

Priority Queues

Data Structures and Algorithms

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Jinho D. Choi



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Priority Queue

- Priority queue
 - Add: insert a comparable key.
 - RemoveMax: find and remove the key with the highest priority.
- Application
 - Scheduling.
Is the incoming order important for scheduling?
- Implementation
 - Lazy vs. Eager.
 - Heap.



Priority Queue Abstract Class

```
public abstract class AbstractPriorityQueue<T extends Comparable<T>>
{
    abstract public void add(T key);
    abstract public T removeMax();
    abstract public int size();

    public boolean isEmpty()
    {
        return size() == 0; defined?
    }

    visibility?

    protected void throwNoSuchElementException()
    {
        if (isEmpty())
            throw new NoSuchElementException("No key exists.");
    }
}
```

These are called?



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Lazy vs. Eager Algorithms

- Lazy algorithm
 - **Add**: insert each key to the back of the list.
 - **RemoveMax**: Find the key with the highest priority and remove it from the list.
- Eager algorithm
 - **Add**: insert each key to the list in ascending order.
 - **RemoveMax**: remove the last key in the list.

| | Add | RemoveMax |
|-------|----------|-----------|
| Lazy | O(1) | O(n) |
| Eager | O(n) | O(1) |

O(log n)



Eager Priority Queue

```
List<T> l_keys = new ArrayList<>();
```

```
public void add(T key)      - (insertion point) - 1
{
    int index = Collections.binarySearch(l_keys, key);
    if (index < 0) index = -(index + 1);
    l_keys.add(index, key);
}
```

Complexity?

```
public T removeMax()
{
    return l_keys.remove(l_keys.size()-1);
}
```



Binary Heap

- What is a heap?
 - A tree where the key of each node is **larger than or equal to** the keys in its **children**.
 - The tree is guaranteed to be **balanced**.
 - What is a **binary heap**?
- Operations
 - Add: **swim**.
 - RemoveMax: **sink**.
 - Both operations can be done in $O(\log n)$.



Binary Heap Using List

Input =

7 3 2 4 6 1 5



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Binary Heap Using List

Input =

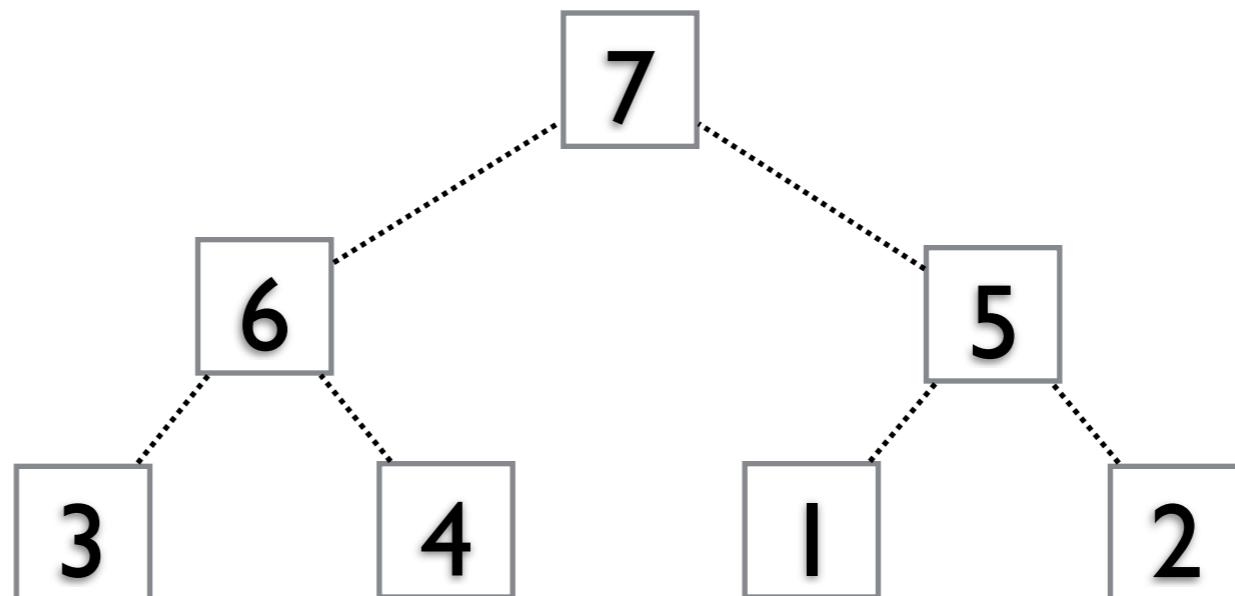
7 3 2 4 6 1 5

Index of the parent?

$k/2$

Index of the children?

$k*2, k*2 + 1$



list =

\emptyset 7 6 5 3 4 1 2

Why?



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Binary Heap Using List

```
List<T> l_keys;
int n_size;

public BinaryHeap()
{
    l_keys = new ArrayList<>();
    l_keys.add(null);
    n_size = 0;
}
```



Add via Swim

Input =

7 3 2 4 6 I 5

Add each key to the **leftmost leaf**.

Keep **swapping** the new node with its parent
until it satisfies the **heap property**.



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Add via Swim

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7



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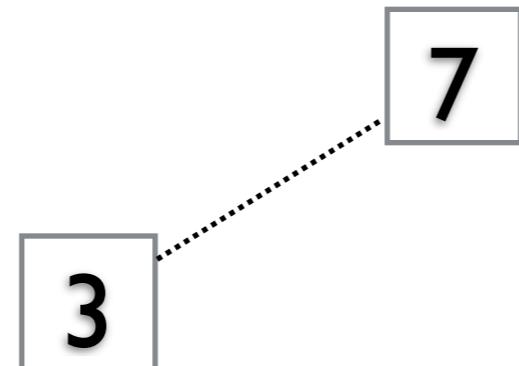
Add via Swim

