

# MORE ON SIMPLICIAL CATS

Recall from last week:

- Monoidal categories, functors, nat. transformations, etc.
- Enriched categories, functors, nat. transformations, etc.

Lemma. We have a 2-functor

$$\underline{\text{MonCat}} \longrightarrow \underline{\text{Cat}}$$

$$\Phi: V \rightarrow V^1 \rightsquigarrow \Phi_+: \underline{\text{Cat}_V} \longrightarrow \underline{\text{Cat}_{V^1}} \text{ a 1-functor}$$

$$\text{ob } \Phi_* \mathcal{C} = \text{ob } \mathcal{C}, \quad \text{Hom}_{\Phi_* \mathcal{C}}(x, y) = \Phi \text{Hom}_{\mathcal{C}}(x, y).$$

This sends monoidal adjunctions to adjunctions.

Def:  $\underline{\text{Cat}_\Delta}$  is the category of simplicial categories, ie. categories enriched over  $sSet$ .

Lemma:  $\underline{\text{Cat}_\Delta}$  is bicomplete.

Recall from earlier: we have adjunctions

$$\begin{array}{ccc} \underline{sSet} & \begin{array}{c} \xleftarrow{c} \\[-1ex] \xrightarrow{\perp} \end{array} & \underline{\text{Set}} \\ & \text{π}_0 - \text{connected components} & \\ & \text{ev}_0 - \text{evaluation at } [0] & \text{constant simplicial set} \end{array}$$

Lemma. These all of these functors are monoidal, and so are the adjunctions, so we get adjunctions

$$\begin{array}{ccc}
 & \text{top} & \\
 \text{Cat}_\Delta & \begin{array}{c} \xleftarrow{\perp} \\[-1ex] \xrightarrow{\perp} \end{array} & \text{Cat} \\
 \text{Horn}(\Delta^0, -)_* & C = C_* & \xleftarrow{\perp} \\
 & \text{bottom} & \\
 & u = (ev_0)_* & \text{underlying category}
 \end{array}$$

$\pi = (\pi_0)_*$  - homotopy category  
 $\perp$   
 $C = C_*$   
 $\perp$   
 $u = (ev_0)_*$  - underlying category

Idea of proof: both sSet and Set are cartesian ( $\otimes = \times$ ), so we have canonical maps

$$\begin{array}{ccccc}
 X \times Y & \xrightarrow{\quad \quad} & X & \rightsquigarrow & F(X \times Y) \\
 \downarrow & & \downarrow & & \downarrow \\
 Y & \xrightarrow{\quad \quad} & F(X \times Y) & \xrightarrow{\quad \quad} & F(X) \times F(Y)
 \end{array}$$

- $C$  and  $ev_0$  are right adjoints, so they preserve products.
- For  $\pi_0$ , consider

$$\begin{array}{ccc}
 (X \times Y)_0 & \xrightarrow{\cong} & X_0 \times Y_0 \\
 \downarrow & & \downarrow \\
 \pi_0(X \times Y) & \longrightarrow & \pi_0 X \times \pi_0 Y
 \end{array}$$

must be surjective  $\nwarrow$  also injective since  
 we can lift relations,  
 because  $ev_0$  is monoidal.  $\square$

Terminology. Let  $\mathcal{C} \in \text{Cat}_\Delta$ , a morphism  $x \rightarrow y$  in  $\mathcal{C}$  is a morphism  $x \rightarrow y$  in  $u\mathcal{C}$ , i.e. a 0-simplex of  $\text{Hom}_{\mathcal{C}}(x, y)$ . A morphism in  $\mathcal{C}$  is an equivalence if its image in  $\pi\mathcal{C}$  is an isomorphism.

Next: defining the nerve of a simplicial category.

We do this by constructing a cosimplicial object in  $\underline{\text{Cat}}_\Delta$ .

$$\Delta \rightarrow \underline{\text{Cat}}_\Delta$$

**Def.** Let  $J$  be a finite non-empty totally ordered set,  $i, j \in J$ .

$$\begin{aligned} P_{i,j} &= \{ I \subseteq J : i, j \in I \text{ and } k \in I \Rightarrow i \leq k \leq j \} \\ &= \{ \text{subsets of } [i,j] \text{ containing the endpoints} \}. \end{aligned}$$

$P_{i,j}$  is a poset under inclusion.

