

# EXERCISES

## CHAPTER 2

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1. Reducted

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**Definition** Some rules for reference.

$$\frac{x : \sigma \in \Gamma}{\Gamma \vdash x : \sigma} \text{ (T-Var)} \quad \frac{\Gamma \vdash M : \sigma \rightarrow \tau \quad \Gamma \vdash N : \sigma}{\Gamma \vdash M N : \tau} \text{ (T-App)}$$
$$\frac{\Gamma, x : \sigma \vdash M : \tau}{\Gamma \vdash \lambda x : \sigma. M : \sigma \rightarrow \tau} \text{ (T-Abst)}$$

In this document, convention is that all type judgements in a proof tree, unless stated otherwise, is derived from a single and unique context per tree.

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### Problem

(2.1) Type the following terms

$x x y \quad x y y \quad x y x \quad x (x y) \quad x (y x)$

*Solution.* The first term cannot be typed.

*Proof.*  $x x y = (x x) y$ . Therefore,  $x$  is a function type, denote it as  $\tau \rightarrow \sigma$ . By the application rule, a subterm applied to  $x$  must be of  $\tau$ , which means that the application  $x x$  is not legally typed. ■

The second one is typable where  $x : \tau \rightarrow \tau \rightarrow \sigma$  and  $y : \tau$ .

1.  $x : \tau \rightarrow \tau \rightarrow \sigma \quad \dashv \Gamma$
2.  $y : \tau \quad \dashv \Gamma$

3.  $x y : \tau \rightarrow \sigma$       **1,2 T-App**
4.  $x y y : \sigma$       **3,2 T-App**

The third term is not typable.

*Proof.* Assume  $x y x = (x y) x$  is typable. Therefore,  $x : \tau$  where  $\tau \equiv \sigma \rightarrow \tau \rightarrow \alpha$  and  $y : \sigma$ . One can construct an infinite chain of function type by substituting  $\tau$ :  $\tau \equiv \sigma \rightarrow (\sigma \rightarrow (\sigma \rightarrow \dots \rightarrow \alpha) \rightarrow \alpha) \rightarrow \alpha$ . By induction, it can be proven that only lambda abstractions can construct function types, meaning that the term is of form

$$(\lambda n : \tau. \lambda m : \tau. \dots. (\lambda a : \sigma. \lambda b : \sigma. \dots)) y (\lambda n : \tau. \lambda m : \tau. \dots. (\lambda a : \sigma. \lambda b : \sigma. \dots))$$

meaning that an infinite reduction path is needed. This is impossible in STLC. ■

The fourth type is typable where  $x : (\tau \rightarrow \tau)$  and  $y : \tau$ .

1.  $x : \tau \rightarrow \tau$        $\dashv \Gamma$
2.  $y : \tau$        $\dashv \Gamma$
3.  $x y : \tau$       **1,2 T-App**
4.  $x (x y) : \tau$       **1,3 T-App**

The fifth term is typable where  $x : (\tau \rightarrow \sigma)$  and  $y : (\tau \rightarrow \sigma) \rightarrow \tau$ :

1.  $x : \tau \rightarrow \sigma$        $\dashv \Gamma$
2.  $y : (\tau \rightarrow \sigma) \rightarrow \tau$        $\dashv \Gamma$
3.  $y x : \tau$       **2,1 T-App**
4.  $x (y x) : \sigma$       **1,3 T-App**

### Problem

(2.2) Find types for zero, one, and two

*Solution.* Term for zero is

$$\text{zero} := \lambda f x . x$$

Here  $x$  is only used as a

$$\text{zero} := \lambda f : \alpha. \lambda x : \beta. x$$

Type derivation shown as below:

1.  $f : \alpha$                               **Bound**
2.      $| x : \beta$                               **Bound**

3.	$x : \beta$	<b>T-Var</b>
4.	$\lambda x . x : \beta \rightarrow \beta$	<b>3 T-Abst</b>
5.	$\lambda f : \alpha . x : \beta . x : \alpha \rightarrow \beta \rightarrow \beta$	<b>4 T-Abst</b>

Term for one is

$$\text{one} := \lambda f x . f x$$

Let  $f$  be an arbitrary function type that consumes  $x$

$$\text{one} := \lambda f : \alpha \rightarrow \beta . x : \alpha . f x$$

Type derivation shown as below

1.	$f : \alpha \rightarrow \beta$	<b>Bound</b>
2.	$x : \alpha$	<b>Bound</b>
3.	$f : \alpha \rightarrow \beta$	<b>T-Var</b>
4.	$x : \alpha$	<b>T-Var</b>
5.	$f x : \beta$	<b>3,4 T-App</b>
6.	$\lambda x . f x : \alpha \rightarrow \beta$	<b>5 T-Abst</b>
7.	$\lambda f : \alpha \rightarrow \beta . x : \alpha . f x : (\alpha \rightarrow \beta) \rightarrow \alpha \rightarrow \beta$	<b>6 T-Abst</b>

Same type signatures can be given to two

$$\text{two} := \lambda f : \alpha \rightarrow \beta . \lambda x : \alpha . f f x$$

Type derivation shown as below

1.	$f : \alpha \rightarrow \beta$	<b>Bound</b>
2.	$x : \alpha$	<b>Bound</b>
3.	$f : \alpha \rightarrow \beta$	<b>T-Var</b>
4.	$x : \alpha$	<b>T-Var</b>
5.	$f x : \beta$	<b>3,4 T-App</b>
6.	$f f x : \beta$	<b>3,5 T-App</b>
7.	$\lambda x . f f x : \alpha \rightarrow \beta$	<b>6 T-Abst</b>
8.	$\lambda f : \alpha \rightarrow \beta . \lambda x : \alpha . f f x : (\alpha \rightarrow \beta) \rightarrow \alpha \rightarrow \beta$	<b>7 T-Abst</b>

### Problem

(2.3) Find types for

$$K := \lambda x y . x$$

$$S := \lambda x y z . x z (y z)$$

*Solution.* There are no occurrences of application in  $K$ 's subterms. Therefore all its binding variables could be given a simple base type.

$$K := \lambda x : \alpha. \lambda y : \beta. x$$

Type derivation shown as below

1.	$x : \alpha$	<b>Bound</b>
2.	$y : \beta$	<b>Bound</b>
3.	$x : \alpha$	<b>T-Var</b>
4.	$\lambda y : \beta. x : \beta \rightarrow \alpha$	<b>3 T-Abst</b>
5.	$\lambda x : \alpha. \lambda y : \beta. x : \alpha \rightarrow \beta \rightarrow \alpha$	<b>4 T-Abst</b>

For the  $S$  combinator, no term was applied to  $z$ . Therefore it can be given a simple base type  $\alpha$ . As  $z$  was applied to  $y$ , it implies that  $y : \alpha \rightarrow \beta$  for some output type  $\beta$ . As  $x$  takes  $z$  and  $(y z)$ , it must be of type  $\alpha \rightarrow \beta \rightarrow \delta$ .

$$S := \lambda x : \alpha \rightarrow \beta \rightarrow \delta. \lambda y : \alpha \rightarrow \beta. \lambda z. \alpha. x z (y z)$$

Complete type derivation shown as below:

