CS321: Computer Networks

Lec 2 : Introduction

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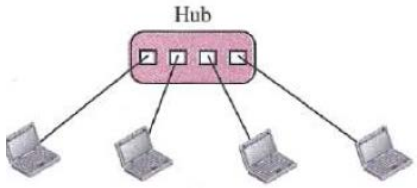
What is Data Communication?

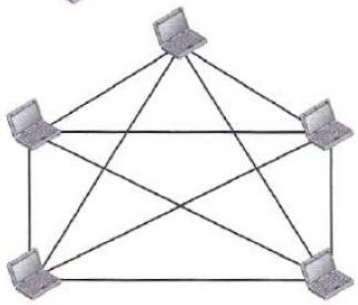
- Data communications are the exchange of data between two devices via some form of transmission medium.

- **Fundamental characteristics**:
  - Delivery: The system must deliver data to the correct destination.
  - Accuracy: The system must deliver the data accurately.
  - Timeliness: The system must deliver data in a timely manner.
  - Jitter: Jitter refers to the variation in the packet arrival time.
- **Components:**
  - Message
  - Sender
  - Receiver
  - Transmission medium
  - Protocol

Networks

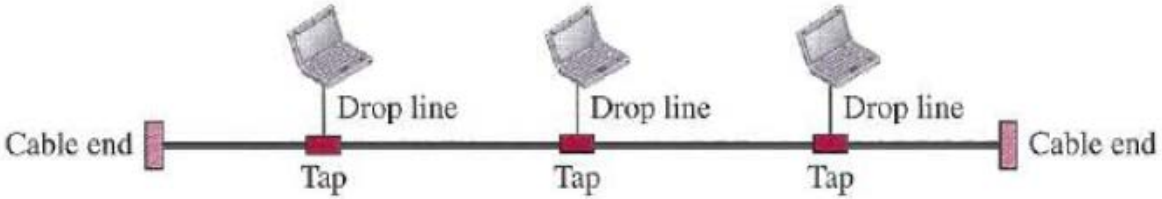
- A network is the interconnection of a set of devices capable of communication.



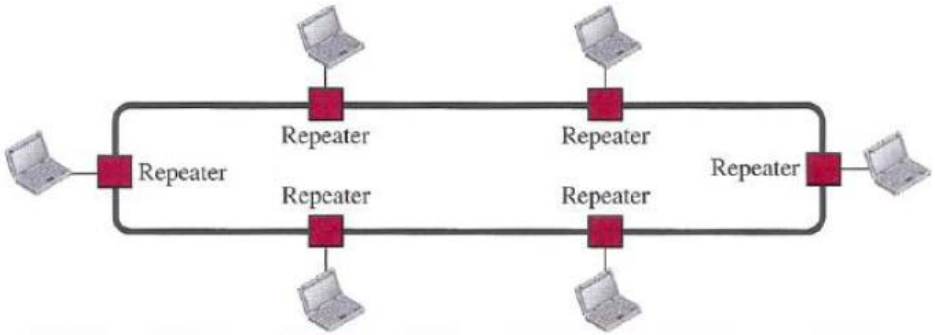


- **Network Criteria:**
  - *Performance:* Throughput, Delay
  - *Reliability:* frequency of failure
  - *Security:* protecting from unauthorized access

- **Type of Connection**
  - Point-to-Point
  - Multipipoint

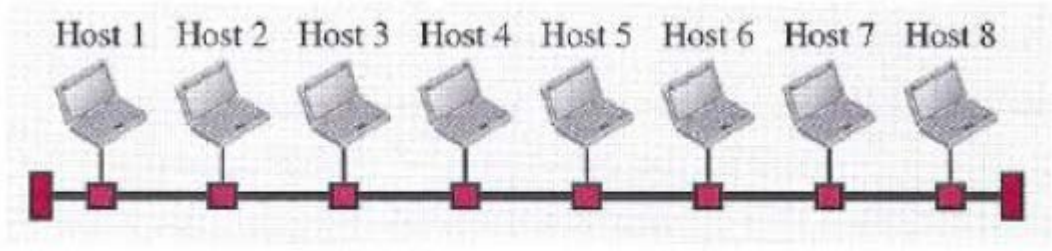


- **Physical Topology**
  - Mesh
  - Star
  - Bus
  - Ring

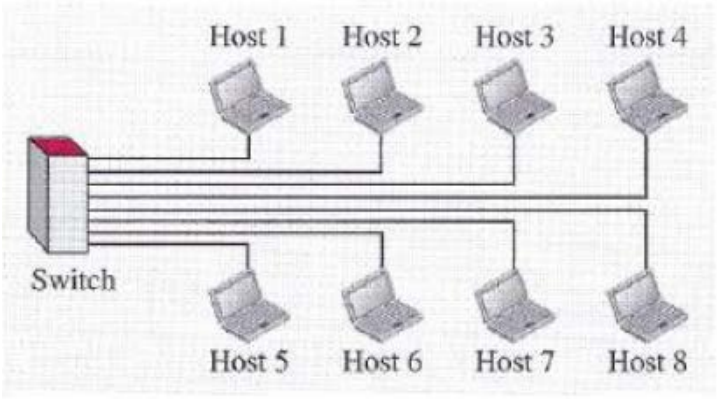


Network Types

- Local Area Networks (LAN)

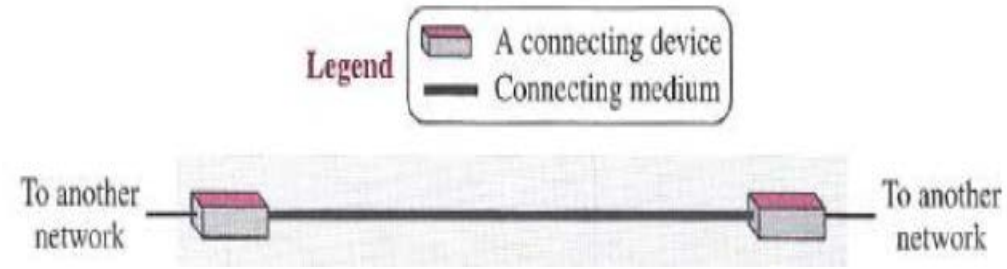


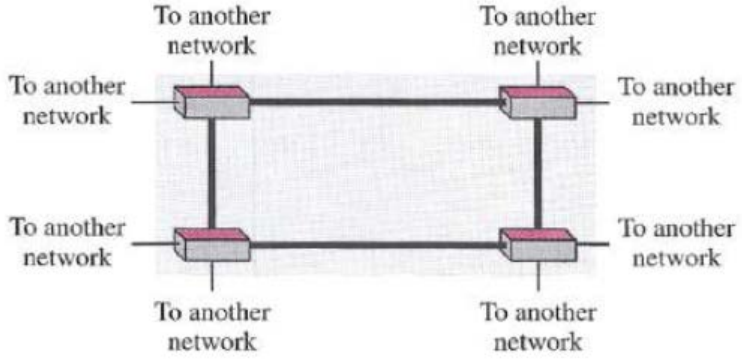
a. LAN with a common cable (past)



b. LAN with a switch (today)

