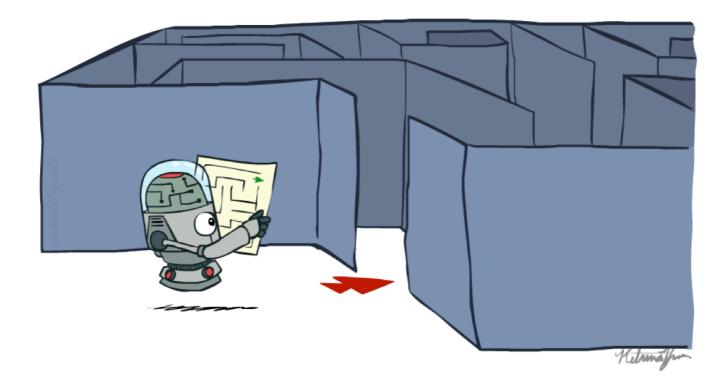
Announcements

- Project 0: Python Tutorial
 - Due today at 11:59pm
- * HW1 will be released today
- * P1 will be released on Wednesday
- * Survey about flipped classroom

Ve492: Introduction to Artificial Intelligence

Search Problem and Uninformed Search



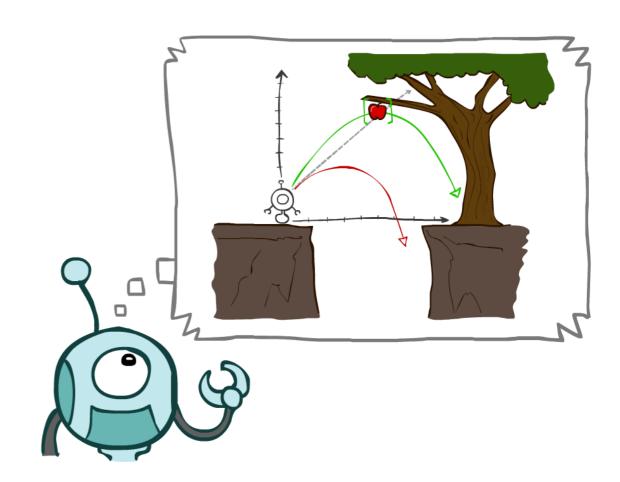
Paul Weng

UM-SJTU Joint Institute

Slides adapted from http://ai.berkeley.edu, AIMA, UM, CMU

Outline

- * Search Problems
- Uninformed Search Methods
 - Depth-First Search
 - * Breadth-First Search
 - * Iterative Deepening Search
 - Uniform-Cost Search



Search Problems

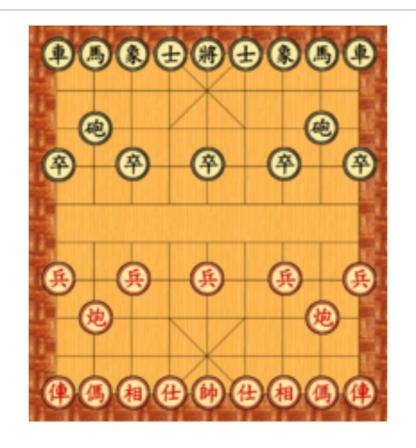


Solving Peg Solitaire is a Search Problem?



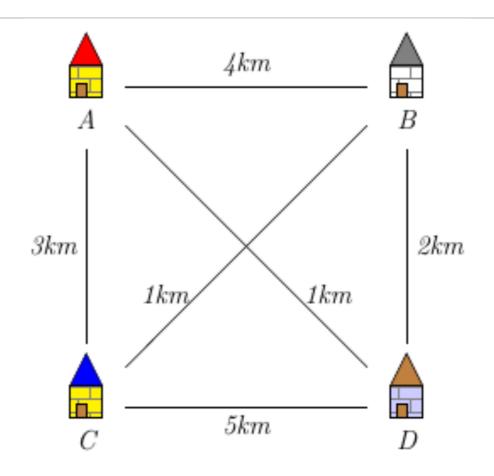
- * True or False?
- * If True, what are the states, actions, initial state, goal test, and costs?
- If True, what is an upper-bound on the size of the state space?
- * If True, what is an upper-bound on the number of actions?