1.	$x : \alpha \rightarrow \beta \rightarrow \delta$	<b>Bound</b>
2.	$y : \alpha \rightarrow \beta$	<b>Bound</b>
3.	$z : \alpha$	<b>Bound</b>
4.	$y : \alpha \rightarrow \beta$	<b>T-Var</b>
5.	$z : \alpha$	<b>T-Var</b>
6.	$y z : \beta$	<b>4,5 T-App</b>
7.	$x : \alpha \rightarrow \beta \rightarrow \delta$	<b>T-Var</b>
8.	$x z : \beta \rightarrow \delta$	<b>7,5 T-App</b>
9.	$x z (y z) : \delta$	<b>8,6 T-App</b>
10.	$\lambda z : \alpha. x z (y z) : \alpha \rightarrow \delta$	<b>9 T-Abstr</b>
11.	$\lambda y : \alpha \rightarrow \beta. \lambda z. \alpha. x z (y z) : (\alpha \rightarrow \beta) \rightarrow \alpha \rightarrow \delta$	<b>10 T-Abstr</b>

12.

$$\lambda x : \alpha \rightarrow \beta \rightarrow \delta. \lambda y : \alpha \rightarrow \beta. \lambda z : \alpha. x z (y z) : (\alpha \rightarrow \beta \rightarrow \delta) \rightarrow (\alpha \rightarrow \beta) \rightarrow \alpha \rightarrow \delta \quad \mathbf{11 \ T-Abst}$$

### Problem

(2.4) Type the bound variables

$$\lambda x y z . x (y z)$$

$$\lambda x y z . y (x z) z$$

*Solution.* For the first term,  $z$  had nothing applied to it. Therefore it could be given a simple base type  $\alpha$ .  $z$  was applied to  $y$ , therefore  $y : \alpha \rightarrow \beta$  to satisfy the application rule. Because the application yielded a type of  $\beta$ , by the application rule  $x : \beta \rightarrow \delta$  for some type  $\delta$ .

$$\lambda x : \beta \rightarrow \delta. \lambda y : \alpha \rightarrow \beta. \lambda z : \alpha. x (y z)$$

Complete type derivation shown below

	<b>Bound</b>
1. $x : \beta \rightarrow \delta$	<b>Bound</b>
2. $y : \alpha \rightarrow \beta$	<b>Bound</b>
3. $z : \alpha$	<b>Bound</b>
4. $y : \alpha \rightarrow \beta$	<b>T-Var</b>
5. $z : \alpha$	<b>T-Var</b>
6. $y z : \beta$	<b>4,5 T-App</b>
7. $x : \beta \rightarrow \delta$	<b>T-Var</b>
8. $x (y z) : \delta$	<b>7,6 T-App</b>
9. $\lambda z : \alpha. x (y z) : \alpha \rightarrow \delta$	<b>8 T-Abst</b>
10. $\lambda y : \alpha \rightarrow \beta. \lambda z : \alpha. x (y z) : (\alpha \rightarrow \beta) \rightarrow \alpha \rightarrow \delta$	<b>9 T-Abst</b>
11.	

$$\lambda x : \beta \rightarrow \delta. \lambda y : \alpha \rightarrow \beta. \lambda z : \alpha. x (y z) : (\beta \rightarrow \delta) \rightarrow (\alpha \rightarrow \beta) \rightarrow \alpha \rightarrow \delta \quad \mathbf{10 \ T-Abst}$$

In the second term  $z$  could still be given a simple base type  $z : \alpha$ . Therefore  $x : \alpha \rightarrow \beta$  for some type  $\beta$ .  $y$  takes  $x z : \beta$  and  $z : \alpha$ , therefore it is of type  $y : \beta \rightarrow \alpha \rightarrow \delta$  for some  $\delta$ .

$$\lambda x : \alpha \rightarrow \beta. \lambda y : \beta \rightarrow \alpha \rightarrow \delta. \lambda z : \alpha. y (x z) z$$

. Complete type derivation shown below

1.	$x : \alpha \rightarrow \beta$	<b>Bound</b>
2.	$y : \beta \rightarrow \alpha \rightarrow \delta$	<b>Bound</b>
3.	$z : \alpha$	<b>Bound</b>
4.	$x : \alpha \rightarrow \beta$	<b>T-Var</b>
5.	$z : \alpha$	<b>T-Var</b>
6.	$x z : \beta$	<b>4,5 T-App</b>
7.	$y : \beta \rightarrow \alpha \rightarrow \delta$	<b>T-Var</b>
8.	$y(x z) : \alpha \rightarrow \delta$	<b>7,6 T-App</b>
9.	$y(x z) z : \delta$	<b>8,5 T-App</b>
10.	$\lambda z : \alpha. y(x z) z : \alpha \rightarrow \delta$	<b>9 T-Abst</b>
11.	$\lambda y : \beta \rightarrow \alpha \rightarrow \delta. \lambda z. y(x z) z : (\beta \rightarrow \alpha \rightarrow \delta) \rightarrow \alpha \rightarrow \delta$	<b>10 T-Abst</b>
12.	$\lambda x : \alpha \rightarrow \beta. \lambda y : \beta \rightarrow \alpha \rightarrow \delta. \lambda z. y(x z) z :$ $(\alpha \rightarrow \beta) \rightarrow (\beta \rightarrow \alpha \rightarrow \delta) \rightarrow \alpha \rightarrow \delta$	<b>11 T-Abst</b>

### Problem

(2.5) Try to type the following terms, and prove if not typable.

$$\lambda x y . x (\lambda z . y) y$$

$$\lambda x y . x (\lambda z . x) y .$$

*Solution.* The first term is trivially typable.

1.	$x : (\delta \rightarrow \alpha) \rightarrow \alpha \rightarrow \beta$	<b>Bound</b>
2.	$y : \alpha$	<b>Bound</b>
3.	$x : (\delta \rightarrow \alpha) \rightarrow \alpha \rightarrow \beta$	<b>T-Var</b>
4.	$z : \delta$	<b>Bound</b>
5.	$y : \alpha$	<b>T-Var</b>
6.	$\lambda z : \delta. y : \delta \rightarrow \alpha$	<b>5 T-Abst</b>
7.	$x(\lambda z : \delta. y) : \alpha \rightarrow \beta$	<b>3,6 T-App</b>
8.	$y : \alpha$	<b>T-Var</b>
9.	$x(\lambda z : \delta. y) y : \beta$	<b>7,8 T-App</b>
10.	$\lambda y : \alpha. x(\lambda z : \delta. y) y : \alpha \rightarrow \beta$	<b>9 T-Abst</b>

11.

$$\begin{aligned} \lambda x : ((\delta \rightarrow \alpha) \rightarrow \alpha \rightarrow \beta) \lambda y : \alpha. x (\lambda z : \delta. y) y \\ : ((\delta \rightarrow \alpha) \rightarrow \alpha \rightarrow \beta) \rightarrow \alpha \rightarrow \beta \end{aligned} \quad \textbf{10 T-Abst}$$

The second term is not typable in STLC.

*Proof.* By induction on the type inference rule that constructed the type judgement for subterm  $x (\lambda z . x)$ . Because the term is an application, the only rule that applies is the application rule.

