Démontrer
$$\vdash \neg(A \lor B) \Rightarrow (\neg A \land \neg B)$$





Démontrer $\vdash \neg (A \lor B) \Rightarrow (\neg A \land \neg B)$

$$\frac{\overline{A \vdash A} \stackrel{Ax}{A}}{A \vdash A \lor B} \stackrel{\bigvee_{intro}}{\bigvee_{intro}} \frac{\neg (A \lor B) \vdash \neg (A \lor B)}{\neg (A \lor B) \vdash \neg (A \lor B)} \stackrel{Ax}{\nearrow_{elim}} \frac{\overline{B \vdash B}}{B \vdash A \lor B} \stackrel{Ax}{\bigvee_{intro}} \frac{\neg (A \lor B) \vdash \neg (A \lor B)}{\neg (A \lor B) \vdash \neg (A \lor B)} \stackrel{Ax}{\nearrow_{elim}} \frac{\neg (A \lor B) \vdash \neg (A \lor B)}{\neg (A \lor B) \vdash \neg (A \lor B)} \stackrel{\neg (A \lor B) \vdash \neg (A \lor B)}{\neg (A \lor B) \vdash \neg (A \lor B)} \stackrel{\neg (A \lor B) \vdash (\neg A \land \neg B)}{\neg (A \lor B) \Rightarrow_{intro}} \Rightarrow_{intro}$$





Démontrer $\vdash \neg(A \land B) \Rightarrow (\neg A \lor \neg B)$





Démontrer $\vdash \neg (A \land B) \Rightarrow (\neg A \lor \neg B)$

$$\frac{\neg A \vdash \neg A}{\neg A \vdash \neg A \lor \neg B} \xrightarrow{\lor_{intro}^{g}} \frac{}{\neg (\neg A \lor \neg B) \vdash \neg (\neg A \lor \neg B)} \xrightarrow{Ax} \frac{}{\neg B \vdash \neg B} \xrightarrow{Ax} \xrightarrow{}_{\neg B \vdash \neg B} \xrightarrow{Ax} \xrightarrow{}_{\neg B \vdash \neg A \lor \neg B} \lor_{intro}^{d} \xrightarrow{}_{\neg (\neg A \lor \neg B) \vdash \neg (\neg A \lor \neg A)} \xrightarrow{}_{\neg (\neg A \lor \neg B) \vdash \neg (\neg A \lor \neg A)} \xrightarrow{}_{\neg (\neg A \lor \neg B) \vdash A} \xrightarrow{}_$$

$$\frac{\vdots}{\neg (\neg A \lor \neg B) \vdash A \land B} \frac{\neg (A \land B) \vdash \neg (A \land B)}{\neg (A \land B) \vdash \neg (A \land B)} \frac{Ax}{\neg elim} \\ \frac{\neg (A \land B), \neg (\neg A \lor \neg B) \vdash \bot}{\neg (A \land B) \vdash (\neg A \lor \neg B)} \xrightarrow{\Rightarrow intro} \bot Abs$$





Règles de base

- Avantage : Avoir choisi un ensemble minimal de règles y sera un atout important.
- Inconvénient : Rendre les démonstrations plus longues.
- Il est donc nécessaire d'introduire des *utilitaires* qui permettront de raccourcir les démonstrations.

Dans cette section, on démontre, à l'aide des règles de bases, des *règles dérivées*.







Règle de coupure

$$\frac{\Gamma, A \vdash B \qquad \Gamma^{'} \vdash A}{\Gamma, \Gamma^{'} \vdash B} \ {}^{\textit{coupure}}$$







Règle de coupure

$$\frac{\Gamma, A \vdash B \qquad \Gamma^{'} \vdash A}{\Gamma, \Gamma^{'} \vdash B} \stackrel{\textit{coupure}}{}$$

$$\frac{\Gamma, A \vdash B}{\Gamma \vdash A \Rightarrow B} \xrightarrow{\Rightarrow_{intro}} \Gamma' \vdash A \xrightarrow{\Rightarrow_{elim}} \Gamma, \Gamma' \vdash B$$







Introduction gauche de la conjonction

$$\frac{\Gamma, A, B \vdash C}{\Gamma, A \land B \vdash C} \, {}^{\land_g}$$







Introduction gauche de la conjonction

$$\frac{\Gamma, A, B \vdash C}{\Gamma, A \land B \vdash C} \stackrel{\wedge_g}{\longrightarrow}$$

$$\begin{array}{c|c} \hline \Gamma, A, B \vdash C & \hline \hline A \land B \vdash A \land B & \land A \land A \land B \vdash A \land B \land A \land B \vdash A \land B \land A \land B \vdash A \land B \land B \land A \land B \vdash A \land B \land B \vdash A \land B \vdash A$$