- Wide Area Networks (WAN)





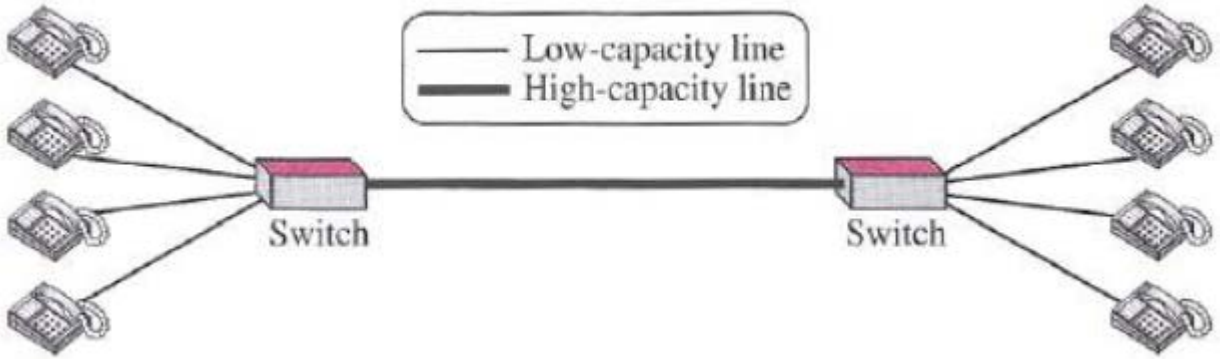
Switching

- An internet is a switched network in which a switch connects at least two links together.

- Common Types:
  - circuit-switched networks
  - packet-switched networks

Circuit-switched Network

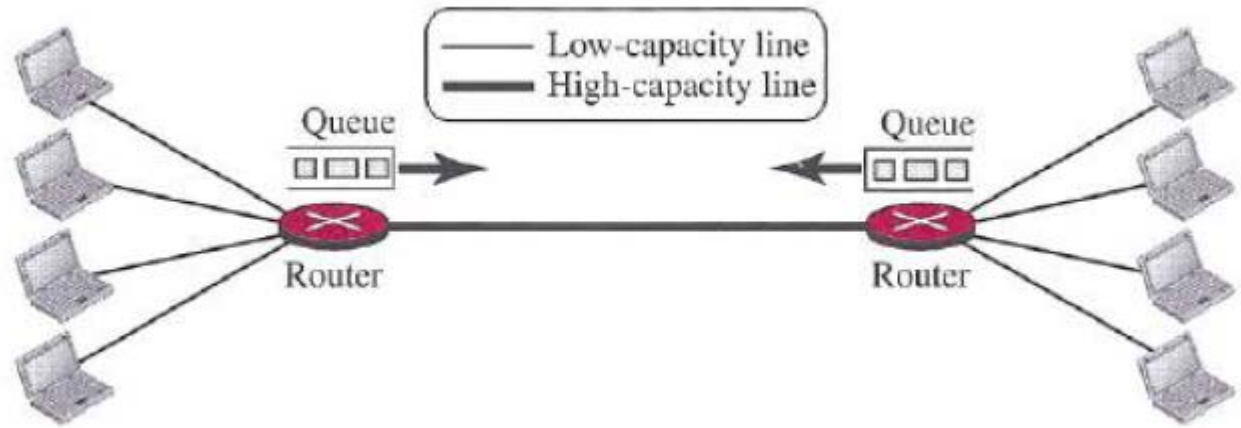
- a dedicated connection, called a circuit, is always available between the two end systems; the switch can only make it active or inactive.
  - E.g., Telephone Network



- Disadvantage:
  - Underutilization of link capacity
  - No storing / holding capacity

Packet-switched Network

- the communication between the two ends is done in blocks of data called packets.
  - E.g., Internet



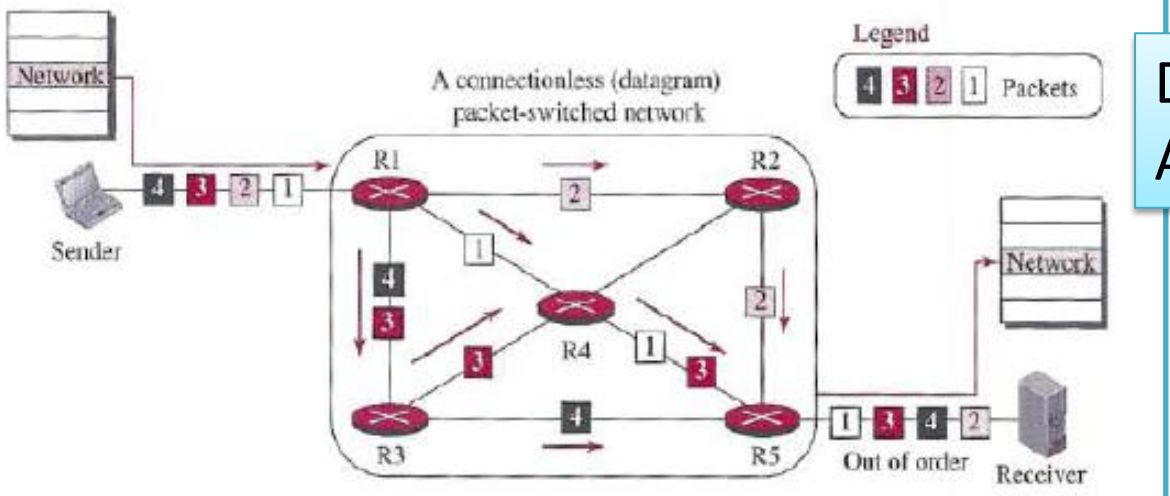
- Advantages:
  - Better utilization of link
  - Router can store and forward data packets