We denote the context inside the abstraction as  $\Gamma'$ . Suppose  $\mathcal{J} \equiv \Gamma' \vdash x (\lambda z . x) : \tau$ . By the inference rule of application,  $x$  must be a function type that accepts the type of  $(\lambda z . x)$ . Let  $\Gamma' \vdash z : \alpha$ , and type of  $x$  as  $\tau_x$ . Therefore,  $\Gamma' \vdash \lambda z : \alpha. x : \alpha \rightarrow \tau_x$ . Therefore,  $\tau_x \equiv (\alpha \rightarrow \tau_x) \rightarrow \tau$ . This is a recursive type, which is not constructable as it requires infinitely nested lambda abstractions that requires infinite reduction paths to reach a normal form. ■

### Problem

(2.6) Prove the pretyped term below is legal.

$$\lambda x : ((\alpha \rightarrow \beta) \rightarrow \alpha). x (\lambda z : \alpha. y)$$

Using the tree format and the flag format.

*Solution.* We suppose a context  $\Gamma \vdash y : \beta$  that obviously exists.

*Proof.*

$$\frac{\frac{\frac{x : (\alpha \rightarrow \beta) \rightarrow \alpha}{\Gamma, z : \alpha \vdash y : \beta} \text{(Bound)} \quad \frac{\Gamma, z : \alpha \vdash y : \beta}{\Gamma \vdash (\lambda z : \alpha. y) : \alpha \rightarrow \beta} \text{(T-Abst)}}{\Gamma, x : (\alpha \rightarrow \beta) \rightarrow \alpha \vdash (x (\lambda z : \alpha. y)) : \alpha} \text{(T-App)}}{\Gamma \vdash \lambda x : ((\alpha \rightarrow \beta) \rightarrow \alpha). x (\lambda z : \alpha. y) : ((\alpha \rightarrow \beta) \rightarrow \alpha) \rightarrow \alpha} \text{(T-Abst)}$$

A valid type could be given to the term. Therefore, the term is typable under an existing context. ■

The flag derivation is given below:

- |  |                 |
|--|-----------------|
| 1. $x : (\alpha \rightarrow \beta) \rightarrow \alpha$   | <b>Bound</b>    |
| 2. $\left  \begin{array}{l} z : \alpha \\ \hline y : \beta \end{array} \right.$  | <b>Bound</b>    |
| 3. $\left  \begin{array}{l} z : \alpha \\ \hline y : \beta \end{array} \right.$  | $\neg \Gamma$   |
| 4. $\left  \begin{array}{l} z : \alpha \\ \hline (\lambda z : \alpha. y) : \alpha \rightarrow \beta \end{array} \right.$ | <b>3 T-Abst</b> |
| 5. $\left  \begin{array}{l} z : \alpha \\ \hline x : (\alpha \rightarrow \beta) \rightarrow \alpha \end{array} \right.$  | <b>T-Var</b>    |

$$6. \quad \boxed{x (\lambda z : \alpha. y) : \beta} \quad \text{5,4 T-App}$$

7.

$$\begin{aligned} \lambda x : ((\alpha \rightarrow \beta) \rightarrow \beta). x (\lambda z : \alpha. y) \\ : (\alpha \rightarrow \beta) \rightarrow \alpha \rightarrow \alpha \end{aligned} \quad \text{6 T-Abst}$$

### Problem

(2.7 a) Derive

$$f : A \rightarrow B \wedge g : B \rightarrow C \Rightarrow g \circ f : A \rightarrow C$$

Using the rules

$$\frac{f : A \rightarrow B, x \in A}{f(x) \in B} \text{ (F-App)} \quad \frac{\forall x \in A, f(x) \in B}{f : A \rightarrow B} \text{ (F-Abst)}$$

*Solution.*

*Proof.*

1.	$f : A \rightarrow B \wedge g : B \rightarrow C$	<b>Assumption</b>
2.	$f : A \rightarrow B$	$1 \wedge E$
3.	$g : B \rightarrow C$	$1 \wedge E$
4.	$a \in A$	
5.	$f(a) \in B$	<b>3, 4 F-App</b>
6.	$g(f(a)) \in C$	<b>5, 4 F-App</b>
7.	$\boxed{(g \circ f)(a) \in C}$	<b>6 Compose Def</b>
8.	$\forall x \in A, (g \circ f)(x) \in C$	$7 \forall E$
9.	$\boxed{g \circ f : A \rightarrow C}$	<b>8 F-Abst</b>
10.	$f : A \rightarrow B, g : B \rightarrow C \Rightarrow g \circ f : A \rightarrow C$	$9 \Rightarrow I$

■

### Problem

(2.7 b) Give a derivation in natural deduction of the following:

$$(A \Rightarrow B) \Rightarrow ((B \Rightarrow C) \Rightarrow (A \Rightarrow C))$$

Using the rules

$$\frac{\frac{A \Rightarrow B \quad A (\Rightarrow E)}{B} \quad \frac{1. \quad A \quad \text{Premise}}{2. \quad \left| \dots \right.} \quad \frac{3. \quad \left| \begin{array}{c} B \\ \hline A \Rightarrow B \end{array} \right.}{\hline A \Rightarrow B} (\Rightarrow I)}{(\Rightarrow I)}$$

*Solution.*

*Proof.*

$$\begin{array}{lll} 1. & A \Rightarrow B & \text{Premise} \\ 2. & \left| \begin{array}{c} B \Rightarrow C \\ \left| \begin{array}{c} A \\ \left| \begin{array}{c} B \\ \left| \begin{array}{c} C \\ \hline A \Rightarrow C \end{array} \right. \end{array} \right. \end{array} \right. & \text{Premise} \\ 3. & \left| \begin{array}{c} A \\ \left| \begin{array}{c} B \\ \left| \begin{array}{c} C \\ \hline A \Rightarrow C \end{array} \right. \end{array} \right. \end{array} \right. & \text{Premise} \\ 4. & \left| \begin{array}{c} B \\ \left| \begin{array}{c} C \\ \hline A \Rightarrow C \end{array} \right. \end{array} \right. & 1, 3 \Rightarrow E \\ 5. & \left| \begin{array}{c} C \\ \hline A \Rightarrow C \end{array} \right. & 2, 4 \Rightarrow E \\ 6. & \left| \begin{array}{c} A \Rightarrow C \\ \hline (B \Rightarrow C) \Rightarrow (A \Rightarrow C) \end{array} \right. & 3-5 \Rightarrow I \\ 7. & (B \Rightarrow C) \Rightarrow (A \Rightarrow C) & 2-6 \Rightarrow I \\ 8. & (A \Rightarrow B) \Rightarrow ((B \Rightarrow C) \Rightarrow (A \Rightarrow C)) & 1-7 \Rightarrow I \end{array}$$

■

### Problem

(2.7 c) Prove the following pre-typed term is legal using flag notation

$$\lambda z : \alpha. y (x z)$$

*Solution.*

*Proof.* Let  $\Gamma \vdash x : \alpha \rightarrow \beta, y : \beta \rightarrow \delta$  for some type  $\beta$  and  $\delta$ .

$$\begin{array}{lll} 1. & z : \alpha & \text{Bound} \\ 2. & \left| \begin{array}{c} x : \alpha \rightarrow \beta \\ z : \alpha \end{array} \right. & \neg \Gamma \\ 3. & \left| \begin{array}{c} x : \alpha \rightarrow \beta \\ z : \alpha \end{array} \right. & \text{T-Var} \end{array}$$

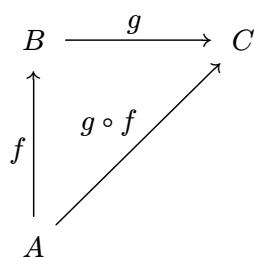
4.	$x z : \beta$	<b>2,3 T-App</b>
5.	$y : \beta \rightarrow \delta$	$\dashv \Gamma$
6.	$y (x z) : \delta$	<b>5,4 T-App</b>
7.	$\lambda z : \alpha. y (x z) : \alpha \rightarrow \delta$	<b>6 T-Abst</b>

■

### Problem

(2.7 d) State the similarity between Q. 2.7 (a), (b), and (c).

