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1 Software Architecture

The solution is divided into four main components:

- The client responsible for interfacing with the server or other methods of data input
- The pathfinding module handling graph search and pathing between two points in any space
- The **simulation** module handling the route selection and the actual "drone control" with the use of pathfinding
- The **visualisation** module responsible for generating visualisations of the output from the simulation module

Most of the inter-module design choices explained below, are fueled by the dependency inversion principle and the rest by good general OOP practice. Due to the size of the solution, a lot of the simpler classes are missed out of the diagrams. The descriptions below will revolve mostly around the key decisions in the architecture and their benefits.

1.1 Client module

Figure 1 shows the structure of the client module classes, it contains Data classes which hold the intermediate, validated and standardised data. These intermediate classes detach the implementation from the exact shape input data. Since any number of formats can map to these classes, the solution is not closely related to the input data this is very desirable.

The ClientService interface is used to actually create the data in the form of Data classes from the input data given. In the case of AQmaps the data is given via http server and so an appropriate implementing class will need to know the base URI/location of the API where sensor/map/w3w data is stored, and have access to the Http protocol.

Any other classes which will "feed" into the rest of the system, but are constructed from input data, will be declared in this module and implement the appropriate external service interfaces.

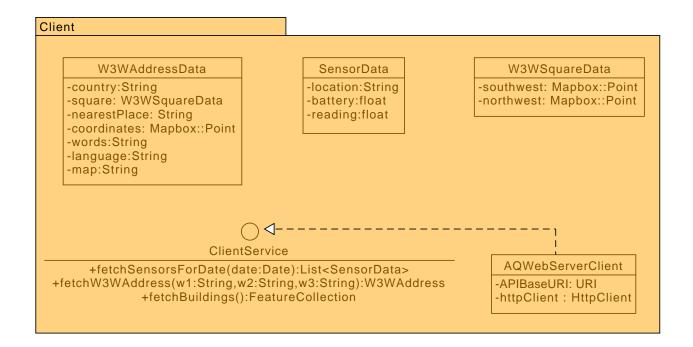


Figure 1: UML diagram of the client module

1.2 Pathfinding module

We separate the problem of pathfinding completely from the AQmap problem. This allows a broader range of techniques to be applied much more easilly.

Figure 3 shows the most important classes from the pathfinding module. The basis of this module is formed by the SearchNode and PathfindingAlgorithm abstract classes as well as the Graph interface.

The PathfindingAlgorithm class defines an abstract findPath method for finding a path to a single goal, leaving the exact way this is done to the concrete implementation. It is marked as abstract since we will always need to apply pathfinding to multiple goals as well, and so this class should deal with that, since this would otherwise always have to be done by the consumer.

Each SearchNode is used to represent a part of the frontier of a graph search algorithm i.e. a **state** in the search space - in the pathfinding case, a path. These are made abstract to allow for the use of generics in such a way that each parent node is of the specific node type needed and also because nodes are expected to be annotated with problem specific data. It is this generic parameter which propagates to the rest of the classes.

Due to the fact we are mostly dealing with multiple goal pathfinding, each node should also hold a deque of goals which can be reached from it - take as an example the problem where you start at a position with 2 goals, no movement is necessary and both are reached.

Pathfinding goals are represented by their own class, alternatively we could simply use coordinates as goals, but this would force the consumer of this functionality to always have to match the coordinates of the path returned to the goals that it needed achieved, as opposed to simply looking through the goals achieved by each node and comparing them via reference to the deque of goals in order.

Finally the Graph interface defines the transition function from any searchNode to all its neighbours. PathfindingAlgorithms always accept the graph as a parameter as it forms the domain of each specific problem. Separating the domain from the algorithm completely makes the algorithms much more flexible and easier to unit test.

Other interfaces shown include the SpatialHash and PathfindingHeuristic, those will be used by implementing algorithms and are defined as interfaces to again further separate the components of algorithms from the domain. This allows us to further change the behaviour and performance of algorithms, by swapping out their components to suit (Composition over Inheritance).

1.3 Simulation module

The simulation module is responsible for: planning the route (using TSP solvers), applying pathfinding to find the detailed path required for collection, and setting the read status in sensors. The name stems from the fact that our sensor data collection is only simulated. This module directly interfaces with the pathfinding module as expected. Principles of dependency inverersion were applied to reduce coupling and increase flexibility.

Figure 4 shows the most important classes of the simulation module. The biggest architectural choice here is the splitting of the drone into components.

This means that we keep the specific pathfinding and route planning behaviour separate from the behaviour of the drone/data collector itself. Since it is not the job of the path or route planner to "read" any sensors, this will be the task of the drone in addition to perhaps applying some strategy to the produced path in case it is not satysfying enough (maybe re-routing from a point with a different collection order).

Another big choice is the fact that this module does not know anything about the input data, it simply defines a contract for the sensor class, letting the consumers of this module deal with data conversion.

The path planner component itself is further composed of the PathfindingAlgorithm and DistanceMatrix which means it only concerns itself with the task of translating a path of individual points it receives from the PathfindingAlgorithm to a path composed of path segments enforcing the move pattern of the drone.

The PathSegment class is necessary since it is a major requirement that the collector must move in a specific pattern, this class can be used to enforce such a contract.

On the other hand CollectionOrderPlanners are only defined by their implementation and the choice of DistanceMatrix as well as the set of route optimisers. This choice decouples the major TSP solving strategy from the distance measure that it uses.

Both the PathPlanner and CollectionOrderPlanner interfaces have a base abstract class implementation since:

- Collection order planners will always apply the set of given optimisers to the final route.
- They will also have to always set up the distance matrix with the given sensors
- Path planners will always have to apply the given pathfinding algorithm before actually performing conversion to path segments
- We also use these abstract bases to enforce problem specific rules, such as the maximum number of moves, this leaves the possibility of straying away from those base classes if the problem changes drastically in the future.