Approaches in Packet-switching

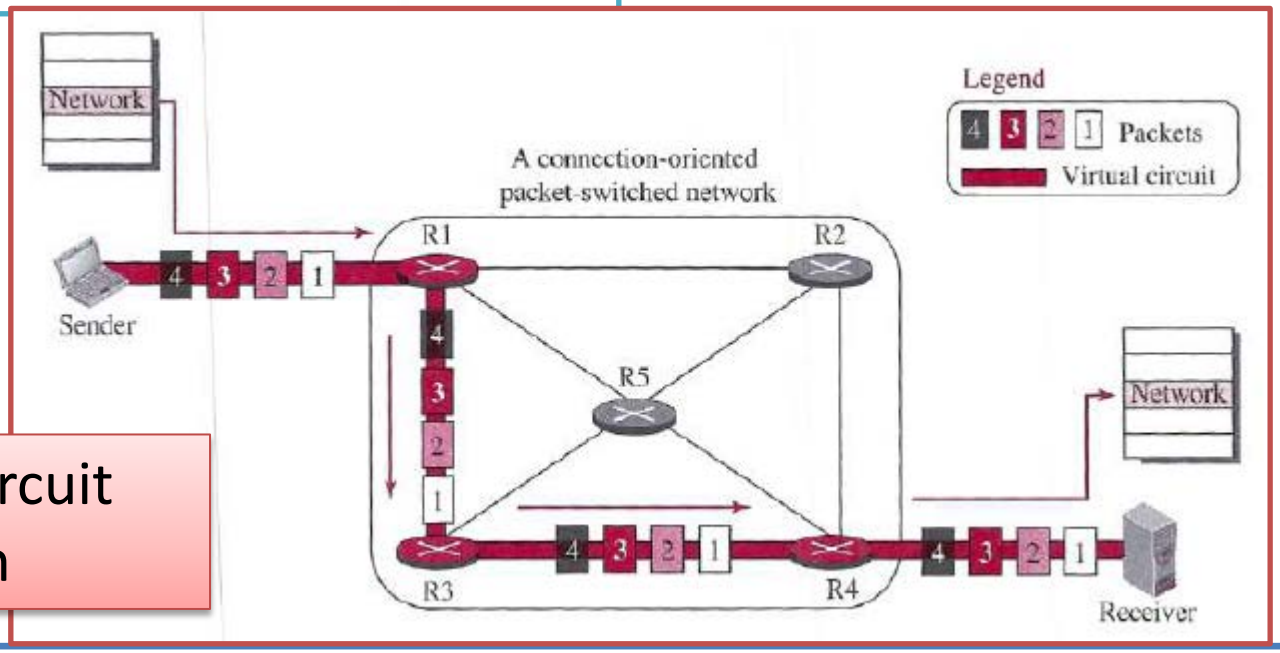
- Datagram approach
  - Connectionless service: each packet is treated independently
  - Packet forwarding decision is based on destination address

- Virtual-Circuit approach
  - Connection-Oriented service: exists relationship among all packets belonging to a message
  - Packet forwarding decision is based on label of the packet

Approaches in Figure



Datagram Approach



Virtual-circuit Approach

How to define network operations?

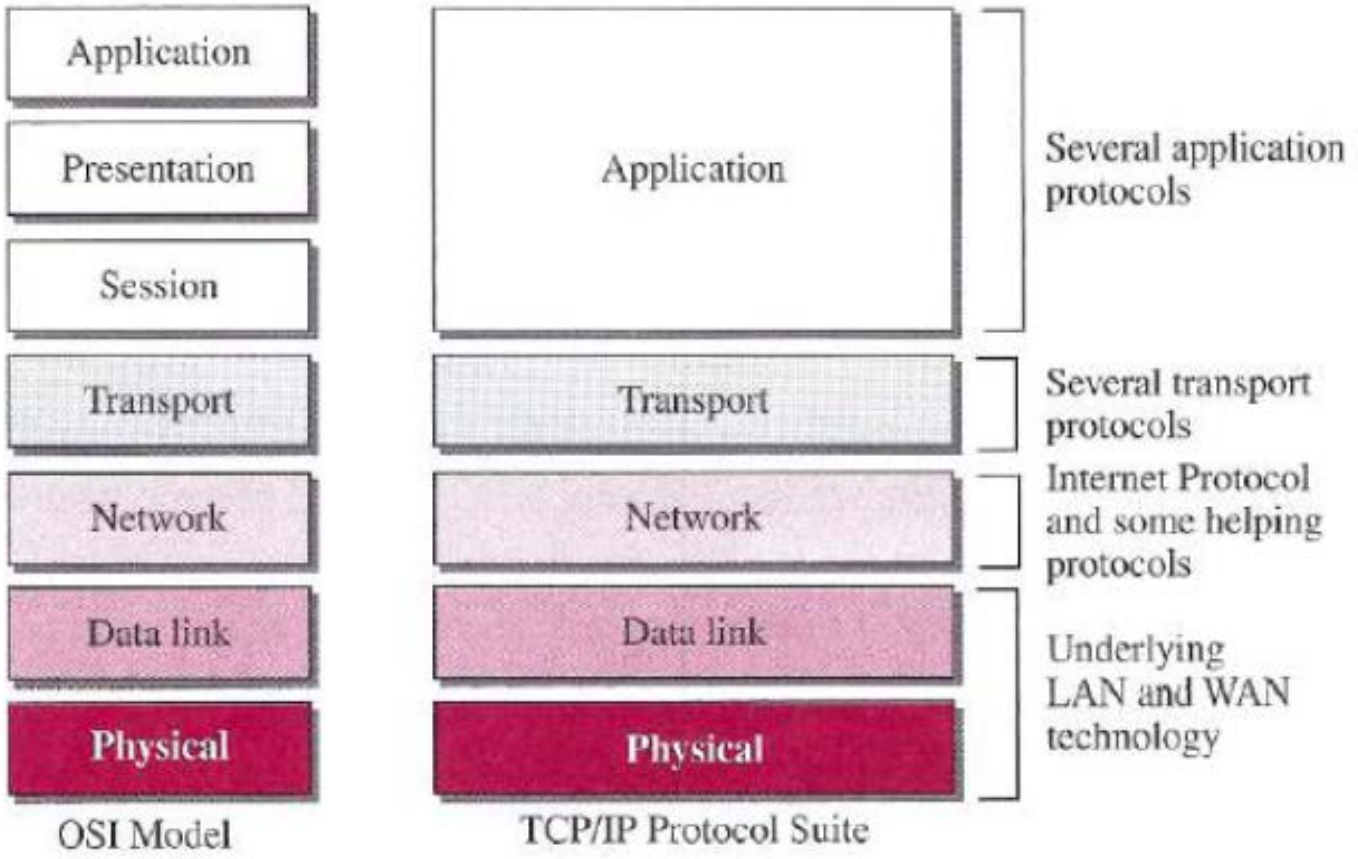
- Ans: Protocol
- Protocol defines the rules that both the sender and receiver and all intermediate devices need to follow to be able to communicate effectively.