Solving Chinese Chess is a Search Problem?



- * True or False?
- * If True, what are the states, actions, initial state, goal test, and costs?
- * If True, what is upper-bound on the size of the state space?
- * If True, what is an upper-bound on the number of actions?

Traveling Salesman Problems are Search Problems?

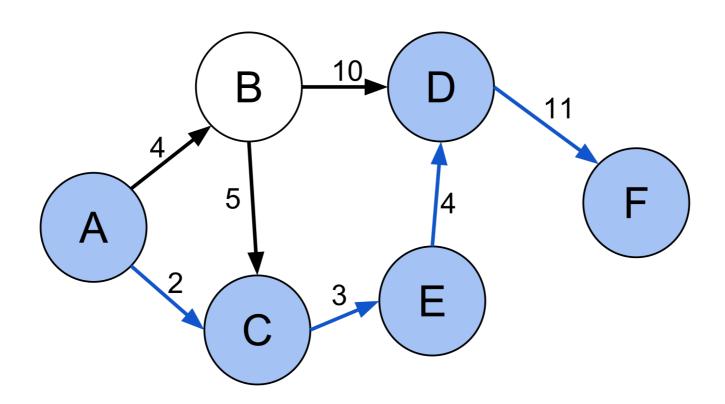


- * True or False?
- * If True, what are the states, actions, initial state, goal test, and costs?
- * If True, what is upper-bound on the size of the state space?
- * If True, what is an upper-bound on the number of actions?

Environment Type for a Search Problem

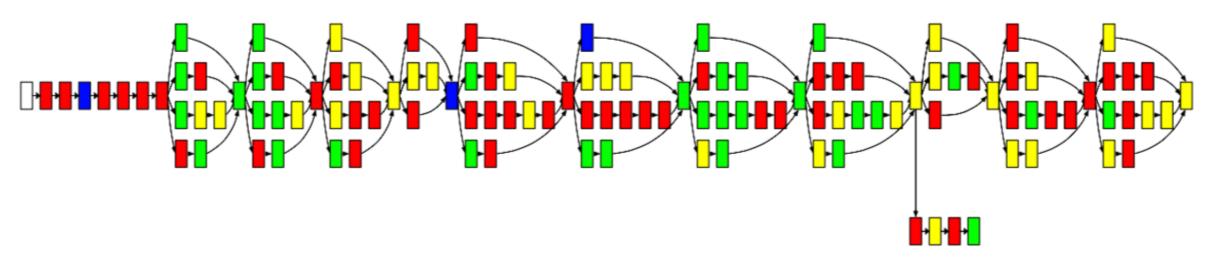
	Search Problem
Fully or partially observable	
Single agent or multi-agent	
Deterministic or non-deterministic	
Static or dynamic	
Discrete or continuous	
Episodic or sequential	

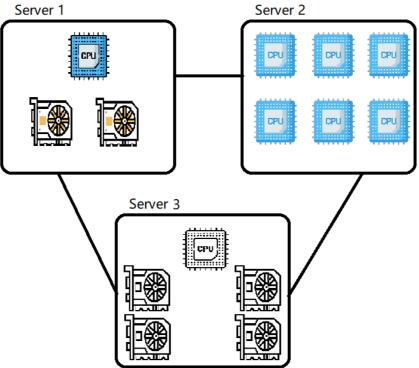
Search Problems are Generalization of Shortest Path Problems



- * Shortest Path Problems are in class P.
- However, state space is generally huge.