*Solution.* All of these examples requires proving something about composing two maps together as like this:



### Problem

(2.8 a) Pre-type the bounding variables for the following term

$$\lambda x y . y (\lambda z . y x) : (\gamma \rightarrow \beta) \rightarrow ((\gamma \rightarrow \beta) \rightarrow \beta) \rightarrow \beta$$

*Solution.*

$$\lambda x : (\gamma \rightarrow \beta). y : ((\gamma \rightarrow \beta) \rightarrow \beta). y (\lambda z : \gamma. y x)$$

### Problem

(2.8 b) Give a derivation in tree format

*Solution.*

$$\begin{array}{c}
\text{(i)} \frac{}{x : (\gamma \rightarrow \beta)} \quad \text{(ii)} \frac{}{y : (\gamma \rightarrow \beta) \rightarrow \beta} \\
\text{(iii)} \frac{x : (\gamma \rightarrow \beta)}{x : (\gamma \rightarrow \beta), y : ((\gamma \rightarrow \beta) \rightarrow \beta), z : \gamma \vdash y x : \beta} \text{T-App} \\
\text{(iv)} \frac{x : (\gamma \rightarrow \beta), y : ((\gamma \rightarrow \beta) \rightarrow \beta)}{x : (\gamma \rightarrow \beta), y : ((\gamma \rightarrow \beta) \rightarrow \beta) \vdash \lambda z : \gamma. y x : \gamma \rightarrow \beta} \text{T-Abst} \\
\text{(v)} \frac{}{y : ((\gamma \rightarrow \beta) \rightarrow \beta)} \text{T-App} \\
\text{(vi)} \frac{}{x : (\gamma \rightarrow \beta)} \text{T-App} \\
\text{(vii)} \frac{x : (\gamma \rightarrow \beta), y : ((\gamma \rightarrow \beta) \rightarrow \beta) \vdash y (\lambda z : \gamma. y x) : \beta}{x : (\gamma \rightarrow \beta) \vdash \lambda y : ((\gamma \rightarrow \beta) \rightarrow \beta). y (\lambda z : \gamma. y x) : ((\gamma \rightarrow \beta) \rightarrow \beta) \rightarrow \beta} \text{T-Abst} \\
\text{(viii)} \frac{x : (\gamma \rightarrow \beta) \vdash \lambda y : ((\gamma \rightarrow \beta) \rightarrow \beta). y (\lambda z : \gamma. y x) : ((\gamma \rightarrow \beta) \rightarrow \beta) \rightarrow \beta}{(\lambda x : (\gamma \rightarrow \beta). y : ((\gamma \rightarrow \beta) \rightarrow \beta). y (\lambda z : \gamma. y x)) : (\gamma \rightarrow \beta) \rightarrow ((\gamma \rightarrow \beta) \rightarrow \beta) \rightarrow \beta} \text{T-Abst}
\end{array}$$

### Problem

(2.8 c) Sketch a diagram of tree structure of derivation

*Solution.* Trivial. See tree above

### Problem

(2.8 d) Transform the derivation into flag notation

*Solution.*

1.	$x : \gamma \rightarrow \beta$	<b>Bound</b>
2.	$y : (\gamma \rightarrow \beta) \rightarrow \beta$	<b>Bound</b>
3.	$z : \gamma$	<b>Bound</b>
(ii)	4. $y : (\gamma \rightarrow \beta) \rightarrow \beta$	<b>T-Var</b>
(i)	5. $x : \gamma \rightarrow \beta$	<b>T-Var</b>
(iii)	6. $y x : \beta$	<b>4,5 T-App</b>
(iv)	7. $\lambda z : \gamma. y x : \gamma \rightarrow \beta$	<b>6 T-Abst</b>
(v)	8. $y : (\gamma \rightarrow \beta) \rightarrow \beta$	<b>T-Var</b>
(vi)	9. $y (\lambda z : \gamma. y x) : \beta$	<b>8,7 T-App</b>
(vii)	10. $\lambda y : (\gamma \rightarrow \beta) \rightarrow \beta. y (\lambda z : \gamma. y x)$	<b>9 T-Abst</b>
(viii)	11. $: ((\gamma \rightarrow \beta) \rightarrow \beta) \rightarrow \beta$	<b>10 T-Abst</b>

$\lambda x : (\gamma \rightarrow \beta). \lambda y : (\gamma \rightarrow \beta) \rightarrow \beta. y (\lambda z : \gamma. y x)$   
 $: (\gamma \rightarrow \beta) \rightarrow ((\gamma \rightarrow \beta) \rightarrow \beta) \rightarrow \beta$

## Problem

(2.9 a) Give derivations of the following judgement

$$x : \delta \rightarrow \delta \rightarrow \alpha, y : \gamma \rightarrow \alpha, z : \alpha \rightarrow \beta \vdash \\ \lambda u : \delta. \lambda v : \gamma. z(yv) : \delta \rightarrow \gamma \rightarrow \beta$$

*Solution.*

1.	$u : \delta$	<b>Bound</b>
2.	$v : \gamma$	<b>Bound</b>
3.	$y : \gamma \rightarrow \alpha$	<b>T-Var</b>
4.	$v : \gamma$	<b>T-Var</b>
5.	$yv : \alpha$	<b>3,4 T-App</b>
6.	$z : \alpha \rightarrow \beta$	<b>T-Var</b>
7.	$z(yv) : \beta$	<b>6,5 T-App</b>
8.	$\lambda v : \gamma. z(yv) : \gamma \rightarrow \beta$	<b>7 T-Abst</b>
9.	$\lambda u : \delta. \lambda v : \gamma. z(yv) : \delta \rightarrow \gamma \rightarrow \beta$	<b>8 T-Abst</b>

## Problem

(2.9 b) Give derivations of the following judgement

$$x : \delta \rightarrow \delta \rightarrow \alpha, y : \gamma \rightarrow \alpha, z : \alpha \rightarrow \beta \vdash \\ \lambda u : \delta. \lambda v : \gamma. z(xuu) : \delta \rightarrow \gamma \rightarrow \beta$$

*Solution.*

1.	$u : \delta$	<b>Bound</b>
2.	$v : \gamma$	<b>Bound</b>
3.	$x : \delta \rightarrow \delta \rightarrow \alpha$	<b>T-Var</b>
4.	$u : \delta$	<b>T-Var</b>
5.	$xu : \delta \rightarrow \alpha$	<b>3,4 T-App</b>
6.	$xuu : \alpha$	<b>5,4 T-App</b>
7.	$z : \alpha \rightarrow \beta$	<b>T-Var</b>
8.	$z(xuu) : \beta$	<b>7,6 T-App</b>
9.	$\lambda v : \gamma. z(xuu) : \gamma \rightarrow \beta$	<b>8 T-Abst</b>
10.	$\lambda u : \delta. \lambda v : \gamma. z(xuu) : \delta \rightarrow \gamma \rightarrow \beta$	<b>9 T-Abst</b>