This module also implements a specific Graph and SearchNode specific to the problem at hand. Here the ConstrainedTreeGraph produces nodes satisfying the angle, move length and obstacle constraints, in effect all logic to do with checking whether the collector is hitting an obstacle, or if the move length is valid etc.. is contained in the graph itself. This is very desirable and intentional. The addition of the direction field to the DirectedSearchNode further emphasizes the angle requirements, and allows for the construction of PathSegments from search nodes without calculating angles twice.

Another important class here is the abstract Sensor class. This is made abstract for two reasons:

• The concrete implementation of the sensor should be done in another module, since it is more convenient for example to closely link it to the input data for convenience. So according to the principle of dependency inversion we do not provide a concrete implementation in this module.

• The sensor cannot show any reading unless it is actually read, this can be enforced by this abstact class using the setHasBeenRead method regardless of the concrete implementation/constructor.

1.4 Visualisation module

This module only needs to interface with the output classes of the simulation module, i.e. the PathSegment and Sensor classes but other than that, it is completely decoupled from the shape of the initial input data of the system. This is very desirable

Figure 2 shows the UML diagram of the most important classes in the visualisation module.

Classes implementing the SensorCollectionVisualiser interface generate geojson visualisations of sensor data collections. These have access to both the flight path and the sensors and hence their readings. The AQMapGenerator is the problem specific example of implementation of such a visualiser.

Using an interface like like this allows for the swapping out of visualisers at will whenever new requirements arise or current ones change.

The OutputFormatter class deals with writing the visualisations and flight paths to a file, these are made static as changes in output format are assumed to be very few in the future, and should such changes be required, new methods can be added to the formatter.

The usage of the AttributeMap interface allows for a lot of flexibility in the way the AQMapGenerator assigns colors and symbols, and should this behaviour need to be changed, it'd be very easy to do.