Example: Scheduling Problem

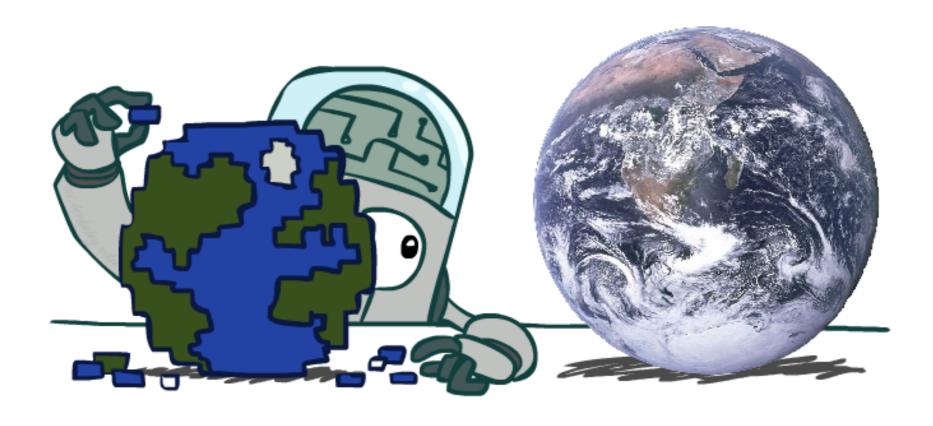




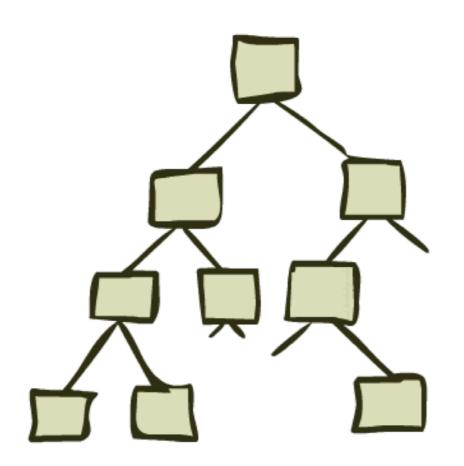
Example: Electronic Design Automation



Search Problems Are Models



State Space Graphs and Search Trees

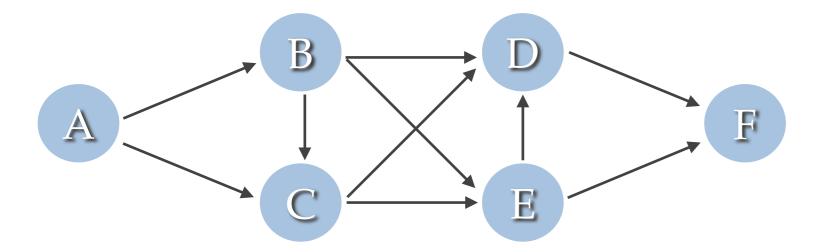


True or False: Search Tree

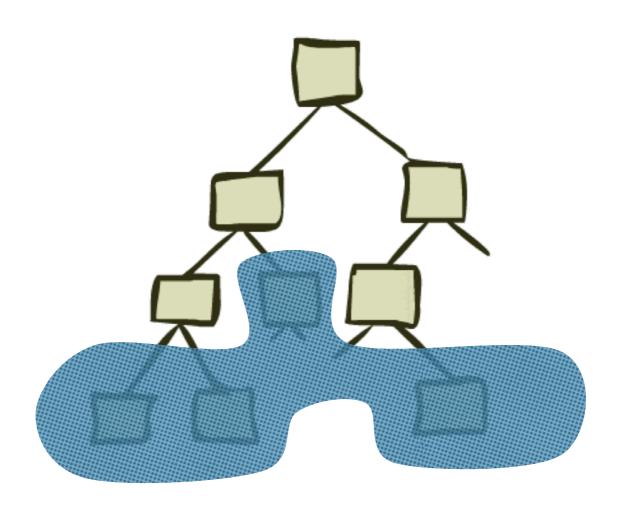
- * A search problem implicitly defines a graph.
- * A graph implicitly defines a search tree.
- * A search problem implicitly defines a search tree.
- * A node of a search tree corresponds to a state in the search problem.
- * A node of a search tree corresponds to a subpath in the graph.
- * A search tree is a compact way to represent a set of path.
- * If a graph has a finite number of states, the search tree has a finite number of nodes.

