

Syntax of switch in C (K&R)

In main text:

```
switch (expression) {  
    case const-expr: statements  
    case const-expr: statements  
    default: statements  
}
```

Grammar Excerpt in Appendix:

```
stmt => labeled stmt | selection stmt | ...  
selection stmt => ... | switch (expression) stmt | ...  
labeled stmt => ... | case const-expr: stmt | ...
```

More C Syntax

- **statement:**

labeled-statement

expression-statement

compound-statement

selection-statement

iteration-statement

jump-statement

- labeled-statement:
 identifier : **statement**
 case **constant-expression** : **statement**
 default : **statement**
- expression-statement: **expression_{opt}** ;
- compound-statement:
 { **declaration-list_{opt}** **statement-list_{opt}** }

- statement-list:
 statement
 statement-list **statement**
- selection-statement:
 if (**expression**) **statement**
 if (**expression**) **statement** **else** **statement**
 switch (**expression**) **statement**
- iteration-statement:
 while (**expression**) **statement**
 do **statement** **while** (**expression**) ;
 for (**expression_{opt}** ; **expression_{opt}** ; **expression_{opt}**) **statement**
- jump-statement:
 goto **identifier** ;
 continue ;
 break ;
 return **expression_{opt}** ;
- expression:
 assignment-expression
 expression , **assignment-expression**

- assignment-expression:
conditional-expression
unary-expression assignment-operator assignment-expression
- assignment-operator: one of
= *= /= %= += -= <<= >>= &= ^= |=
- conditional-expression:
logical-OR-expression
logical-OR-expression ? expression : conditional-expression
- constant-expression: conditional-expression
- logical-OR-expression:
logical-AND-expression
logical-OR-expression || logical-AND-expression
- logical-AND-expression:
inclusive-OR-expression
logical-AND-expression && inclusive-OR-expression
- inclusive-OR-expression:
exclusive-OR-expression
inclusive-OR-expression | exclusive-OR-expression
- exclusive-OR-expression:
AND-expression
exclusive-OR-expression ^ AND-expression

...

- AND-expression:
equality-expression
AND-expression & equality-expression
- equality-expression:
relational-expression
equality-expression == relational-expression equality-expression != relational-expression
- relational-expression:
shift-expression
relational-expression < shift-expression
relational-expression > shift-expression
relational-expression <= shift-expression
relational-expression >= shift-expression
- shift-expression:
additive-expression
shift-expression << additive-expression
shift-expression >> additive-expression
- additive-expression:
multiplicative-expression
additive-expression + multiplicative-expression
additive-expression - multiplicative-expression
- multiplicative-expression:
multiplicative-expression * cast-expression
multiplicative-expression / cast-expression
multiplicative-expression % cast-expression

C Compiler accepts “wrong” code?

```
// switch.c
```

```
int main() {  
    switch(1) { if (2) 3+2; else 1+1;};  
    return(1);  
}
```

```
cc -S -save-temps switch.c
```

```
switch.c:4:10: warning: no case matching constant switch condition '1'
```

```
    switch(1) { if (2) 3+2; else 1+1;};
```

^

1 warning generated.

```
        .section      __TEXT,__text,regular,pure_instructions  
        .build_version macos, 12, 0      sdk_version 13, 1  
        .globl _main  
function main  
        .p2align      2  
_main:                                     ; @main  
        .cfi_startproc  
; %bb.0:  
        sub      sp, sp, #16  
        .cfi_def_cfa_offset 16  
        str      wzr, [sp, #12]  
        mov      w0, #1  
        add      sp, sp, #16  
        ret  
        .cfi_endproc  
                                           ; -- End function  
        .subsections_via_symbols
```

Language, Grammar, Derivations, Recognizers

- $a^n b^n$

- $a^n b^n c^n$

In a **leftmost derivation**, at each step, the leftmost nonterminal is replaced.

- $a^n b^n c^n d^n$

In a **rightmost derivation**, at each step, the rightmost nonterminal is replaced.