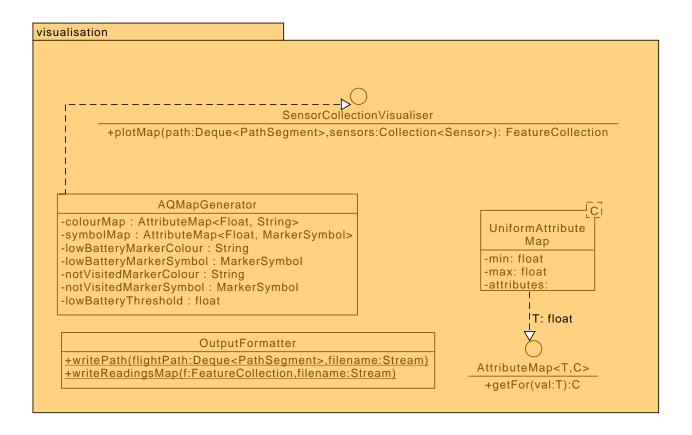


Figure 2: UML diagram of the Visualisation module

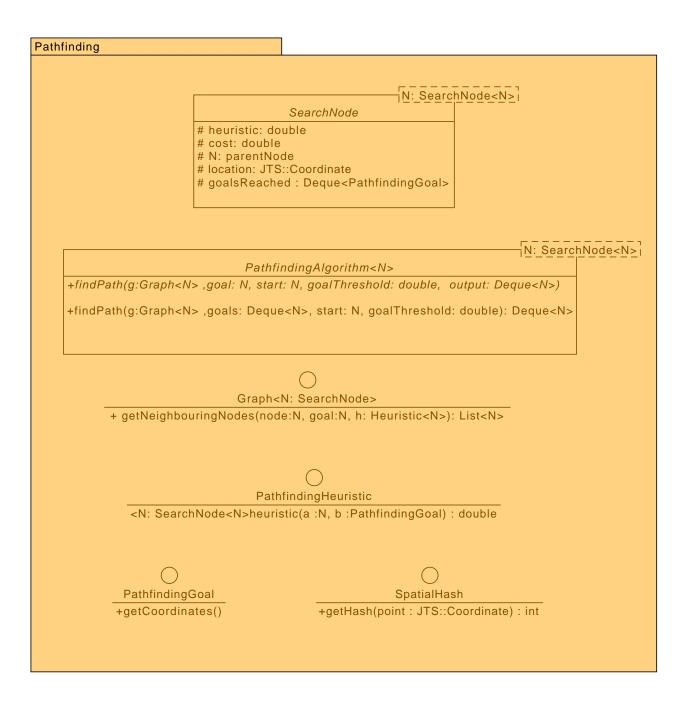


Figure 3: UML diagram of the pathfinding module

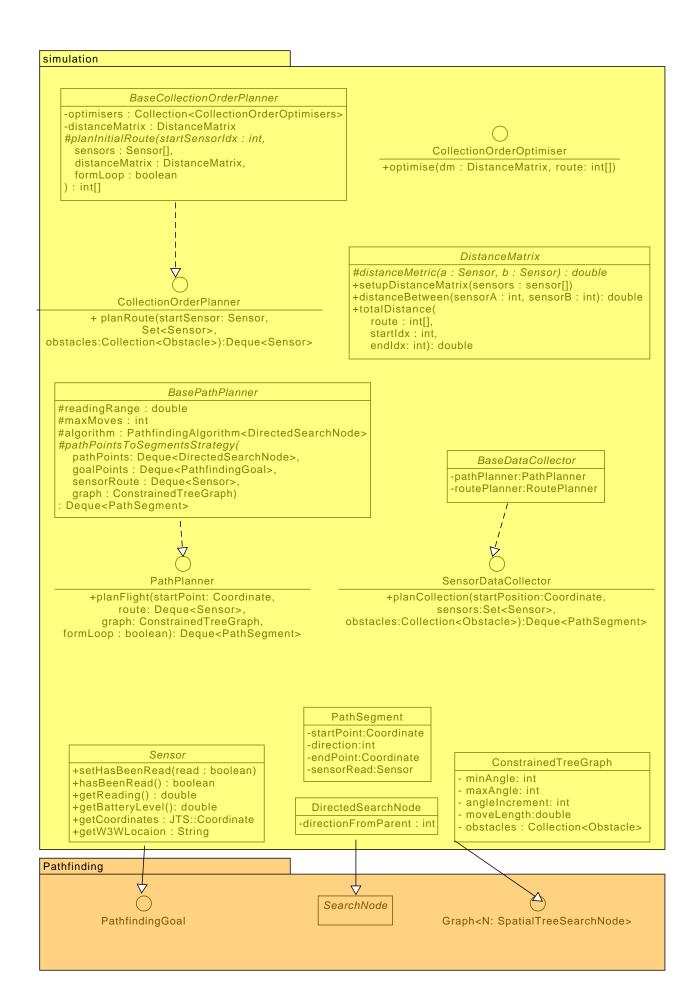


Figure 4: UML diagram of the Simulation module

2 Drone Control Algorithm

2.1 TSP Solver

```
Algorithm NearestInsertion:
     Let R be the sensor nearest to the starting point
     Let U \leftarrow unvisited sensors
     Let T \leftarrow [R] be the current tour (implicitly looping)
     While U \neq \emptyset:
           R \leftarrow \operatorname{arg\,min}_{s \in R} \operatorname{distanceToTour}(s)
           i \leftarrow \operatorname{arg\,min}_{i \in I(T)} \operatorname{insertion\,Cost}(i, R)
           T \leftarrow T with R inserted at i
           U \leftarrow U - \{R\}
     Return T
Function distanceToTour(s):
     \min \min \leftarrow \infty
     For t in T:
           If dist(s,t) < minimum:
                 minimum \leftarrow dist(s,t)
     Return minimum
Function insertionCost(i,s):
     N \leftarrow T with s inserted at i
     Return euclidian length of N
```

Figure 5: Nearest Insertion TSP solver

The more general part of the problem is selecting the collection order of the sensors on any given day. Assuming the drone already has the data for the appropriate collection day, it must decide on a visiting order of the sensors.

For this part many techniques were tried, and experimentation proved that applying the Nearest Insertion heuristic together with a euclidian distance matrix was an adequate solution. At each step we chose the sensor which is closest to one of the sensors present in the tour already - the first sensor is the one nearest to the starting point -, and insert it in a way which minimises the cost. If U is the set of unvisited sensors and T is set of sensors in the tour, at each step we pick the sensor s such that:

$$\exists s, \exists k \in T, \forall o \in U.distance(s, k) < distance(o, k)$$
 (1)

and insert it into the tour between consecutive sensors $i, j \in T$ such that the cost of the tour is minimized. Following this initial pass, 2-opt optimisations are applied to the route to remove crossings in the graph and further reduce the cost. The 2-opt algorithm in each pass checks if reversing any sub-segment of all possible sub-segments of a full tour reduces it's cost, and if so keeps the reversal. The passes are repetead until improvements fall under a threshold of 0.00003 degrees

This alrogithm turns out to be pareto-optimal [1] among a family of cutting-edge algorithms for problem sizes of between 30-50 vertices. Pareto-optimality means that this algorithm was either finding the best path compared to the other algorithms, or was finding one the quickest. It also happens that this algorithm is just a variation of Prim's algorithm as in this process we find the minimum spanning tree of the implicit graph, therefore Nearest insertion is actually a 2-approximation algorithm [2] of the optimal solution.

2.2 Generating neighbour nodes

All the other sections of the algorithm make use of the "neighbour" function, i.e. the function which generates neighbour nodes for any node representing a point on the map.

This is done using the Bounding Volume Hierarchy data structure. The structure allows for logarithmic time lookups of possibly colliding obstacles (with any shape). We do this by creating a binary tree whose nodes are defined by an Axis Aligned Bounding Volume enveloping all the obstacles present in the leaf nodes underneath the node. The root node then envelops all the obstacles present in the hierarchy.

When creating the tree we find the axis along which the difference between the extremal coordinates of the AABB's is the largest, i.e. the "longest axis". We then pick a splitting point on the axis and partition the shapes into the left and right sub trees according to which side of the splitting point they're on.

With this setup, the structure can tell us which obstacles are possibly colliding with any given shape by checking for collisions (cheaply) with the AABB's and only returning those leafs whose AABB's were collided with (possibly from both subtrees).

2.3 Pathfinding

```
Algorithm A*:
     Let R \leftarrow starting node
     let G \leftarrow goal coordinates
     Let O \leftarrow \{R\} // open set
     Let V \leftarrow \{\} // approximately visited
     While O \neq \emptyset:
           \mathbf{R} \leftarrow \, \arg \min\nolimits_{o \in O} \,\, \mathbf{Fvalue} \, (\, \mathbf{o} \, )
           O \leftarrow O - R
           If isNearGoal(R):
                R.\ goalsReached\ \leftarrow\ R.\ goalsReached\ +\ G
                Return reconstructPath(R) // using parent references
           // generate neighbour nodes with appropriate costs and parent set to R
           // nodes colliding with obstacles or outside of boundary are excluded
           // this is done via Bounding Volume Hierarchies
           N \leftarrow neighbours(R)
           For n in N:
                hash \leftarrow cantor Hash(n.x, n.y)
                 If hash in V:
                      Skip n
                V \leftarrow V + hash
                O \leftarrow O + n
     // no path found
     Return []
Function Fvalue(n):
     Return dist (n,G)*1.5 + cost(n)
Function cantorHash(x,y):
     \begin{array}{lll} \textbf{Let} & gridWidth \leftarrow \frac{1}{75}*0.0003 \\ \textbf{Let} & gridX\,, gridY \leftarrow coordinates & of center & of drone & confinement & area \\ \end{array}
     nx, ny \leftarrow (x - gridX)/gridWidth, (y - gridY)/gridWidth
     nx, ny \leftarrow makePositive(|nx|), makePositive(|ny|)
     return |(((0.5 * (nx + ny)) * (nx + ny + 1)) + ny)|;
Function makePositive(n):
     If n \geq 0:
          Return 2n
     Otherwise:
           Return -2n-1
```

Figure 6: Custom A* pathfinding algorithm

Pathfinding between any two points on the plane was carried out using the A* tree search algorithm with the euclidian distance heuristic. The algorithm is just a modified breadth-first search, where the search nodes are

picked in order of least f value which is defined as:

$$f(n) = h(n) + c(n) \tag{2}$$

with c(n) is the cost of reaching the node n from the start state, and h(n) is the approximation of the cost of reaching the goal state from the node n. Assuming that the node in the search frontier with the smallest f value is always expanded and that the heuristic is admissible A* will return the optimal path.

