(x) \Rightarrow Q(x) \wedge R(x)))$ | 7 \forallI |

■

λC .

| | | |
|----|--|--------------------------------|
| 1. | $S : *, P : S \rightarrow *, Q : S \rightarrow *, R : S \rightarrow *$ | |
| 2. | $x : S$ | |
| 3. | $nhp : \neg(P x)$ | |
| 4. | $hp : P x$ | |
| 5. | $nhp \ hp : \perp$ | 3,4 App (Contradiction) |
| 6. | $Q x \wedge R x : *$ | Form |
| 7. | $nhp \ hp \ (Q x \wedge R x) : Q x \wedge R x$ | 5,6 App (Ex Falso) |
| 8. | $\lambda hp : P x . nhp \ hp \ (Q x \wedge R x)$ | |
| | $: P x \rightarrow Q x \wedge R x$ | 7 Abst |

| | | |
|-----|---|---------------|
| 9. | $\frac{\lambda nhp : \neg P x . \lambda hp : P x . nhp hp (Q x \wedge R x) : \neg P x \rightarrow P x \rightarrow Q x \wedge R x}{\lambda x : S . \lambda nhp : \neg P x . \lambda hp : P x . nhp hp (Q x \wedge R x)}$ | 8 Abst |
| 10. | $\frac{\lambda x : S . \lambda nhp : \neg P x . \lambda hp : P x . nhp hp (Q x \wedge R x) : \Pi x : S . \neg P x \rightarrow P x \rightarrow Q x \wedge R x}{\lambda x : S . \lambda nhp : \neg P x . \lambda hp : P x . nhp hp (Q x \wedge R x)}$ | 9 Abst |

■

Problem

(7.9 b) Verify that the following expression is a tautology in constructive logic by giving a proof in first-order natural deduction and a flag styled derivation.

$$\forall x \in S, P(x) \Rightarrow \forall y \in S, P(y) \vee Q(y)$$

Solution.

Natural Deduction.

| | | |
|----|---|-------------------------------------|
| 1. | $\forall x \in S, P(x)$ | |
| 2. | $y : S$ | |
| 3. | $P(y)$ | 2,1 $\forall E$ |
| 4. | $P(y) \vee Q(y)$ | 3 $\vee I$ |
| 5. | $\forall y : S, (P(y) \vee Q(y))$ | 4 $\forall I$ |
| 6. | $\forall x \in S, P(x) \Rightarrow \forall y : S, (P(y) \vee Q(y))$ | 5 $\Rightarrow I$ |

■

λC .

| | | |
|----|---|----------------|
| 1. | $S : *, P : S \rightarrow *, Q : S \rightarrow *$ | |
| 2. | $h : \Pi x : S . P x$ | |
| 3. | $y : S$ | |
| 4. | $h y : P y$ | 2,3 App |
| 5. | $C : *$ | |
| 6. | $hp : P y \rightarrow C$ | |
| 7. | $hq : Q y \rightarrow C$ | |
| 8. | $hp (h y) : C$ | 6,4 App |
| 9. | $\lambda hq : Q y \rightarrow C . hp (h y) : (Q y \rightarrow C) \rightarrow C$ | 8 Abst |

| | | |
|-----|---|----------------|
| 10. | $\lambda hp : P\ y \rightarrow C . \lambda hq : Q\ y \rightarrow C . hp\ (h\ y)$ $: (P\ y \rightarrow C) \rightarrow (Q\ y \rightarrow C) \rightarrow C$ | 9 Abst |
| 11. | $\lambda C : * . \lambda hp : P\ y \rightarrow C . \lambda hq : Q\ y \rightarrow C . hp\ (h\ y)$ $: \Pi C : * . (P\ y \rightarrow C) \rightarrow (Q\ y \rightarrow C) \rightarrow C$ | 10 Abst |
| 12. | $\lambda y : S . \lambda C : * .$ $\lambda hp : P\ y \rightarrow C . \lambda hq : Q\ y \rightarrow C . hp\ (h\ y)$ $: \Pi y : S . P\ y \vee Q\ y$ | 11 Abst |
| 13. | $\lambda h : \Pi x : S . P\ x .$ $\lambda y : S . \lambda C : * .$ $\lambda hp : P\ y \rightarrow C . \lambda hq : Q\ y \rightarrow C . hp\ (h\ y)$ $: (\Pi x : S . P\ x) \rightarrow \Pi y : S . P\ y \vee Q\ y$ | 12 Abst |

■