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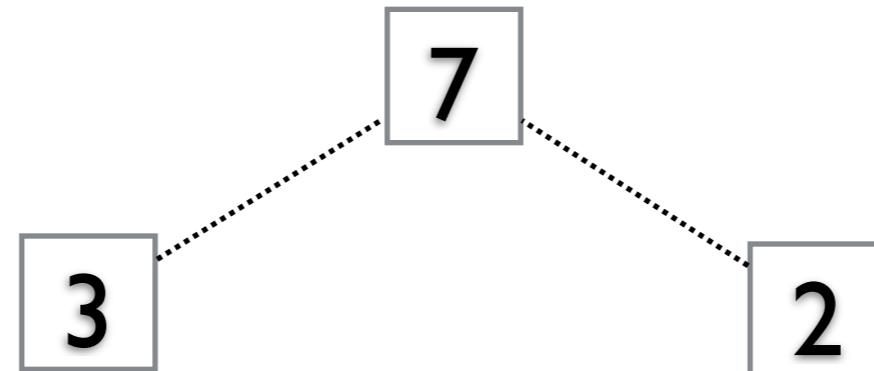
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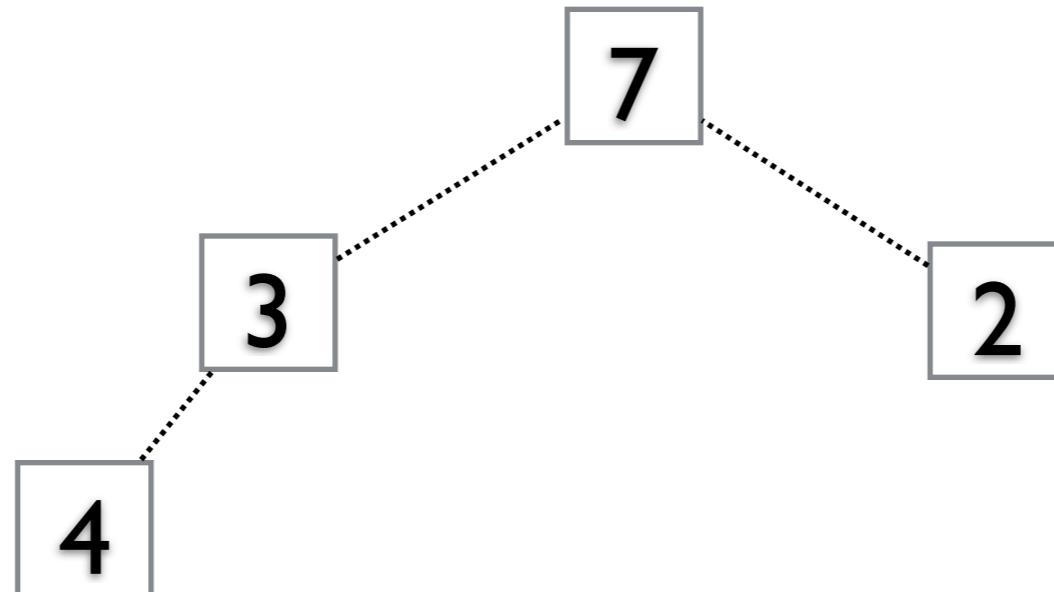


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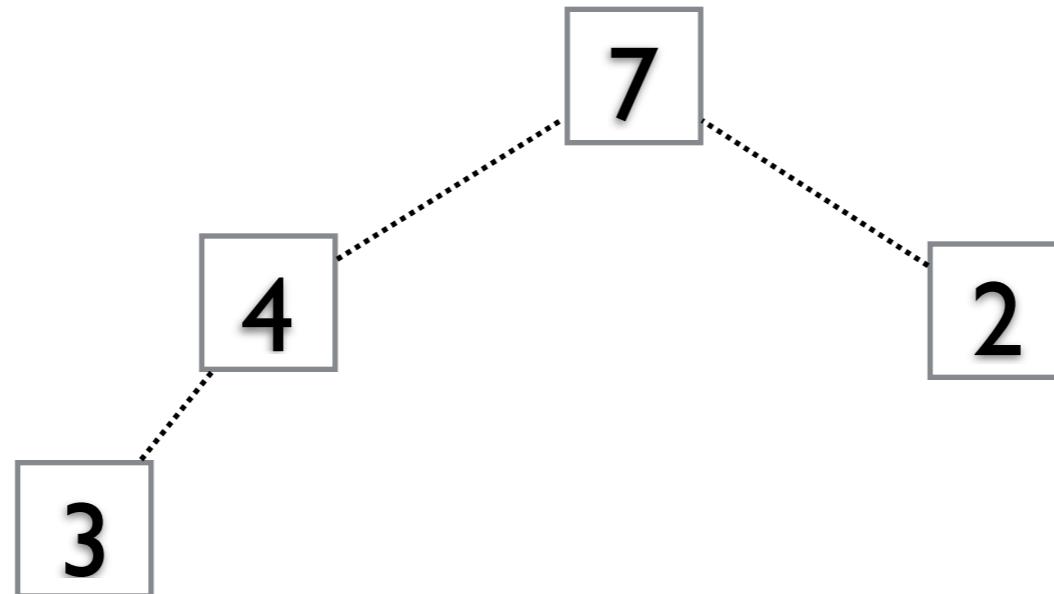


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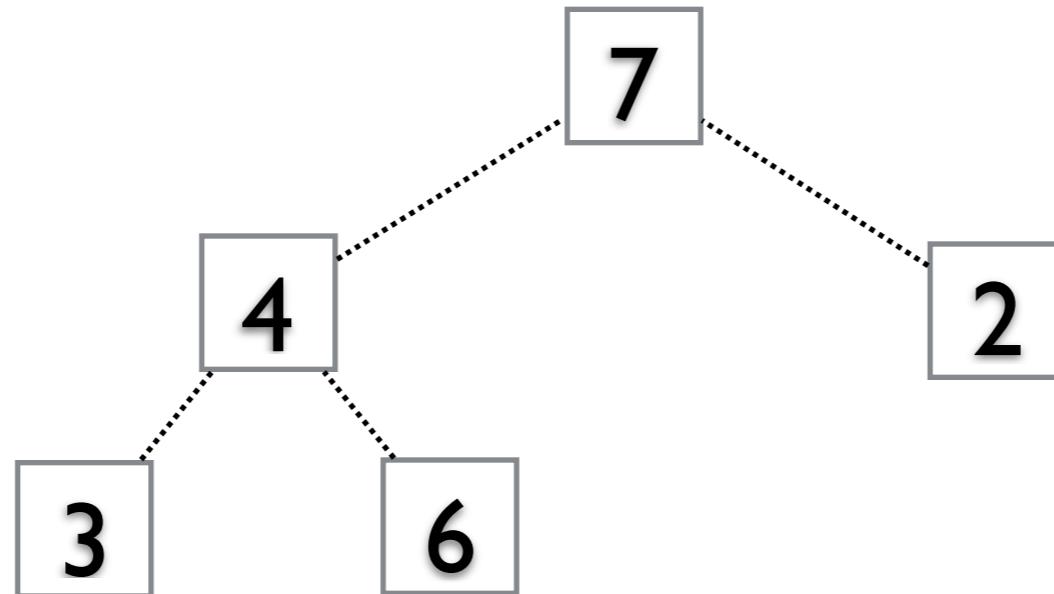