Introduction gauche de la disjonction

$$\frac{\Gamma, A \vdash C \qquad \Gamma^{'}, B \vdash C}{\Gamma, \Gamma^{'}, A \lor B \vdash C} \lor_{\mathbf{g}}$$







Introduction gauche de la disjonction

$$\frac{\Gamma, A \vdash C \qquad \Gamma^{'}, B \vdash C}{\Gamma, \Gamma^{'}, A \lor B \vdash C} \, {}_{\mathbf{V_g}}$$







Introduction gauche de l'implication

$$\frac{\Gamma, A, B \vdash C}{\Gamma, A, A \Rightarrow B \vdash C} \Rightarrow_{g}$$







Introduction gauche de l'implication

$$\frac{\Gamma, A, B \vdash C}{\Gamma, A, A \Rightarrow B \vdash C} \Rightarrow_{g}$$







Introduction gauche de la négation

$$\Gamma, A, \neg A \vdash \bot$$







Introduction gauche de la négation

$$\Gamma, A, \neg A \vdash \bot$$







Loi de Peirce

$$\frac{\Gamma, \neg A \vdash A}{\Gamma \vdash A} \underset{\textbf{l.p.}}{\underline{\textbf{l.p.}}}$$







A Loi de Peirce

$$\frac{\Gamma, \neg A \vdash A}{\Gamma \vdash A} \, {}_{l.p.}$$

$$\frac{ \begin{array}{c|c} \hline \Gamma, \neg A \vdash \neg A & Ax & \hline \Gamma, \neg A \vdash A \\ \hline \hline \begin{array}{c} \hline \Gamma, \neg A \vdash \bot \\ \hline \hline \begin{array}{c} \Gamma, \neg A \vdash \bot \\ \hline \hline \end{array} \\ \hline \end{array} } \begin{array}{c} Premisse \\ \hline \begin{array}{c} elim \\ \hline \end{array} }$$







Tiers exclu

$$\frac{\Gamma, A \vdash B \qquad \Gamma^{'}, \neg A \vdash B}{\Gamma, \Gamma^{'} \vdash B} \text{ $t.e.$}$$







Tiers exclu

$$\frac{\Gamma, A \vdash B \qquad \Gamma', \neg A \vdash B}{\Gamma, \Gamma' \vdash B} \text{ t.e.}$$

Preuve:

$$\frac{A \vdash A \qquad Ax}{A \vdash A \lor \neg A} \bigvee_{intro}^{g} \qquad \frac{Ax}{\neg (A \lor \neg A) \vdash \neg (A \lor \neg A)} \qquad Ax}{\neg (A \lor \neg A) \vdash \bot} \\
\frac{A, \neg (A \lor \neg A) \vdash \bot}{\neg (A \lor \neg A) \vdash \neg A} \gamma_{intro} \qquad \gamma_{dintro} \\
\frac{\neg (A \lor \neg A) \vdash A \lor \neg A}{\vdash A \lor \neg A} \qquad l.p.$$



M. Ledmi



Contraposition

$$\frac{\Gamma \vdash \neg B \Rightarrow \neg A}{\Gamma \vdash A \Rightarrow B} \overset{c.c.}{\sim}$$







Contraposition

$$\frac{\Gamma \vdash \neg B \Rightarrow \neg A}{\Gamma \vdash A \Rightarrow B} \stackrel{c.c.}{\leftarrow}$$







Lois de Morgan

$$\frac{\Gamma, \neg A \vee \neg B \vdash C}{\Gamma, \neg (A \wedge B) \vdash C} \stackrel{\wedge_{m}}{\longrightarrow} \frac{\Gamma, \neg A, \neg B \vdash C}{\Gamma, \neg (A \vee B) \vdash C} \stackrel{\vee_{m}}{\longrightarrow}$$

$$\Gamma, A, \neg B \vdash C$$

$$\frac{\Gamma, A, \neg B \vdash C}{\Gamma, \neg (A \Rightarrow B) \vdash C} \Rightarrow_{m}$$







Lois de Morgan

$$\frac{\Gamma, \neg A \vee \neg B \vdash C}{\Gamma, \neg (A \wedge B) \vdash C} \stackrel{\wedge_{m}}{\longrightarrow} \frac{\Gamma, \neg A, \neg B \vdash C}{\Gamma, \neg (A \vee B) \vdash C} \stackrel{\vee_{m}}{\longrightarrow} \frac{\Gamma, A, \neg B \vdash C}{\Gamma, \neg (A \Rightarrow B) \vdash C} \stackrel{\Rightarrow_{m}}{\longrightarrow}$$

Preuve:

Laissées en TD