A naive application of A* could not work however due to floating-point error causing the expansion of a lot of poor nodes and the sheer amount of nodes present in the frontier at any time. The straight-line-distance heuristic was relaxed and the new definition was the following:

$$h(n) = distance(n, q) * 1.5$$
(3)

This makes nodes closer to the goal receive much smaller heuristic values and so more direct paths to the goal are expanded first. This means that the returned paths are no longer optimal as the heuristic is no longer admissible, but since most paths would not be obstucted by buildings - this change does not actually impact the length of the paths all that much.

Another change was the addition of a spatial hash to try and prune nodes which were approximately visited, by finding out their coordinates relative to a grid centered around the center of the boundary (with integer width square size), and hashing those coordinates using a modified cantor pairing allowing for negative values

2.4 Path segmenting

Once we have a path of individual points, we need to convert it to a path of path segments with information about which sensor is read at each segment. We can do this by "sliding" a window over the points and always looking at two points at a time in order. We call the points at each position of the windo P and N, where P is the point before N. The pathfinding algorithm also attaches a deque of sensors reached by each P/N. Multiple conversion problems arise when naively pairing up the points to create segments:

- P or N may reach multiple sensors which is invalid, each P-N segment must only allow for one reading
- The list of points might be empty if no path is found
- P might reach a sensor which is illegal, only N's in each segment are allowed to read sensors (since an N point always occurs at the end of a move)

The empty list of points can be prevented by simply returning no path segments, the rest of the problems can be defeated if we apply a number of rules in order (portrayed in Figure 7):

- While P attains a goal, create a proxy segment to any neighbouring node, and one back, assign the first sensor attained by P to the newest back-facing segment's N node
- If N attains no goals, see if the next segment's P node attains any sensors, if so "steal" one away from it
- while N attains more than one sensor reduce the number of attained sensors by creating proxy segments as above for each
- If N only attains one sensor or none, create P-N path segment as normal

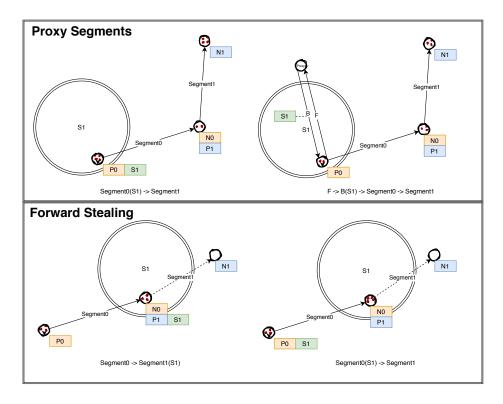


Figure 7: Segmenting behaviours

2.5 Examples

The algorithm was tested on over 35000 configurations over the data provided. The hardest and easiest collection days' geojson visualisations are shown below (with the algorithm set to optimal parameters as given above)

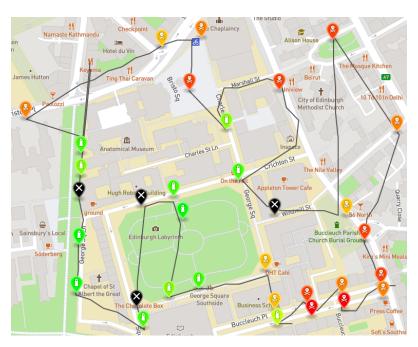


Figure 8: geojson.io rendering of hardest collection at day 9-2-2020 with a starting point of -3.19087,55.945778, with 111 moves

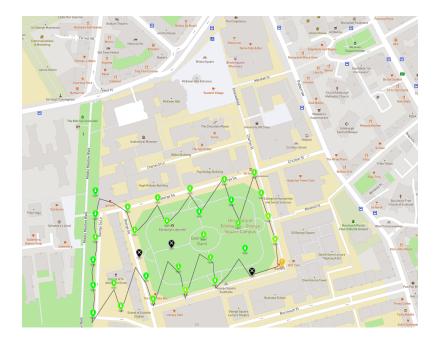


Figure 9: geojson.io rendering of easiest collection at day 2-1-2021 with a starting point of -3.1878,55.9444, with 46 moves

The artifacts of path segmenting are evident in the harder day, where the drone wanders near closely spaced sensors on the bottom right it creates a horizontal proxy segment. But overall the algorithm manages very well, never carrying out the collection in more than 111 moves and averaging at 90 moves \pm 6.7. Execution times averaged at 103ms.

3 Class Documentation

3.1 Quicklinks

3.2 Documentation

3.2.1 uk.ac.ed.inf.aqmaps.visualisation

enum MarkerSymbol

Available symbols for use with geojson

${\bf class} \ {\bf UniformAttribute Map}$

A general attribute map which maps a range of values from a (min,max) range to attribute buckets of size (max-min)/buckets number uniformly.

Implements AttributeMap

Members

public UniformAttributeMap (
 Float min,
 Float max,
 C... attributes)

public C getFor (
 Float value)

Retrieve attribute for the given input

3.2.2 uk.ac.ed.inf.aqmaps.utilities

class BVHNode

Bounding Volume Hierarchy Node. This class forms a tree of AABB (Axis aligned bounding boxes) for internal nodes

and of any shapes at the leaf nodes. Allows for quick broad phase collision checks between objects. Will never return a false negative but might return

false positives. I.e. this structure only tells you which objects are possibly colliding (whose AABB's intersect).

Members public BVHNode (Collection<T> shapes) Construct a new bvh hierarchy with the given shapes at the leaf nodes Construct a new bvh hierarchy with the given shapes at the leaf nodes public Collection<T> getPossibleCollisions (org.locationtech.jts.geom.Geometry other) Retrieves all possibly coliding objects (their bounding boxes intersect) from within the tree.