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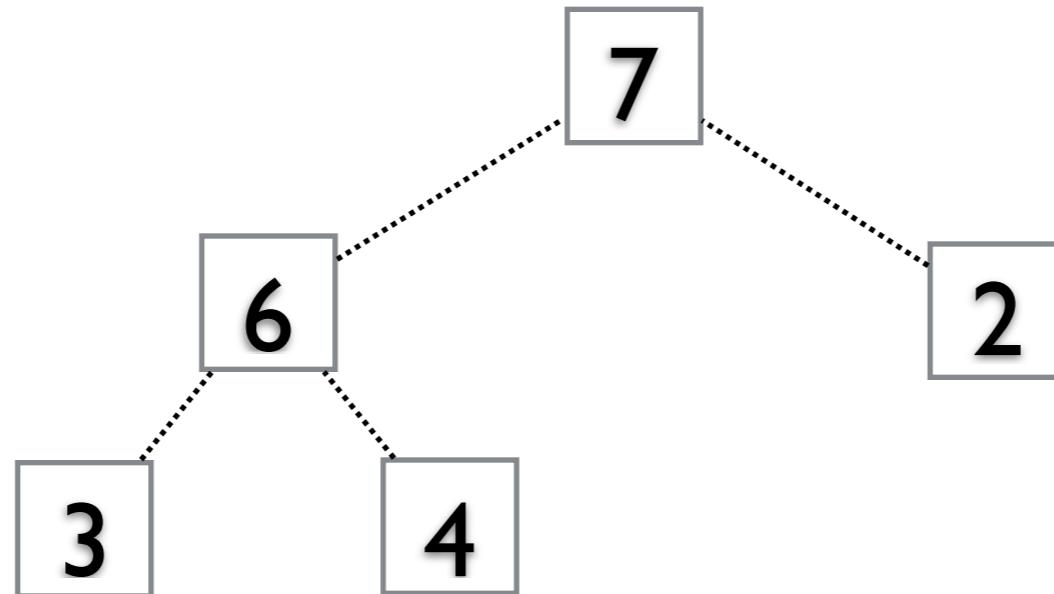


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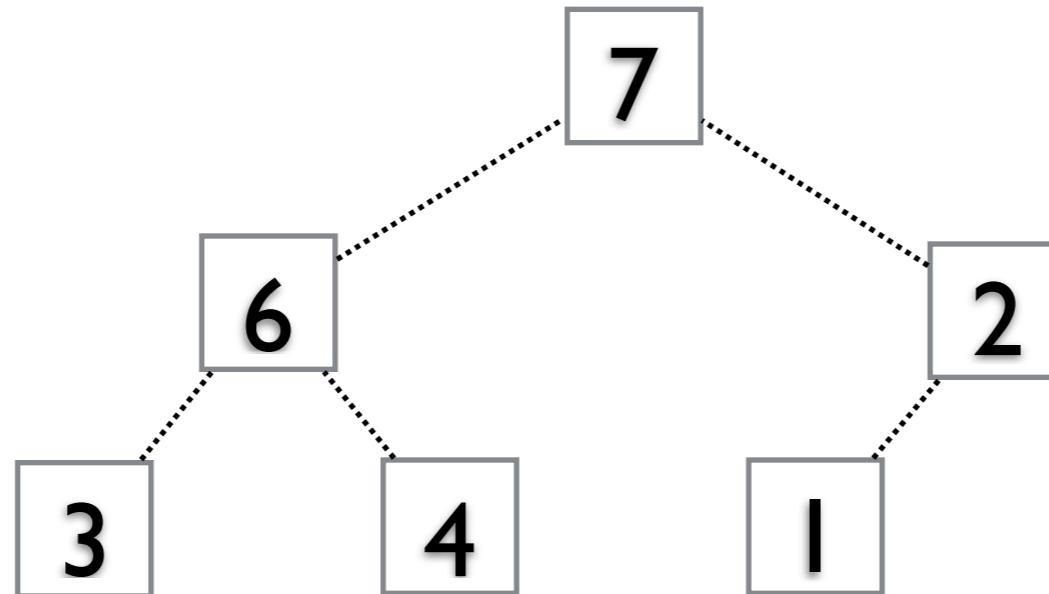


