

**M A S A R Y K
U N I V E R S I T Y**

FACULTY OF INFORMATICS

**Disk sector content analysis and
visualization**

Bachelor's Thesis

JAKUB MALOŠTÍK

Brno, Spring 2022

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Advisor: Ing. Milan Brož, Ph.D.

Department of Computer Systems and Communications

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Declaration

Hereby I declare that this paper is my original authorial work, which I have worked out on my own. All sources, references, and literature used or excerpted during elaboration of this work are properly cited and listed in complete reference to the due source.

Jakub Malošík

Advisor: Ing. Milan Brož, Ph.D.

Acknowledgements

These are the acknowledgements for my thesis, which can span multiple paragraphs.

Abstract

This is the abstract of my thesis, which can span multiple paragraphs.

Keywords

keyword1, keyword2, ...

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Introduction

Disks (e.g., hard drives, SSDs, Flash drives) are usually divided into atomic parts named sectors, which are represented as blocks in the software layer. Sectors store a fixed amount of data, usually 512 bytes and 4KiB, but other sector sizes can be used. Sectors may contain partition tables, file system information, files or be empty.

Some of the possible contents may contain specific byte patterns which can be analyzed and used to identify the type of content stored in the sector. When a byte pattern is not present, sector content can be analyzed for entropy to estimate whether it is encrypted. A good way to get an idea about which parts of the disk are encrypted and where filesystem data is stored is to visualize the data. This visualization will allow humans to distinguish between different data encryption methods such as filesystem-level and full-disk encryption and even uncover faulty encryption. Visualizing can also be very useful as an illustration while teaching.

The utility introduced in this bachelor's thesis analyzes the sectors of a user-specified size of a provided disk image and visualizes the result using the Pillow Python library. The utility is also easily extensible by other output methods.

The text of this thesis is structured into five chapters. Chapter number one explains the foundations of the thesis and examines prior work. Chapter number two lists some byte patterns of sectors and discusses algorithms for their detection. Chapter number three discusses algorithms used to calculate entropy and possible issues with their accuracy. Chapter number four discusses ways of visualization and their advantages and disadvantages. The last chapter concludes with an evaluation of the resulting utility.

The resulting utility is available on GitHub¹ under the MIT License.

1. <https://github.com/malon43/entropy-visualization>

1 Prior work

This review focuses on works on the topics of block detection, entropy calculation, and ways of visualization.

1.1 Analysis

This section describes parts of analyzing the disk sectors.

1.1.1 Block Patterns

Each disk is divided into tens, even hundreds of millions of sectors. Each disk sector stores some data. Sectors of empty new drives would be mostly initialized with a pattern of zeroes, except for partitioning tables and file system metadata.

Most recent drives use 4KiB sized sectors, also known as Advanced Format, but still provide backward compatibility with older systems which expect 512B sector size with 512B sector size emulation.[1]

Sector byte pattern is a specific configuration of bytes, which would indicate what this sector is used for. For example, a repeated pattern of byte x00 often signals that this sector has not been used yet, that the blocks have been freed by the TRIM command. The TRIM command is used by the software to inform the drive which sectors no longer contain user data in order to increase performance.[2] Or, bytes x55xAA at the end of the sector would signalize a block containing master boot record (MBR). However, while in many cases, analysis for positions of bytes is not as time-intensive as analysis of randomness or single-byte patterns, multiple problems show up:

- Testing for many positions and byte configuration will add up.
- Files with magic bytes may be contained in the first sector where the file is stored, but there is no easy way of telling whether the file simply ends, continues on the next sector, or is placed in a completely different sector.
- If the file is unencrypted, it will mostly get picked up by the randomness analysis.

Most works focusing on detecting patterns of bytes on sectors[3, 4] do it through the lens of forensic analysis and use the filesystem metadata in combination with magic bytes of files to allow the user to find information faster. These, while up an abstraction layer from what this thesis focuses on, can provide beneficial information when identifying common patterns of entire sectors or repeating portions of bytes in a single sector.[3]

1.1.2 Randomness

In order to properly classify all disk sectors, one cannot rely exclusively on byte patterns since files can span multiple sectors and can even be encrypted. In this case, it is possible to check the predictability of byte values or even of single bits.