- For simple communication (between 2 computer): One protocol is enough

- For complex communication (multiple computers): Modularization / Protocol Layering is essential

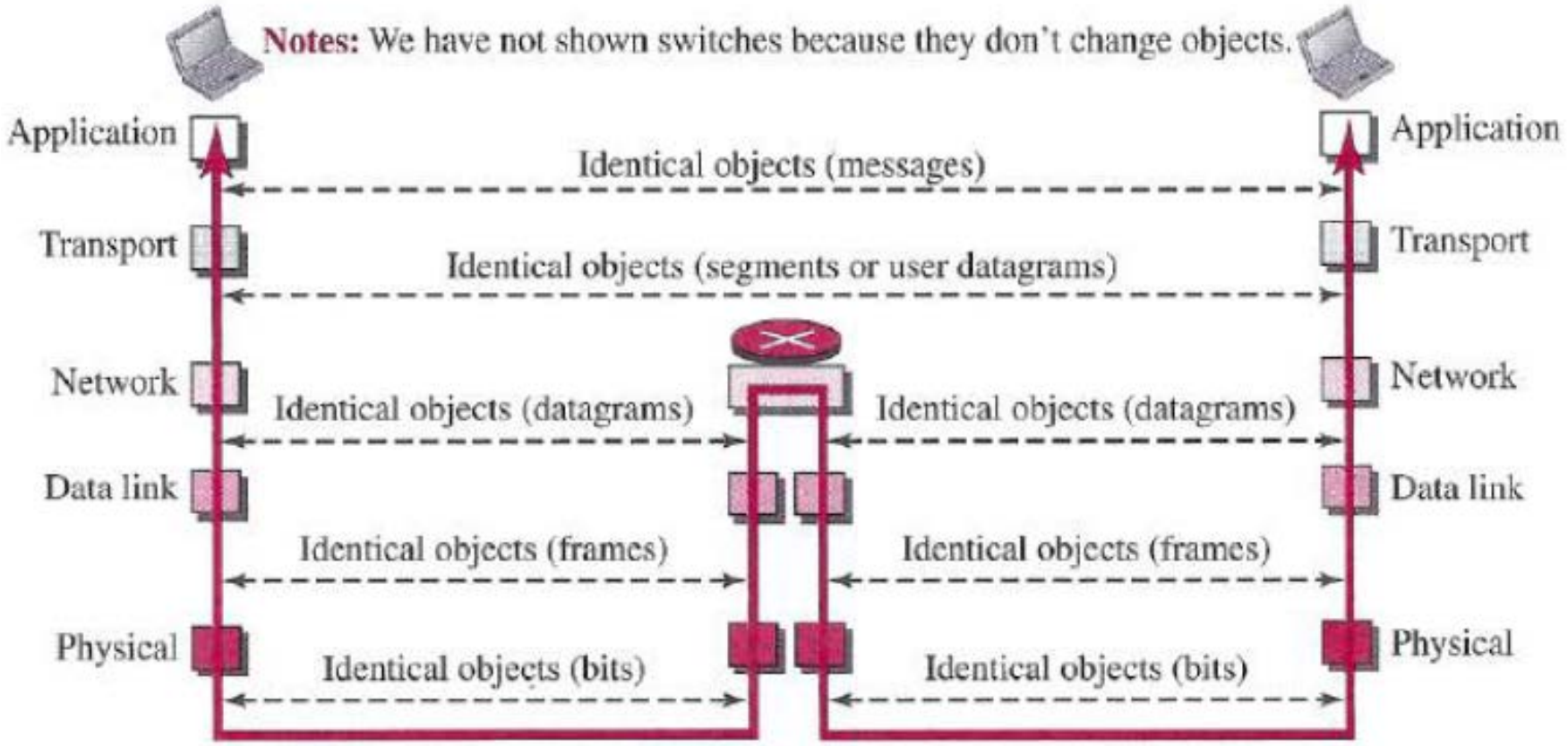
- Principles of Protocol Layering:
  - each layer performs two opposite tasks, one in each direction
  - Two objects under each layer at both sites should be identical

Protocol Layering Models



- Allegory:: Please Do Not Touch Steve's Pet Alligator

TCP/IP Protocol Suite



Application Layer

- The application layer consists of what most users think of as programs.
- Although each application is different, some applications are so useful that they have become standardized.

- **File transfer (FTP)**: Connect to a remote machine and send or fetch an arbitrary file. FTP deals with authentication, listing a directory contents, ASCII or binary files, etc.
- **Remote login (telnet)**: A remote terminal protocol that allows a user at one site to establish a TCP connection to another site, and then pass keystrokes from the local host to the remote host.
- **Mail (SMTP)**: Allow a mail delivery agent on a local machine to connect to a mail delivery agent on a remote machine and deliver mail.
- **Web (HTTP)**: Base protocol for communication on the World Wide Web.

Presentation and Session Layers

- Presentation layer:
  - This layer is concerned with Syntax and Semantics of the information transmitted
  - Service: Encoding / Decoding, Abstract data structures and conversion among them, etc.

- Session layer:
  - This layer allows users on different machines to establish session between them.
  - Service: dialogue control, token management, synchronization, etc.

Transport Layer

- The transport layer provides end-to-end communication between processes executing on different machines.
  - Connection establishment
  - Segmentation
  - Quality and Type of Services: throughput, delay, protection, priority, reliability, etc.
  - Flow control
  - Congestion control
  - Etc.

Network Layer

- The network layer provides an end-to-end communication between machines belonging to different/same network.
  - Routing
  - Internetworking
  - Interface between the host and the network
  - Etc.