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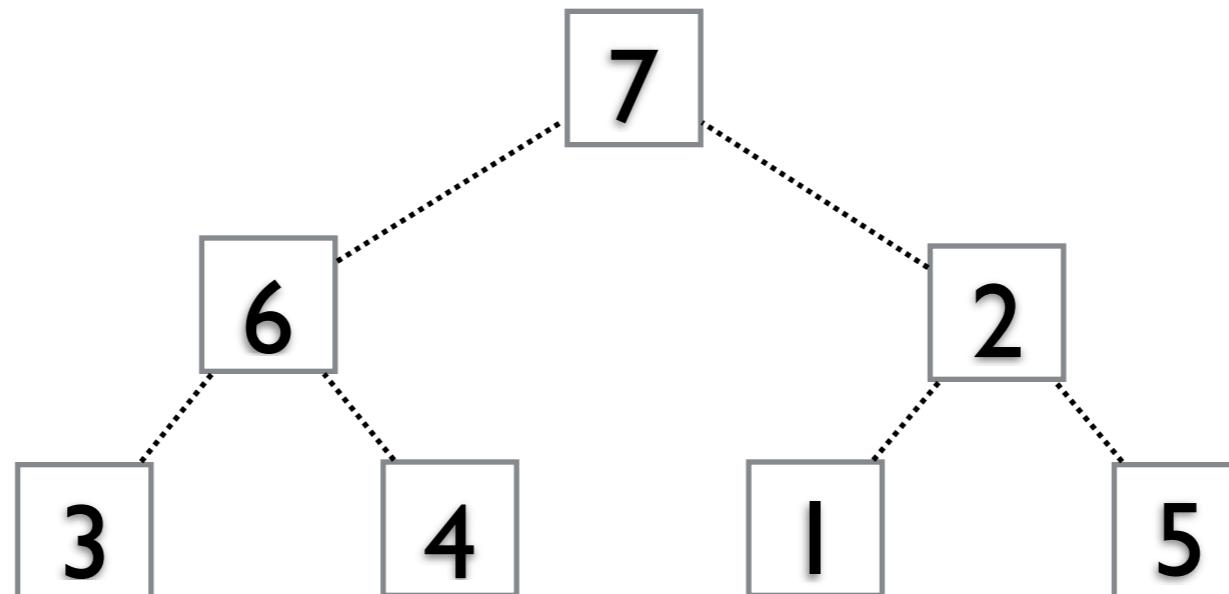


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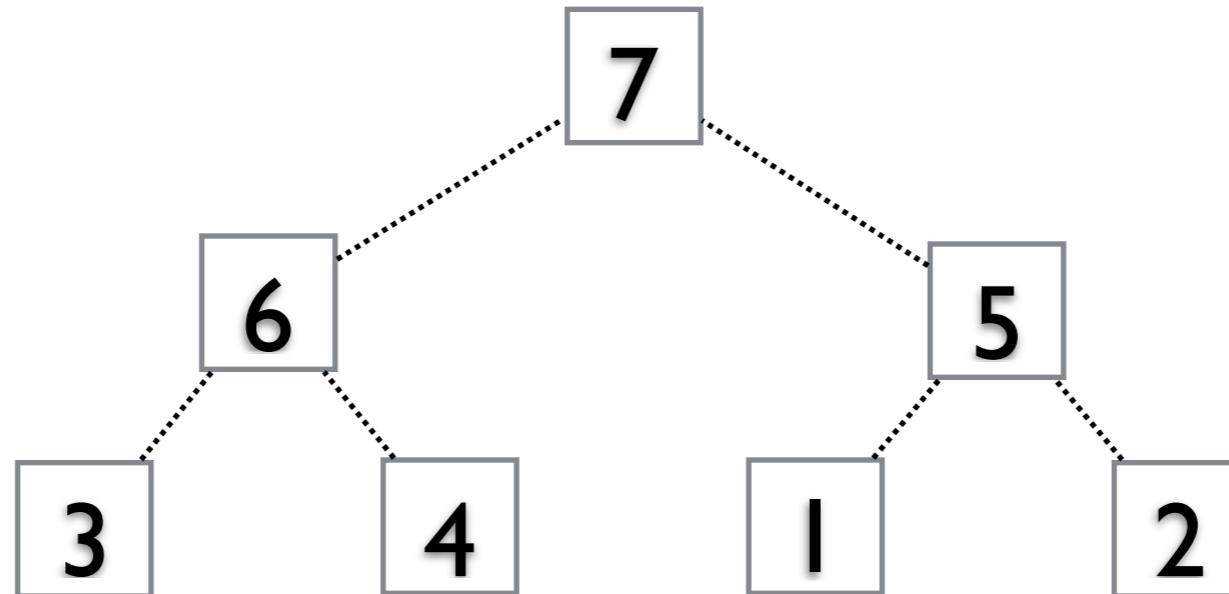


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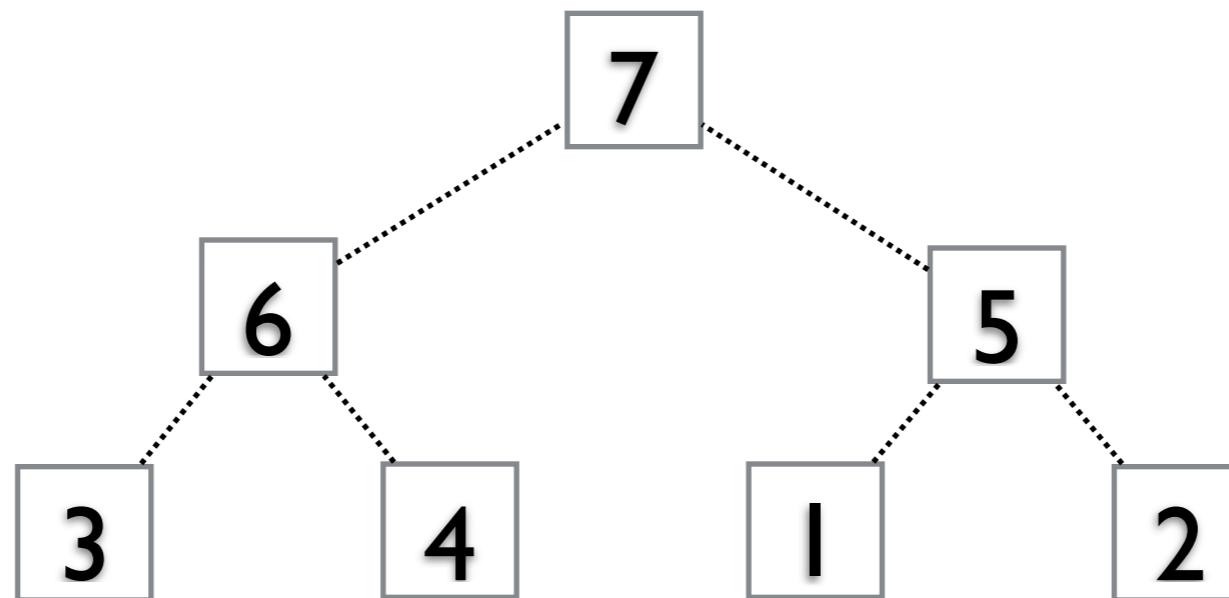


Add via Swim

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Add each key to the **leftmost leaf**.

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Complexity?