In order to precisely differentiate random data, the provided samples would need to be in the order of gigabytes, which is far from the provided 512 or 4096 Bytes. However, we can at least get an estimate using the techniques described in this subsection.

Entropy

Shannon's entropy calculates the amount of information in bits provided by each byte value in the sector.[5] For example, the entropy of 8 bits means that every byte value is contained the same number of times (i.e., exactly $\frac{s}{256}$ times). Whereas the entropy value of 0 means that only a single byte value is contained and is repeated through the whole sector. Shannon's entropy can be calculated using:

$$H(S) = - \sum_{i=0}^{255} (P(x_i) \log_2(P(x_i)))$$

Where $P(x_i)$ represents the probability of byte value i (i.e., number of times value i appears in the sector divided by the number of all bytes in the sector). Which can be then normalized:

$$\mu(S) = \frac{H}{H_{max}} = -\frac{1}{8} \sum_{i=0}^{255} (P(x_i) \log_2(P(x_i)))$$

Where s is equal to the sector size in bytes. Normalized Shannon's entropy ranges from 0, the least random (a single repeated byte value), to 1, the most random (every byte value is contained in the sector an equal amount of times). Using this value, one can estimate whether the sector contains encrypted data.

However, multiple problems arise when using Shannon's entropy. There is no simple line where all sectors with a higher entropy are encrypted, and all with lower entropy are not. That means that most sectors containing compressed file formats like videos, jpeg images, or zip files will be almost indistinguishable from encrypted sectors by entropy. Another problem is that Shannon's entropy completely disregards the order of values. For example, simple counting up (x00 x01 ... xFE xFF) repeatedly, which is often part of files, results in the entropy of 1, despite this clearly not being random.

Most works I found that attempted to use entropy calculation to classify small data samples used Shannon's entropy despite its drawbacks mentioned above. However, each work aimed to use the calculated entropy differently. Some used[3] or tried to use[4] it to classify blocks for use in file carving and not encryption detection.

Other works used[6] or tried to use[7] entropy calculation as input or part of the input for machine learning trained to classify network packets. Work[6] also suggested using Tsallis entropy for calculation. However, the work did not attempt to calculate Tsallis entropy and instead decided to focus on Shannon's entropy.

Another work worthy of consideration[8] compared multiple entropy estimation algorithms. The work concluded by recommending the Miller-Madow method for uniform byte value distributions to estimate entropy. Entropy estimation will be helpful when considering the efficiency and speed of the entropy calculation.

Chi-squared test

The chi-squared test or χ^2 test is used to determine whether or not the data fit our expectations.[9] For example, consider flipping five fair coins and counting flipped heads. The probability distribution of the results (assuming that the coins cannot land on their side) would look like this:

number of heads	0	1	2	3	4	5
probability	$\frac{1}{32}$	$\frac{5}{32}$	$\frac{10}{32}$	$\frac{10}{32}$	$\frac{5}{32}$	$\frac{1}{32}$

First, we select a significance level (e.g. $\alpha = 0.05$). Then, after repeating the experiment of flipping five coins 160 times and adding up the results, we get the following table:

number of heads	0	1	2	3	4	5
number of flips (X)	2	8	34	64	44	8
expected number of flips (E)	5	25	50	50	25	5

Given counts of variables X_i , expected counts of variables $E_i = 160 * \text{probability}$, and the number of columns n chi-square test statistic can be calculated using:

$$\chi^2 = \sum_{i=0}^{n-1} \left(\frac{(X_i - E_i)^2}{E_i} \right)$$

After getting the value approximation, the corresponding value from the cumulative chi-squared distribution for $n - 1$ degrees of freedom represents how likely was the measured data is from the distribution of our null hypothesis.