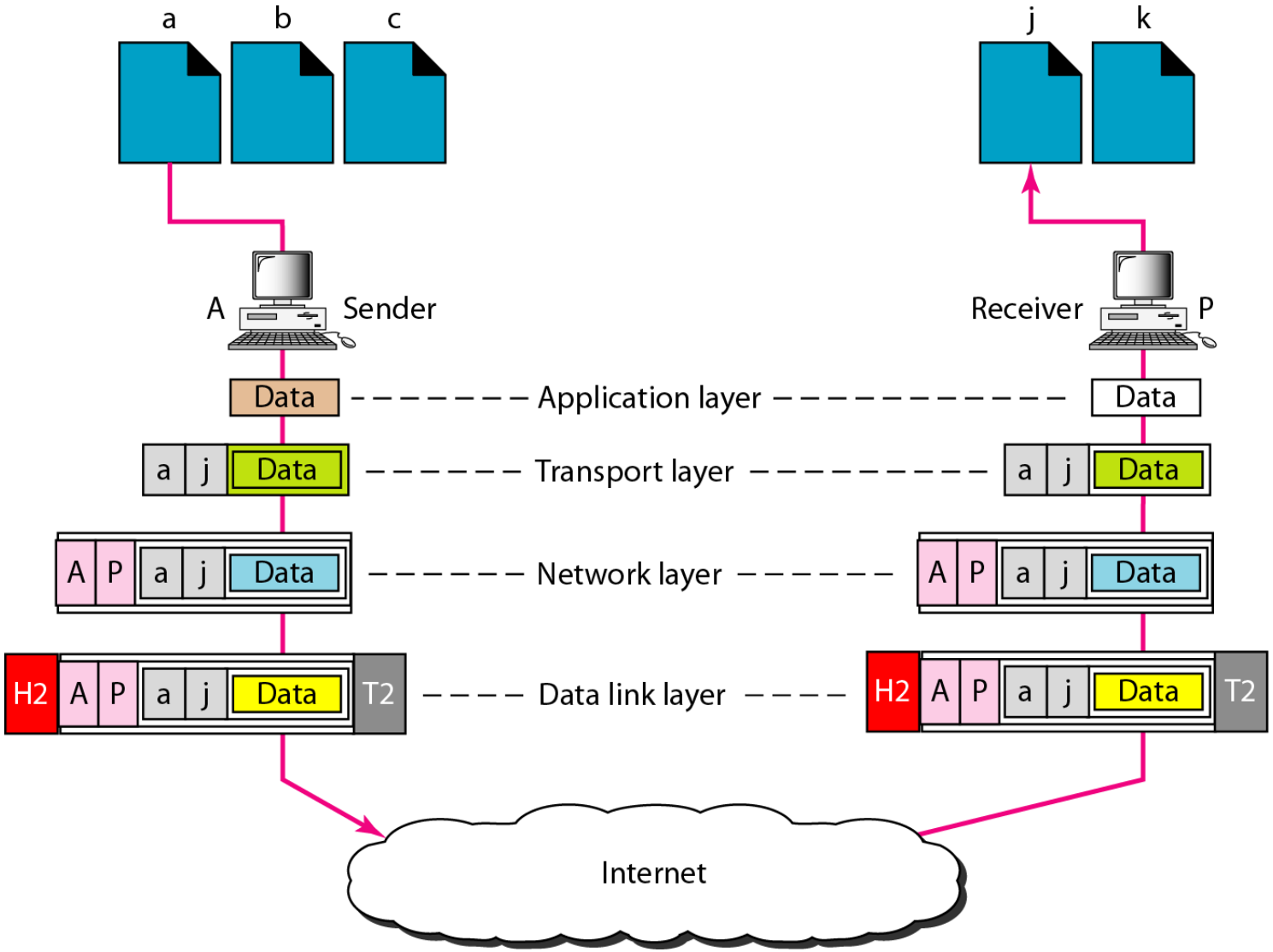
Datalink Layer

- provides reliable, efficient communication between adjacent machines connected by a single communication channel.
  - Framing
  - Reliable delivery
  - Acknowledged delivery
  - Error control
  - Flow control
  - Etc.

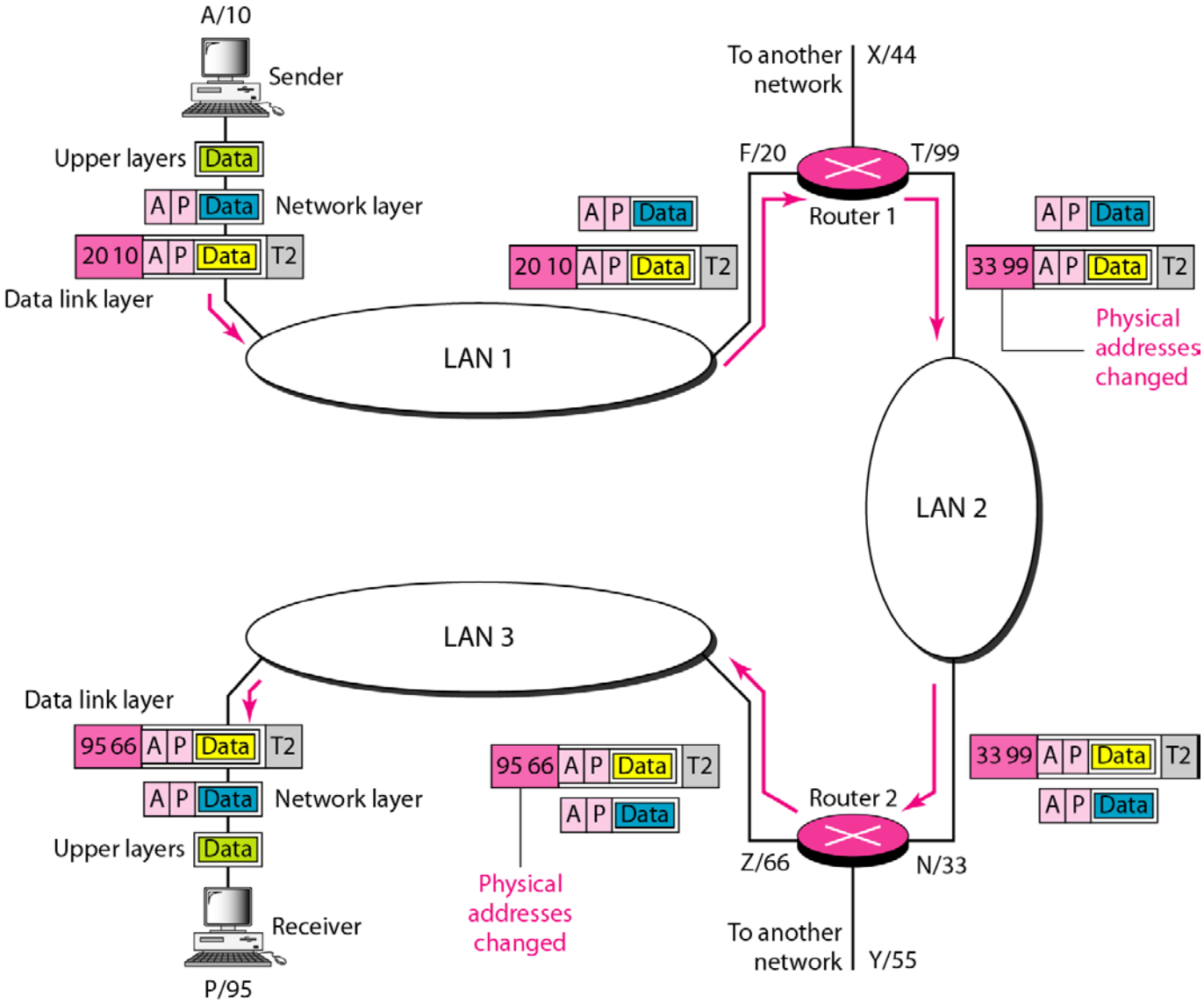
Physical Layer

- This layer is concerned with transmission of raw bits over a communication channel.
  - Deals with physical transmission
  - Establishing and breaking of connection
  - Signal Level, Data rate
  - Etc.

