

EXERCISES

CHAPTER 3

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1. Redacted

Definition Some rules for reference.

$$\begin{array}{c} \frac{x : \sigma \in \Gamma \quad \Gamma \text{ is a } \lambda 2 \text{ context}}{\Gamma \vdash x : \sigma} \text{ (T-Var)} \qquad \frac{\Gamma \vdash M : \sigma \rightarrow \tau \quad \Gamma \vdash N : \sigma}{\Gamma \vdash M N : \tau} \text{ (T-App)} \\[10pt] \frac{\Gamma, x : \sigma \vdash M : \tau}{\Gamma \vdash \lambda x : \sigma. M : \sigma \rightarrow \tau} \text{ (T-Abst)} \qquad \frac{\alpha \in \mathbb{T}_2 \quad \text{FV } \alpha \subseteq \text{dom } \Gamma}{\alpha : * \in \Gamma} \text{ (T-Form)} \\[10pt] \frac{\Gamma \vdash M : \Pi_{\alpha:*.} A \quad \Gamma \vdash B : *}{\Gamma \vdash M B : A [\alpha := B]} \text{ (T2-App)} \qquad \frac{\Gamma, \alpha : * \vdash M : A}{\Gamma \vdash \lambda \alpha : *. M : \Pi_{\alpha:*.} A} \text{ (T2-Abst)} \end{array}$$

In this document, convention is that all type judgements in a proof tree, unless stated otherwise, is derived from a single and unique $\lambda 2$ context per tree. Multiple conclusions might be drawn on a single line from usage of the same inference rule for compactness. Eg:

ex 1. $\alpha, \beta : *$ **T-Form**

Is shorthand for

ex 1. $\Gamma \vdash \alpha : *$ **T-Form**

ex 2. $\Gamma \vdash \beta : *$ **T-Form**

Problem

(3.1) How many $\lambda 2$ contexts consisting of four and only four declarations

- (1) $\Gamma \vdash \alpha : *$ (2) $\Gamma \vdash \beta : *$
 (3) $\Gamma \vdash f : \alpha \rightarrow \beta$ (4) $\Gamma \vdash x : \alpha$

Solution. The last two declarations depend on the first two. Therefore this is an easy combinatorics problem: $2! \times 2! = 4$ contexts:

- 1 – 2 – 3 – 4 1 – 2 – 4 – 3
 2 – 1 – 3 – 4 2 – 1 – 4 – 3

Problem

(3.2) Give a full derivation in $\lambda 2$ to show the following type term is legal:

$$M \equiv \lambda \alpha, \beta, \gamma : *. \lambda f : \alpha \rightarrow \beta. \lambda g : \beta \rightarrow \gamma. \lambda x : \alpha. g (f x)$$

Solution.

1.	$\alpha : *$	Bound
2.	$\beta : *$	Bound
3.	$\gamma : *$	Bound
4.	$f : \alpha \rightarrow \beta$	Bound
5.	$g : \beta \rightarrow \gamma$	Bound
6.	$x : \alpha$	Bound
7.	$\alpha, \beta, \gamma : *$	T-Form
8.	$\alpha \rightarrow \beta, \beta \rightarrow \gamma : *$	T-Form
9.	$f : \alpha \rightarrow \beta, x : \alpha$	T-Var
10.	$f x : \beta$	8,8 T-App
11.	$g : \beta \rightarrow \gamma$	T-Var
12.	$g (f x) : \gamma$	11,10 T-App
13.	$\lambda x : \alpha. g (f x) : \alpha \rightarrow \gamma$	12 T-Abst
14.	$\lambda g : \beta \rightarrow \gamma. \lambda x : \alpha. g (f x) : (\beta \rightarrow \gamma) \rightarrow \alpha \rightarrow \gamma$	13 T-Abst
15.	$\lambda f : \alpha \rightarrow \beta. \lambda g : \beta \rightarrow \gamma. \lambda x : \alpha. g (f x) : (\alpha \rightarrow \beta) \rightarrow (\beta \rightarrow \gamma) \rightarrow \alpha \rightarrow \gamma$	14 T-Abst

16.	$\frac{\lambda\gamma : *. \lambda f : \alpha \rightarrow \beta. \lambda g : \beta \rightarrow \gamma. \lambda x : \alpha. g (f x) : \Pi\gamma : *. (\alpha \rightarrow \beta) \rightarrow (\beta \rightarrow \gamma) \rightarrow \alpha \rightarrow \gamma}{\lambda\beta, \gamma : *. \lambda f : \alpha \rightarrow \beta. \lambda g : \beta \rightarrow \gamma. \lambda x : \alpha. g (f x) : \Pi\beta, \gamma : *. (\alpha \rightarrow \beta) \rightarrow (\beta \rightarrow \gamma) \rightarrow \alpha \rightarrow \gamma}}$	15 T2-Abst
17.	$\lambda\beta, \gamma : *. \lambda f : \alpha \rightarrow \beta. \lambda g : \beta \rightarrow \gamma. \lambda x : \alpha. g (f x) : \Pi\beta, \gamma : *. (\alpha \rightarrow \beta) \rightarrow (\beta \rightarrow \gamma) \rightarrow \alpha \rightarrow \gamma$	16 T2-Abst
18.	$\lambda\alpha, \beta, \gamma : *. \lambda f : \alpha \rightarrow \beta. \lambda g : \beta \rightarrow \gamma. \lambda x : \alpha. g (f x) : \Pi\alpha, \beta, \gamma : *. (\alpha \rightarrow \beta) \rightarrow (\beta \rightarrow \gamma) \rightarrow \alpha \rightarrow \gamma$	17 T2-Abst

Problem

(3.3 a) Given M in 3.2, and a context Γ such that

$$\Gamma \vdash \text{nat} : *$$

$$\Gamma \vdash \text{bool} : *$$

$$\Gamma \vdash \text{succ} : \text{nat} \rightarrow \text{nat}$$

$$\Gamma \vdash \text{even} : \text{nat} \rightarrow \text{bool}$$

Prove $M \text{ nat nat bool succ even}$ is legal.

Solution. Proof by deriving the term's type.

Proof.

1.	$M : \Pi\alpha, \beta, \gamma : *. (\alpha \rightarrow \beta) \rightarrow (\beta \rightarrow \gamma) \rightarrow \alpha \rightarrow \gamma$	T-Var
2.	$\text{nat}, \text{bool} : *$	T-Form
3.	$M \text{ nat} : \Pi\beta, \gamma : *. (\text{nat} \rightarrow \beta) \rightarrow (\beta \rightarrow \gamma) \rightarrow \text{nat} \rightarrow \gamma$	2,1 T2-App
4.	$M \text{ nat nat} : \Pi\gamma : *. (\text{nat} \rightarrow \text{nat}) \rightarrow (\text{nat} \rightarrow \gamma) \rightarrow \text{nat} \rightarrow \gamma$	2,3 T2-App
5.	$M \text{ nat nat bool} : (\text{nat} \rightarrow \text{nat}) \rightarrow (\text{nat} \rightarrow \text{bool}) \rightarrow \text{nat} \rightarrow \text{bool}$	2,3 T2-App
6.	$\text{succ} : \text{nat} \rightarrow \text{nat}, \text{even} : \text{nat} \rightarrow \text{bool}$	T-Var
7.	$M \text{ nat nat bool succ} : (\text{nat} \rightarrow \text{bool}) \rightarrow \text{nat} \rightarrow \text{bool}$	6,5 T-App
8.	$M \text{ nat nat bool succ even} : \text{nat} \rightarrow \text{bool}$	6,7 T-App

■

Problem

(3.3 b.i) Prove $\lambda x : \text{nat}. \text{even} (\text{succ } x)$ via 3.3 a.