${\bf class}\ {\bf Geometry Factory Singleton}$

The geometry factory containing the precision model to be used when generating geometries with JTS

public GeometryFactorySingleton () Retrieve the geometry factory, containing the precision model getGeometryFactory ()

class GeometryUtilities

A collection of utility geometry methods

Members public GeometryUtilities () public static org.locationtech.jts.geom.Coordinate MapboxPointToJTSCoordinate (com.mapbox.geojson.Point p) Convert a mapbox point to a jts coordinate public static com.mapbox.geojson.Point JTSCoordinateToMapboxPoint (Convert a jts coordinate to a mapbox point public static org.locationtech.jts.geom.Coordinate p) Convert a jts coordinate to a mapbox point public static org.locationtech.jts.geom.Polygon MapboxPolygonToJTSPolygon (Convert mapbox polygon to jts polygon com.mapbox.geojson.Polygon p) Convert mapbox polygon to jts polygon

class MathUtilities		
Members		
public MathUtilities ()		
public static double angleFromEast (
${\it org.location tech.jts.math.} Vector 2D\ a)$	returns the angle from the eastern direction clockwise between 0 and 360 of the vector	
public static double oppositeAngleFromEast (
double angle)	Return the angle representing the given angle incremented by 180 degrees (but witin 360 degrees)	
public static org.locationtech.jts.math.Vector2D getH-		
eadingVector (
double angle)	gets unit vector in the direction of angle	
public static boolean thresholdEquals (
double a,		
double b,	Determine the Court of the second telegraph of the sec	
double epsilon)	Returns true if the given values are within a threshold of each other	
public static boolean thresholdEquals (double a.		
double b)	Returns true if the given values are within a threshold of each other. Uses default epsilon	
public static boolean thresholdEquals (
org.locationtech.jts.geom.Coordinate a,		
org.locationtech.jts.geom.Coordinate b,		
double epsilon)	Returns true if the given coordinates are within a threshold of each other	
public static boolean thresholdEquals (
org.locationtech.jts.geom.Coordinate a,		
${\it org.} location tech. jts. geom. Coordinate\ b)$	Returns true if the given coordinates are within a threshold of each other. uses default epsilon	

3.2.3 uk.ac.ed.inf.aqmaps.simulation.planning.path

class BasePathPlanner

Base class for planners with a limited number of maximum moves and a

minimum reading range. All inheriting path planners must make sure that each sensor is read at the endpoint of some path segment

and that that sensor is within reading range of the endpoint. They also must make sure that the path is under the maximum move limit

Implements PathPlanner

Members		
public BasePathPlanner (double readingRange, int maxMoves, PathfindingAlgorithm Public Deque <pathsegment> planPath (org.locationtech.jts.geom.Coordinate startCoordinate, Deque<sensor> route, ConstrainedTreeGraph graph,</sensor></pathsegment>		
boolean formLoop)	Plans the exact path required to reach all the given sensors, the specific constraints on placed on the route are decided by the specific implementation of the planner itself. The ConstrainedPathPlanner class adds a maximum move, reading range and also the move sequence constraint to the path planning problem, a valid path will allow the collector to come within READING_RANGE of each sensor (the distance between each sensor and the collector at some path segment's endpoint will be less than or equal to READING_RANGE). The number of segments returned will always be <= MAX_MOVES. The move sequence in each path segment requires that a sensor be read at the end of each path segment only, i.e. the collector must move before collecting any reading, and only one reading must be made per path segment.	
protected abstract Deque <pathsegment> path-PointsToSegmentsStrategy (Deque<directedsearchnode> pathPoints, Deque<pathfindinggoal> goalsRoute, Deque<sensor> sensorRoute,</sensor></pathfindinggoal></directedsearchnode></pathsegment>		
ConstrainedTreeGraph graph)	The main defining characteristic of a constrained path planner. Converts a path of points to a path of path segments needs to make sure that each pathfinding goal is visited only in the end segment of some path segment in range. the passed deque arguments will be consumed	

class PathSegment

represents a singular move made by the sensor data collector each move follows the pattern of: move->collect reading, we cannot collect a reading in a move unless we have moved

| Description |

${\bf class\ Simple Path Planner}$

This path planner will apply some simple optimisations in order to produce a smaller number of path segments than the naive implementation.

Implements PathPlanner

$\overline{\text{Members}}$

Liv Girls D. U.D.			
public SimplePathPlanner (
double readingRange,			
int maxMoves,			
Path finding Algorithm < Directed Search Node > algorithm)			
protected Deque <pathsegment> pathPointsToSeg-</pathsegment>			
mentsStrategy (
Deque <directedsearchnode> pathPoints,</directedsearchnode>			
Deque < Pathfinding Goal > goals Route,			
Deque < Sensor > sensor Route,			
ConstrainedTreeGraph graph)	The main defining characteristic of a constrained path planner. Converts a path of points to a path of path segments needs to make sure that each pathfinding goal is visited only in the end segment of some path segment in range. the passed deque arguments will be consumed. This planner will try to perform some simple optimisations in order to shorten the route. In order to produce a valid route this planner will introduce proxy segments which go back and forth between the nearest neighbour whenever a sensor is read at the start point of a segment or if more than one sensor is read at the endpoint. The optimisations currently include: 1) if the current segment does not read anything at the end point and the next reads a sensor at its start point, we "absorb" that sensor into the current segment.		