Add via Swim

```
public void add(T key)
{
    l_keys.add(key); ← Leftmost leaf
    swim(++n_size);
}
```

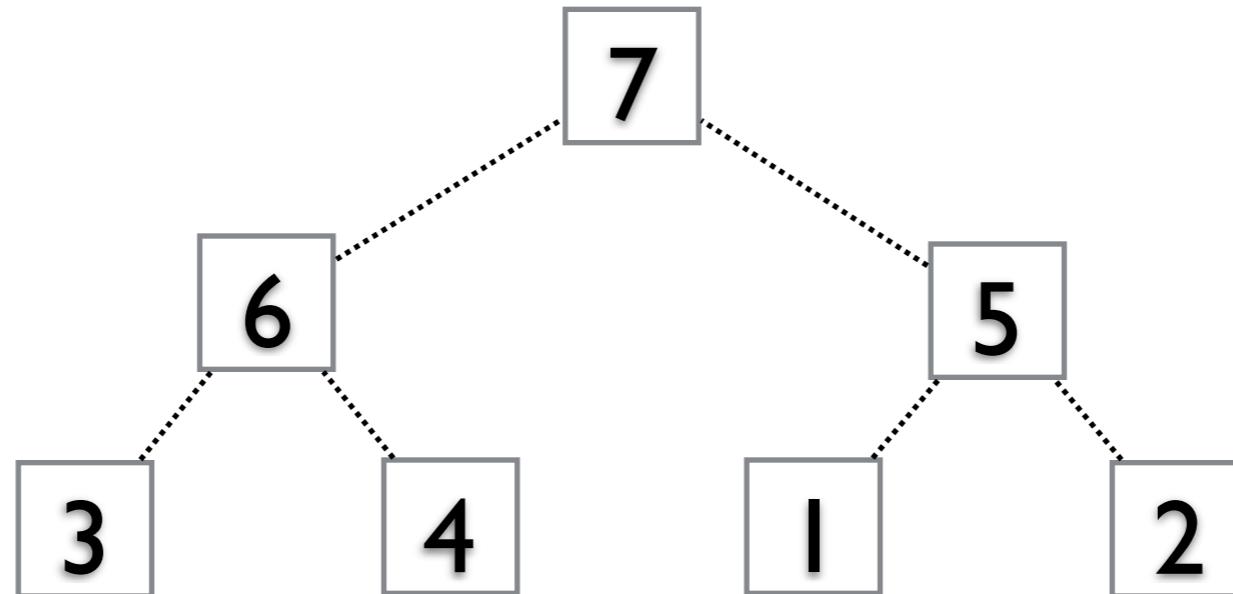
```
private void swim(int k)
{
    while (k > 1 && compareTo(k/2, k) < 0)
    {
        swap(k/2, k); ↑ Parent
        k /= 2;
    }
}
```



RemoveMax vis Sink

Replace the root with the **rightmost leaf**.

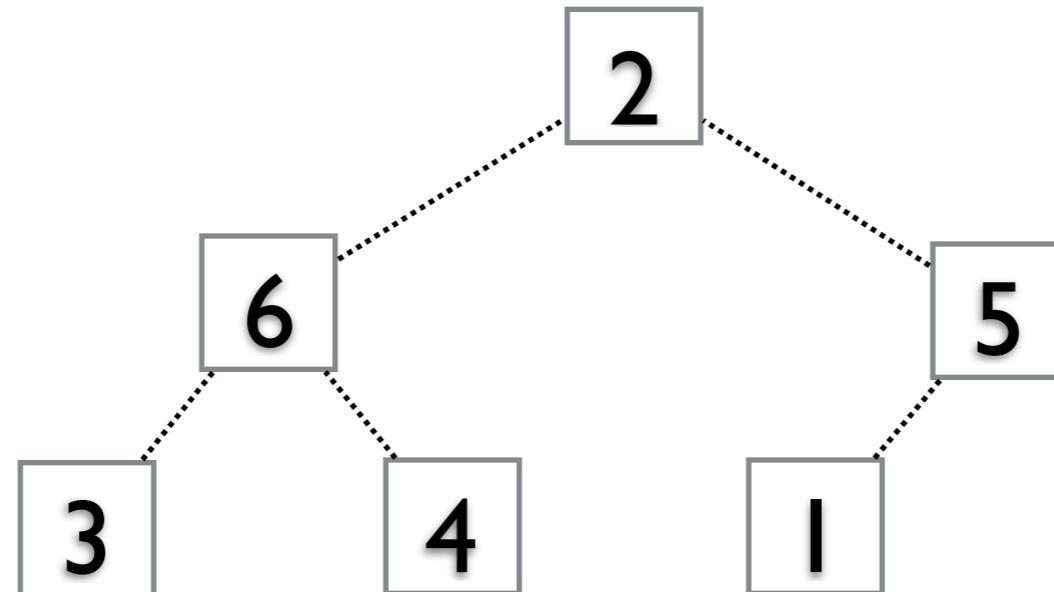
Keep **swapping** the root node with its child until it satisfies the **heap property**.



RemoveMax vis Sink

Replace the root with the **rightmost leaf**.