If  $i > j$ , then  $P_{i,j} = \emptyset$ . If  $i \leq j \leq k \in J$ , then

$$N(P_{i,j}) \times N(P_{j,k}) \longrightarrow N(P_{i,k})$$

$$(I, I') \mapsto I \cup I'$$

is an associative binary operation.

[n]

For each  $J$ , we have a simplicial category  $C[\Delta^J]$ :

$$\text{ob } C[\Delta^J] = J, \quad \text{Hom}_{C[\Delta^J]}(i, j) = \begin{cases} \emptyset & \text{if } i > j \\ N(P_{i,j}) & \text{o/w.} \end{cases}$$

Composition is induced by the binary operation above.

**Lemma.** Let  $n \geq 1$ ,  $J = [n]$ . Then

$$(a) N(P_{0,n}) \cong (\Delta^n)^{n-1}$$

$$(b) P_{0,j} \cong P_{0,j-1}.$$

**Proof.** (b) is clear. For (a), suffice to show

$$\begin{aligned} P_{0,n} &\cong \underbrace{[1] \times \cdots \times [1]}_{n-1 \text{ times}} && \text{an element of this represents} \\ &\uparrow && \text{a subset of } (0, n) \text{ in } [n]. \\ s &\mapsto (x_s(1), \dots, x_s(n-1)) \end{aligned}$$

**Lemma.** If  $\mathcal{C}$  is a category with an initial or terminal object, then  $N(\mathcal{C})$  is contractible.

$$\emptyset \in \mathcal{C} \text{ initial} \Rightarrow \emptyset \rightarrow 1_{\mathcal{C}}$$

$$\mathcal{C} \times [1] \rightarrow \mathcal{C} \rightsquigarrow N(\mathcal{C}) \times \Delta^1 \rightarrow N(\mathcal{C})$$

**Cor.** For  $i \leq j$ ,  $\text{Hom}_{\mathcal{C}[\Delta^J]}(i, j)$  is contractible.

**Lemma** There is a unique isomorphism  $\pi: \mathcal{C}[\Delta^n] \cong [n]$  which is the identity on objects. Hence, we have a canonical functor  $\mathcal{C}[\Delta^n] \rightarrow \mathcal{C}[n]$ .

$$\mathcal{C}[\Delta^n]$$

**Lemma.** The assignment  $J \mapsto \mathcal{C}[\Delta^J]$  defines a functor

$$\text{LinordSet} \longrightarrow \underline{\text{Cat}}_{\Delta}$$

In particular, we have a cosimplicial object  $\Delta \rightarrow \underline{\text{Cat}}_{\Delta}$  given by  $[n] \rightarrow \mathcal{C}[\Delta^n]$ .

**Ideas of proof:** Unfolding definitions, it suffices to construct for any monotone map  $f: J \rightarrow J'$  and  $i \leq j \in J$ , a monotone map  $P_{i,j} \rightarrow P_{f(i), f(j)}$ . This is  $I \mapsto f(I)$ .  $\square$

**Def.** For  $\mathcal{C} \in \underline{\text{Cat}}_{\Delta}$ , its simplicial nerve (or homotopy-coherent nerve) is the sSet.

$$N(\mathcal{C})_n = \text{Hom}_{\underline{\text{Cat}}_{\Delta}}(\mathcal{C}[\Delta^n], \mathcal{C})$$

**Lemma.** If  $\mathcal{C} \in \underline{\text{Cat}}$ , then  $N(\mathcal{C}) \cong N(c\mathcal{C})$ .

**Proof:**  $N(c\mathcal{C})_n = \text{Hom}_{\underline{\text{Cat}}_{\Delta}}(\mathcal{C}[\Delta^n], c\mathcal{C})$   
 $\cong \text{Hom}_{\underline{\text{Cat}}}(\pi_* \mathcal{C}[\Delta^n], \mathcal{C})$   
 $\cong \text{Hom}_{\underline{\text{Cat}}}([n], \mathcal{C}) = N(\mathcal{C})_n$ .

Discussion. What do (low dimensional) simplices in  $N(\mathcal{C})$  look like?