${\bf 3.2.4} \quad \hbox{uk.ac.ed.} in f. a quaps. simulation. planning. collection Order. optimisers$

class Optimiser2Opt

An optimiser which performs the 2-opt algorithm to remove crossings in the path

Implements CollectionOrderOptimiser

Members

public Optimiser2Opt (double epsilon)	Construct a 2 opt optimiser with the given epsilon threshold. The threshold determines the minimum decrease in path cost required for the optimiser to keep optimising each loop.	
public void optimise (DistanceMatrix distanceMatrix, int[] path)		

3.2.5 uk.ac.ed.inf.aqmaps.simulation.planning.collectionOrder

${\bf class\ Base Collection Order Planner}$

Collection order planners generate good traversal orders between the given set of sensors, where "good" criteria are defined by each implementation of the collection planner.

public BaseCollectionOrderPlanner (
Collection <collectionorderoptimiser> optimisers,</collectionorderoptimiser>	
DistanceMatrix distMat)	Creates a collection order planner with the given opimisers and distance matrix method
public Deque <sensor> planRoute (</sensor>	
Sensor startSensor,	
Set < Sensor > sensors,	
boolean formLoop)	Generates a collection order over the sensors.
protected abstract int[] planInitialRoute (
int startSensorIdx,	
Sensor[] sensors,	
DistanceMatrix distanceMatrix,	
	Drafts a route between the given sensors, using the given matrix. If
boolean formLoop)	form loop is true then the route will also begin and end on the same
- /	sensor

class GreedyCollectionOrderPlanner

Plans a collection of sensor data in a greedy order and in a way that forms a loop, i.e. by picking the closest sensor at each step.

Members

```
public GreedyCollectionOrderPlanner (
  Collection<CollectionOrderOptimiser> optimiser,
  DistanceMatrix distMat)
protected int[] planInitialRoute (
  int startSensorIdx,
  Sensor[] sensors,
  DistanceMatrix distanceMatrix,
                                                                {\bf Description\ copied\ from\ class:\ BaseCollectionOrderPlanner}
  boolean formLoop)
```

${\bf class\ Nearest Insertion Collection Order Planner}$

Collection order planner which employs the nearest insertion method to try and pick the best route.

Members ${\bf public\ Nearest Insertion Collection Order Planner\ } \\$ Collection
 CollectionOrderOptimiser> optimisers,
 DistanceMatrix distMat) protected int[] planInitialRoute (int startSensorIdx, Sensor[] sensors, DistanceMatrix distanceMatrix, Drafts a route between the given sensors, using the given matrix. If form loop is true then the route will also begin and end on the same $\frac{1}{2}$

sensor. The route is planned using the nearest insertion heuristic.

3.2.6 uk.ac.ed.inf.aqmaps.simulation.planning

$class\ Constrained Tree Graph$

boolean formLoop)

A graph which imposes angle, move length and boundary (+ obstacle) constraints for the nodes, and does not keep track of already produced nodes (tree search) i.e. a new node is returned each time

Implements SearchGraph

public ConstrainedTreeGraph (
int minAngle,	
int maxAngle,	
int angleIncrement,	
double moveLength,	
Collection < Obstacle > obstacles,	
org.locationtech.jts.geom.Polygon boundary)	The angle system needs to allow for each possible angle to have a "complement angle" which takes you back to where you started if you moved in its direction after steping in any possible angle. the min and max angle need to cover a range of 360 degrees - the angle increment.
public ConstrainedTreeGraph (
int minAngle,	
int maxAngle,	
int angleIncrement,	
double moveLength,	
Collection < Obstacle > obstacles)	The angle system needs to allow for each possible angle to have a "complement angle" which takes you back to where you started if you moved in its direction after steping in any possible angle. the min and max angle need to cover a range of 360 degrees - the angle increment.
public Collection <obstacle> getObstacles ()</obstacle>	Retrieve the obstacles present on the map
public org.locationtech.jts.geom.Polygon getBoundary	Retrieve the bounds of the map
public double getMoveLength ()	get the distance between a node and its neighbour
public List <directedsearchnode> getNeighbouringN-</directedsearchnode>	
odes (
DirectedSearchNode node)	Returns the neighbours of the given node within the graph
public int getClosestValidAngle (
double direction)	returns the closest valid angle in this graph to the given double
public int[] getValidDirectionsBetween (
int lowAngle,	
int highAngle)	Retrieves all valid directions between the given low and high angle

class DistanceMatrix

Class which stores distance information between sensors

Members		
public DistanceMatrix ()	Creates new blank distance matrix	
public void setupDistanceMatrix (
Sensor[] sensors)	Fills in the distance matrix with distance data for the given sensors	
public double distanceBetween (
int a,		
int b)	returs the distance from sensor a to sensor b at the given indices in the sensor list	
public double totalDistance (
int[] route,		
int startIdx,		
int endIdx)	calculates the total distance of the subpath specified with start and end indices within the given path of sensor indices.	
public double totalDistance (
int[] route)		
protected abstract double distanceMetric (
Sensor a,		
Sensor b)	the specific distance measure used to calculate distances. Does NOT	

class EuclidianDistanceMatrix

Distance matrix using euclidian distance as the value for distances

Members

_public EuclidianDistanceMatrix ()	
protected double distanceMetric (
Sensor a,	
Sensor b)	Description copied from class: DistanceMatrix

${\bf class}\ {\bf Greatest Avoidance Distance Matrix}$

Distance matrix using the greatest avoidance distance as the distance metric. This distance is calculated by forming a minimum bounding circle around all obstacles

between any two sensors and calculating the length of the path which "wraps" around the circle (approximated as a triangle)

Members

public GreatestAvoidanceDistanceMatrix (
Collection < Obstacle > obstacles)	initialize blank distance matrix with the given obstacles
protected double distanceMetric (
Sensor a,	
Sensor b)	Description copied from class: DistanceMatrix

${\bf 3.2.7} \quad \hbox{ uk.ac.ed.} \hbox{inf.aqmaps.simulation.collection}$

${\bf class\ Base Data Collector}$

Each data collector follows the same pattern, it uses a path planner to find a way between two points, as well as a collection order planner which

sets out the route around all the sensors. Each collector may use this data differently, for example it may discard the path given by a path planner under certain circumstances, or change the route mid-way.

Implements SensorDataCollector

public BaseDataCollector (PathPlanner fp, BaseCollectionOrderPlanner rp)	

class Drone

the drone collector is not constrained by the map layout, if the graph (or map) allows a node to be reached the drone can fly through it, the graph itself may impose constraints indirectly, but the drone assumes absolutely no restrictions in its movements.