So, for the coin example, the chi-square statistic is calculated:

$$\chi^2 = \frac{(2 - 5)^2}{5} + \frac{(8 - 25)^2}{25} + \dots + \frac{(8 - 5)^2}{5} = 38.64$$

After calculating the image of 38.64 under the chi-square cumulative distribution function for $n - 1$ degrees of freedom, we can see that $F_5(38.64) = 0.9999997$, which means that the p-value $= 1 - 0.9999997$ is smaller than our α , and we can therefore reject our null hypothesis, meaning that the fact that the in the measured data fair coins were used is less than 5%. And indeed, the obtained counts are from tests using two fair and three rigged coins with the probability of getting heads of $\frac{2}{3}$. After calculating the chi-square statistic for the hypothesis with rigged coins, we get $\chi^2 = 4.141$ and $F_5(4.141) = 0.4707$ and since $1 - 0.4707 > \alpha$, this hypothesis cannot be rejected.

As the chi-square test is only an approximation and gets more precise with more data, expected values should be at least 5, and it is preferable that they are much higher. [10]

For detection of random numbers, it is possible, for example, to create a column for each possible number, a column for ranges of numbers, or create a column for each remainder after division by a preselected number. Since the distribution of truly random numbers should be uniform, the expected value (E_i) should be the same for all columns. When the null hypothesis of uniformity of the numbers with sufficiently small α gets rejected, we can assume that the numbers are not random enough.

However, the same counting-up problem as with Shannon's entropy arises. When each number is represented the same number of times, the chi-squared statistic for uniform distribution is equal to zero. The image of 0 for any number of degrees of freedom is 0, giving the p-value of 1, meaning the hypothesis cannot be rejected. As having the random numbers very close to the expected distribution is very unlikely, it is possible to counteract this by also rejecting when the $pvalue > 1 - \alpha$. [10] This way, all number sequences the chi-squared statistics of which would be less likely than α will be rejected.

1.2 Visualization

After classifying all disk sectors based on byte patterns and entropy, it all comes down to visualizing the gathered data. While it would be certainly possible to draw a histogram of all sectors' entropy values or a pie graph based on detected patterns, this would not be as illustrative as the chosen approach, and much of the information about sector position in the disk would be lost. That is why the resulting utility visualizes the data using a bitmap, where each pixel represents a single sector on a disk.

Many works which were visualizing data used the most straightforward technique of *sweeping*. [11, 12, 13] This means that the first pixel is placed in the top-left corner, and each following pixel is placed to the right of the previous one except for when the position exceeds the fixed width of the image. In that case, the pixel is placed on the left-most position on the following line. This technique can be very illustrative in cases when the disk contains long sequences of equally classified sectors. However, when the disk would contain a shorter

sequence, this would produce only a horizontal line with a single-pixel width, which could be hard to see and easily overlooked.

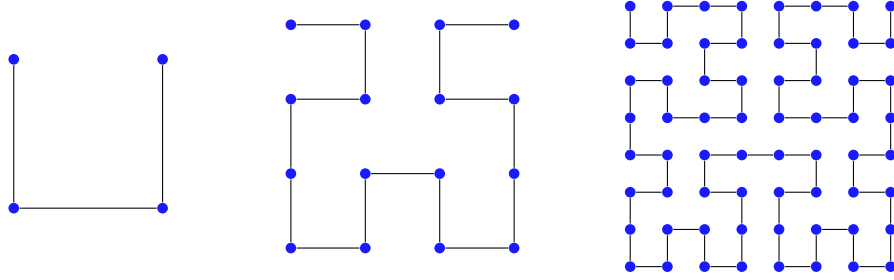


Figure 1.1: the first three Hilbert curve iterations

That is why the work [14] used the more complex Hilbert space-filling curve. The Hilbert curve passes through every pixel in a square exactly once in such a recursive pattern which always keeps consecutive pixels next to each other.[15] Moreover, placing pixels in these specific ways ensures that the shorter sequences are expanded into multiple lines and aggregated into clusters which makes them more easily visible. The curve covers a square with the side length of 2^i pixels (i.e., a total of 4^i pixels) where i is the number of iterations.