Solution. The result of beta reduction on the term in 3.3 a is what we are proving.

Proof.

$$\begin{aligned}
 & M \text{ nat nat bool succ even} \\
 & \equiv (\lambda \alpha, \beta, \gamma, f, g. \lambda x : \alpha. g (f x)) \text{ nat nat bool succ even} \\
 & \xrightarrow[\beta]{\rightarrow} (\lambda f : \text{nat} \rightarrow \text{nat}. \lambda g : \text{nat} \rightarrow \text{bool}. \lambda x : \text{nat}. g (f x)) \text{ succ even} \\
 & \xrightarrow[\beta]{\rightarrow} (\lambda x : \text{nat}. \text{even} (\text{succ } x))
 \end{aligned}$$

By the subject reduction lemma, $\lambda x : \text{nat}. \text{even} (\text{succ } x) : \text{nat} \rightarrow \text{bool}$, thus is legal. ■

Problem

(3.3 b.ii) Prove $\lambda x : \text{nat}. \text{even} (\text{succ } x)$ via derivation in the context provided in 3.3 a.

Solution.

Proof.

1.	$\text{nat}, \text{bool} : *$	T-Form
2.	$x : \text{nat}$	Bound
3.	$\text{succ} : \text{nat} \rightarrow \text{nat}$	T-Var
4.	$x : \text{nat}$	T-Var
5.	$\text{succ } x : \text{nat}$	3,4 T-App
6.	$\text{even} : \text{nat} \rightarrow \text{bool}$	T-Var
7.	$\text{even} (\text{succ } x) : \text{bool}$	6,5 T-App
8.	$\lambda x : \text{nat}. \text{even} (\text{succ } x) : \text{nat} \rightarrow \text{bool}$	7 T-Abst

■

Problem

(3.4) Give a shorthand (omit T-Var and T-Form) derivation in $\lambda 2$ to show the following term is valid in $\Gamma \equiv \text{nat} : *, \text{bool} : *$

$$(\lambda \alpha, \beta : *. \lambda f : \alpha \rightarrow \alpha. \lambda g : \alpha \rightarrow \beta. \lambda x : \alpha. g (f (f x))) \text{ nat bool}$$

Solution.

Proof.

1.	$\alpha, \beta : *$	Bound
2.	$f : \alpha \rightarrow \alpha$	Bound
3.	$g : \alpha \rightarrow \beta$	Bound
4.	$x : \alpha$	Bound
5.	$f x : \alpha$	*, T-App
6.	$f (f x) : \alpha$	*,5 T-App
7.	$g (f (f x)) : \beta$	*,6 T-App
8.	$\lambda x : \alpha. g (f (f x)) : \alpha \rightarrow \beta$	7 T-Abst
9.	$\lambda g : \alpha \rightarrow \beta. \lambda x : \alpha. g (f (f x)) : (\alpha \rightarrow \beta) \rightarrow \alpha \rightarrow \beta$	8 T-Abst
10.	$\lambda f : \alpha \rightarrow \alpha. \lambda g : \alpha \rightarrow \beta. \lambda x : \alpha. g (f (f x))$ $: (\alpha \rightarrow \alpha) \rightarrow (\alpha \rightarrow \beta) \rightarrow \alpha \rightarrow \beta$	9 T-Abst
11.		
	$\lambda \alpha, \beta : *. \lambda f : \alpha \rightarrow \alpha. \lambda g : \alpha \rightarrow \beta. \lambda x : \alpha. g (f (f x))$ $: \Pi \alpha, \beta : *. (\alpha \rightarrow \alpha) \rightarrow (\alpha \rightarrow \beta) \rightarrow \alpha \rightarrow \beta$	10 T2-Abst
12.		
	$(\lambda \alpha, \beta : *. \lambda f : \alpha \rightarrow \alpha. \lambda g : \alpha \rightarrow \beta. \lambda x : \alpha. g (f (f x))) \text{ nat}$ $: \Pi \beta : *. (\text{nat} \rightarrow \text{nat}) \rightarrow (\text{nat} \rightarrow \beta) \rightarrow \text{nat} \rightarrow \beta$	*,11 T2-App
13.		
	$(\lambda \alpha, \beta : *. \lambda f : \alpha \rightarrow \alpha. \lambda g : \alpha \rightarrow \beta. \lambda x : \alpha. g (f (f x))) \text{ nat bool}$ $: \Pi \beta : *. (\text{nat} \rightarrow \text{nat}) \rightarrow (\text{nat} \rightarrow \text{bool}) \rightarrow \text{nat} \rightarrow \text{bool}$	*,12 T2-App

■

Problem

(3.5 a) Let $\perp \equiv \Pi \alpha : *. \alpha$. Prove \perp is legal.

Solution. Here a notion called kind checking is introduced. This has not yet been discussed in this book (?)

Proof.

1.	$\alpha : *$	Bound
2.	$\alpha : *$	T-Form

3. $\Pi\alpha : * . \alpha : * \rightarrow *$ **T2-Abst**

■

Problem

(3.5 b) Consider the context $\Gamma \equiv \beta : *, x : \perp$. Find an inhabitant of type β under Γ .

Solution. $x \beta$ is. Because x is of second-order type, it must be parametric to a type, thus x is of form $\lambda\alpha : * . M$ where $\Gamma, \alpha : * \vdash M : \alpha$.

Proof.

1. $x : \Pi\alpha : * . \alpha$ **T-Var**
2. $\beta : *$ **T-Form**
3. $x \beta : \beta$ **1,2 T2-App**

■

Problem

(3.5 c) Give three inhabitants of $\beta \rightarrow \beta$ in β -nf under Γ in 3.5 b.

Solution.

1. $\lambda y : \beta . y$.

Proof.

1. $y : \beta$ **Bound**
2. $\lambda y : \beta . y$ **T-Var**
3. $\lambda y : \beta . y : \beta \rightarrow \beta$ **2 T-Abst**

■

2. $\lambda y : \beta . x \beta$.

Proof.

1. $y : \beta$ **Bound**
2. $x : \Pi\alpha : * . \alpha$ **T-Var**
3. $\beta : *$ **T-Form**
4. $x \beta : \beta$ **2,3 T2-App**

5. $\lambda y : \beta. x \beta : \beta \rightarrow \beta$ **4 T-Abst**

■

3. $\lambda y : \beta. x (\beta \rightarrow \beta) y.$

Proof.

1.	$y : \beta$	Bound
2.	$x : \Pi \alpha : * . \alpha$	T-Var
3.	$\beta \rightarrow \beta : *$	T-Form
4.	$x (\beta \rightarrow \beta) : \beta \rightarrow \beta$	2,3 T2-App
5.	$y : \beta$	T-Var
6.	$x (\beta \rightarrow \beta) y : \beta$	4,5 T-App
7.	$\lambda y : \beta. x (\beta \rightarrow \beta) y : \beta \rightarrow \beta$	5 T-Abst

■

Problem

(3.5 d) Prove that the following terms inhabit the same type in Γ :

$$\lambda f : \beta \rightarrow \beta \rightarrow \beta. f (x \beta)(x \beta)$$

$$x ((\beta \rightarrow \beta \rightarrow \beta) \rightarrow \beta)$$

Solution. We simply derive the types.

