

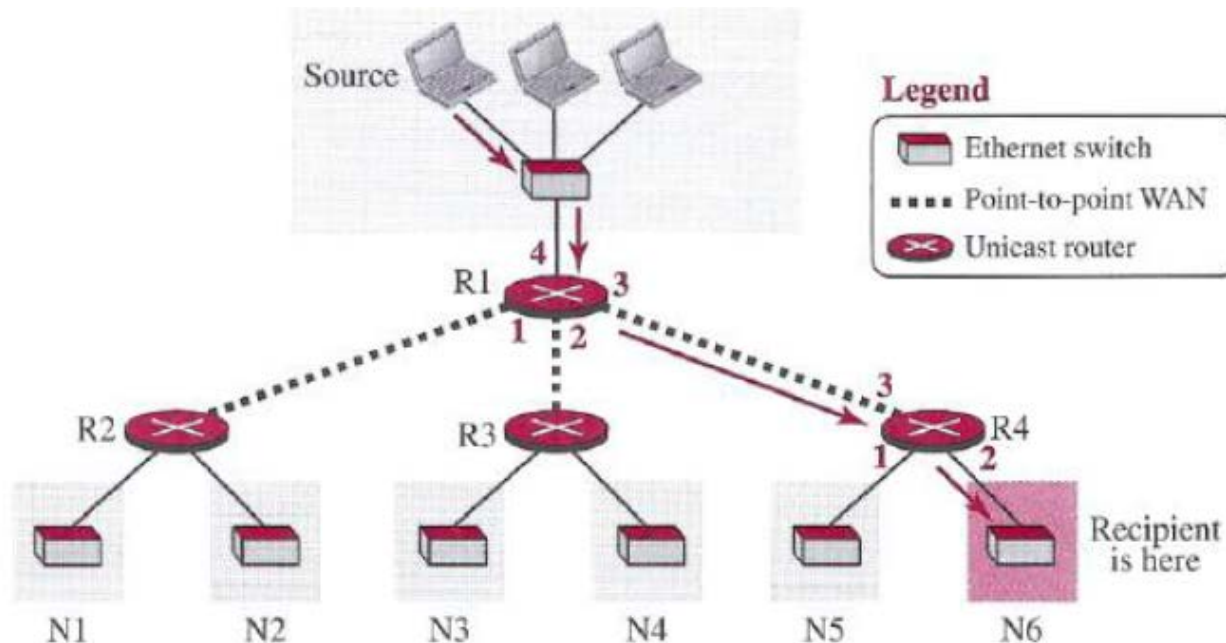
## Multicast Routing

Dr. Manas Khatua  
Assistant Professor  
Dept. of CSE  
IIT Jodhpur

E-mail: [manaskhatua@iitj.ac.in](mailto:manaskhatua@iitj.ac.in)

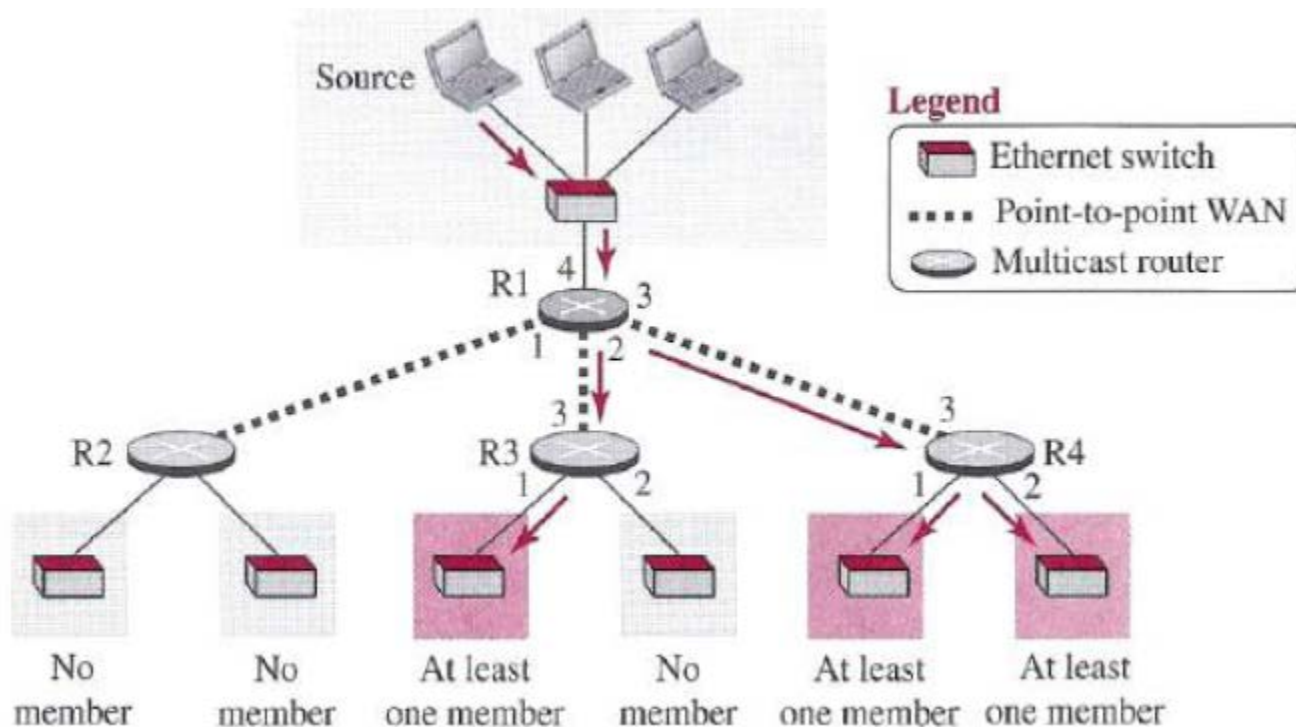
# Unicasting

- There is **one source** and **one destination network**.
- The **relationship** between the source and the destination network is **one to one**.
- Each router in the path tries to **forward** the packet **to one** and only one of its **interfaces**.