However, when using the Hilbert curve, another problem arises. There is no intuitive way to tell where the visualized sectors are located in the source image.

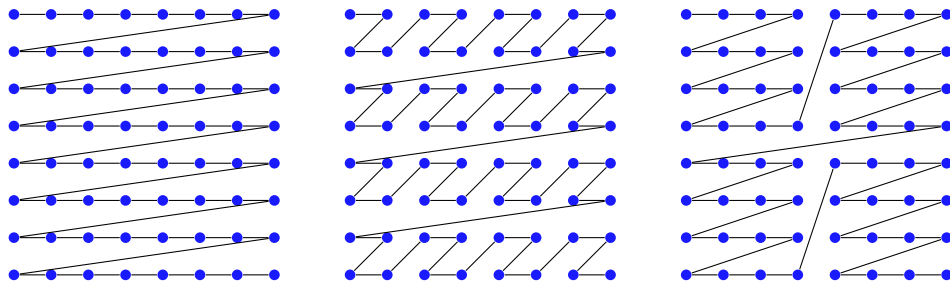


Figure 1.2: sweeping, 2x2 block sweeping, and 4x4 block sweeping

A middle ground between simple sweeping and the Hilbert curve would be block-sweeping. Block-sweeping uses the sweeping method to fill up a square $N \times N$ pixels in size, then continues to another $N \times N$ pixel block and places these pixel blocs in the same way simple sweeping would place individual pixels. This means that sweeping is just a version of block-sweeping, with pixel blocks of 1×1 pixels. By employing this technique, most shorter sequences are still more pronounced by getting expanded into multiple lines, and the position of pixels in the image more closely resembles the sector position in the source image than in the Hilbert curve. However, same as with sweeping, consecutive sectors are not always guaranteed to be right next to each other.

2 Used tools

2.1 Pillow

Pillow[16] is an image manipulation library for Python, which is a fork of the discontinued library PIL[17]. This library is used to visualize the analysis results. Since the results should be visualized as an image, where each pixel represents a single disk sector, statistical visualization libraries like Matplotlib[18], seaborn[19], or Gnuplot[20] were not good choices as they were not created with this exact type of visualization in mind. While they provide the means to create such visualizations, they are not as straightforward as the means provided by Pillow. While Pillow offers many more features beyond the very basics needed, it still keeps the interface for drawing one pixel at a time very simple.

2.1.1 Image

The module `PIL.Image` provides an essential toolkit for manipulating images. Given the image mode (e.g., RGB or RGBA) and image size, function `Image.new` creates an instance of the `PIL.Image.Image` class. The `Image` class stores the state of the resulting image and can be modified using its methods.

The method `Image.putpixel` modifies the state of the `Image` object and changes the color of the pixel on the given coordinates to the given color. `Image.save` tries to store the image on the provided path, and the method `Image.close` releases allocated memory. [21]

2.1.2 ImageDraw

The module `PIL.ImageDraw` is used to modify the `Image` class from `PIL.Image` in more powerful ways than just changing single pixels. Function `ImageDraw.Draw` creates a special context object for the given `Image` object, which can be used for further in-place modifications. The class of the context object `ImageDraw.ImageDraw` provides a wide range of shape drawing methods.

The method `ImageDraw.rectangle`, allows for drawing rectangle for provided coordinates and colors. It is also possible to specify the

width and color of the rectangle outline. The method `ImageDraw.text` allows for writing a provided string in a font on the image. Both of these methods can be used for drawing the image legend. [22]

2.1.3 ImageFont

The module `ImageFont` is used for working with fonts with the Pillow library. It provides the means to load installed fonts by name or from path using the function `ImageFont.load` or to load a fallback font in case no other font is found with function `ImageFont.load_default`. Both of these functions return an `ImageFont.ImageFont` object, which contains a method `ImageFont.getsize` for calculating the dimensions in pixels of the box occupied by provided text written in this font. [23]

2.2 Scipy

SciPy[24] is a Python library for scientific computing. This library is used to calculate the inverse cumulative distribution function for the chi-squared distribution. SciPy provides a couple of different modules, but only one was used in the implementation.