First Term.

1.	$f : \beta \rightarrow \beta \rightarrow \beta$	Bound
2.	$f : \beta \rightarrow \beta \rightarrow \beta$	T-Var
3.	$x : \Pi \alpha : * . \alpha$	T-Var
4.	$\beta : *$	T-Form
5.	$x \beta : \beta$	3,4 T2-App
6.	$f (x \beta) : \beta \rightarrow \beta$	2,5 T-App
7.	$f (x \beta)(x \beta) : \beta$	6,5 T-App
8.	$\lambda f : \beta \rightarrow \beta \rightarrow \beta. f (x \beta)(x \beta) : (\beta \rightarrow \beta \rightarrow \beta) \rightarrow \beta$	6 T-Abst

■

Second Term.

1. $(\beta \rightarrow \beta \rightarrow \beta) \rightarrow \beta : *$ **T-Form**
2. $x : \Pi \alpha : *. \alpha$ **T-Var**
3. $x ((\beta \rightarrow \beta \rightarrow \beta) \rightarrow \beta) : (\beta \rightarrow \beta \rightarrow \beta) \rightarrow \beta$ **2,1 T2-App**

■

The two terms were shown to both inhabit $(\beta \rightarrow \beta \rightarrow \beta) \rightarrow \beta$.

Problem

(3.6 a) Find inhabitant of type

$$\Pi \alpha, \beta : *. (\text{nat} \rightarrow \alpha) \rightarrow (\alpha \rightarrow \text{nat} \rightarrow \beta) \rightarrow \text{nat} \rightarrow \beta$$

In context $\Gamma \equiv \text{nat} : *$.

Solution.

$$\lambda \alpha, \beta : *. \lambda x : \text{nat} \rightarrow \alpha. \lambda y : (\alpha \rightarrow \text{nat} \rightarrow \beta). \lambda z : \text{nat}. y (x z) z$$

Proof.

- | | | |
|-----|---|--------------------|
| 1. | $\alpha, \beta : *$ | Bound |
| 2. | $\text{nat} \rightarrow \alpha : *$ | T-Form |
| 3. | $x : \text{nat} \rightarrow \alpha$ | Bound |
| 4. | $\alpha \rightarrow \text{nat} \rightarrow \beta : *$ | T-Form |
| 5. | $y : \alpha \rightarrow \text{nat} \rightarrow \beta$ | Bound |
| 6. | $\text{nat} : *$ | Bound |
| 7. | $z : \text{nat}$ | Bound |
| 8. | $y : \alpha \rightarrow \text{nat} \rightarrow \beta$ | T-Var |
| 9. | $x : \text{nat} \rightarrow \alpha$ | T-Var |
| 10. | $z : \text{nat}$ | T-Var |
| 11. | $x z : \alpha$ | 9,10 T-App |
| 12. | $y (x z) : \text{nat} \rightarrow \beta$ | 8,11 T-App |
| 13. | $y (x z) z : \beta$ | 12,10 T-App |
| 14. | $\lambda z : \text{nat}. y (x z) z : \text{nat} \rightarrow \beta$ | 13 T-Abst |
| 15. | $\lambda y : \alpha \rightarrow \text{nat} \rightarrow \beta. \lambda z : \text{nat}. y (x z) z$
$: (\alpha \rightarrow \text{nat} \rightarrow \beta) \rightarrow \text{nat} \rightarrow \beta$ | 14 T-Abst |
| 16. | $\lambda x : \text{nat} \rightarrow \alpha. \lambda y : \alpha \rightarrow \text{nat} \rightarrow \beta. \lambda z : \text{nat}. y (x z) z$
$: (\text{nat} \rightarrow \alpha) \rightarrow (\alpha \rightarrow \text{nat} \rightarrow \beta) \rightarrow \text{nat} \rightarrow \beta$ | 15 T2-Abst |

17.

$\lambda\alpha, \beta : *. x : \mathbf{nat} \rightarrow \alpha. y : \alpha \rightarrow \mathbf{nat} \rightarrow \beta. \lambda z : \mathbf{nat}. y (x z) z$
 $: \Pi\alpha, \beta : *. (\mathbf{nat} \rightarrow \alpha) \rightarrow (\alpha \rightarrow \mathbf{nat} \rightarrow \beta) \rightarrow \mathbf{nat} \rightarrow \beta$

16 T2-Abst

■

Problem

(3.6 b) Find inhabitant of type

$$\Pi\delta : *. ((\alpha \rightarrow \gamma) \rightarrow \delta) \rightarrow (\alpha \rightarrow \beta) \rightarrow (\beta \rightarrow \gamma) \rightarrow \delta$$

In context $\Gamma \equiv \alpha : *, \beta : *, \gamma : *$

Solution.

$$\lambda\delta : *. \lambda x : (\alpha \rightarrow \gamma) \rightarrow \delta. \lambda y : \alpha \rightarrow \beta. \lambda z : \beta \rightarrow \gamma. x (\lambda u : \alpha. z (y u))$$

A derivation in shorthand will be given (omitting T-Form / T-Var)

Proof.

1.	$\delta : *$	Bound
2.	$x : (\alpha \rightarrow \gamma) \rightarrow \delta$	Bound
3.	$y : \alpha \rightarrow \beta$	Bound
4.	$z : \beta \rightarrow \gamma$	Bound
5.	$u : \alpha$	Bound
6.	$y u : \beta$	*, T-App
7.	$z (y u) : \gamma$	*, 6 T-App
8.	$\lambda u : \alpha. z (y u) : \alpha \rightarrow \gamma$	7 T-Abst
9.	$x (\lambda u : \alpha. z (y u)) : \delta$	8 T-Abst
10.	$\lambda z : \beta \rightarrow \gamma. x (\lambda u : \alpha. z (y u)) : (\beta \rightarrow \gamma) \rightarrow \delta$	9 T-Abst
	$\lambda y : \alpha \rightarrow \beta. \lambda z : \beta \rightarrow \gamma. x (\lambda u : \alpha. z (y u))$	
11.	$: (\alpha \rightarrow \beta) \rightarrow (\beta \rightarrow \gamma) \rightarrow \delta$	10 T-Abst
	$\lambda x : (\alpha \rightarrow \gamma) \rightarrow \delta. \lambda y : \alpha \rightarrow \beta. \lambda z : \beta \rightarrow \gamma. x (\lambda u : \alpha. z (y u))$	
12.	$: ((\alpha \rightarrow \gamma) \rightarrow \delta) \rightarrow (\alpha \rightarrow \beta) \rightarrow (\beta \rightarrow \gamma) \rightarrow \delta$	11 T-Abst
13.		

$$\lambda\delta : *. \lambda x : (\alpha \rightarrow \gamma) \rightarrow \delta. \lambda y : \alpha \rightarrow \beta. \lambda z : \beta \rightarrow \gamma. x (\lambda u : \alpha. z (y u))$$

$$: \Pi\delta : *. ((\alpha \rightarrow \gamma) \rightarrow \delta) \rightarrow (\alpha \rightarrow \beta) \rightarrow (\beta \rightarrow \gamma) \rightarrow \delta$$

12 T2-Abst

■

Problem

(3.6 c) Find inhabitant of type

$$\Pi\alpha, \beta, \gamma : *. (\alpha \rightarrow (\beta \rightarrow \alpha) \rightarrow \gamma) \rightarrow \alpha \rightarrow \gamma$$

In the empty context

Solution.

$$\lambda\alpha, \beta, \gamma : *. \lambda f : (\alpha \rightarrow (\beta \rightarrow \alpha) \rightarrow \gamma). \lambda x : \alpha. f x (\lambda u : \beta. x)$$

Proof.

