Decomposition of Graphs: Strongly connected components

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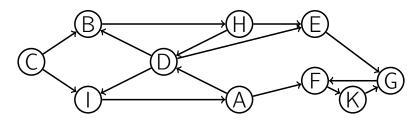
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Graph Algorithms

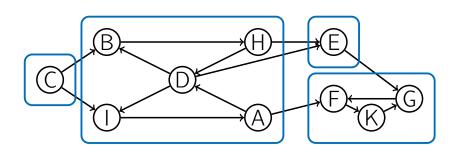
Data Structures and Algorithms

Strongly connected components

A strongly connected component of a directed graph is an inclusion-wise maximal subset of vertices such that there is a (directed) path between any two of them in both directions.



Metagraph





Lemma

If C and C' are strongly connected components, and there is an edge from a node in C to a node in C', then the highest post number in C is bigger than the highest post number in C'.

Proof

Case 1. If DFS visits C before C' then all of C and C' will be examined before we end processing a vertex from C,

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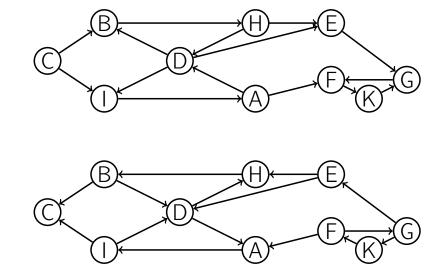
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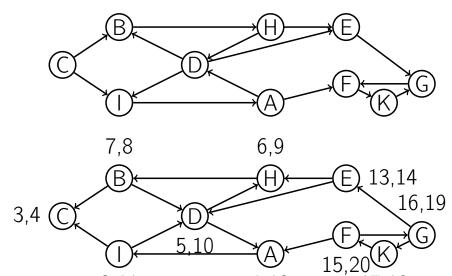
Corollaries

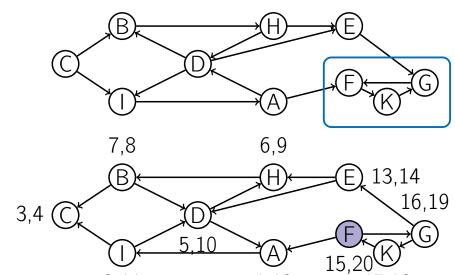
1 The vertex with the highest post number lies in a source SCC.

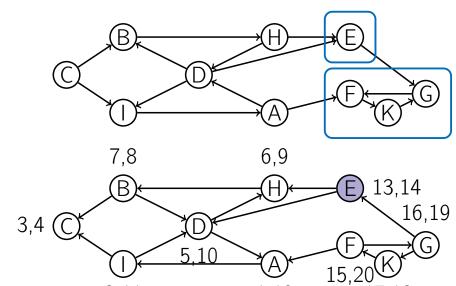
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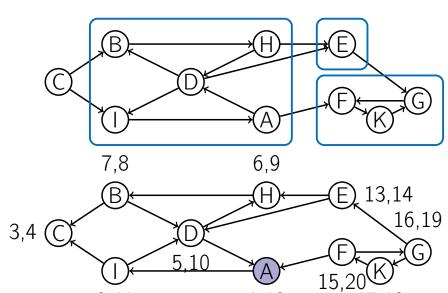
- 1 The vertex with the highest post number lies in a source SCC.
- 2 SCC's can be linearized by arranging them in decreasing order of their highest post numbers.











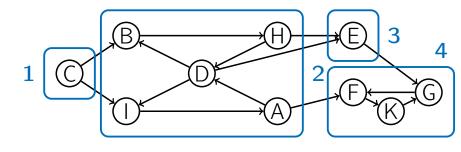
Computing SCC's

SCC(G)

run DFS on G^R

run the undirected connected components algorithm on G processing the vertices in decreasing order of their post numbers from step 1

Constructing metagraph in linear time



Turn the list of edges $\{(A, D), (C, B), \ldots\}$ into $\{(2, 2), (1, 2), \ldots\}$, sort it by calling counting sort twice, remove all duplicates