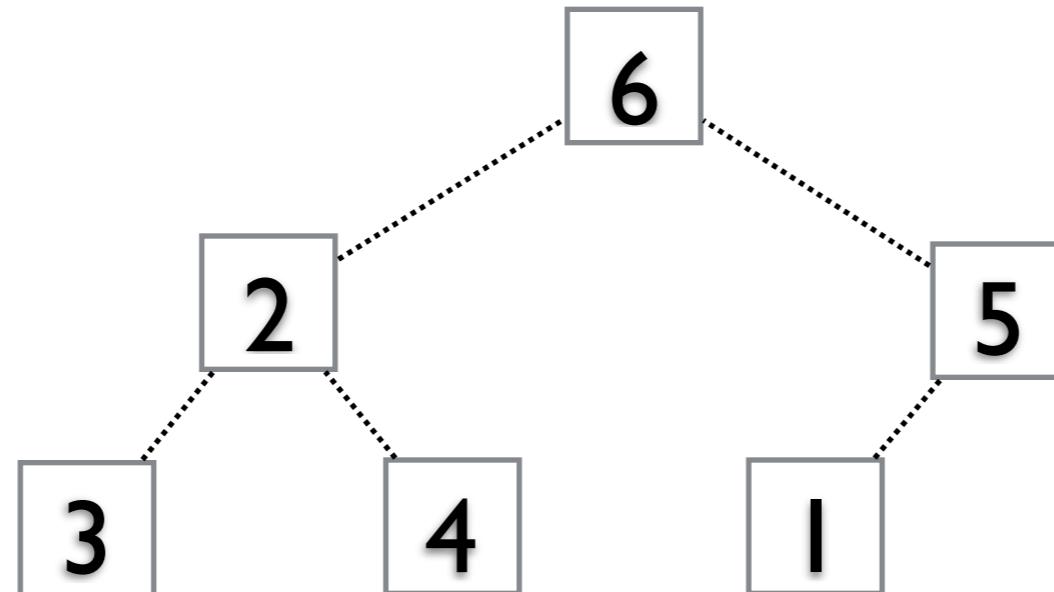
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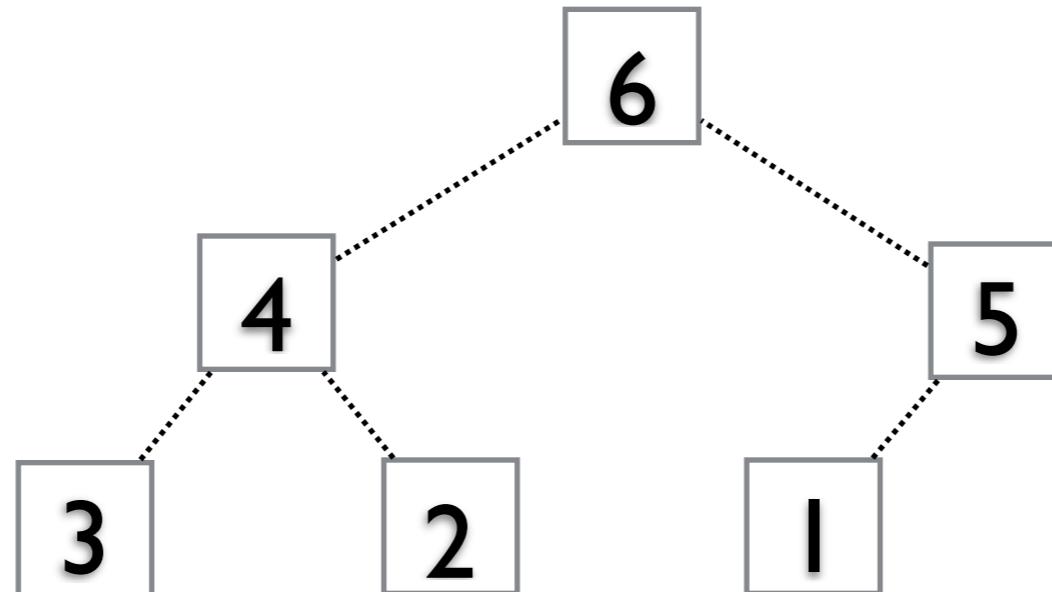
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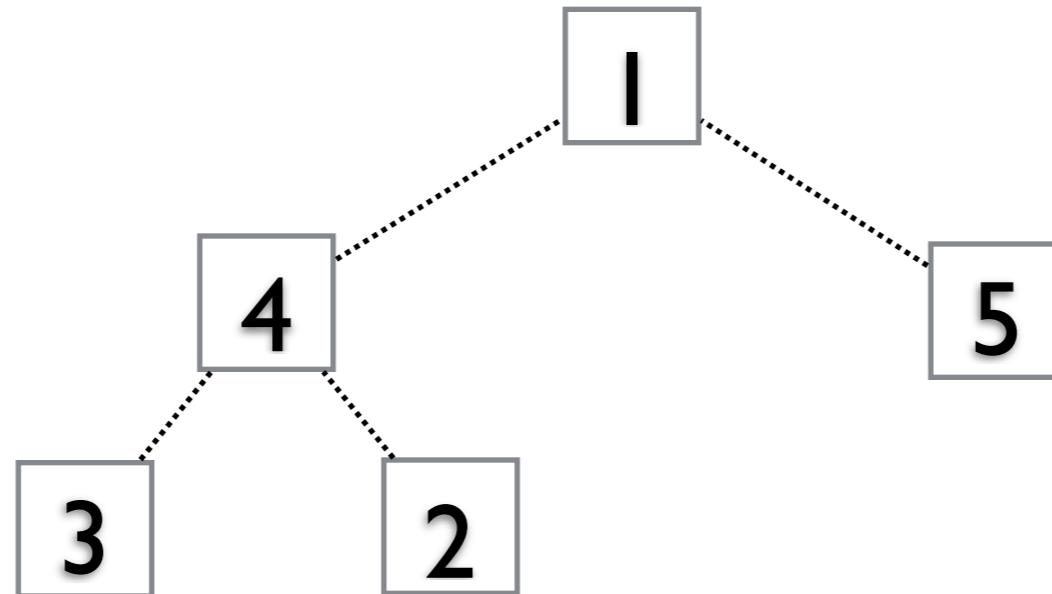
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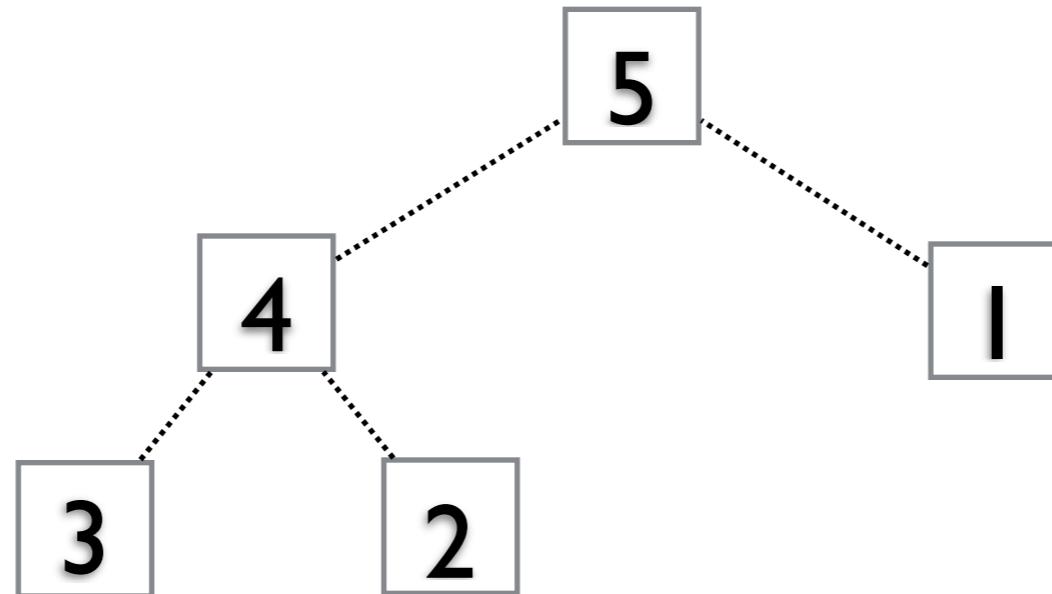
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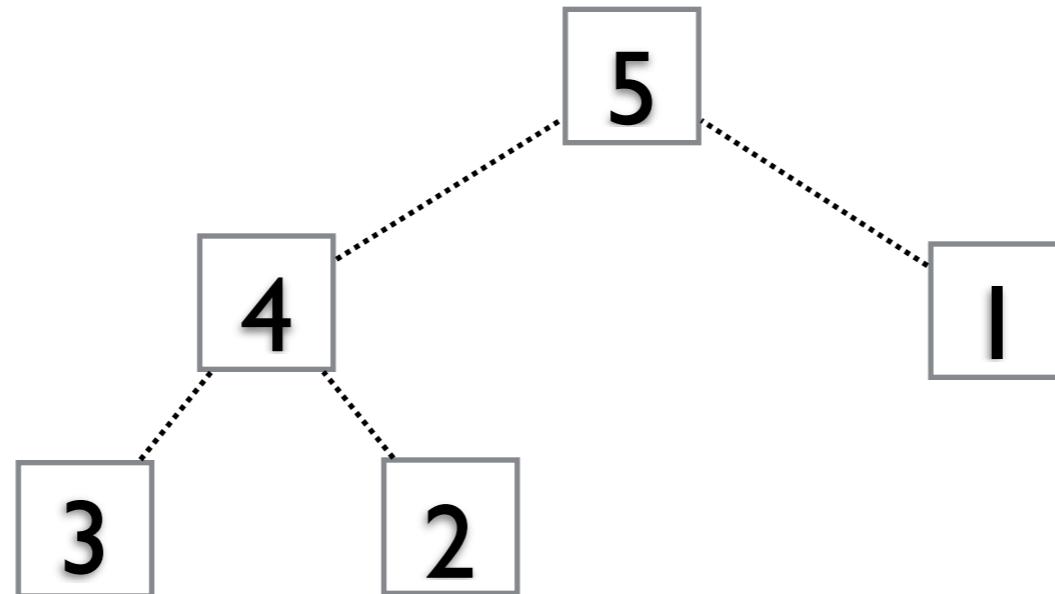
Keep **swapping** the root node with its child until it satisfies the **heap property**.