1.	α, β, γ	Bound
2.	$f : \alpha \rightarrow (\beta \rightarrow \alpha) \rightarrow \gamma$	Bound
3.	$x : \alpha$	Bound
4.	$f x : (\beta \rightarrow \alpha) \rightarrow \gamma$	*, T-App
5.	$u : \beta$	Bound
6.	$x : \alpha$	T-Var
7.	$\lambda u : \beta. x : \beta \rightarrow \alpha$	6 T-Abst
8.	$f x (\lambda u : \beta. x) : \gamma$	4,7 T-App
9.	$\lambda x : \alpha. f x (\lambda u : \beta. x) : \alpha \rightarrow \gamma$	8 T-Abst
	$\lambda f : \alpha \rightarrow (\beta \rightarrow \alpha) \rightarrow \gamma. \lambda x : \alpha. f x (\lambda u : \beta. x)$	
10.	$: (\alpha \rightarrow (\beta \rightarrow \alpha) \rightarrow \gamma) \rightarrow \alpha \rightarrow \gamma$	9 T-Abst
11.		
	$\lambda\alpha, \beta, \gamma : *. \lambda f : \alpha \rightarrow (\beta \rightarrow \alpha) \rightarrow \gamma. \lambda x : \alpha. f x (\lambda u : \beta. x)$	
	$: \Pi\alpha, \beta, \gamma : *. (\alpha \rightarrow (\beta \rightarrow \alpha) \rightarrow \gamma) \rightarrow \alpha \rightarrow \gamma$	10 T2-Abst

■

Problem

(3.7) Let $\perp \equiv \Pi\alpha : *. \alpha$ and context $\Gamma \equiv \alpha : *, \beta : *, x : \alpha \rightarrow \perp, f : (\alpha \rightarrow \alpha) \rightarrow \alpha$. Give a derivation that succesively calculate an inhabitant of α and β , both in context Γ .

Solution. Have $M : \alpha := f (\lambda n : \alpha. n)$. Then $\Gamma \vdash x M \beta : \beta$.

Typing M.

1. $f : (\alpha \rightarrow \alpha) \rightarrow \alpha$ **T-Var**
2. $n : \alpha$ **Bound**
3. $\boxed{n : \alpha}$ **T-Var**
4. $\lambda n : \alpha. n : \alpha \rightarrow \alpha$ **3 T-Abst**
5. $f (\lambda n : \alpha. n) : \alpha$ **1,4 T-App**

■

Typing $x M \beta$.

1. $M : \alpha$ **T-Var**
2. $x : \alpha \rightarrow \Pi \alpha : * . \alpha$ **T-Var**
3. $M x : \Pi \alpha : * . \alpha$ **2,1 T-App**
4. $M x \beta : \beta$ **3,* T2-App**

■

Problem

(3.8) Recall $K \equiv \lambda x y . x \in \Lambda$ from untyped lambda calculus. Consider the following types

$$T_1 \equiv \Pi \alpha, \beta : * . \alpha \rightarrow \beta \rightarrow \alpha \quad T_2 \equiv \Pi \alpha : * . \alpha \rightarrow (\Pi \beta : * . \beta \rightarrow \alpha)$$

Find inhabitants of both type $t_1 : T_1$ and $t_2 : T_2$ under the empty context, which may be considered the closed $\lambda 2$ form of $K \in \Lambda_{\mathbb{T}2}$.

Solution.

$$\lambda \alpha, \beta : * . \lambda x : \alpha . \lambda y : \beta . x$$

$$\lambda \alpha : * . \lambda x : \alpha . \lambda \beta : * . \lambda y : \beta . x$$

First Form.

1. $\alpha, \beta : *$ **Bound**
2. $\boxed{x : \alpha}$ **Bound**
3. $\boxed{\boxed{y : \beta}}$ **Bound**
4. $\boxed{\boxed{\boxed{x : \alpha}}}$ **T-Var**
5. $\boxed{\boxed{\lambda y : \beta. x : \beta \rightarrow \alpha}}$ **4 T-Abst**
6. $\boxed{\lambda x : \alpha. \lambda y : \beta. x : \alpha \rightarrow \beta \rightarrow \alpha}$ **5 T-Abst**

7. $\lambda\alpha, \beta : * . \lambda x : \alpha . \lambda y : \beta . x : \Pi\alpha, \beta : * . \alpha \rightarrow \beta \rightarrow \alpha$ **5 T2-Abst**

■

Second Form.

1.	$\alpha : *$	Bound
2.	$x : \alpha$	Bound
3.	$\beta : *$	Bound
4.	$y : \beta$	Bound
5.	$x : \alpha$	T-Var
6.	$\lambda y : \beta . x : \beta \rightarrow \alpha$	5 T-Abst
7.	$\lambda\beta : * . \lambda y : \beta . x : \Pi\beta : * . \beta \rightarrow \alpha$	6 T2-Abst
8.	$\lambda x : \alpha . \lambda\beta : * . \lambda y : \beta . x : \alpha \rightarrow (\Pi\beta : * . \beta \rightarrow \alpha)$	7 T-Abst
9.	$\lambda\alpha : * . \lambda x : \alpha . \lambda\beta : * . \lambda y : \beta . x : \Pi\alpha : * . \alpha \rightarrow (\Pi\beta : * . \beta \rightarrow \alpha)$	8 T2-Abst

■

Problem

(3.9) Pretype the combinator

$$S \equiv \lambda x y z . x z (y z)$$

In closed form (typable in an empty context) in Λ_{T2} .

Solution.

$$S \equiv \lambda\alpha, \beta, \gamma : * . \lambda x : \alpha \rightarrow \beta \rightarrow \gamma . \lambda y : \alpha \rightarrow \beta . \lambda z : \alpha . x z (y z)$$

Proof.

1.	$\alpha, \beta, \gamma : *$	Bound
2.	$x : \alpha \rightarrow \beta \rightarrow \gamma$	Bound
3.	$y : \alpha \rightarrow \beta$	Bound
4.	$z : \alpha$	Bound
5.	$x z : \beta \rightarrow \gamma$	*,* T-App
6.	$y x : \beta$	*,* T-App
7.	$x z (y x) : \gamma$	5,6 T-App
8.	$\lambda z : \alpha . x z (y x) : \alpha \rightarrow \gamma$	7 T-Abst
9.	$\lambda y : \alpha \rightarrow \beta . \lambda z : \alpha . x z (y x) : (\alpha \rightarrow \beta) \rightarrow \alpha \rightarrow \gamma$	8 T-Abst

10.
$$\frac{\lambda x : \alpha \rightarrow \beta \rightarrow \gamma. \lambda y : \alpha \rightarrow \beta. \lambda z : \alpha. x z (y x)}{ : (\alpha \rightarrow \beta \rightarrow \gamma) \rightarrow (\alpha \rightarrow \beta) \rightarrow \alpha \rightarrow \gamma} \quad \mathbf{9\ T\text{-}Abst}$$
11.
$$\lambda \alpha, \beta, \gamma : *. \lambda x : \alpha \rightarrow \beta \rightarrow \gamma. \lambda y : \alpha \rightarrow \beta. \lambda z : \alpha. x z (y x)$$

$$: \Pi \alpha, \beta, \gamma : *. (\alpha \rightarrow \beta \rightarrow \gamma) \rightarrow (\alpha \rightarrow \beta) \rightarrow \alpha \rightarrow \gamma \quad \mathbf{10\ T2\text{-}Abst}$$

■

Problem

(3.10 a) Consider the term

$$M \equiv \lambda x : \Pi \alpha : *. \alpha \rightarrow \alpha. x (\sigma \rightarrow \sigma)(x \sigma)$$

Prove that M is legal.