- $\mathbb{C}[\Delta^0]$  has a single object and  $\text{Hom}_{\mathbb{C}[\Delta^0]}(0,0) = N(P_{0,0}) = \Delta^0$ . C1
- $\mathbb{C}[\Delta^1]$  has two objects and all hom-sets are  $\Delta^0$ . P\_{0,1}

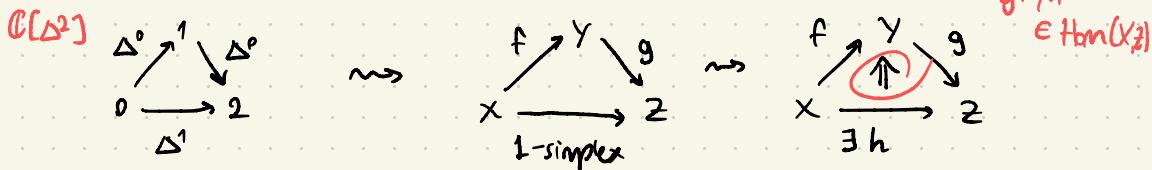
Hence,  $\mathbb{C}[\Delta^0] = \mathbb{C}[0]$  and  $\mathbb{C}[\Delta^1] = \mathbb{C}[1]$ .

$$\bullet N(\mathcal{C})_0 = \text{Hom}_{\underline{\text{Cat}}_\Delta}(\mathbb{C}[\Delta^0], \mathcal{C}) \cong \text{Hom}_{\underline{\text{Cat}}}([0], \mathcal{C}) \cong \text{ob } \mathcal{C}$$

$$\bullet N(\mathcal{C})_1 = \text{Hom}_{\underline{\text{Cat}}_\Delta}(\mathbb{C}[\Delta^1], \mathcal{C}) \cong \text{Hom}_{\underline{\text{Cat}}}([1], \mathcal{C}) \cong \text{mor } \mathcal{C}$$

What is  $\mathbb{C}[\Delta^2]$ ? It has three objects. All hom-sets are  $\Delta^0$  except  $\text{Hom}_{\mathbb{C}[\Delta^2]}(0,2) = N(P_{0,2}) \cong \Delta^1$ . to 23, to 1,23

Hence, a simplicial functor  $\mathbb{C}[\Delta^2] \rightarrow \mathcal{C}$  picks out:



Lemma. There is a unique colimit preserving functor

$$\mathbb{C}[-]: \underline{sSet} \rightarrow \underline{\text{Cat}}_\Delta$$

which sends  $\Delta^n \mapsto \mathbb{C}[\Delta^n]$ . It is left-adjoint to  $N$ .

Fact.  $\mathbb{C}[-]$  preserves monomorphisms, i.e. if  $A \hookrightarrow X$  is an inclusion of sets then  $\mathbb{C}[A] \hookrightarrow \mathbb{C}[X]$ .

Lemma. Let  $0 < j < n$ .  $\mathbb{C}[\Delta_j^n] \subseteq \mathbb{C}[\Delta^n]$  is given by:

(1)  $\text{ob } \mathbb{C}[\Delta_j^n] = \text{ob } \mathbb{C}[\Delta^n]$ .

(2)  $\text{Hom}_{\mathbb{C}[\Delta_j^n]}(i, k) = \text{Hom}_{\mathbb{C}[\Delta^n]}(i, k)$  except

$\text{Hom}_{\mathbb{C}[\Delta_j^n]}(0, n) \subseteq \text{Hom}_{\mathbb{C}[\Delta^n]}(0, n)$

which is given by the subsimplicial set of  $(\Delta^1)^{n-1}$  obtained by deleting the interior of the bottom  $j$ -face.

Proof idea:  $\Delta_j^n = \bigcup_{i \neq j} \Delta^{n-1, i}$  so look at  $\mathbb{C}[\Delta^{n-1, i}] \subseteq \mathbb{C}[\Delta^n]$ .

Lemma.  $N(\mathcal{C})$  is a composet. Further, if all hom-sets in  $\mathcal{C}$  are Kan complexes then  $N(\mathcal{C})$  is an  $\infty$ -category.

Proof:

$$\begin{array}{ccc} I^n & \longrightarrow & N(\mathcal{C}) \\ \downarrow & \curvearrowright & \downarrow \\ \Delta^n & & \end{array} \quad \Leftrightarrow \quad \begin{array}{ccc} \mathbb{C}[I^n] & \longrightarrow & \mathcal{C} \\ \downarrow & & \downarrow \\ \mathbb{C}[\Delta^n] & & \end{array}$$

But we have a retraction of  $\mathbb{C}[I^n] \rightarrow \mathbb{C}[\Delta^n]$ .

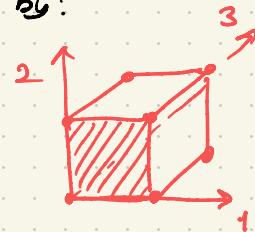
To see this, note  $\mathbb{C}[I^n] = c[n]$  since  $I^n = I^{n-1} \cup_{\Delta^0} I^1$ .