RemoveMax vis Sink

Replace the root with the **rightmost leaf**.

Keep **swapping** the root node with its child until it satisfies the **heap property**.



Complexity?



RemoveMax vis Sink

```
public T removeMax()
{
    throwNoSuchElementException();
    swap(1, n_size);
    T max = l_keys.remove(n_size--);
    sink(1);
    return max;
}

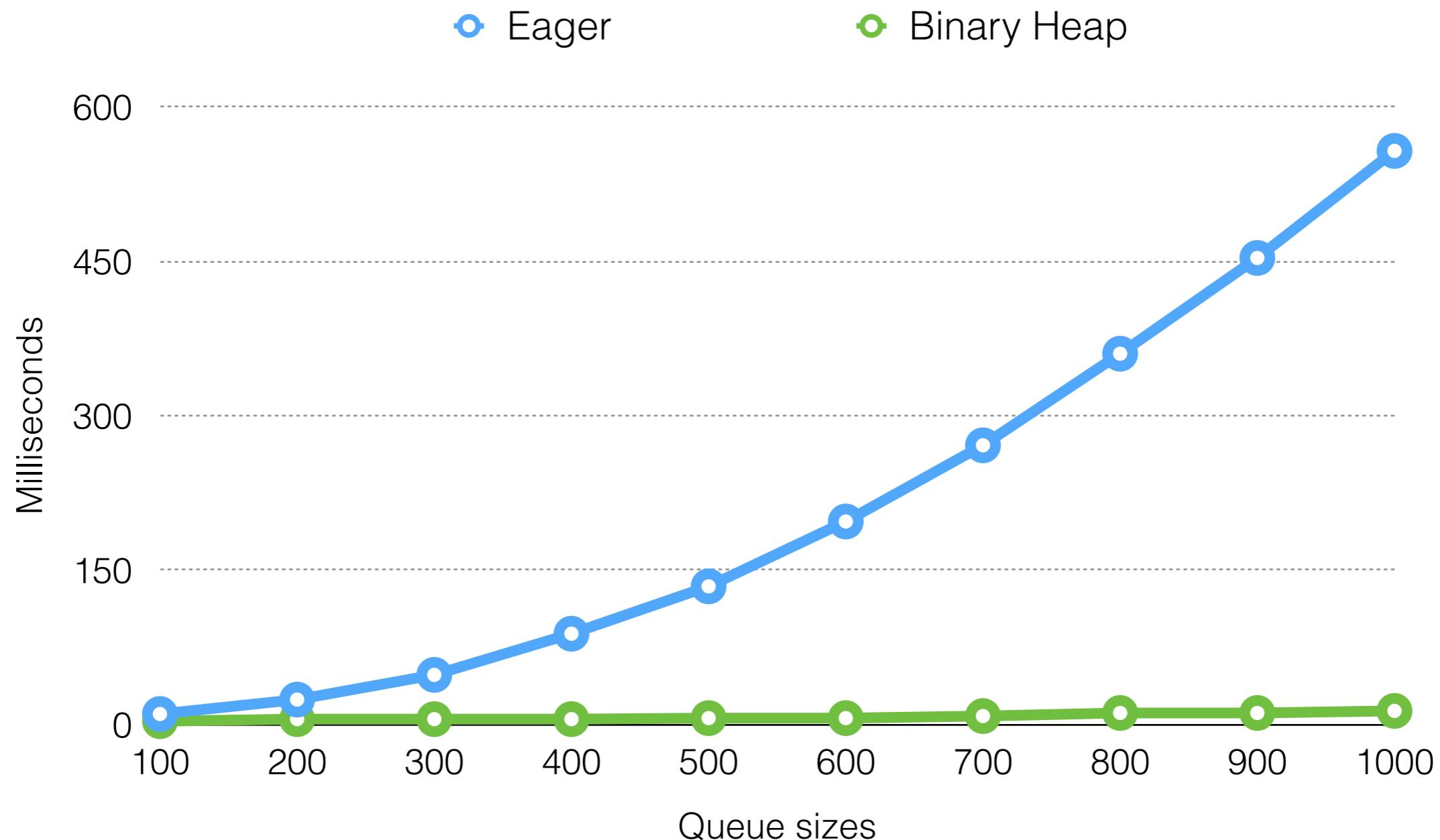
private void sink(int k)
{
    for (int i=k*2; i<=n_size; k=i, i=k*2) ← Left child
    {
        if (i < n_size && compareTo(i, i+1) < 0) i++;
        if (compareTo(k, i) >= 0) break;
        swap(k, i);
    }
}
```

