

Please give clear and rigorous answers.
Be to the point. Show your work.

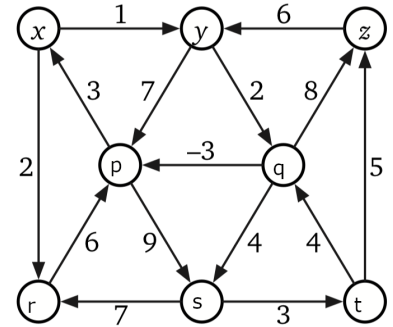
Name: _____

ERP: _____

Question 1: Graphs 11 marks

(a) [3 marks] Clearly indicate the following structures in the directed graph below.

- i. A depth-first tree rooted at x .
- ii. A breadth-first tree rooted at y .
- iii. The shortest directed cycle.



(b) [4 marks] Run Bellman-Ford algorithm on the above graph starting from node z . Show **dist** values after each iteration of the main loop. Draw the computed shortest-path tree.

(c) [2 marks] Consider the execution of depth-first search on a directed graph G from vertex s , beginning with the function call $dfs(G, s)$. Suppose that $dfs(G, v)$ is called during the depth-first search. Which of the following statements can you infer at the moment when $dfs(G, v)$ is called? Mark each as True or False.

- i. G contains a directed path from s to v .
- ii. The function-call stack contains a directed path from s to v .
- iii. If G includes an edge $v \rightarrow w$ for which w has been previously marked, then G has a directed cycle containing v .
- iv. If G includes an edge $v \rightarrow w$ for which w is currently a vertex on the function-call stack, then G has a directed cycle containing v .

(d) [2 marks] Consider the execution of breadth-first search on a directed graph G , starting from vertex s . Suppose that vertex v is removed from the queue during the breadth-first search. Which of the following statements can you infer at the moment when v is removed from the queue? Mark each as True or False.

- i. G contains a directed path from s to v .
- ii. The queue contains a directed path from s to v .
- iii. If G includes an edge $v \rightarrow w$ for which w has been previously marked, then G has a directed cycle containing v .
- iv. If G includes an edge $v \rightarrow w$ for which w is currently a vertex on the queue, then G has a directed cycle containing v .

Question 2: Running Time 7 marks

(a) For each of the following Java methods, choose the best matching running time from:

$O(\log N)$, $O(N)$, $O(N \log N)$, $O(N^2)$, $O(2^N)$, $O(N!)$

i.

```
public static int f1(int N) {
    int x = 0;
    for (int i = 0; i < N; i++)
        x++;
    return x;
}
```

iv.

```
public static int f4(int N) {
    if (N == 0) return 0;
    return f4(N/2) + f1(N) +
        f4(N/2);
}
```

ii.

```
public static int f2(int N) {
    int x = 0;
    for (int i = 0; i < N; i++)
        for (int j = 0; j < i; j++)
            x++;
    return x;
}
```

v.

```
public static int f5(int N) {
    int x = 0;
    for (int i = N; i > 0; i = i/2)
        x += f1(i);
    return x;
}
```

iii.

```
public static int f3(int N) {
    if (N == 0) return 1;
    int x = 0;
    for (int i = 0; i < N; i++)
        x += f3(N-1);
    return x;
}
```

vi.

```
public static int f6(int N) {
    if (N == 0) return 1;
    return f6(N-1) + f6(N-1);
}
```

vii.

```
public static int f7(int N) {
    if (N == 1) return 0;
    return 1 + f7(N/2);
}
```

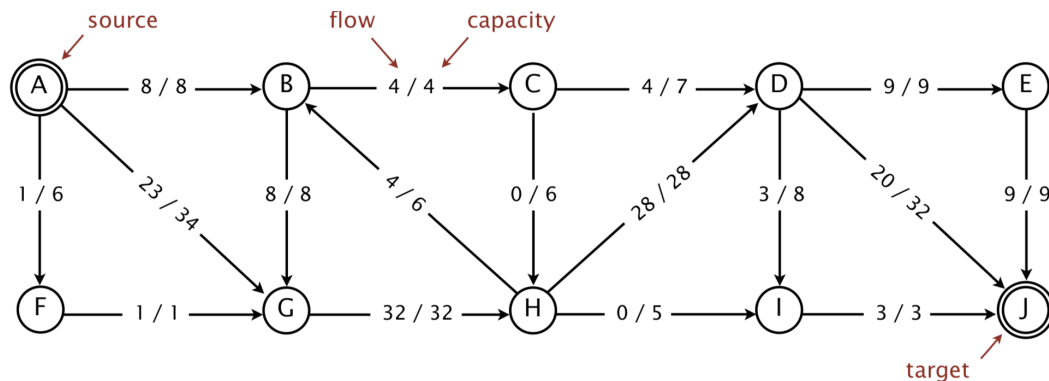
Question 3: Reductions 3 marks

Suppose that Problem X poly-time reduces to Problem Y . Mark whether each of the following statements are True or False.

