

CS202

Data Structures and Algorithms

Assignment 1

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Report for sorting algorithms

Question1.

1) **Insertion Sort**

1) **Pseudo Code:**

1. for $j=2$ to n
2. $key=A[j]$
3. $i=j-1$
4. while($A[i]>key$ and $i>0$)
5. $A[i+1]=A[i]$
6. $i=i-1$
7. $A[i+1]=key$

2) **Time Complexity analysis:**

Best Case Running Time: $O(n)$

Worst Case Running Time: $O(n^2)$

Average Case: $O(n^2)$

Space Complexity: $O(1)$

3) Remarks:

- 1) Good for sorting small arrays.
- 2) In place and stable
- 3) Not good for very large numbers as n^2 complexity.

2. Merge Sort

A) Pseudo Code:

```

1. Merge_Sort(A,p,r)
2. If p<r
3.   q=(p+r)/2
4.   Merge_Sort(A,p,q)
5.   Merge_Sort(A,q+1,r)
6.   Merge(A,p,q,r)
7. Merge(A,p,q,r)
8.   i=p,j=q+1,k=0,tem[r-p+1]
9.   while(i<q and j<=r)
10.    if(A[i]<A[j])
11.        tem[k]=A[i]
12.        k+=1
13.        i+=1
14.    else()
15.        tem[k]=A[j]
16.        j+=1
17.        k+=1
18.    for(g=0 to n)
19.        A[g]=tem[g]
```

B) Complexity Analysis:

Best Case Running Time: $O(n \log n)$
Worst Case Running Time: $O(n \log n)$
Average Case: $O(n \log n)$
Space Complexity: $O(n)$

C) Remarks:

- 1) Not for large arrays as space complexity $O(n)$.
- 2) Stable algorithm

3) Good point is it always works in $O(n \log n)$ time complexity.

3. Quick Sort

A) Pseudo Code:

- 1) quickSort(A[],low,high)
- 2) if(low<high)
- 3) q=partition(A,low,high)
- 4) quickSort(A,low,q-1)
- 5) quickSort(A,q+1,high)
- 6) partition(A,low,high)
 - a) pivot=A[high]
 - b) i=low-1
 - c) j=high+1
 - d) while(true)
 - e) j=j-1
 - f) while(A[j]>pivot)
 - g) i=i+1
 - h) while(A[i]<pivot)
 - i) if(j>i)
 - j) exchange(A[i],A[j])
 - k) Else if j=i
 - l) Return j-1
 - m) Else
 - n) Return j

B) Complexity Analysis:

Best Case Running Time: $O(n \log n)$
Worst Case Running Time: $O(n^2)$
Average Case: $O(n \log n)$
Space Complexity: $O(\log n)$

C) Remarks:

- 1) In place but not stable.
- 2) Better than merge sort as it takes less place
- 3) Can be time taking for large n.

3. Heap Sort

A)Pseudo Code:

- 1) HeapSort(A[])
- 2) BuildHeap(A)
- 3) For i= length downto 2
- 4) Swap A[i] and A[1]
- 5) heapsize-=1
- 6) Heapify(A,1,3)
- 7) BuildHeap(A)
- 8) heapsize=length(A)
- 9) For i=length/2 downto 1
- 10) Heapify(A,i)

Heapify(A,i)

 - 1) left=2i+1
 - 2) right=2i+2
 - 3) If left<=heapsize and A[left]>A[i]
 - 4) largest=left
 - 5) Else
 - 6) largest=i
 - 7) If right<=heapsize and A[right]>A[i]
 - 8) largest=right

- 9) Else
- 10) largest=i
- 11) If largest!=i
- 12) Exchange A[i] and A[largest]
- 13) Heapify(A,largest)

B) Complexity Analysis:

Best Case Running Time: $O(n \log n)$

Worst Case Running Time: $O(n \log n)$

Average Case: $O(n \log n)$

Space Complexity: $O(1)$

C) Remarks:

- 1) Always time complexity is $n \log n$.
- 2) In place but unstable
- 3) Memory efficient as no extra space required