### Problem

(2.10 a) Give derivation for

$$x z (y z)$$

*Solution.* Assume an context

$$\Gamma \vdash x : \alpha \rightarrow \beta \rightarrow \gamma$$

$$\Gamma \vdash y : \alpha \rightarrow \beta$$

$$\Gamma \vdash z : \alpha$$

- |  |                  |
|--|------------------|
| 1. $x : \alpha \rightarrow \beta \rightarrow \gamma$ | <b>T-Var</b>     |
| 2. $y : \alpha \rightarrow \beta$                    | <b>T-Var</b>     |
| 3. $z : \alpha$                                      | <b>T-Var</b>     |
| 4. $x z : \beta \rightarrow \gamma$                  | <b>1,3 T-App</b> |
| 5. $y z : \beta$                                     | <b>2,3 T-App</b> |
| 6. $x z (y z) : \gamma$                              | <b>4,5 T-App</b> |

### Problem

(2.10 b) Give derivation for

$$\lambda x : (\alpha \rightarrow \beta) \rightarrow \beta. x (y z)$$

*Solution.* Assume an context

$$\Gamma \vdash y : \gamma \rightarrow (\alpha \rightarrow \beta)$$

$$\Gamma \vdash z : \gamma$$

- |   |                  |
|---|------------------|
| 1. $x : (\alpha \rightarrow \beta) \rightarrow \beta$   | <b>Bound</b>     |
| 2. $x : (\alpha \rightarrow \beta) \rightarrow \beta$   | <b>T-Var</b>     |
| 3. $y : \gamma \rightarrow \alpha \rightarrow \beta$  | <b>T-Var</b>     |
| 4. $z : \gamma$   | <b>T-Var</b>     |
| 5. $y z : \alpha \rightarrow \beta$   | <b>3,4 T-App</b> |
| 6. $x (y z) : \beta$  | <b>2,5 T-App</b> |
| 7. $\lambda x : (\alpha \rightarrow \beta) \rightarrow \beta. x (y z) : ((\alpha \rightarrow \beta) \rightarrow \beta) \rightarrow \beta$ | <b>6 T-Abst</b>  |

### Problem

(2.10 c) Give derivation for

$$\lambda y : \alpha. \lambda z : \beta \rightarrow \gamma. z(x y y)$$

*Solution.* Assume a context

$$\Gamma \vdash x : \alpha \rightarrow \alpha \rightarrow \beta$$

1.	$y : \alpha$	<b>Bound</b>
2.	$z : \beta \rightarrow \gamma$	<b>Bound</b>
3.	$z : \beta \rightarrow \gamma$	<b>T-Var</b>
4.	$x : \alpha \rightarrow \alpha \rightarrow \beta$	<b>T-Var</b>
5.	$y : \alpha$	<b>T-Var</b>
6.	$x y : \alpha \rightarrow \beta$	<b>4,5 T-App</b>
7.	$x y y : \beta$	<b>6,5 T-App</b>
8.	$z(x y y) : \gamma$	<b>3,6 T-App</b>
9.	$\lambda z : \beta \rightarrow \gamma. z(x y y) : (\beta \rightarrow \gamma) \rightarrow \gamma$	<b>8 T-Abst</b>
10.	$\lambda y : \alpha. \lambda z : \beta \rightarrow \gamma. z(x y y) : \alpha \rightarrow (\beta \rightarrow \gamma) \rightarrow \gamma$	<b>8 T-Abst</b>

### Problem

(2.10 d) Give derivation for

$$\lambda x : \alpha \rightarrow \beta. y(x z) z$$

*Solution.* Consider a context

$$\begin{aligned} \Gamma &\vdash z : \alpha \\ \Gamma &\vdash y : \beta \rightarrow \alpha \rightarrow \gamma \end{aligned}$$

1.	$x : \alpha \rightarrow \beta$	<b>Bound</b>
2.	$x : \alpha \rightarrow \beta$	<b>T-Var</b>
3.	$z : \alpha$	<b>T-Var</b>
4.	$x z : \beta$	<b>2,3 T-App</b>
5.	$y : \beta \rightarrow \alpha \rightarrow \gamma$	<b>T-Var</b>
6.	$y(x z) : \alpha \rightarrow \gamma$	<b>5,4 T-App</b>
7.	$y(x z) z : \gamma$	<b>3,5 T-App</b>

$$8. \quad \lambda x : \alpha \rightarrow \beta. y (x z) z : (\alpha \rightarrow \beta) \rightarrow \gamma \quad 7 \text{ T-Abst}$$

### Problem

(2.11 a) Find an inhabitant of type and prove through derivation

$$(\alpha \rightarrow \alpha \rightarrow \gamma) \rightarrow \alpha \rightarrow \beta \rightarrow \gamma$$

*Solution.*

$$\lambda x : (\alpha \rightarrow \alpha \rightarrow \gamma). \lambda y : (\alpha). \lambda z : (\beta). x y y$$

*Proof.*

1.	$x : \alpha \rightarrow \alpha \rightarrow \gamma$	<b>Bound</b>
2.	$y : \alpha$	<b>Bound</b>
3.	$z : \beta$	<b>Bound</b>
4.	$x : \alpha \rightarrow \alpha \rightarrow \gamma$	T-Var
5.	$y : \alpha$	<b>Bound</b>
6.	$x y : \alpha \rightarrow \gamma$	4,5 T-App
7.	$\underline{x y y : \gamma}$	6,5 T-App
8.	$\lambda z : \beta. x y y : \beta \rightarrow \gamma$	7 T-Abst
9.	$\lambda y : \alpha. \lambda z : \beta. x y y : \alpha \rightarrow \beta \rightarrow \gamma$	8 T-Abst
10.		

$$\begin{aligned} & \lambda x : \alpha \rightarrow \alpha \rightarrow \gamma. \lambda y : \alpha. \lambda z : \beta. x y y \\ & : (\alpha \rightarrow \alpha \rightarrow \gamma) \rightarrow \alpha \rightarrow \beta \rightarrow \gamma \quad \text{8 T-Abst} \end{aligned}$$

■

### Problem

(2.11 b) Find an inhabitant of type and prove through derivation

$$((\alpha \rightarrow \gamma) \rightarrow \alpha) \rightarrow (\alpha \rightarrow \gamma) \rightarrow \beta \rightarrow \gamma$$

*Solution.*

$$\lambda x : (\alpha \rightarrow \gamma) \rightarrow \alpha. \lambda y : (\alpha \rightarrow \gamma). \lambda z : \beta. y (x y)$$