Solution. For a term to be legal there must exist a context so that the term could be typed. Here, a witness context is $\Gamma \equiv \sigma : *$.

Proof.

1. $x : \Pi \alpha : *. \alpha \rightarrow \alpha$ **Bound**
2. $x (\sigma \rightarrow \sigma) : (\sigma \rightarrow \sigma) \rightarrow (\sigma \rightarrow \sigma)$ ***, T2-App**
3. $x \sigma : \sigma \rightarrow \sigma$ ***, T2-App**
4. $x (\sigma \rightarrow \sigma)(x \sigma) : \sigma \rightarrow \sigma$ **2,3 T-App**
5.
$$\lambda x : \Pi \alpha : *. \alpha \rightarrow \alpha. x (\sigma \rightarrow \sigma)(x \sigma)$$

$$: (\Pi \alpha : *. \alpha \rightarrow \alpha) \rightarrow \sigma \rightarrow \sigma \quad \mathbf{4\ T\text{-}Abst}$$

■

Problem

(3.10 b) Find a term N such that $M N$ is legal and may be considered to be a way to add type information to $(\lambda x . x x)(\lambda y . y)$

Solution.

$$M \sigma N \equiv (\lambda x : \Pi \alpha : *. \alpha \rightarrow \alpha. x (\sigma \rightarrow \sigma)(x \sigma))\sigma(\lambda y : \sigma. y)$$

Is the same as $(\lambda x . x x)(\lambda y . y)$ modulo type annotations.

Proof.

1. $M : (\Pi \alpha : * . \alpha \rightarrow \alpha) \rightarrow \sigma \rightarrow \sigma$ **T-Var**
2. $M \sigma : (\sigma \rightarrow \sigma) \rightarrow \sigma \rightarrow \sigma$ **1,* T2-App**
3. $y : \sigma$ **Bound**
4. $\boxed{y : \sigma}$ **T-Var**
5. have $N := \lambda y : \sigma . y : \sigma \rightarrow \sigma$ **4 T-Abst**
6. $M \sigma N : \sigma \rightarrow \sigma$ **2,5 T-Abst**

■

Problem

(3.11) Recall $\perp \equiv \Pi \alpha : * . \alpha$ from 3.5. Type and prove the following term legal:

$$\lambda x : \perp . x (\perp \rightarrow \perp \rightarrow \perp)(x (\perp \rightarrow \perp) x)(x (\perp \rightarrow \perp \rightarrow \perp) x x)$$

Solution.

Proof. The type $\perp \rightarrow \perp$ is closed and well formed. Therefore, the term is legal.

1. $\perp : * \equiv \Pi \alpha : * . \alpha$ **T-Form**
2. $x : \perp$ **Bound**
3. $\boxed{x (\perp \rightarrow \perp \rightarrow \perp) : \perp \rightarrow \perp \rightarrow \perp}$ ***,* T2-App**
4. $\boxed{x (\perp \rightarrow \perp) : \perp \rightarrow \perp}$ ***,* T2-App**
5. $\boxed{x (\perp \rightarrow \perp) x : \perp}$ **4,* T-App**
6. $\boxed{x (\perp \rightarrow \perp \rightarrow \perp)(x (\perp \rightarrow \perp) x) : \perp \rightarrow \perp}$ **3,5 T-App**
7. $\boxed{x (\perp \rightarrow \perp \rightarrow \perp) x : \perp \rightarrow \perp}$ **3,* T-App**
8. $\boxed{x (\perp \rightarrow \perp \rightarrow \perp) x x : \perp}$ **7,* T-App**
9. $\boxed{x (\perp \rightarrow \perp \rightarrow \perp)(x (\perp \rightarrow \perp) x)(x (\perp \rightarrow \perp \rightarrow \perp) x x) : \perp}$ **6,8 T-App**
10. $\lambda x : \perp . x (\perp \rightarrow \perp \rightarrow \perp)(x (\perp \rightarrow \perp) x)$
- $(x (\perp \rightarrow \perp \rightarrow \perp) x x) : \perp \rightarrow \perp$ **9 T-Abst**

■

Problem

(3.12) Given the Polymorphic Church Numerals:

$$\mathbf{nat} \in \mathbb{T}_2 := \Pi \alpha : *. (\alpha \rightarrow \alpha) \rightarrow \alpha \rightarrow \alpha$$

$$\bar{0} \equiv \lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. x : \mathbf{nat}$$

$$\bar{1} \equiv \lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. f x : \mathbf{nat}$$

$$\bar{2} \equiv \lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. f (f x) : \mathbf{nat}$$

$$\mathbf{succ} \equiv \lambda n : \mathbf{nat}. \lambda \beta : *. \lambda f : \beta \rightarrow \beta. \lambda x : \beta. f (n \beta f x)$$

Prove that

$$\mathbf{succ} \bar{0} \underset{\beta}{=} \bar{1}$$

$$\mathbf{succ} \bar{1} \underset{\beta}{=} \bar{2}$$

Solution.

$$\begin{aligned} & \mathbf{succ} \bar{0} \\ & \equiv (\lambda n : \mathbf{nat}. \lambda \beta : *. \lambda f : \beta \rightarrow \beta. \lambda x : \beta. f (n \beta f x)) (\lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. x) \\ & \xrightarrow[\beta]{} (\lambda \beta : *. \lambda f : \beta \rightarrow \beta. \lambda x : \beta. f ((\lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. x) \beta f x)) \\ & \xrightarrow[\beta_{T2}]{} (\lambda \beta : *. \lambda f : \beta \rightarrow \beta. \lambda x : \beta. f ((\lambda f : \beta \rightarrow \beta. \lambda x : \beta. x) f x)) \\ & \xrightarrow[\beta]{} (\lambda \beta : *. \lambda f : \beta \rightarrow \beta. \lambda x : \beta. f x) \underset{\alpha_{T2}}{\overset{\beta \rightarrow \alpha}{\equiv}} \bar{1} \\ & \mathbf{succ} \bar{1} \\ & \equiv (\lambda n : \mathbf{nat}. \lambda \beta : *. \lambda f : \beta \rightarrow \beta. \lambda x : \beta. f (n \beta f x)) (\lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. f x) \\ & \xrightarrow[\beta]{} (\lambda \beta : *. \lambda f : \beta \rightarrow \beta. \lambda x : \beta. f ((\lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. f x) \beta f x)) \\ & \xrightarrow[\beta_{T2}]{} (\lambda \beta : *. \lambda f : \beta \rightarrow \beta. \lambda x : \beta. f ((\lambda f : \beta \rightarrow \beta. \lambda x : \beta. f x) f x)) \\ & \xrightarrow[\beta]{} (\lambda \beta : *. \lambda f : \beta \rightarrow \beta. \lambda x : \beta. f (f x)) \underset{\alpha_{T2}}{\overset{\beta \rightarrow \alpha}{\equiv}} \bar{2} \end{aligned}$$

Problem

(3.13 a) We define addition in Polymorphic Church Numerals as

$$\text{add} \equiv \lambda m, n : \mathbf{nat}. \lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \mathbf{nat}. m \alpha f (n \alpha f x)$$

Show that

$$\text{add } \bar{1} \quad \bar{1} \underset{\beta}{=} \bar{2}$$

Solution.

$$\begin{aligned} & \text{add } \bar{1} \quad \bar{1} \\ & \equiv (\lambda m, n : \mathbf{nat}. \lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. m \alpha f (n \alpha f x)) \bar{1} \quad \bar{1} \\ & \xrightarrow[\beta]{\lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. \bar{1} \alpha f (\bar{1} \alpha f x)} \\ & \equiv \lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. \bar{1} \alpha f ((\lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. f x) \alpha f x) \\ & \xrightarrow[\beta]{\lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. \bar{1} \alpha f (f x)} \\ & \equiv \lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. (\lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. f x) \alpha f (f x) \\ & \xrightarrow[\beta]{\lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. f (f x)} \underset{\alpha}{=} \bar{2} \end{aligned}$$

Problem

(3.13 b) Find a term `mul` simulates multiplication on `nat`.