But we saw there is a unique simplicial functor

$\mathbb{C}[\Delta^n] \rightarrow c[n]$  that is the identity on objects. Then

$$c[n] \cong \mathbb{C}[I^n] \hookrightarrow \mathbb{C}[\Delta^n] \rightarrow c[n]$$

is the identity.



For Kan complexes, need  $0 < j < n$

$$\begin{array}{ccc} \Delta_j^n & \longrightarrow & N(\mathcal{C}) \\ \downarrow & \nearrow & \Downarrow \\ \Delta^n & & C[\Delta^n] \end{array} \quad \Leftrightarrow \quad \begin{array}{ccc} C[\Delta_j^n] & \xrightarrow{f} & \mathcal{C} \\ \downarrow & \nearrow & \Downarrow \\ C[\Delta^n] & & \end{array}$$

From the previous lemma, we only need to worry about

$$\begin{array}{ccc} \text{Hom}_{C[\Delta_j^n]}(0, n) & \xrightarrow{f} & \text{Hom}_{\mathcal{C}}(f(0), f(n)) \\ \downarrow & \nearrow & \Downarrow \\ \text{Hom}_{C[\Delta^n]}(0, n) & & \end{array}$$

But this follows from the fact that the vertical maps is **anodyne**. One checks that this extension in fact defines a simplicial functor. □

**Remark.** Nerves of categories enriched over Kan complexes will be  $(\infty, 1)$ -categories?

**Def.** Consider the simplicial category of CW-complexes and hom-sets the singular set of the mapping space. Its simplicial nerve is the  **$\infty$ -category of spaces**, Spec.

**Lemma.** The product and coproduct of  $\infty$ -cats is an  $\infty$ -cat.

**Proof.** For product, find extension in each factor and then put together. For coproduct, note that  $\Delta_j^n$  and  $\Delta^n$  are connected. □

**Def.** A sub- $\infty$ -cat  $\mathcal{C}' \subseteq \mathcal{C}$  is a subsimplicial set determined by subsets  $X \subseteq \mathcal{C}_0$  and  $S \subseteq \mathcal{C}$ , between objects in  $X$  and closed under composition and equivalences. An  $n$ -simplex of  $\mathcal{C}$  belongs in  $\mathcal{C}'$  iff the edges of its restriction to the spine  $I^n$  are in  $S$ . A subcategory is full if  $S$  contains all 1-simplices whose boundary is in  $X$ .

**Lemma.** A sub- $\infty$ -category is an  $\infty$ -cat. Its homotopy category is the subcategory of  $\text{hc}$  on the image of  $S$ .

$$\begin{array}{ccc} \mathcal{C}' & \longrightarrow & \mathcal{C} \\ \downarrow & & \downarrow \\ N(\text{hc}') & \longrightarrow & N(\text{hc}) \end{array}$$

is a pullback. For any subcategory  $\mathcal{D} \subseteq \text{hc}$ , this pullback gives a sub- $\infty$ -cat.

**Cor.** There is a 1-to-1 correspondence between sub- $\infty$ -cats of  $\mathcal{C}$  and subcategories of  $\text{hc}$ . Full correspond to full.

**Def.** A natural transformation between functors  $f, g: \mathcal{C} \rightarrow \mathcal{D}$  is a simplicial map  $\mathcal{C} \times \Delta^1 \rightarrow \mathcal{D}$  that restricts to  $f$  and  $g$ .

**Remark.** This exactly a 1-simplex in  $\text{Hom}(\mathcal{C}, \mathcal{D})$ .

We will see that  $\text{Hom}(\mathcal{C}, \mathcal{D})$  is also an  $\infty$ -cat and  $N(\text{Fun}(\mathcal{C}, \mathcal{D})) = \text{Hom}(N\mathcal{C}, N\mathcal{D})$  since  $N$  is fully faithful.

$$-\times A \dashv \underline{\text{Hom}}(A, -)$$

# Anodyne & (co)Fibration C-sets

Def: A (left/right/inner) fibration  $p: K \rightarrow S$  has the RLP wrt  $\mathbb{A}$  left/right/inner horns

right lifting      eg  $\begin{array}{ccc} \mathbb{A}^n & \xrightarrow{i^n} & X \\ & \downarrow & \curvearrowright \\ \mathbb{A}^m & \xrightarrow{\quad} & * \end{array}$  is an inner fibration iff  $X$  is  $\infty$ -cat