Search Tree

* Draw the search tree corresponding to this graph where the initial state is A and the goal test is F:



Tree Search



True or False: General Tree Search

- * The fringe corresponds to the nodes that have already been visited.
- * Expanding a node means visiting all its neighbors.
- * A state that is visited cannot be visited again.

General Tree Search

```
function TREE-SEARCH( problem, strategy) returns a solution, or failure initialize the search tree using the initial state of problem loop do

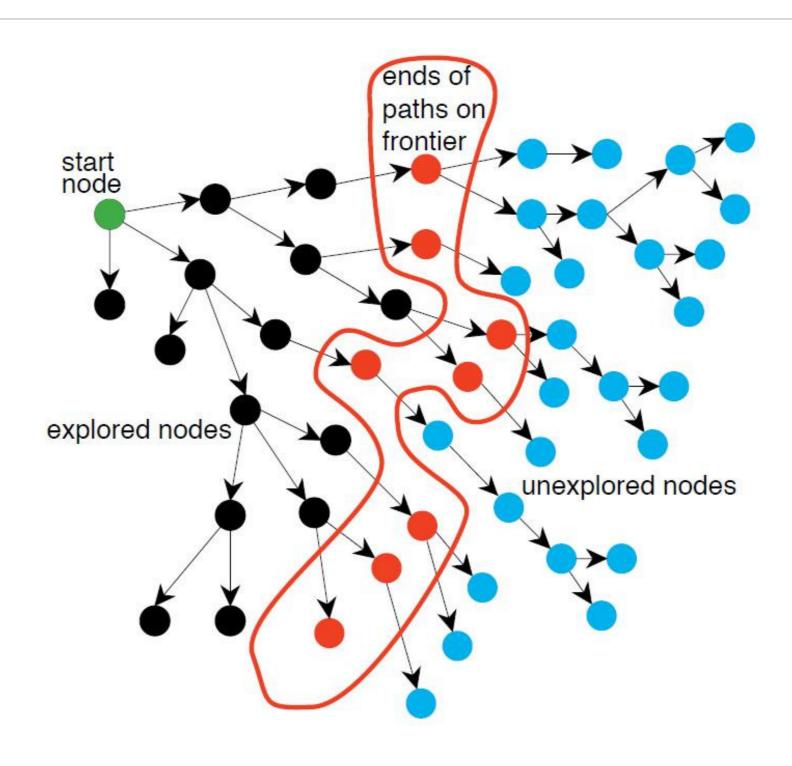
if there are no candidates for expansion then return failure choose a leaf node for expansion according to strategy

if the node contains a goal state then return the corresponding solution else expand the node and add the resulting nodes to the search tree end
```

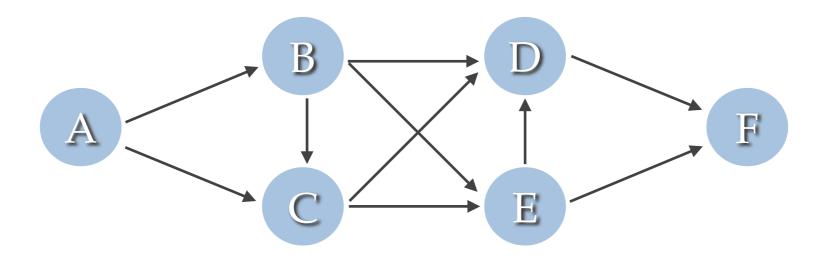
Important ideas:

- * Fringe
- * Expansion
- Exploration strategy
- * Main question: which fringe nodes to explore?