Solution.

$$\text{mul} \equiv \lambda m, n : \mathbf{nat}. \lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. m \alpha (n \alpha f) x$$

Proof. We derive the type first to prove a legal term.

1.	$m, n : \mathbf{nat}$	Bound
2.	$\alpha : *$	Bound
3.	$f : \alpha \rightarrow \alpha$	Bound
4.	$x : \alpha$	Bound
5.	$m \alpha : (\alpha \rightarrow \alpha) \rightarrow \alpha \rightarrow \alpha$	*, T2-App
6.	$n \alpha : (\alpha \rightarrow \alpha) \rightarrow \alpha \rightarrow \alpha$	*, T2-App
7.	$n \alpha f : \alpha \rightarrow \alpha$	6, T-App
8.	$m \alpha (n \alpha f) : \alpha \rightarrow \alpha$	5,7 T-App

9.	$\frac{}{m \alpha(n \alpha f) x : \alpha}$	8,* T-App
10.	$\frac{}{\lambda x : \alpha. m \alpha(n \alpha f) x : \alpha \rightarrow \alpha}$	9 T-Abst
11.	$\frac{}{\lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. m \alpha(n \alpha f) x : (\alpha \rightarrow \alpha) \rightarrow \alpha \rightarrow \alpha}$	10 T-Abst
12.	$\frac{}{\lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. m \alpha(n \alpha f) x : \mathbf{nat}}$	11 T2-Abst
13.		

$\lambda m, n : \mathbf{nat}. \lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. m \alpha(n \alpha f) x$
 $: \mathbf{nat} \rightarrow \mathbf{nat} \rightarrow \mathbf{nat}$

12 T-Abst

This proves that the term does indeed produce a natural number from two. Next let's prove that

$$\forall \bar{n}, \bar{m} : \mathbf{nat} \quad \text{mul } \bar{n} \ \bar{m} = \overline{n \times m}$$

It could be proven by induction that

$$\forall \bar{a} : \mathbf{nat} \quad \bar{a} \equiv \lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. f^a x$$

$$\begin{aligned} \text{mul } \bar{n} \ \bar{m} &\equiv (\lambda m, n : \mathbf{nat}. \lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. m \alpha(n \alpha f) x)(\bar{n})(\bar{m}) \\ &\xrightarrow[\beta]{} \lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. \bar{n} \alpha(\bar{m} \alpha f) x \\ &\xrightarrow[\beta]{} \lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. \bar{n} \alpha(\lambda u : \alpha. f^m u) x \\ &\xrightarrow[\beta]{} \lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. (\lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. f^n x)(\lambda u : \alpha. f^m u) x \\ &\xrightarrow[\beta]{} \lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. (\lambda u : \alpha. f^m u)^n x \end{aligned}$$

By induction this can be further beta-reduced to

$$\xrightarrow[\beta]{} \lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. f^{mn} x \equiv \overline{m n}$$

■

Problem

(3.14) We present the Church-Encoded Boolean:

$$\begin{aligned} \mathbf{bool} &\in \mathbb{T}_2 := \Pi \alpha : *. \alpha \rightarrow \alpha \rightarrow \alpha \\ \mathbf{true} &\equiv \lambda \alpha : *. \lambda x, y : \alpha. x \\ \mathbf{false} &\equiv \lambda \alpha : *. \lambda x, y : \alpha. y \end{aligned}$$

Construct a $\lambda 2$ term \mathbf{neg} such that $\mathbf{neg} \ \mathbf{true} \equiv_{\beta} \mathbf{false}$ and $\mathbf{neg} \ \mathbf{false} \equiv_{\beta} \mathbf{true}$.

Solution.

$$\mathbf{neg} \equiv \lambda b : \mathbf{bool}. \lambda \alpha : *. b \alpha(\mathbf{false} \ \alpha)(\mathbf{true} \ \alpha)$$

Neg True.

$$\begin{aligned}
 \text{neg true} &\equiv (\lambda b : \text{bool}. \lambda \alpha : * . b \alpha (\text{false } \alpha) (\text{true } \alpha)) \text{ true} \\
 &\xrightarrow[\beta]{\rightarrow} \lambda \alpha : * . (\lambda x, y : \alpha. x) (\text{false } \alpha) (\text{true } \alpha) \\
 &\xrightarrow[\beta]{\rightarrow} \lambda \alpha : * . \text{false } \alpha \xrightarrow[\eta]{\rightarrow} \text{false}
 \end{aligned}$$

■

Neg False.

$$\begin{aligned}
 \text{neg false} &\equiv (\lambda b : \text{bool}. \lambda \alpha : * . b \alpha (\text{false } \alpha) (\text{true } \alpha)) \text{ false} \\
 &\xrightarrow[\beta]{\rightarrow} \lambda \alpha : * . (\lambda x, y : \alpha. y) (\text{false } \alpha) (\text{true } \alpha) \\
 &\xrightarrow[\beta]{\rightarrow} \lambda \alpha : * . \text{true } \alpha \xrightarrow[\eta]{\rightarrow} \text{true}
 \end{aligned}$$

■

Problem

(3.15) Define

$$M \equiv \lambda u, v : \text{bool}. \lambda \beta : * . \lambda x, y : \beta. u \beta (v \beta x y) (v \beta y y)$$

And reduce $M \text{ true true}$, $M \text{ true false}$, $M \text{ false true}$, $M \text{ false false}$, and decide which logical operator is represented by M .