Example of Data Delivery



Example of Data Delivery



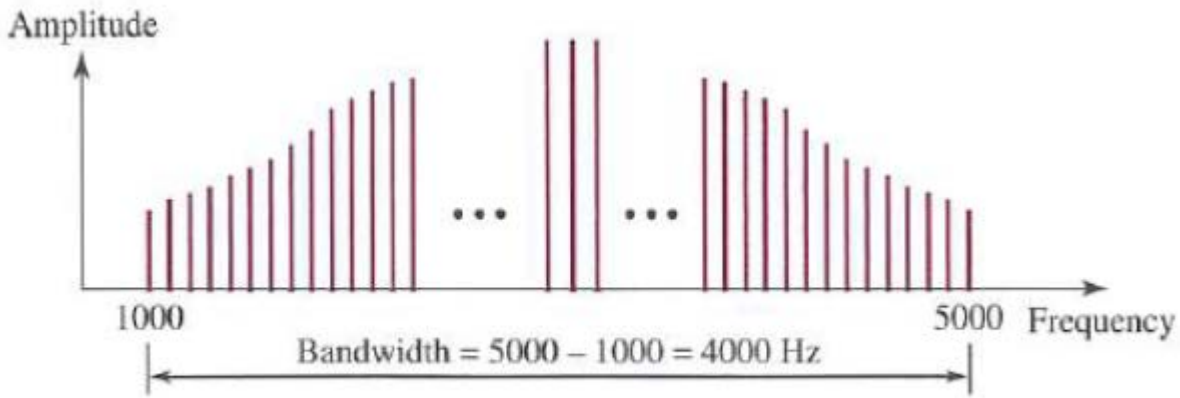
Performance Metrics

- Bandwidth
- Latency (delay)
- Bandwidth-delay product
- Throughput

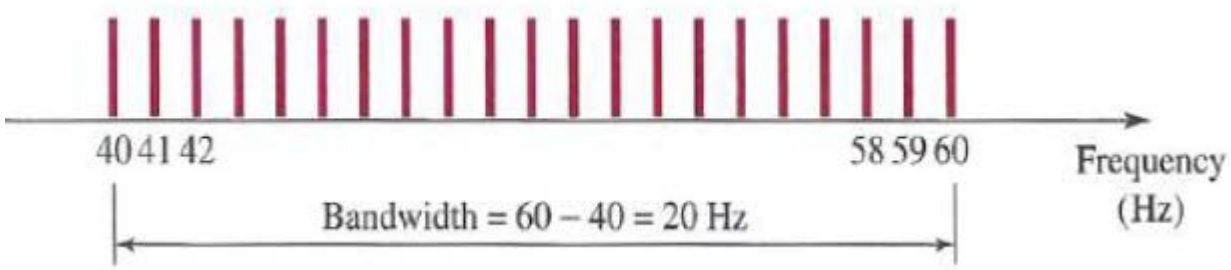
- Bandwidth:
  - Number of bits transmitted per second
  - Difference between the highest and lowest frequencies contained in that signal

Bandwidth (in Hertz)

- For Analog Signal

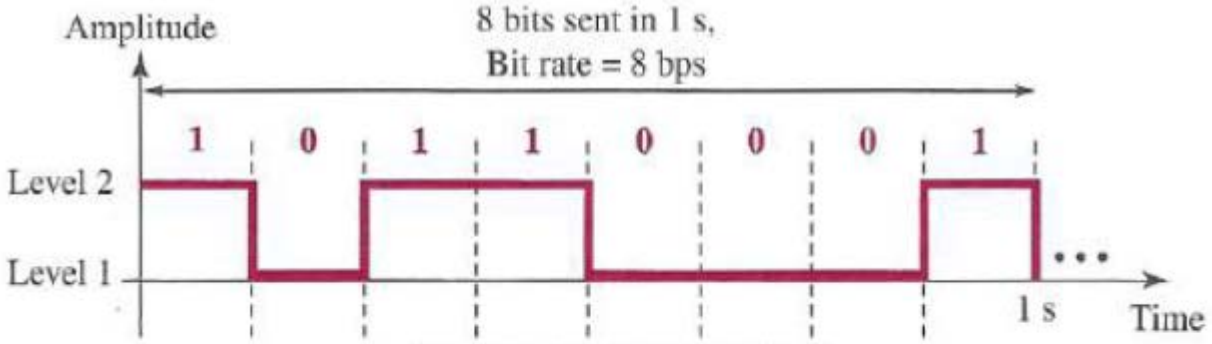


a. Bandwidth of a periodic signal

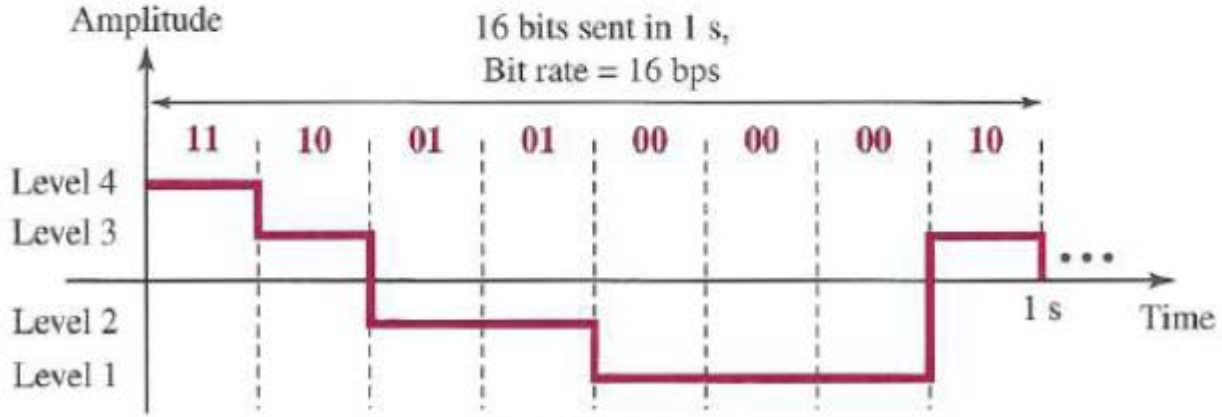


Bandwidth (in Bits per Sec.)

