Kathmandu University

Department of Computer Science and Engineering

Dhulikhel, Kavre



Lab Report 1 COMP 314

(For partial fulfillment of 3rd Year/ 2nd Semester in Computer Engineering)

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- 1. Implement linear and binary search algorithms.
 - → Filename: search.py

```
def linear_search(data, target):
    for i in range(len(data)):
        if data[i] == target:
def binary_search(data, target):
    first = 0
    while first <= last:</pre>
        mid = (first + last)//2
        if data[mid] == target:
            return mid
        if data[mid] < target:</pre>
            first = mid + 1
            last = mid - 1
```

- 2. Write some test cases to test your program.
 - → Filename: searchTest.py

```
import unittest
from search import linear search
from search import binary search
class TestSearch(unittest.TestCase):
   def test search(self):
       self.assertEqual(linear search(data, 6), 4)
       self.assertEqual(linear search(data, 10), -1)
       self.assertEqual(binary search(data, 6), 4)
       self.assertEqual(binary search(data, 10), -1)
       data = ['t', 'a', 'b', 'l', 'e']
       self.assertEqual(linear search(data, 'a'), 1)
       self.assertEqual(binary search(data, 'a'), 1)
if name == " main ":
```

- 3. Generate some random inputs for your program and apply both linear and binary search algorithms to find a particular element on the generated input. Record the execution times of both algorithms for best and worst cases on inputs of different size (e.g. from 10000 to 100000 with step size as 10000). Plot an input-size vs execution-time graph.
 - → Filename: searchMain.py

```
from search import linear search
from search import binary search
from time import time
import random
import matplotlib.pyplot as plt
def partition(arr, low, high):
    i = (low - 1)
   pivot = arr[high]
    for j in range(low , high):
        if arr[j] <= pivot:</pre>
            arr[i], arr[j] = arr[j], arr[i]
    arr[i+1], arr[high] = arr[high], arr[i+1]
    return i+1
def quickSort(arr, low, high):
    if low < high:
        q = partition(arr, low, high)
        quickSort(arr, low, q - 1)
        quickSort(arr, q + 1, high)
```

```
linearBest = []
linearWorst = []
binaryBest = []
\overline{\text{binaryWorst}} = \overline{[]}
i = 10000
dataSize = []
while i <= 100000:
    dataSize.append(i)
    data = random.sample(range(i), i)
    indexLB = linear search(data, data[0])
    end = time()
    linearBest.append(end - start)
    start = time()
    end = time()
    linearWorst.append(end - start)
```

```
indexBB = binary search(data, data[(len(data) - 1)//2])
    binaryBest.append(end - start)
    start = time()
    indexBW = binary_search(data, data[-1])
    binaryWorst.append(end - start)
    i += 10000
print(indexLB)
print(indexLW)
print(indexBB)
print(indexBW)
print(linearBest)
print(linearWorst)
print(binaryBest)
print(binaryWorst)
plt.plot(dataSize,linearBest,"g")
plt.show()
plt.plot(dataSize,linearWorst,"r")
plt.show()
```

```
plt.plot(dataSize, binaryBest, "g")
plt.show()

plt.plot(dataSize, binaryWorst, "r")
plt.show()
```

4. Explain your observations.

→ Output:

Indexes for searched data value:

Linear Best Case: 0

Linear Worst Case: 99999

Binary Best Case: 49999

Binary Worst Case: 99999

Time taken:

Linear Worst Case: [0.0, 0.000997304916381836, 0.0019524097442626953, 0.001995086669921875, 0.003900766372680664, 0.003239154815673828, 0.003956317901611328, 0.005024433135986328, 0.0059163570404052734, 0.006980180740356445]

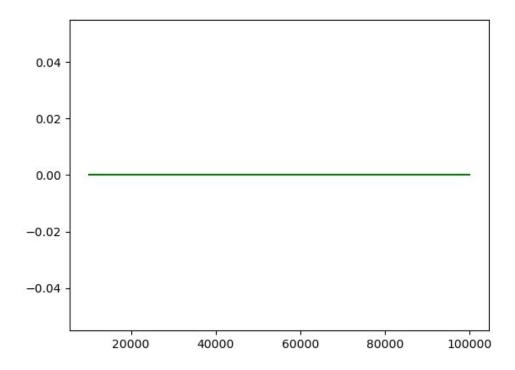


Figure: Linear Search Best Case Time Complexity

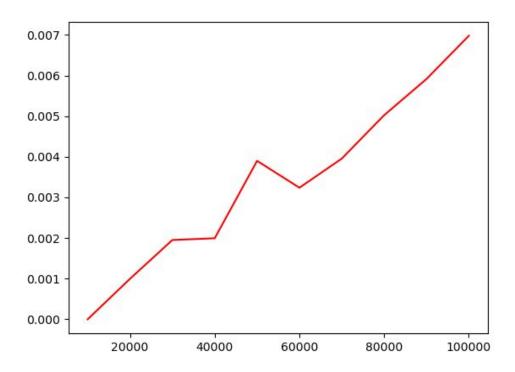


Figure: Linear Search Worst Case Time Complexity

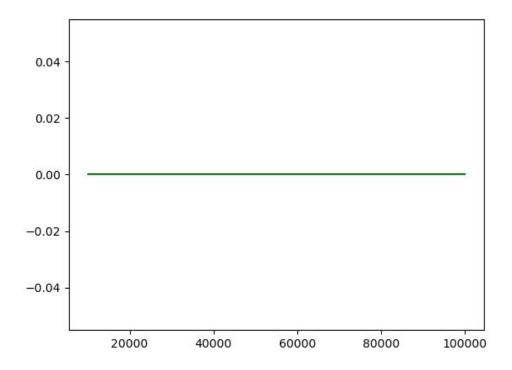


Figure: Binary Search Best Case Time Complexity

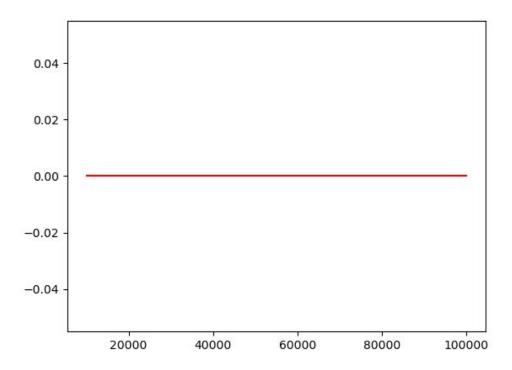


Figure: Binary Search Worst Case Time Complexity

Linear Search Best Case Time Complexity:

For this, the first element of the list is taken and searched. So the time complexity is O(1). Hence, a graph is a straight line parallel to x-axis as time is constant even when data size is increased

Linear Search Worst Case Time Complexity:

For this, the last element of the list is taken and searched. So the time complexity is O(n). Hence, a graph is linear as time increases when data size is increased.

Binary Search Best Case Time Complexity:

For this, the middle element of the list is taken, as it is the root node in the binary tree and searched. So the time complexity is O(1). Hence, a graph is a straight line parallel to x-axis as time is constant even when data size is increased.

Binary Search Worst Case Time Complexity:

For this, the last element of the list is taken and searched. So the time complexity is $O(\log_2 n)$. Since the data size is small, a graph with constant time complexity is obtained, even when data size is increased.