Implements SensorDataCollector

viembers		
public Drone (PathPlanner fp, BaseCollectionOrderPlanner rp)		
public Deque <pathsegment> planCollection (</pathsegment>		

org.locationtech.jts.geom.Coordinate startCoordinate, Set<Sensor> sensors, ConstrainedTreeGraph graph, boolean formLoop,

Description copied from interface: SensorDataCollector

3.2.8 uk.ac.ed.inf.aqmaps.simulation

class Building

int randomSeed)

An obstacle which has a polygonal shape

Implements Obstacle Implements Shape

Members

public Building (
org.locationtech.jts.geom.Polygon shape)	initialize new building with the given shape
public org.locationtech.jts.geom.Polygon getShape ()	Retrieve the shape of the obstacle
public boolean intersectsPath (
org.locationtech.jts.geom.Coordinate a,	
org.locationtech.jts.geom.Coordinate b)	Returns true if the line formed from a to b intersects this obstacle

${\bf class\ Directed Search Node}$

A data structure representing a tree search node for spatial pathfinding problems with integer angles.

public DirectedSearchNode (
org.locationtech.jts.geom.Coordinate location,	
DirectedSearchNode parent,	
int directionFromParent,	
	Creates a new spatial tree search node which is fully specified apart
double cost)	from
	the heuristic value
public DirectedSearchNode (
org.locationtech.jts.geom.Coordinate location,	
DirectedSearchNode parent,	
int directionFromParent,	
double heuristic,	
double cost)	Creates fully specified tree search node
public int getDirectionFromParent ()	

class Sensor

Sensors contain a reading value and a battery level as well as the what 3 words location and coordinates of the

This sensor will contain placeholder values for readings and battery status untill it is set to have been read.

Implements PathfindingGoal

Members public Sensor org.locationtech.jts.geom.Coordinate coordinates, float reading, float batteryLevel, String W3WLocation) public boolean equals Object o) public int hashCode () public String toString (public void setHaveBeenRead (Sets the sensor state to "read". If the sensor is not set to have been boolean read) read, asking it for readings and battery levels will return placeholder values $public\ org. location tech. jts. geom. Coordinate\ get Coordinate$ Return the coordinates of the sensor nates () Return the reading at this sensor. Will return NaN untill the setHave- ${\bf public\ float\ getReading}\ ()$ BeenRead(true) has been called Return the battery level at this sensor. Will return NaN untill the public float getBatteryLevel () setHaveBeenRead(true) has been called Returns true if this sensor has been read (i.e. setHaveBeenRead(true)

has been called on this object)

Retrieves the 3 word address (w3w) of this sensor

class SensorDataCollectorFactory

public boolean hasBeenRead ()

public String getW3WLocation ()

A utility class for instantiating sensor data collectors with correct parameter values in one call

Members

_public SensorDataCollectorFactory ()	
public static SensorDataCollector createCollector (
ConstrainedTreeGraph g,	
double readingRange,	
int maxMoves,	
SensorDataCollectorFactory.CollectorType collectorType,	
SensorDataCollectorFactory.PathfindingHeuristicType	
heuristicType,	
SensorDataCollectorFactory.CollectionOrderPlannerType	
plannerType,	
SensorDataCollectorFactory.DistanceMatrixType matrix-	Create a new collector within the given domain and with the given
Type)	parameters.
public static SensorDataCollector createCollector (
ConstrainedTreeGraph g,	
double readingRange,	
int maxMoves,	
SensorDataCollectorFactory.CollectorType collectorType,	
SensorDataCollectorFactory.PathfindingHeuristicType	
heuristicType,	
SensorDataCollectorFactory.CollectionOrderPlannerType	
plannerType,	
SensorDataCollectorFactory.DistanceMatrixType matrix-	
Type,	
float relaxationFactor,	
double hashingGridWidth,	
double opt2Epsilon)	Create a new collector within the given domain and with the given parameters

$enum\ Sensor Data Collector Factory. Collection Order Planner Type$

enum SensorDataCollectorFactory.CollectorType

$enum\ Sensor Data Collector Factory. Distance Matrix Type$

$enum\ Sensor Data Collector Factory. Path finding Heuristic Type$

${\bf 3.2.9} \quad \hbox{uk.ac.ed.inf.} a quaps. path finding. heuristics$

$class\ StraightLineDistance$

The simplest heuristic, uses the euclidian distance between the node and its goal as the value of its heuristic.

Implements PathfindingHeuristic

Members	
public StraightLineDistance ()	Initializes a new instance of the straight line heuristic with a relaxation factor of 1 (i.e. no relaxation)
public StraightLineDistance (
double relaxationFactor)	Initializes a new instance of the straight line heuristic with the given relaxation factor. Some pathfinding algorithms such as Astar will run much faster with a relaxed heuristic, but the relaxed solution might not be optimal (path cost <= relaxationFactor * optimal path cost)
public <t extends="" searchnode<t="">> double heuristic</t>	
Та,	
PathfindingGoal b)	Returns the value of the heuristic between the given node and the goal. A heuristic must be a) admissible - i.e. never overestimate the real cost of reaching the goal b) consistent - i.e. follow the triangle inequality Otherwise pathfinding algorithms might struggle to find a good solution or to even find one at all

3.2.10 uk.ac.ed.inf.aqmaps.pathfinding.hashing

${\bf class}~{\bf Grid Snapping Spatial Hash}$

hashes real coordinates by scaling them and "snapping" them to a grid..

Implements SpatialHash

Members

public GridSnappingSpatialHash (
double gridSize,	
org.locationtech.jts.geom.Coordinate gridCenter)	Create new instance of grid snapping hash
public int getHash (
org.locationtech.jts.geom.Coordinate a)	Returns a hash for the given coordinate. A good hash will be equal for points which are near each other according to some metric, and different for ones that are not.

${\bf 3.2.11} \quad \hbox{ uk.ac.ed.inf.} a quaps. path finding. goals$

class PointGoal

A point goal is the simplest kind of goal, a single location in space.