- CFL

- Recursive descent parsers
- SLR(k) with k lookahead tokens, $k=0, 1, \dots$
- LL(k), $k=0, 1, \dots$
 - Left to Right, Leftmost derivation (“left recursive”)
- LR(k), $k=0, 1, \dots$
 - Left to Right, Rightmost derivation (“right recursive”)
- LALR(1) same as LR(0) with 1 lookahead token
 - “optimized” LR(0)

Dragon book

Leftmost derivation parse

$$\begin{aligned} E &\rightarrow T E' \\ E' &\rightarrow + T E' \mid \epsilon \\ T &\rightarrow F T' \\ T' &\rightarrow * F T' \mid \epsilon \\ F &\rightarrow (E) \mid \text{id} \end{aligned}$$

Parsing table

NON - TERMINAL	INPUT SYMBOL					
	id	+	*	()	\$
E	$E \rightarrow T E'$			$E \rightarrow T E'$		
E'		$E' \rightarrow + T E'$			$E' \rightarrow \epsilon$	$E' \rightarrow \epsilon$
T	$T \rightarrow F T'$			$T \rightarrow F T'$		
T'		$T' \rightarrow \epsilon$	$T' \rightarrow * F T'$		$T' \rightarrow \epsilon$	$T' \rightarrow \epsilon$
F	$F \rightarrow \text{id}$			$F \rightarrow (E)$		

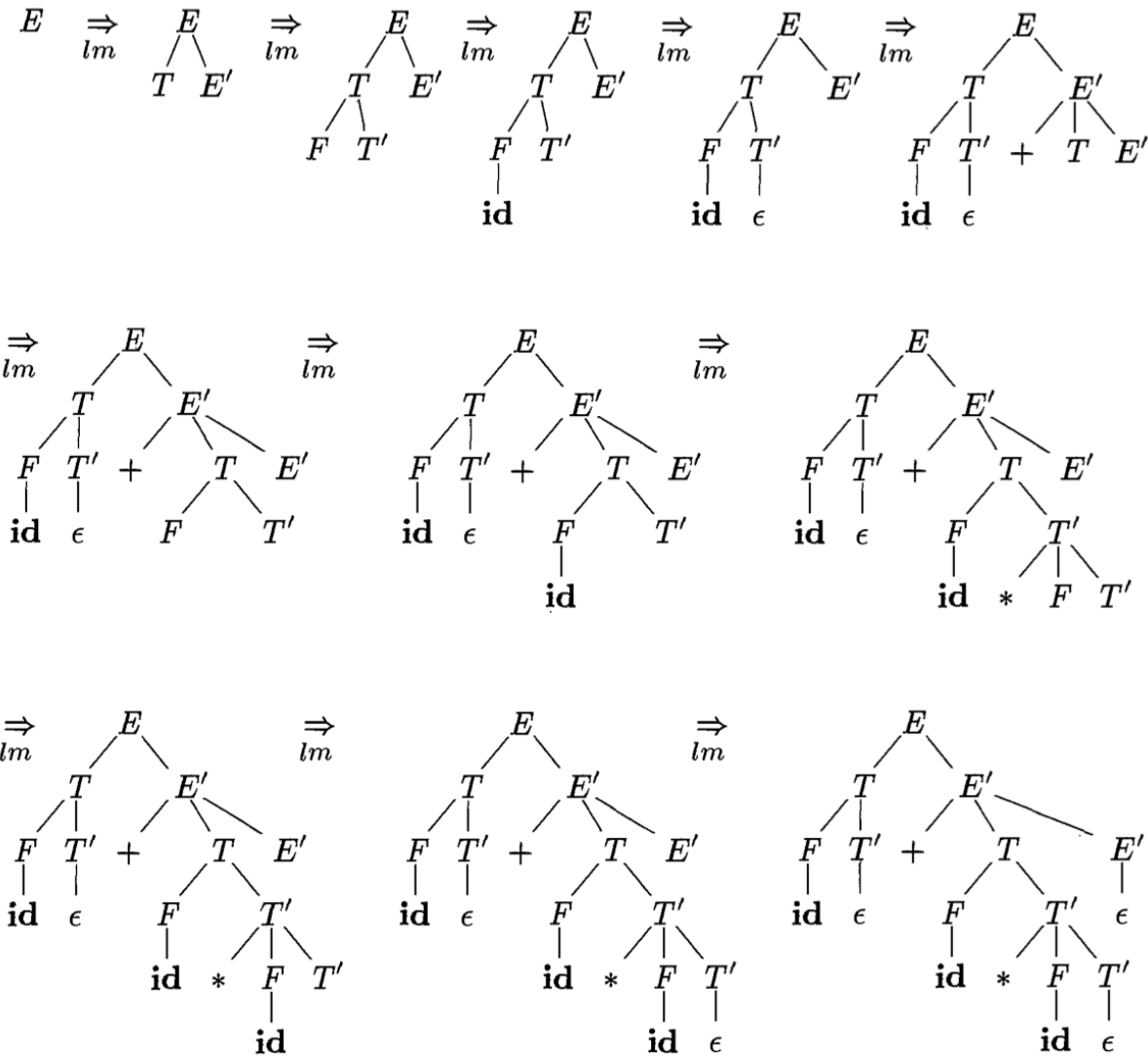


Figure 4.12: Top-down parse for **id + id * id**