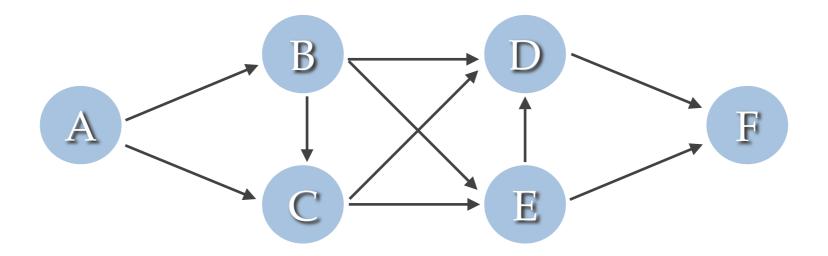
Search



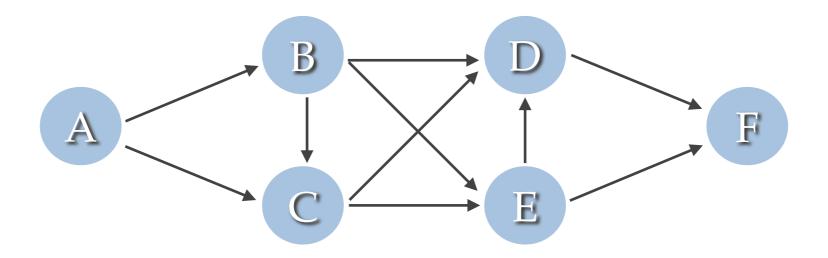
Run Depth First Search



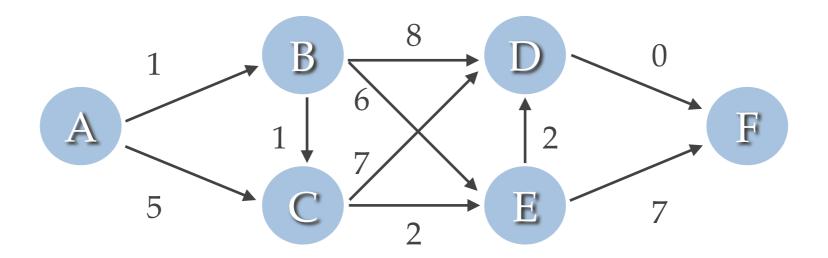
Run Breadth First Search



Run Iterative Deepening Search



Run Uniform Cost Search

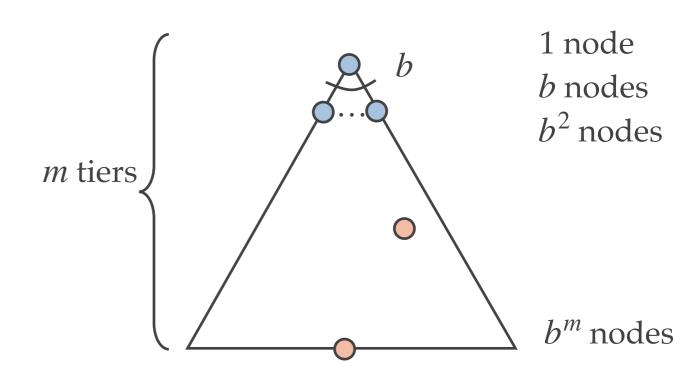


True or False: DFS, BFS, UCS

- * Consider a search problem where for every action, the cost is equal to ϵ , with ϵ >0.
 - Depth-first search is guaranteed to return an optimal solution.
 - Breadth-first search is guaranteed to return an optimal solution.
 - Uniform-cost search is guaranteed to return an optimal solution.
- * Consider a search problem where for every action, the cost is at least ϵ , with ϵ >0.
 - Depth-first search is guaranteed to return an optimal solution.
 - Breadth-first search is guaranteed to return an optimal solution.
 - * Uniform-cost search is guaranteed to return an optimal solution.