Solution.

$$\begin{aligned}
 M \text{ true true} &\equiv (\lambda u, v : \text{bool}. \lambda \beta : * . \lambda x, y : \beta. u \beta (v \beta x y) (v \beta y y)) \text{ true true} \\
 &\xrightarrow[\beta]{\rightarrow} \lambda \beta : * . \lambda x, y : \beta. \text{true } \beta (\text{true } \beta x y) (\text{true } \beta y y) \\
 &\xrightarrow[\beta]{\rightarrow} \lambda \beta : * . \lambda x, y : \beta. (\text{true } \beta x y) \\
 &\xrightarrow[\beta]{\rightarrow} \lambda \beta : * . \lambda x, y : \beta. x \equiv_{\alpha} \text{true}
 \end{aligned}$$

$$\begin{aligned}
 M \text{ true false} &\equiv (\lambda u, v : \text{bool}. \lambda \beta : * . \lambda x, y : \beta. u \beta (v \beta x y) (v \beta y y)) \text{ true false} \\
 &\xrightarrow[\beta]{\rightarrow} \lambda \beta : * . \lambda x, y : \beta. \text{true } \beta (\text{false } \beta x y) (\text{false } \beta y y) \\
 &\xrightarrow[\beta]{\rightarrow} \lambda \beta : * . \lambda x, y : \beta. (\text{false } \beta x y) \\
 &\xrightarrow[\beta]{\rightarrow} \lambda \beta : * . \lambda x, y : \beta. y \equiv_{\alpha} \text{false}
 \end{aligned}$$

$$\begin{aligned}
M \text{ false true} &\equiv (\lambda u, v : \text{bool}. \lambda \beta : *. \lambda x, y : \beta. u \beta (v \beta x y) (v \beta y y)) \text{ false true} \\
&\xrightarrow[\beta]{} \lambda \beta : *. \lambda x, y : \beta. \text{false} \beta (\text{true} \beta x y) (\text{true} \beta y y) \\
&\xrightarrow[\beta]{} \lambda \beta : *. \lambda x, y : \beta. (\text{true} \beta y y) \\
&\xrightarrow[\beta]{} \lambda \beta : *. \lambda x, y : \beta. y \equiv_{\alpha} \text{false}
\end{aligned}$$

$$\begin{aligned}
M \text{ false false} &\equiv (\lambda u, v : \text{bool}. \lambda \beta : *. \lambda x, y : \beta. u \beta (v \beta x y) (v \beta y y)) \text{ false false} \\
&\xrightarrow[\beta]{} \lambda \beta : *. \lambda x, y : \beta. \text{false} \beta (\text{false} \beta x y) (\text{false} \beta y y) \\
&\xrightarrow[\beta]{} \lambda \beta : *. \lambda x, y : \beta. (\text{false} \beta y y) \\
&\xrightarrow[\beta]{} \lambda \beta : *. \lambda x, y : \beta. y \equiv_{\alpha} \text{false}
\end{aligned}$$

Therefore M is equivalent to logical AND.

Problem

(3.16) Find $\lambda 2$ term representing the logical OR, XOR, IMP.

Solution.

$$\begin{aligned}
\text{OR} &\equiv \lambda u, v : \text{bool}. \lambda \beta : *. \lambda x, y : \beta. u \beta x (v \beta x y) \\
\text{XOR} &\equiv \lambda u, v : \text{bool}. \lambda \beta : *. \lambda x, y : \beta. u \beta (v \beta y x) (v \beta x y) \\
\text{IMP} &\equiv \lambda u, v : \text{bool}. \lambda \beta : *. \lambda x, y : \beta. u \beta (v \beta x y) x
\end{aligned}$$

All of them could be checked by finite enumeration over $\text{bool} \times \text{bool}$.

Problem

(3.17) Find $\text{isZero} : \text{nat} \rightarrow \text{bool}$ such that $\forall n : \text{nat}, \text{isZero } n \equiv_{\beta} \text{false}$ except when $n \equiv \bar{0}$.

Solution.

$$\text{isZero} \equiv \lambda n : \text{nat}. n \text{ bool } (\lambda u : \text{bool}. \text{false}) \text{ true}$$

Proof.

$$\begin{aligned}
\text{isZero } \bar{0} &\equiv (\lambda n : \mathbf{nat}. n \text{ bool } (\lambda u : \mathbf{bool}. \text{false}) \text{ true}) \bar{0} \\
&\xrightarrow[\beta]{} (\lambda \alpha : *. \lambda f : \alpha \rightarrow \alpha. \lambda x : \alpha. x) \text{ bool } (\lambda u : \mathbf{bool}. \text{false}) \text{ true} \\
&\xrightarrow[\beta]{} (\lambda f : \mathbf{bool} \rightarrow \mathbf{bool}. \lambda x : \mathbf{bool}. x) (\lambda u : \mathbf{bool}. \text{false}) \text{ true} \\
&\xrightarrow[\beta]{} \text{true}
\end{aligned}$$

■

By induction it could be proven that any other natural numbers must be applied $\lambda u : \mathbf{bool}. \text{false}$ to the body, making the result false, except for $\bar{0}$, where the function $f : \alpha \rightarrow \alpha$ never got applied.

Problem

(3.18 a) Define type

$$\mathbf{tree} \equiv \Pi \alpha : *. (\mathbf{bool} \rightarrow \alpha) \rightarrow (\mathbf{bool} \rightarrow \alpha \rightarrow \alpha \rightarrow \alpha) \rightarrow \alpha$$

Then we construct an inhabitant

$$\lambda \alpha. *. \lambda u : \mathbf{bool} \rightarrow \alpha. \lambda v : \mathbf{bool} \rightarrow \alpha \rightarrow \alpha \rightarrow \alpha. M$$

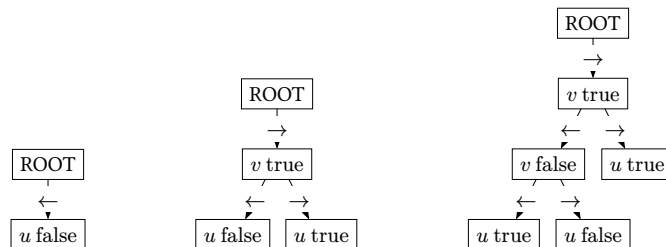
Is a node of a binary tree. Sketch graphs of trees where M is

$$\begin{aligned}
&u \text{ false} \\
&v \text{ true } (u \text{ false})(u \text{ true}) \\
&v \text{ true } (u \text{ true})(v \text{ false } (u \text{ true})(u \text{ false}))
\end{aligned}$$

Solution. A binary tree is usually defined as this:

inductive **Tree** (**a** : **Type**) where
 | **leaf** (**value** : **a**) : **Tree a**
 | **node** (**left right** : **Tree a**) : **Tree a**

With two constructors: a leaf or a node. Here α is the type of the payload at each node. There are two constructors: u is the left constructor, taking a \mathbf{bool} value for the direction of the node. The v term is the node constructor, taking a \mathbf{bool} as the direction, two α -typed terms as it's children.



Problem

(3.18 b) Give a $\lambda 2$ term, which, given input a polymorphic boolean p and two trees s and t , delivers the combined tree with p on top, left subtree s and right subtree t .

Solution.

$$\begin{aligned} \text{leaf} &:= \lambda p : \text{bool} . \lambda s, t : \text{tree} . \\ &\quad \lambda \alpha . \lambda u . \text{bool} \rightarrow \alpha . \lambda v : \text{bool} \rightarrow \alpha \rightarrow \alpha \rightarrow \alpha . \\ &\quad v p (s \alpha u v) (t \alpha u v) \end{aligned}$$

Proof. We suppose

$$\begin{aligned} s &\equiv \lambda \beta : * . \lambda u_s : \text{bool} \rightarrow \beta . \lambda v_s : \text{bool} \rightarrow \beta \rightarrow \beta \rightarrow \beta . S \\ t &\equiv \lambda \gamma : * . \lambda u_t : \text{bool} \rightarrow \gamma . \lambda v_t : \text{bool} \rightarrow \gamma \rightarrow \gamma \rightarrow \gamma . T \end{aligned}$$

We want

$$\text{leaf } p \ s \ t \underset{\beta}{=} \lambda \alpha : * . \lambda u : \text{bool} \rightarrow \alpha . \lambda v : \text{bool} \rightarrow \alpha \rightarrow \alpha \rightarrow \alpha . v \ p \ S \ T$$