Implements PathfindingGoal

public PointGoal (
org.locationtech.jts.geom.Coordinate goal)	Creates a new point goal
public org.locationtech.jts.geom.Coordinate getCoordi-	
nates ()	

${\bf 3.2.12} \quad \hbox{uk.ac.ed.inf.aqmaps.pathfinding}$

class AstarTreeSearch

Classic pathfinding algorithm, modified BFS which uses both the cost to reach a node and the predicted cost from that node to the goal

to chose the nodes to be expanded next. This version of Astar treats the search as a tree search and so uses a hashing function to determine if a node has been visited yet.

Members	
public AstarTreeSearch (PathfindingHeuristic heuristic,	
SpatialHash hash)	Creates a new instance of Astar search with the given heuristic and spatial hashing function.
public void findPath (
SearchGraph <t> g,</t>	
PathfindingGoal goal,	
T start,	
double goalThreshold,	
Deque <t> output)</t>	Description copied from class: PathfindingAlgorithm

class PathfindingAlgorithm

Pathfinding algorithms operate over any valid search nodes and graph definitions. The graph defines the transition function between one node and its neighbours while the search nodes are used as

the path constructing object. Any number of pathfinding algorithms can be defines in these terms, both tree and graph searches are possible with the correct set of graph and node definitions.

Members	
public PathfindingAlgorithm ()	
public abstract void findPath (
SearchGraph <t> g,</t>	
PathfindingGoal goal,	
T start,	
double goalThreshold,	
double goal i meshold,	Finds a path from the start to the goal node and outputs the result
	into the provided deque.
	If a path doesn't exist no nodes will be added to the output, if it does at
Deque <t> output)</t>	least one node will be added. The nodes will have their goalsReached
Deque (1) Gutput)	deque's set
	to the corresponding goals that can be reached from their locations
	within the goal threshold. One node can reach multiple nodes
public Deque <t> findPath (</t>	within the goal threshold one node can reach material nedes
SearchGraph <t> g,</t>	
Deque < Pathfinding Goal > route,	
T start,	
1 50010,	Finds a path from the start node through the provided route. Modifies
	the goal nodes' locations to accommodate for the goal threshold.
double goalThreshold)	if at any point there exists no path between the given goals, the route
	will halt before that segment (so might be empty).
protected boolean isAtGoal (3 (0 17)
double threshold,	
T node,	
PathfindingGoal goal)	Checks whether the given node is within a threshold away from the
r atililiding Goal goal)	goal
protected void reconstructPathUpToIncluding (
T node,	
Deque < T > out,	
T limitNode)	reconstructs the path to node and deposits it in the given queue

${\bf class} \,\, {\bf SearchNode}$

Search Nodes are used to hold the path information in pathfinding algorithms including heuristic, cost and parent node values. Search nodes also contain information about which pathfinding goals can be achieved from their position (can be multiple).

Members	
<pre>public SearchNode (org.locationtech.jts.geom.Coordinate location, T parent,</pre>	Creates a new spatial tree search node which is fully specified apart
double cost)	from the heuristic value
public SearchNode (
org.locationtech.jts.geom.Coordinate location,	
T parent,	
double heuristic,	
double cost)	Creates a fully specified search node
public Deque <pathfindinggoal> getGoalsReached ()</pathfindinggoal>	Retrieve the deque of goals achievable from this search node's position
public void setGoalsReached (
Deque <pathfindinggoal> goalReached)</pathfindinggoal>	Set the goals achievable from this search node's position
public void addGoalReached (
PathfindingGoal goalReached)	Add a goal achievable from this node's position to the tail of its deque
_public int getNumberOfGoalsReached ()	
public void removeGoalReached (
PathfindingGoal goalReached)	pop the first goal achievable at this node's position from its deque
public double getHeuristic ()	Retrieve the heuristic value of this node
public org.locationtech.jts.geom.Coordinate getCoordi-	Retrieve the coordinates in space of this node
_ nates ()	
public void setHeuristic (
double heuristic)	Set the heuristic value of this node
public double getCost ()	Retrieve the cost of reaching this search node
public T getParentNode ()	Retrieve the parent node of this node
public int hashCode ()	
public String toString ()	

3.2.13 uk.ac.ed.inf.aqmaps.client.data

${\bf class\ Sensor Data}$

Object representing the reading, battery and location data of a sensor

Members	
public SensorData ()	Empty constructor, necessary for proper de-serialization
public SensorData (
String W3WLocation,	
float battery,	
float reading)	Constructs a new sensor data object
public String getW3WLocation ()	Returns the w3w location of this sensor
public Float getBattery ()	Returns the battery level of the sensor, always a normal float value
public float getReading ()	Returns the reading value of the sensor, can be NaN or a float
public boolean equals (
Object o)	
public int hashCode ()	
public String toString ()	

${\bf class~W3WAddressData}$

An object containing information about a what-3-word square including it's coordinates and meta-data

Members	
public W3WAddressData ()	Empty constructor, necessary for proper de-serialization
public W3WAddressData (
String country,	
W3WSquareData square,	
String nearestPlace,	
com.mapbox.geojson.Point coordinates,	
String words,	
String language,	
String map)	
public String getCountry ()	
public W3WSquareData getSquare ()	
public String getNearestPlace ()	
public com.mapbox.geojson.Point getCoordinates ()	
public String getWords ()	
public String getLanguage ()	
public String getMap ()	
public boolean equals (
Object o)	
public int hashCode ()	
public String toString ()	

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public String toString ()

${\bf class~AQWebServerClient}$	
{The web server client communicates with a server which contains the necessary information and retrieves it}	
Implements ClientService	
Members	
public AQWebServerClient (
HttpClient client,	
URI APIBaseURI)	
public List <sensordata> fetchSensorsForDate (</sensordata>	
LocalDate date)	
public W3WAddressData fetchW3WAddress (
String wordAddress)	
public W3WAddressData fetchW3WAddress (
String w1,	
String w2,	
String w3)	
public com.mapbox.geojson.FeatureCollection fetch-	
Buildings ()	

class HTTPException	
An exception thrown whenever an unexpected http status is returned	
Implements Serializable	
Members	
public HTTPException (
int errorStatusCode, String message)	

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class App	
Sensor data pick up path planner and visualiser	
Members	
public App () public static void main (String[] args)	

References

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