- (a) If X can be solved in polynomial time, then so can Y .
- (b) If Y can be solved in polynomial time, then so can X .
- (c) If X cannot be solved in polynomial time, then neither can Y .
- (d) If Y cannot be solved in polynomial time, then neither can X .
- (e) If Y is NP-complete, then so is X .

Question 4: Maximum Flows 5 marks

Consider the following flow network and feasible flow f from the source vertex A to the sink vertex J .

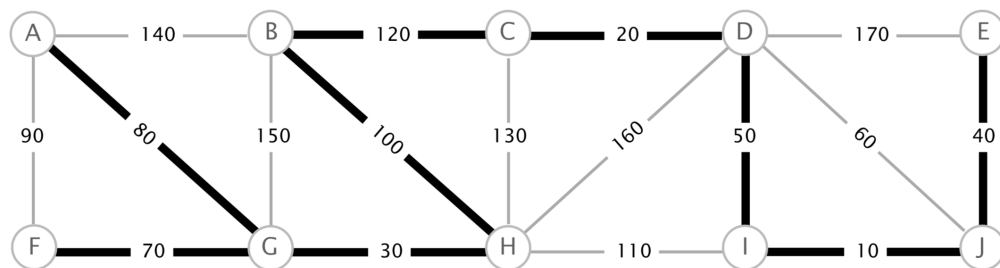


- [1 mark] What is the size of the flow f ?
- [2 marks] Starting from the flow f , perform one iteration of the Ford-Fulkerson algorithm. List all vertices that are on the (unique) augmenting path. What is the bottleneck capacity of the augmenting path?
- [1 mark] List the vertices on the source side of a minimum cut.
- [1 mark] For which of the following edges, doubling the capacity would increase the value of the maximum flow?

$A \rightarrow G, B \rightarrow C, D \rightarrow I, G \rightarrow H, I \rightarrow J, H \rightarrow D$

Question 5: Minimum Spanning Trees 6 marks

- Consider the following edge-weighted graph G containing 10 vertices and 17 edges. The thick black edges T define a spanning tree of G but not a minimum spanning tree of G .



- [1 mark] Find a cut in G whose minimum weight crossing edge is not an edge in T .
 - [2 marks] Run Prim's algorithm on G starting from vertex A to find an MST T' in G . In what order the edges of T' are found?
- [3 marks] Let G be any simple graph (no self-loops or parallel edges) with positive and distinct edge weights. For each of the following statements, make a short proof if the statement is true; otherwise provide a counterexample to show that it is false.
 - Any MST of G must include the edge of minimum weight.
 - Any MST of G must exclude the edge of maximum weight.
 - If the weights of all edges in G are increased by 17, then any MST in G is an MST in the modified edge-weighted graph.

Question 6: Divide & Conquer 6 marks

- (a) [3 marks] You are given a sorted array of numbers where every value except one appears exactly twice; the remaining value appears only once. Design a $O(\log n)$ algorithm for finding which value appears only once.

Here are some example inputs to the problem:

```

1 1 2 2 3 4 4 5 5 6 6 7 7 8 8
10 10 17 17 18 18 19 19 21 21 23
1 3 3 5 5 7 7 8 8 9 9 10 10

```

- (b) [2 marks] State Master theorem.
- (c) [1 mark] Use master theorem to solve the recurrence: $T(n) = 3T(n/2) + O(n)$

Question 7: Dynamic Programming 7 marks

- (a) [3 marks] In this question we will compute edit-distance between the strings $X = \text{your-firstname}$ and $Y = \text{lastname}$. (If your first-name or last-name have more than 5 letters then consider the first 5 letters only).
- Write the recurrence for computing the optimal cost of a problem given the optimal solution of relevant subproblems. How many subproblems we get?
 - Fill in the appropriate table using the recurrence from previous part.
- (b) [4 marks] Given a set of n positive integers, find if we can partition it into two subsets such that the sum of elements in both the subsets is equal. E.g., the set $\{1, 1, 3, 4, 7\}$ can be partitioned into two subsets with equal sum: $\{1, 3, 4\}$ & $\{1, 7\}$. While the set $\{2, 3, 4, 6\}$ cannot be partitioned into two subsets with equal sum.

Your algorithm should run in $O(nS)$ time, where S is the sum of all elements in the given input set.

Question 8: Multithreaded Algorithms 5 marks

- (a) [3 marks] Consider the following multithreaded pseudocode for transposing an $n \times n$ matrix A in place:

TRANSPOSE(A)

```

1:  $n = A.rows$ 
2: parallel for  $j = 2$  to  $n$  do
3:   parallel for  $i = 1$  to  $j - 1$  do
4:     exchange  $a_{ij}$  with  $a_{ji}$ 

```

Analyze the work, span, and parallelism of this algorithm.

- (b) [2 marks] Suppose that we replace the parallel for loop in line 3 of TRANSPOSE with an ordinary for loop. Analyze the work, and span of the resulting algorithm.