Search Algorithm Properties

- Complete: Guaranteed to find a solution if one exists?
- * Optimal: Guaranteed to find the least cost path?
- * Time complexity?
- * Space complexity?
- Cartoon of search tree:
 - * *b* is the branching factor
 - * *m* is the maximum depth
 - solutions at various depths



Number of nodes in entire tree?

$$* 1 + b + b^2 + ... + b^m = O(b^m)$$

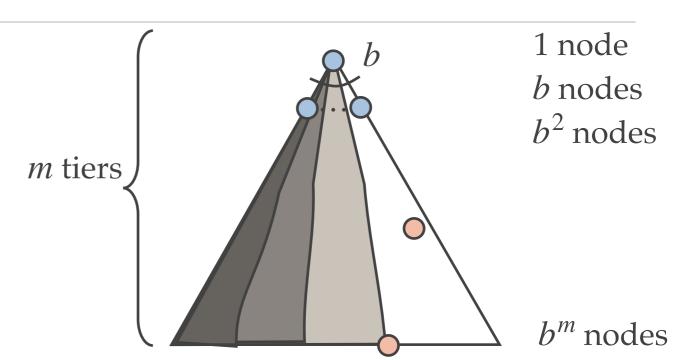
True or False: Complexity

- * The time complexity of a search algorithm depends on the number of nodes that are visited.
- * The space complexity of a search algorithm depends on the number of nodes that are visited.
- * The time complexity of IDS is the same as DFS.
- * The space complexity of IDS is the same as BFS.
- * The ratio of the time complexity of IDS and BFS tends to 1 when s tends to infinity.

Depth-First Search (DFS) Properties

* What nodes DFS expand?

- Some left prefix of the tree.
- Could process the whole tree!
- * If m is finite, takes time $O(b^m)$



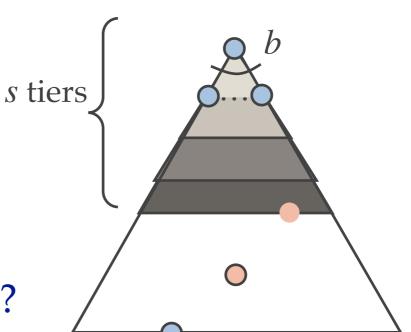
Now much space does the fringe take?

- * Only has siblings on path to root, so O(bm)
- * Is it complete?
 - * *m* could be infinite, so only if we prevent cycles (more later)
- * Is it optimal?
 - No, it finds the "leftmost" solution, regardless of depth or cost

Breadth-First Search (BFS) Properties

* What nodes does BFS expand?

- Processes all nodes above shallowest solution
- * Let depth of shallowest solution be *s*
- * Search takes time $O(b^s)$
- * How much space does the fringe take?
 - * Has roughly the last tier, so (b^s)
- * Is it complete?
 - * *s* must be finite if a solution exists, so yes!
- * Is it optimal?
 - * Only if costs are all 1 (more on costs later)



1 node b nodes b^2 nodes

 b^m nodes

Iterative Deepening

- Idea: get DFS's space advantage with BFS's time / shallow-solution advantages
 - * Run a DFS with depth limit 1. If no solution...
 - * Run a DFS with depth limit 2. If no solution...
 - * Run a DFS with depth limit 3.
- Isn't that wastefully redundant?
 - Generally most work happens in the lowest level searched, so not so bad!

Uniform Cost Search (UCS) Properties

* What nodes does UCS expand?

- Processes all nodes with cost less than cheapest solution!
- * If that solution costs C^* and arcs cost at least ε , then the "effective depth" is roughly C^*/ε
- * Takes time $O(b^{\frac{C^*}{\varepsilon}})$ (exponential in effective depth)
- * How much space does the fringe take?
 - * Has roughly the last tier, so $O(b^{\frac{C^*}{\varepsilon}})$
- * Is it complete?
 - Assuming best solution has a finite cost and minimum arc cost is positive, yes!
- * Is it optimal?
 - Yes! (Proof next lecture via A*)