*Proof.*

1.	$x : (\alpha \rightarrow \gamma) \rightarrow \alpha$	<b>Bound</b>
2.	$y : \alpha \rightarrow \gamma$	<b>Bound</b>
3.	$z : \beta$	<b>Bound</b>
4.	$x : (\alpha \rightarrow \gamma) \rightarrow \alpha$	<b>T-Var</b>
5.	$y : \alpha \rightarrow \gamma$	<b>T-Var</b>
6.	$x y : \alpha$	<b>4,5 T-App</b>
7.	$y(x y) : \gamma$	<b>5,6 T-App</b>
8.	$\lambda z : \beta. x(x y) : \beta \rightarrow \gamma$	<b>7 T-Abst</b>
9.	$\lambda y : \alpha \rightarrow \gamma. \lambda z : \beta. x(x y) : (\alpha \rightarrow \gamma) \rightarrow \beta \rightarrow \gamma$	<b>7 T-Abst</b>
10.	$\lambda x : (\alpha \rightarrow \gamma) \rightarrow \alpha. \lambda y : \alpha \rightarrow \gamma. \lambda z : \beta. x(x y)$	
	$: ((\alpha \rightarrow \gamma) \rightarrow \alpha) \rightarrow (\alpha \rightarrow \gamma) \rightarrow \beta \rightarrow \gamma$	<b>7 T-Abst</b>

■

### Problem

(2.12 a) Construct a term of type

$$((\alpha \rightarrow \beta) \rightarrow \alpha) \rightarrow (\alpha \rightarrow \alpha \rightarrow \beta) \rightarrow \alpha$$

*Solution.*

$$\lambda x : (\alpha \rightarrow \beta) \rightarrow \alpha. \lambda y : \alpha \rightarrow \alpha \rightarrow \beta. x(\lambda z : \alpha. y z z)$$

*Proof.*

1.	$x : (\alpha \rightarrow \beta) \rightarrow \alpha$	<b>Bound</b>
2.	$y : \alpha \rightarrow \alpha \rightarrow \beta$	<b>Bound</b>
3.	$x : (\alpha \rightarrow \beta) \rightarrow \alpha$	<b>T-Var</b>
4.	$z : \alpha$	<b>Bound</b>
5.	$y : \alpha \rightarrow \alpha \rightarrow \beta$	<b>T-Var</b>
6.	$z : \alpha$	<b>Bound</b>
7.	$y z : \alpha \rightarrow \beta$	<b>5,6 T-App</b>
8.	$y z z : \beta$	<b>7,6 T-App</b>
9.	$\lambda z : \alpha. y z z : \alpha \rightarrow \beta$	<b>8 T-Abst</b>
10.	$x(\lambda z : \alpha. y z z) : \alpha$	<b>3,9 T-App</b>
11.	$\lambda y : \alpha \rightarrow \alpha \rightarrow \beta. x(\lambda z : \alpha. y z z) : (\alpha \rightarrow \alpha \rightarrow \beta) \rightarrow \alpha$	<b>3,9 T-App</b>

12.

$$\lambda x : (\alpha \rightarrow \beta) \rightarrow \alpha. \lambda y : \alpha \rightarrow \alpha \rightarrow \beta. x (\lambda z : \alpha. y z z)$$

$$: ((\alpha \rightarrow \beta) \rightarrow \alpha) \rightarrow (\alpha \rightarrow \alpha \rightarrow \beta) \rightarrow \alpha$$

**3,9 T-App**

■

### Problem

(2.12 b) Construct a term of type

$$((\alpha \rightarrow \beta) \rightarrow \alpha) \rightarrow (\alpha \rightarrow \alpha \rightarrow \beta) \rightarrow \beta$$

*Solution.*

$$\lambda x : (\alpha \rightarrow \beta) \rightarrow \alpha. \lambda y : \alpha \rightarrow \alpha \rightarrow \beta. (\lambda z : \alpha. y z z)(x (\lambda z : \alpha. y z z))$$

*Proof.*

1.	$x : (\alpha \rightarrow \beta) \rightarrow \alpha$	<b>Bound</b>
2.	$y : \alpha \rightarrow \alpha \rightarrow \beta$	<b>Bound</b>
3.	$x : (\alpha \rightarrow \beta) \rightarrow \alpha$	<b>T-Var</b>
4.	$z : \alpha$	<b>Bound</b>
5.	$y : \alpha \rightarrow \alpha \rightarrow \beta$	<b>T-Var</b>
6.	$z : \alpha$	<b>Bound</b>
7.	$y z : \alpha \rightarrow \beta$	<b>5,6 T-App</b>
8.	$\underline{y z z : \beta}$	<b>7,6 T-App</b>
9.	have $f := \lambda z : \alpha. y z z : \alpha \rightarrow \beta$	<b>8 T-Abst</b>
10.	$x f : \alpha$	<b>3,9 T-App</b>
11.	$\underline{f(x f) : \beta}$	<b>9,10 T-App</b>
12.	$\lambda y : \alpha \rightarrow \alpha \rightarrow \beta. f(x f) : (\alpha \rightarrow \alpha \rightarrow \beta) \rightarrow \beta$	<b>9,10 T-App</b>
13.	$\lambda x : (\alpha \rightarrow \beta) \rightarrow \alpha. \lambda y : \alpha \rightarrow \alpha \rightarrow \beta. f(x f)$ $: (\alpha \rightarrow \beta) \rightarrow \alpha \rightarrow (\alpha \rightarrow \alpha \rightarrow \beta) \rightarrow \beta$	<b>9,10 T-App</b>

■

### Problem

(2.13 a) Find a term of type

$$(\alpha \rightarrow \beta) \rightarrow \alpha \rightarrow \gamma$$

in context  $\Gamma$

$$x : \alpha \rightarrow \beta \rightarrow \gamma \in \Gamma$$

*Solution.*

$$\lambda u : \alpha \rightarrow \beta. \lambda v : \alpha. x v (u v)$$

*Proof.*

1.	$u : \alpha \rightarrow \beta$	<b>Bound</b>
2.	$v : \alpha$	<b>Bound</b>
3.	$x : \alpha \rightarrow \beta \rightarrow \gamma$	<b>T-Var</b>
4.	$v : \alpha$	<b>T-Var</b>
5.	$x v : \beta \rightarrow \gamma$	<b>3,4 T-App</b>
6.	$u : \alpha \rightarrow \beta$	<b>T-Var</b>
7.	$u v : \beta$	<b>6,4 T-App</b>
8.	$x v (u v) : \gamma$	<b>5,7 T-App</b>
9.	$\lambda v : \alpha. x v (u v) : \alpha \rightarrow \gamma$	<b>8 T-Abst</b>
10.	$\lambda u : \alpha \rightarrow \beta. \lambda v : \alpha. x v (u v) : (\alpha \rightarrow \beta) \rightarrow \alpha \rightarrow \gamma$	<b>9 T-Abst</b>

■

### Problem

(2.13 b) Find a term of type

$$\alpha \rightarrow (\alpha \rightarrow \beta) \rightarrow \gamma$$

in context  $\Gamma$

$$x : \alpha \rightarrow \beta \rightarrow \alpha \rightarrow \gamma \in \Gamma$$