↑ Compare
between children



Speed Comparison - Add

Both $O(\log n)$?

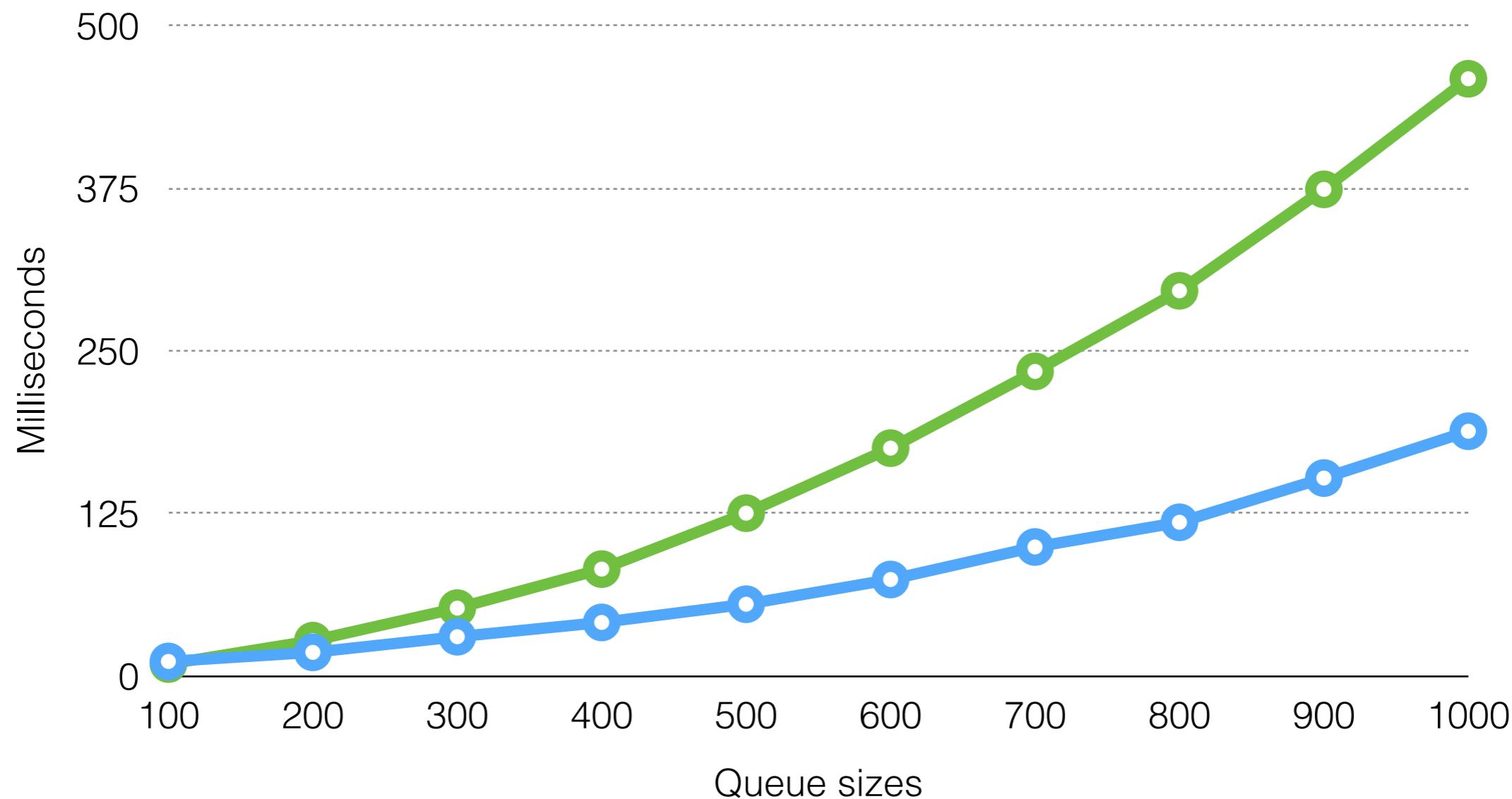


Speed Comparison - RemoveMax

$O(1)$ vs. $O(\log n)$?

● Eager

● Binary Heap



Exercise - Extra Credit

- HeapSelect
 - Create a class called **HeapSelect** under the **select** package.
 - Use **BinaryHeap** for finding the k'th maximum key.
 - Compare the speeds between **SmartSelect** and **HeapSelect**.
 - <https://github.com/jdchoi77/emory-courses/wiki/CS323:-Priority-Queues>
 - Submit your answer by **Sep. 11th** (Thus).
- Next class reading
 - Sections 2.1, 2.2, 2.4 - insertion sort, selection sort, merge sort.