- For Digital Signal



a. A digital signal with two levels



b. A digital signal with four levels

Latency (Delay)

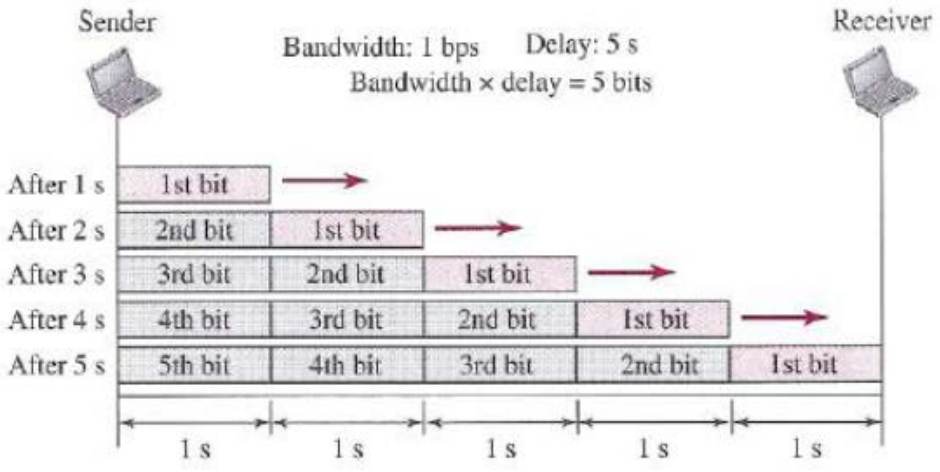
- The latency or delay defines how long it takes for an entire message to completely arrive at the destination from the time the first bit is sent out from the source.
  - Latency = propagation time + transmission time + queuing time (+ processing delay)
  - Propagation time = Distance / (Propagation Speed)
    - Speed of Light in:
      - Vacuum $3x10^8$ m/s
      - Copper $2.3x10^8$ m/s
      - Fiber $2x10^8$ m/s
  - Transmission time = Message size / Datarate
  - queuing time= the time needed for each intermediate or end device to hold the message before it can be processed
  - Processing delay= time needed to process by devices
  - Round Trip Time: roughly two times of Latency

Example

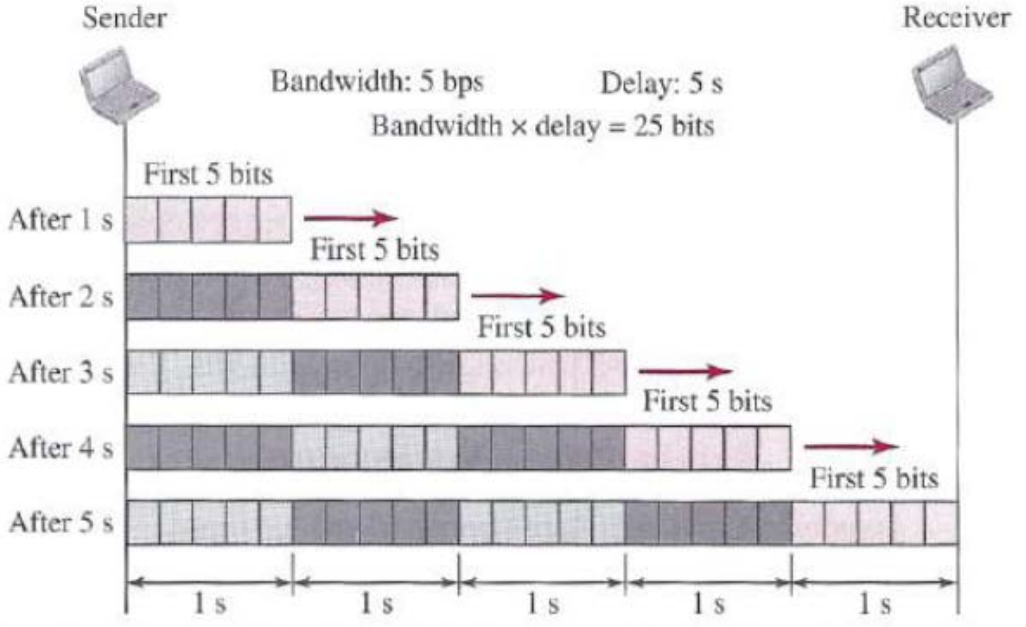
- A network with bandwidth of 10 Mbps can pass only an average of 12,000 frames per minute with each frame carrying an average of 10,000 bits. What is the throughput of this network?
- Solution:
  - Throughput = (12,000 x 10,000) / 60 = 2 Mbps

- What are the propagation time and the transmission time for a 5-MB (megabyte) message (an image) if the bandwidth of the network is 1 Mbps? Assume that the distance between the sender and the receiver is 12,000 km and that light travels at 2.4×10^8 m/s.
- Solution:
  - Propagation time = (12,000 x 1000) / (2.4×10^8) = 50 ms
  - Transmission time == (5,000,000 x 8) /10^6 = 40 sec

Bandwidth-Delay Product



The bandwidth-delay product is a measure of the number of bits a sender can transmit through the system while waiting for an ACK from the receiver.