*Solution.*

$$\lambda u : \alpha. \lambda v : \alpha \rightarrow \beta. x u (v u) u$$

*Proof.*

1.	$u : \alpha$	<b>Bound</b>
2.	$v : \alpha \rightarrow \beta$	<b>Bound</b>
3.	$x : \alpha \rightarrow \beta \rightarrow \alpha \rightarrow \gamma$	<b>T-Var</b>
4.	$u : \alpha$	<b>Bound</b>
5.	$v : \alpha \rightarrow \beta$	<b>Bound</b>
6.	$v u : \beta$	<b>5,4 T-App</b>
7.	$x u : \beta \rightarrow \alpha \rightarrow \gamma$	<b>3,4 T-App</b>
8.	$x u (v u) : \alpha \rightarrow \gamma$	<b>7,6 T-App</b>
9.	$x u (v u) u : \gamma$	<b>8,4 T-App</b>
10.	$\lambda v : \alpha \rightarrow \beta. x u (v u) u : (\alpha \rightarrow \beta) \gamma$	<b>9 T-Abst</b>
11.	$\lambda u : \alpha. \lambda v : \alpha \rightarrow \beta. x u (v u) u : \alpha \rightarrow (\alpha \rightarrow \beta) \gamma$	<b>10 T-Abst</b>

■

### Problem

(2.13 c) Find a term of type

$$(\alpha \rightarrow \gamma) \rightarrow (\beta \rightarrow \alpha) \rightarrow \gamma$$

in context  $\Gamma$

$$x : (\beta \rightarrow \gamma) \rightarrow \gamma \in \Gamma$$

*Solution.*

$$\lambda u : \alpha \rightarrow \gamma. \lambda v : \beta \rightarrow \alpha. x (\lambda y : \beta. u (v y))$$

*Proof.*

1.	$u : \alpha \rightarrow \gamma$	<b>Bound</b>
2.	$v : \beta \rightarrow \alpha$	<b>Bound</b>
3.	$x : (\beta \rightarrow \gamma) \rightarrow \gamma$	<b>T-Var</b>
4.	$y : \beta$	<b>Bound</b>
5.	$v : \beta \rightarrow \alpha$	<b>T-Var</b>
6.	$y : \beta$	<b>T-Var</b>
7.	$v y : \alpha$	<b>5,6 T-App</b>
8.	$u : \alpha \rightarrow \gamma$	<b>T-Var</b>
9.	$u (v y) : \gamma$	<b>8,7 T-App</b>
10.	$\lambda y : \beta. u (v y) : \beta \rightarrow \gamma$	<b>9 T-Abst</b>

11.	$x (\lambda y : \beta. u (v y)) : \gamma$	<b>10,3 T-App</b>
12.	$\lambda v : \beta \rightarrow \alpha. x (\lambda y : \beta. u (v y)) : (\beta \rightarrow \alpha) \rightarrow \gamma$	<b>11 T-Abst</b>
13.		
	$\lambda u : \alpha \rightarrow \gamma. \lambda v : \beta \rightarrow \alpha. x (\lambda y : \beta. u (v y))$ $: (\alpha \rightarrow \gamma) \rightarrow (\beta \rightarrow \alpha) \rightarrow \gamma$	<b>12 T-Abst</b>

■

### Problem

(2.14) Find an inhabitant of type  $\alpha \rightarrow \beta \rightarrow \gamma$  in the context  $\Gamma$

$$x : (\gamma \rightarrow \beta) \rightarrow \alpha \rightarrow \gamma \in \Gamma$$

*Solution.*

$$\lambda u : \alpha. \lambda v : \beta. x (\lambda z : \gamma. v) u$$

*Proof.*

1.	$u : \alpha$	<b>Bound</b>
2.	$v : \beta$	<b>Bound</b>
3.	$x : (\gamma \rightarrow \beta) \rightarrow \alpha \rightarrow \gamma$	<b>T-Var</b>
4.	$z : \gamma$	<b>Bound</b>
5.	$v : \beta$	<b>T-Var</b>
6.	$\lambda z : \gamma. v : \gamma \rightarrow \beta$	<b>5 T-Abst</b>
7.	$x (\lambda z : \gamma. v) : \alpha \rightarrow \gamma$	<b>3,6 T-App</b>
8.	$u : \alpha$	<b>T-Var</b>
9.	$x (\lambda z : \gamma. v) u : \gamma$	<b>7,8 T-App</b>
10.	$\lambda v : \beta. x (\lambda z : \gamma. v) u : \beta \rightarrow \gamma$	<b>9 T-Abst</b>
11.	$\lambda u : \alpha. \lambda v : \beta. x (\lambda z : \gamma. v) u : \alpha \rightarrow \beta \rightarrow \gamma$	<b>10 T-Abst</b>

■

### Problem

(2.15) Prove the thinning lemma via induction.

*Solution.*

---

*Lemma 1.* **Thinning Lemma.** Given two contexts  $\Gamma' \subseteq \Gamma$ ,

$$\Gamma' \vdash M : \sigma \Rightarrow \Gamma \vdash M : \sigma$$

---

We can do structural induction over the structure of  $M$  by the finite syntactic constructors of  $\lambda$  terms.

*Case 1 : Variable.* By the generation lemma  $M : \sigma \in \Gamma'$ . By the definition of subcontexts,  $M : \sigma \in \Gamma$ . It follows from the T-Var rule that  $\Gamma \vdash M : \sigma$ . ■

*Case 2 : Application.* Therefore the judgement must be of form

$$\Gamma' \vdash A B : \sigma$$

Where  $M = A B$ . By the generation lemma, there exists a type  $\tau$  such that  $\Gamma' \vdash A : \tau \rightarrow \sigma$  and  $\Gamma' \vdash B : \tau$ . By the principle of induction the thinning lemma holds for the terms  $A$  and  $B$ . By plugging in  $\Gamma \vdash A : \tau \rightarrow \sigma, B : \tau$ , the T-App rule proves  $\Gamma \vdash M : \sigma$ . ■

*Case 3 : Abstraction.* Therefore  $M$  is of a function type. We denote  $M$  as a abstraction term

$$\lambda x : \alpha. N : \alpha \rightarrow \beta$$

where  $x \notin \text{FR } N$ . By the generation lemma,  $\Gamma', x : \alpha \vdash N : \beta$ .

By the principle of induction, the thinning lemma already holds for  $N : \beta$ . Because  $\Gamma', x : \alpha \subseteq \Gamma, x : \alpha$ , the thinning lemma proves  $\Gamma, x : \alpha \vdash N : \beta$ . By the T-Abst rule,  $\Gamma \vdash \lambda x : \alpha. N : \alpha \rightarrow \beta$ . ■

### Problem

(2.16) Prove the subterm lemma.

*Solution.*

---

*Lemma 2.* **Subterm Lemma.** If  $M \in \Lambda_T$  is legal, then all subterms of  $M$  are.

---

Let's formalize the lemma. The definition of a legal term is a term  $M$  with the existence of a context  $\Gamma$  and a type  $\tau$  such that  $\Gamma \vdash M : \tau$ , in other words,  $M$  is typable. Because  $M$  is typable iff there's an applicable inference rule and each term appearing in the premise of each inference rule is an immediate subterm of  $M$  that

could construct  $M$ , structural induction could be done down the typing tree as it is also an induction down the term's structure, inducting over each subterm.