2.2.1 stats

The module `stats` can be used to calculate values of many discrete[25] or continuous[26] statistical distributions, including the chi-squared distribution. The module provides callable object `stats.chi2.ppf`, which calculates the percentile point function (inverse cumulative distribution function) for given probability and degrees of freedom.[27] The resulting value is then used as the threshold or limit for when the results of the chi-square statistic are significant and should therefore be marked as *too random* or *not random*.

3 Implementation

...description of the chapter...

The utility reads the provided disk image in blocks of the provided block size, analyzes the read data, and passes the results alongside information about the position of the read block to the visualization class.

3.1 Analysis

...describe section and the fact that it is a class...

3.1.1 Entropy

The provided buffer's Shannon entropy can be calculated using the formula mentioned in [Entropy](#). However, since $P(x_i) = \frac{c_i}{s}$ where c_i is the number of times the value i was in the buffer and s is the unchanging sector size, the formula for normalized entropy can be simplified to

$$-\frac{1}{8s} \sum_{i=0}^{255} (c_i * \log_2(c_i)) + \log_2(s)$$

The counts c_i are calculated using `Counter` from the `collections` module[28]. The fact that the `Counter` object only iterates over nonzero counts does not matter. By the entropy definition, the impossible results should not affect it.

However, Shannon's entropy is not a statistical test, and its results are not used to distinguish between random and not random. Therefore there are no levels of significance, and the analysis result is just a number from the range 0-1. While this analysis is not helpful to recognize random from non-random data at a glance, it can illustrate whether sectors with similar entropy (and therefore with similar file contents) are stored next to each other. Image generated with Shannon's entropy can be found [HERE](#).

3.1.2 Chi-square

The implementation includes multiple options for analyzing by chi-square test. Each option takes n consecutive bits and treats it as a separate number. If the sector size is not a multiple of n , the remaining bits are ignored.

The chi-square statistic can be calculated using the formula mentioned in **Chi-squared test**. As the numbers should be from the uniform distribution and the expected count for each value should be the same, the formula can be modified to

$$\frac{1}{E} \sum_{i=0}^{2^n-1} ((X_i - E)^2)$$

where E is the expected count of each number and can be calculated using

$$E = \frac{\lfloor \frac{8s}{n} \rfloor}{2^n}$$

where s is the sector size in bytes.

Each option, while easily generalizable, is implemented separately due to possible speed improvements. (i.e., there is no need for reading a byte bit per bit and then merging the bits into an eight-bit number again, however for $n = 3$, there is no such easy workaround) `chi2-8` simply counts each byte value, `chi2-4` separates the byte values into halves using the python bitwise `&` and `»` operators, `chi2-3` goes through each byte bit-by-bit, and `chi2-1` uses the `int.bit_count` method to get the number of bits set to 1 in the byte value.

Each method first precalculates limits for the chi-square statistic, for which the sector is marked as *suspiciously random*, *random*, or *not random*. Then for each sector, its statistic is compared to these limits, and a matching output flag is selected. All methods also check for single-byte patterns. The most common ones (i.e., `0x00` and `0xFF`) are always checked. However, all other single-byte patterns are checked only by `chi2-8` and `chi2-4`. These two implementations were already counting the number of bytes, and this check could be done with a minimal additional performance penalty. Finding the other single-byte patterns is not essential as they are not likely to be contained unless they are part of an unencrypted stored file, which will almost

always be marked as *not random* by the analysis. The only exception is with `chi2-1` for the seventy byte patterns, which contain an equal number of zeroes and ones, in which case the sector will be marked as *suspiciously random*.

4 Results

5 Conclusion

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