Transmission of Digital Signal

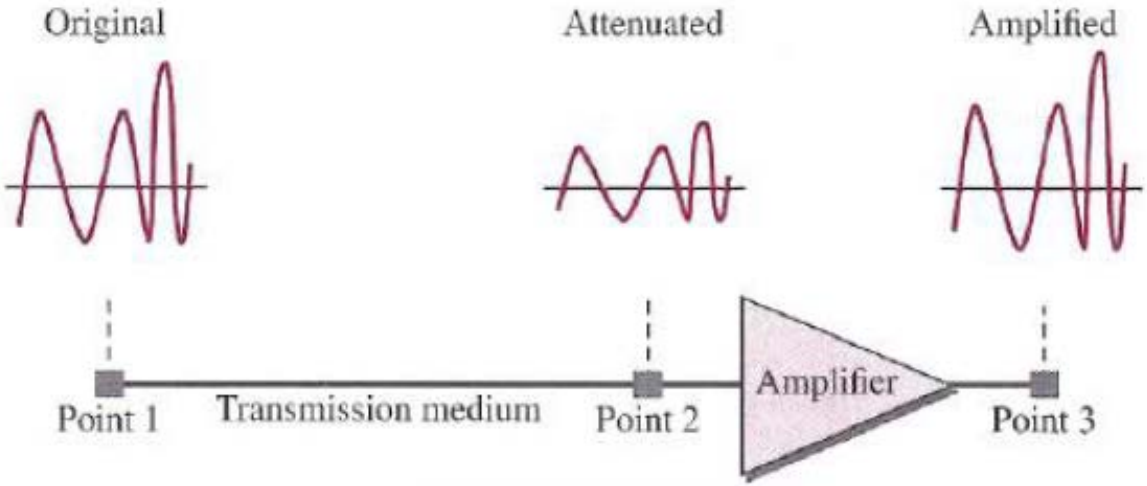
- Digital signal is treated as composite analog signal
- Two different approaches:
  - Baseband Transmission: sending digital signal without changing to analog signal
    - Need dedicated medium as frequency starts from zero
    - Low-pass channel
    - Required bandwidth is proportional to the bit rate; approximately bitrate/2.
  - Broadband Transmission: changing the digital signal to analog signal for transmission
    - Need modulation & de-modulation
    - Bandpass channel

Transmission Impairment

- **Attenuation** means a loss of energy.
- **Distortion** means that the signal changes its form or shape.
- Several types of **noise**, such as thermal noise, induced noise, crosstalk, and impulse noise, may corrupt the signal.
- Thermal noise: the random motion of electrons in a wire
- Induced noise: comes from sources such as motors and appliances
- Crosstalk: is the effect of one wire on the other.
- Impulse noise: is a spike (a signal with high energy in a very short time) that comes from power lines, lightning, and so on.

Decibel

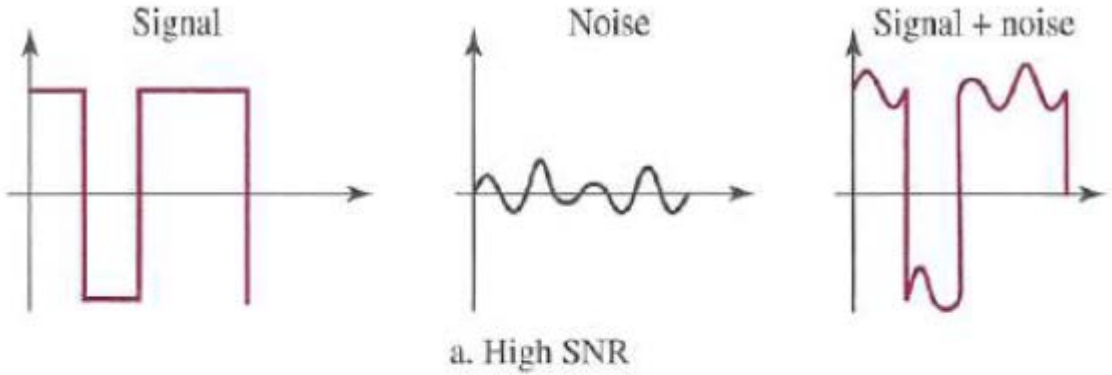
- The **decibel (dB)** measures the relative strengths of two signals or one signal at two different points.



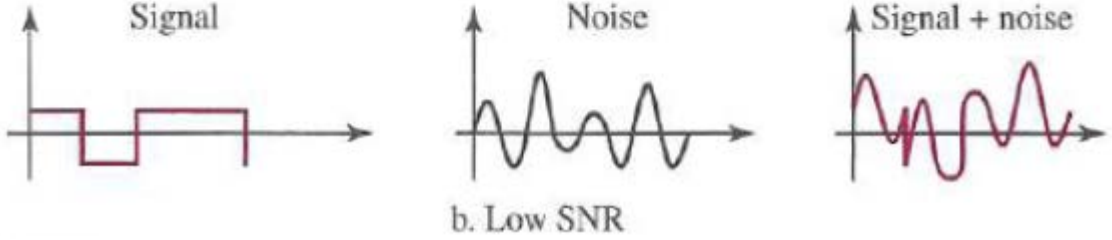
$$\mathrm{dB} = 10 \log_{10} \frac{P_2}{P_1}$$

Signal-to-Noise Ratio (SNR)

- It is the ratio of the signal power to the noise power.



a. High SNR



b. Low SNR

$$\text{SNR}_{\text{dB}} = 10 \log_{10} \text{SNR}$$

The values of SNR and SNR_{dB} for a noiseless channel are

$$\text{SNR} = (\text{signal power}) / 0 = \infty \longrightarrow \text{SNR}_{\text{dB}} = 10 \log_{10} \infty = \infty$$

Throughput

- how fast we can actually send data through a network.
  - Bits per Sec.
  - A link may have a bandwidth of *B* bps, but we can only send *T* bps through this link with $T < B$.
  - *Corollary*:: Imagine a highway designed to transmit 1000 cars per minute from one point to another. However, if there is congestion on the road, this figure may be reduced to 100 cars per minute. The bandwidth is 1000 cars per minute; the throughput is 100 cars per minute.

Data Rate

- how fast we can send data, in bits per second, over a channel.
  - Bits per sec.

- Depends on:
  - Bandwidth available
  - Level of signals we use
  - Quality of the channel

Noiseless Channel: Nyquist Bit Rate

- Nyquist Capacity Theorem:: For a noiseless channel, the Nyquist bit rate formula defines the theoretical maximum bit rate.
  - **BitRate** = 2 x Bandwidth x $\log_2(L)$;

    where L is the number of signal levels used to represent data.
- Note: Increasing the levels of a signal may reduce the reliability of the system.

- Sampling Theorem: Any signal whose bandwidth is *B* can be completely recovered by the sampled data at rate *2B* samples per second.

Noisy Channel: Shannon Capacity

- In reality, we cannot have a noiseless channel; the channel is always noisy.
- the Shannon capacity determines the theoretical highest data rate for a noisy channel.
  - Capacity = Bandwidth x $\log_2(1+SNR)$

- The signal-to-noise ratio is often given in decibels. Assume that $SNR_{dB} = 36$ and the channel bandwidth is 2 MHz. The theoretical channel capacity can be calculated as:

$$SNR_{dB} = 10 \log_{10} SNR \longrightarrow SNR = 10^{SNR_{dB}/10} \longrightarrow SNR = 10^{3.6} = 3981$$

$$C = B \log_2(1 + SNR) = 2 \times 10^6 \times \log_2 3982 = 24 \text{ Mbps}$$

Thanks!