*Case 1 : Variable.*  $M$  is the only subterm of  $M$  thus is trivially typable and legal. ■

*Case 2 : Application.* Therefore the derivation must be of form

$$\frac{\Gamma \vdash A : \tau \rightarrow \sigma \quad \Gamma \vdash B : \tau}{\Gamma \vdash A B : \sigma} \text{T-App}$$

Where  $M = A B$ . From the generation lemma, those  $A$  and  $B$  must exist typable, thus legal. By the principle of induction, the subterm lemma holds for  $A$  and  $B$ . Denote the lemma as  $P(\Lambda_{\mathbb{T}})$ :

$$\begin{aligned} & P(A B) \wedge P(A) \wedge P(B) \wedge (\forall a \in \text{Sub } A, P(a)) \wedge (\forall b \in \text{Sub } B, P(b)) \\ &= \forall m \in \{A B, A, B\} \cup \text{Sub } A \cup \text{Sub } B, P(m) \\ &= \forall m \in \{A B\} \cup \text{Sub } A \cup \text{Sub } B, P(m) \\ &= \forall m \in \text{Sub } A B, P(m) = P(A B) = P(M) \end{aligned}$$

■

*Case 3 : Abstraction.* Therefore the derivation must be of form

$$\frac{\Gamma, x : \alpha \vdash N : \beta}{\Gamma \vdash \lambda x : \alpha. N : \alpha \rightarrow \beta} \text{T-Abst}$$

Where  $M = \lambda x : \alpha. N$  and  $x \notin \text{FR } N$ . By the generation lemma,  $N$  is typable under the context  $\Gamma, x : \alpha$ , thus valid. By the inductive hypothesis, the subterm lemma holds for  $N$ . Denote the lemma as  $P(\Lambda_{\mathbb{T}})$

$$\begin{aligned} & P(\lambda x : \alpha. N) \wedge (P(N) \wedge \forall n \in \text{Sub } N, P(n)) \\ &= \forall m \in \{\lambda x : \alpha. N, N\} \cup \text{Sub } N, P(m) \\ &= \forall m \in \{\lambda x : \alpha. N\} \cup \text{Sub } N, P(m) = P(\lambda x : \alpha. N) = P(M) \end{aligned}$$

■

### Problem

(2.17) Prove the uniqueness of types lemma.

*Solution.*

*Lemma 3.* Assume  $\Gamma \vdash M : \sigma$  and  $\Gamma \vdash M : \tau$ . Then  $\sigma = \tau$ .

Again do induction on the construction of  $M$ .

*Case 1 : Variable.* By the generation lemma,  $M : \sigma \in \Gamma$  and  $M : \tau \in \Gamma$  By the definition of contexts. Because the context could only provide one judgement per variable,  $M$  must have a unique type. ■

*Case 2 : Application.*  $M$  must be of form  $A B$ . By the generation lemma, two judgements could be obtained where  $\alpha, \beta \in \mathbb{T}$ :

$$\Gamma \vdash A : \alpha \rightarrow \sigma, B : \alpha$$

$$\Gamma \vdash A : \beta \rightarrow \tau, B : \beta$$

By the inductive hypothesis, the uniqueness lemma already holds for  $A$  and  $B$ . Therefore  $\alpha \equiv \beta$ . Because  $A$  is a function type, it could be expanded into two forms:

$$\Gamma \vdash (\lambda x : \alpha. N) : \alpha \rightarrow \sigma$$

$$\Gamma \vdash (\lambda x : \beta. N) : \beta \rightarrow \tau$$

By the generation lemma and the equivalence  $\alpha \equiv \beta$ , the above further deduce to

$$\Gamma, x : \alpha \vdash N : \sigma$$

$$\Gamma, x : \alpha \vdash N : \tau$$

By the principle of induction, the uniqueness lemma holds for  $N$ , therefore  $\sigma \equiv \tau$ . By substituting  $\tau$  for  $\sigma$ , the final type for  $M$  should be the same in the conclusion for both judgements. ■

*Case 3 : Abstraction.* Because  $M$  is a function type, two judgements could be obtained for some  $\alpha, \beta \in \mathbb{T}$

$$\Gamma \vdash \lambda x : \gamma. N : \sigma \text{ where } \sigma \equiv \gamma \rightarrow \alpha$$

$$\Gamma \vdash \lambda x : \gamma. N : \tau \text{ where } \tau \equiv \gamma \rightarrow \beta$$

By applying the generation lemma

$$\Gamma, x : \gamma \vdash N : \alpha$$

$$\Gamma, x : \gamma \vdash N : \beta$$

By the principle of induction,  $N$  already conform to the lemma, thus  $\alpha \equiv \beta$ . Therefore,  $\gamma \rightarrow \alpha \equiv \gamma \rightarrow \beta$ . ■

### Problem

(2.18) Prove the *Subject Reduction* lemma.

---

*Lemma 4.* If  $\Gamma \vdash L : \rho$  and  $L \xrightarrow{\beta} L'$ , then  $\Gamma \vdash L' : \rho$

---

*Solution.* Proof by generation of  $L'$ : which is by establishing all possibilities of how  $L'$  was derived.

*Basis : Direct Reduction.*

$$L \equiv (\lambda x : \sigma. M) N \text{ and } L' \equiv M [x := N]$$

By the generation lemma and some deconstructing we have  $M : \gamma$  under  $\Gamma, x : \sigma$  and  $N : \sigma$  under  $\Gamma$ . By applying the T-App rule,  $L : \gamma$ . By the uniqueness of types lemma,  $\gamma \equiv \rho$ . Now we define  $\Gamma' \equiv \emptyset$  and  $\Gamma'' \equiv \Gamma$ , by the substitution lemma we have  $\Gamma \vdash M [x := N] : \gamma$ . We substitue  $\gamma \equiv \rho$  back and get  $\Gamma \vdash M [x := N] : \rho$ . ■

*Case 1 : Left Term Reduction.*

$$L \equiv M N \text{ and } L' \equiv M' N \text{ where } M \xrightarrow{\beta} M'$$

We type the terms  $M : \sigma \rightarrow \rho$  and  $N : \sigma$ . By the inductive hypothesis,  $\Gamma \vdash M' : \sigma \rightarrow \rho$ . By the T-App rule,  $\Gamma \vdash M' N : \rho$ . ■

*Case 2 : Left Term Reduction.*

$$L \equiv M N \text{ and } L' \equiv M N' \text{ where } N \xrightarrow{\beta} N'$$

We type the terms  $M : \sigma \rightarrow \rho$  and  $N : \sigma$ . By the inductive hypothesis,  $\Gamma \vdash N' : \sigma$ . By the T-App rule,  $\Gamma \vdash M N' : \rho$ . ■

*Case 3 : Reduction In Body.*

$$L \equiv (\lambda x : \tau. M) \text{ and } L' \equiv (\lambda x : \tau. M') \text{ where } M \xrightarrow{\beta} M'$$

Let  $\sigma \in \mathbb{T}$  where  $\rho \equiv \tau \rightarrow \sigma$ . By the generation lemma  $\Gamma, x : \tau \vdash M : \sigma$ . By the inductive hypothesis,  $\Gamma, x : \tau \vdash M' : \sigma$ , and by the T-Abst rule,  $\Gamma \vdash (\lambda x : \tau. M') : \tau \rightarrow \sigma$ . Therefore,  $\Gamma \vdash L' : \rho$  by substituting  $\tau \rightarrow \sigma$  for  $\rho$ . ■

—

Completed Dec 15 2:24 am.