For compactness we denote

$$\begin{aligned} \alpha : * \vdash \tau_{\text{mkleaf}} &\equiv \text{bool} \rightarrow \alpha \\ \alpha : * \vdash \tau_{\text{mknode}} &\equiv \text{bool} \rightarrow \alpha \rightarrow \alpha \rightarrow \alpha \end{aligned}$$

By beta reduction we have

$$\begin{aligned} \text{leaf } p \ s \ t &\equiv (\lambda p : \text{bool} . \lambda s, t : \text{tree} . \dots) \ p \ s \ t \\ &\underset{\beta}{\rightarrow} (\lambda \alpha : * . \lambda u . \tau_{\text{mkleaf}} . \lambda v : \tau_{\text{mknode}} . v \ p \ (s \alpha u v) (t \alpha u v)) \\ &\underset{\beta}{\rightarrow} (\lambda \alpha : * . \lambda u . \tau_{\text{mkleaf}} . \lambda v : \tau_{\text{mknode}} . v \ p \ (s \alpha u v) (t \alpha u v)) \\ &\underset{\beta}{\rightarrow} \lambda \alpha : * . \lambda u . \tau_{\text{mkleaf}} . \lambda v : \tau_{\text{mknode}} . v \ p \ S \ T \end{aligned}$$

■

Problem

(3.19) If $\Gamma \vdash L : \sigma$, then Γ is a valid $\lambda 2$ context.

Solution. The definition of “valid” here would be taken as relative to a judgement as the fact to be able to derive the judgement. Thus, it meant a complete inference path could be made using only statements and judgements derived from the context. Now proof by induction on inference rule that deducted $L : \sigma$.

Case 1 : T-Var. The premise is that Γ is a $\lambda 2$ context. ■

Case 2 : T-App. Therefore $L \equiv M N$ for some $M, N \in \Lambda_{\mathbb{T}_2}$. Therefore,

$$\Gamma \vdash M : \tau \rightarrow \sigma \quad \Gamma \vdash N : \tau$$

for some $\tau \in \mathbb{T}_2$. By the inductive hypothesis on any premise Γ is a valid λ_2 context. ■

Case 3 : T-Abst. Therefore $L \equiv \lambda x : \alpha. N$ such that

$$\Gamma, x : \alpha \vdash N : \beta$$

for some $\alpha, \beta \in \mathbb{T}_2$ such that $\sigma \equiv \alpha \rightarrow \beta$. By the inductive hypothesis, $\Gamma, x : \alpha$. By the recursive definition of λ_2 contexts, for some valid context Δ , $\forall n \in \text{dom } \Delta$, n could not depend on statement declared after n in the context. Therefore, no statement in Γ could depend on $x : \alpha$. Therefore, Γ is a valid context. ■

Case 4 : T-Form. The premise is that Γ is a valid λ_2 context. ■

Case 5 : T2-App. Therefore $L \equiv N B$ for some $N, B \in \mathbb{V}_2$ such that

$$\Gamma \vdash N : \Pi\alpha : *. \sigma \quad \Gamma \vdash B : *$$

By the inductive hypothesis on the any premise Γ is a valid λ_2 context. ■

Case 6 : T2-Abst. Therefore $L \equiv \lambda\alpha : *. M$ for some $M \in \Lambda_{\mathbb{T}_2}$ such that

$$\Gamma, \alpha : * \vdash M : \beta$$

and $\sigma \equiv \Pi\alpha : *. \beta$. By reasoning in *Case 3* no statement in the context could depend on any statement before the latter's declaration. Therefore no statement in Γ could depend on $\alpha : *$, making it a valid context. ■

Problem

(3.20) Prove the free variable lemma for λ_2 .

$$\Gamma \vdash L : \sigma \Rightarrow \text{FV } L \subseteq \text{dom } \Gamma$$

Solution. Proof by induction on inference rules that deduced $L : \sigma$. The only rule not considered is T-Form since all terms apparent in \mathbb{T}_2 .

Case 1 : T-Var. Therefore L is the only free variable in L . By the generation lemma $L : \sigma \in \Gamma$, so $\{L\} \subseteq \text{dom } \Gamma$ ■

Case 2 : T-App. Therefore by the generation lemma $L \equiv M N$ for some $M, N \in \Lambda_{\mathbb{T}_2}$ such that

$$\Gamma \vdash M : \tau \rightarrow \sigma \quad \Gamma \vdash N : \tau$$

For some type τ . By the inductive hypothesis $\text{FV } M \subseteq \Gamma$ and $\text{FV } N \subseteq \text{dom } \Gamma$. Therefore $\text{FV } L = (\text{FV } M) \cup (\text{FV } N) \subseteq \text{dom } \Gamma$. ■

Case 3 : T-Abst. Therefore by the generation lemma $L \equiv \lambda x : \alpha. M$ for some $M \in \Lambda_{\mathbb{T}_2}$ such that

$$\Gamma, x : \alpha \vdash M : \beta$$

and $\sigma \equiv \alpha \rightarrow \beta$. By the inductive hypothesis $\text{FV } M \subseteq \text{dom } \Gamma \cup \{x\}$. Therefore

$$\text{FV } L = \text{FV } M \setminus \{x\} \subseteq (\text{dom } \Gamma \cup \{x\}) \setminus \{x\} = \text{dom } \Gamma$$

■

Case 4 : T2-App. Therefore $L \equiv B$ for some $N, B \in \mathbb{V}_2$ such that

$$\Gamma \vdash N : \Pi\alpha : * . \sigma \quad \Gamma \vdash B : *$$

By the inductive hypothesis $\text{FV } N \subseteq \text{dom } \Gamma$. Since $B \in \mathbb{T}_2$ then $\text{FV } B = \emptyset$. Therefore $\text{FV } L = \text{FV } N \subseteq \text{dom } \Gamma$.

■

Case 5 : T2-Abst. Therefore $L \equiv \lambda\alpha : * . M$ for some $M \in \Lambda_{\mathbb{T}_2}$ such that

$$\Gamma, \alpha : * \vdash M : \beta$$

and $\sigma \equiv \Pi\alpha : * . \beta$. Because $\alpha \in \mathbb{T}_2$ so $\alpha \notin \text{FV } M$ and $\text{FV } L = \text{FV } M$, thus $\text{FV } M = \text{FV } L \subseteq \Gamma$.

■

Problem

(3.21) Give a recursive definition for $\text{FTV} : \mathbb{T}_2 \cup \Lambda_{\mathbb{T}_2} \rightarrow \mathbb{V}_2$

Solution. Here $\alpha \in \mathbb{V}_2$, $A, B \in \mathbb{T}_2$, $x \in V$ and $M \in \Lambda_{\mathbb{T}_2}$.

Form

$\text{FTV } \alpha$

$\text{FTV } x$

$\text{FTV } (A \rightarrow B)$

$\text{FTV } (\Pi\alpha : * . A)$

$\text{FTV } (M N)$

$\text{FTV } (\lambda x : A . M)$

$\text{FTV } (M B)$

$\text{FTV } (\lambda\alpha : * . M)$

Value

$\{\alpha\}$

\emptyset

$\text{FTV}(A) \cup \text{FTV}(B)$

$\text{FTV}(A) \setminus \{\alpha\}$

$\text{FTV}(M) \cup \text{FTV}(N)$

$\text{FTV}(A) \cup \text{FTV}(M)$

$\text{FV}(M) \cup \text{FTV}(B)$

$\text{FTV}(M) \setminus \{\alpha\}$

—

Completed Dec 16 10:11 pm.