$S \subset \text{Mor}(C)$

$$\chi_R(S) = \{ f: X \rightarrow C \mid f \text{ have RLP wrt } S \} \quad \chi_L(S) \text{ (LRP)}$$

$\chi_R(\{\text{left horns}\}) = \text{Left Fibration.}$

Def: A map  $f: A \rightarrow B$  is (left/right/inner) anodyne if it has LLP wrt  $(\dots)$  fibrations. i.e.  $\chi_L(\chi_R(\{\text{left/right/inner horns}\}))$

$$S \rightarrow S \subset \chi_L(\chi_R(S)) =: \chi(S)$$

Def: A  $S \subset \text{Mor}(C)$  is saturated if:

1) it contains all RPs

2) Stable under pushouts:

$$\begin{array}{ccc} A & \xrightarrow{} & A' \\ S \ni i & \downarrow & \downarrow i' \\ B & \xrightarrow{} & B' \end{array} \Rightarrow i' \in S$$

3) Stable retracts: id

$$\begin{array}{c} A' \xrightarrow{\quad} A \xrightarrow{\quad} A' \\ \downarrow i \quad \downarrow i' \quad \downarrow i'' \\ B' \xrightarrow{\quad} B \xrightarrow{\quad} B' \end{array} \quad \text{if } i \in S, \text{ so is } i'. \\ \text{if } i \in S, \text{ so is } i''. \quad \text{if } i \in S, \text{ so is } i''.$$

4) Stable under countable composition (N-diagram)

$$\begin{array}{ccccccc} A_0 & \xrightarrow{i_0} & A_1 & \xrightarrow{i_1} & A_2 & \rightarrow \cdots & i_j \in S \quad k_j \\ & \nearrow j_0 & \downarrow j_1 & \downarrow j_2 & & & \\ & & & & & & \\ \Rightarrow \text{es} & \xrightarrow{\quad} & \text{claim } A_k & \xrightarrow{\quad} & & & \end{array}$$

5) Coprode: If  $\{A_j \xrightarrow{i_j} B_j\}_{j \in J} \subset S$

$$\Rightarrow \coprod_i A_j \xrightarrow{\coprod i_j} \coprod_i B_j \in S$$

Def. S<sub>char(C)</sub>. The saturated closure of  $S$  is the smallest saturated set containing  $S$ , i.e.  $\bar{S} = \bigcap_{\substack{T \text{ saturated} \\ T \supseteq S}} T$

Lem  $S \in \text{Schar}(C)$ . Then,  $\mathcal{X}_L(S)$  is saturated.

Pr:

$$\begin{array}{ccc} A & \xrightarrow{\quad} & X \\ \downarrow \sim & \nearrow \sim & \downarrow f \in S \\ B & \xrightarrow{\quad} & Y \end{array}$$

$$\begin{array}{ccc} A & \xrightarrow{\quad} & A' \xrightarrow{\quad} X \\ \downarrow \sim & \nearrow \sim & \downarrow \sim \\ B & \xrightarrow{\quad} & B' \xrightarrow{\quad} Y \end{array}$$

$$\begin{array}{ccccc} A & \xrightarrow{id} & A' & \xrightarrow{\quad} & X \\ \downarrow \sim & \nearrow \sim & \downarrow \sim & \nearrow \sim & \downarrow f \in S \\ B & \xrightarrow{id} & B' & \xrightarrow{\quad} & Y \end{array}$$

$$\begin{array}{ccc} A_0 & \xrightarrow{i_0} & X \\ \downarrow \sim & \nearrow \sim & \downarrow \sim \\ (A_1 \xrightarrow{i_1} \text{colim } A_i \xrightarrow{y}) & \xrightarrow{\quad} & Y \\ \downarrow \sim & \nearrow \sim & \downarrow \sim \\ A_2 & \xrightarrow{i_2} & \vdots \end{array}$$

• similar to coprode

Prop: (Small Object Argument)  $c = \text{ssets}$

□

Let  $S = \{A_i \rightarrow B_j\}_{i \in I, j \in J}^{\text{small}} \subset \text{Mor}(C)$  s.t.  $A_i$ 's have finitely many non-deg. simplices  
Then any  $F: X \rightarrow Y$  in sssets can be factorized.

$$X \xrightarrow{F} Y$$

$$\xrightarrow{w} Z \xrightarrow{g}$$

s.t.  $w \in \bar{S}$ ,  $g \in \mathcal{X}_R(S)$ .