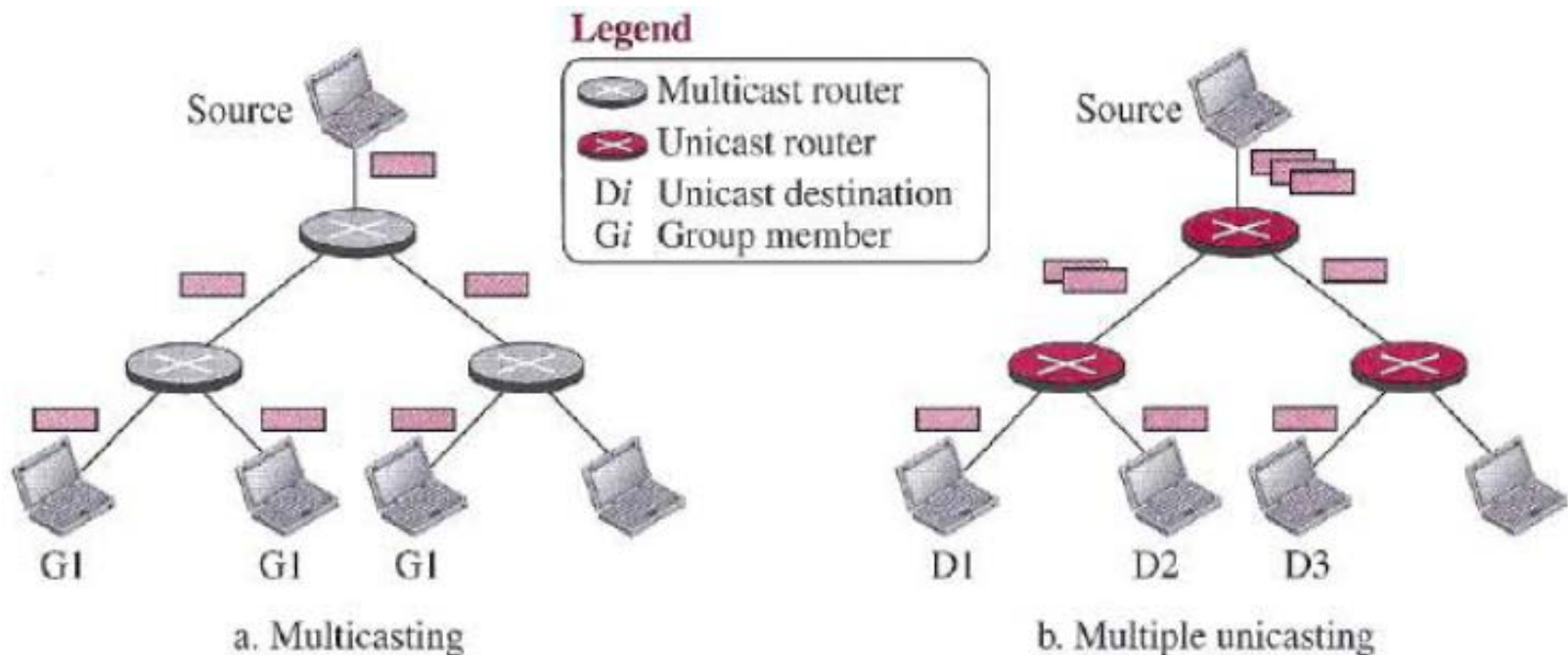
# Multicasting

- There is **one source** and a **group of destinations**.
- The relationship is **one to many**.
- The source address is a **unicast address**, but the destination address is a **group address**, in which there is **at least one member of the group** that is interested in receiving the multicast datagram.



# Multicast vs Multiple Unicast

- Multicasting starts with a **single packet** from the source that is **duplicated by the routers**.
- The **destination address** in each packet is the **same** for all duplicates.
- **Only a single copy** of the packet travels between any two routers.



# Cont...



- In multiple unicasting, several packets start from the source.
- If there are three destinations, for example, the source sends three packets, each with a different unicast destination address.
- Note that there may be multiple copies traveling between two routers.
- Example:
  - When a person sends an e-mail message to a group of people, this is multiple unicasting.
  - Teleconferencing: A group of workstations form a multicast group such that a transmission from any member is received by all other group members.

# Why multicasting?



- Two main reasons:
  - Multicasting requires **less bandwidth** than multiple unicasting.
  - In multiple unicasting, the packets are created by the source with a **relative delay between packets**. In multicasting, there is no delay because only one packet is created by the source.
- Why group e-mail is multiple unicast?
  - Multicast involves a **subscription from the receiver's** side.
  - But, multiple unicast is a **decision from the sender's** side.
  - Usually, sender manage the group of multiple unicast, but a receiver is associated with a multicast group.

# Multicast Applications

- Teleconferencing
- Distance Learning
- Information Dissemination
- Access to Distributed Databases
- etc.

## Broadcasting:

- one-to-all communication: a host sends a packet to all hosts in an internet.

# Multicast Address

- In IP datagram, we can only write **one destination address**. So, we need multicast address for sending the datagram to many destinations.
- a multicast address is an identifier for a group.
- If a new group is formed with some active members, an **authority can assign** an unused multicast address to this group to uniquely define it



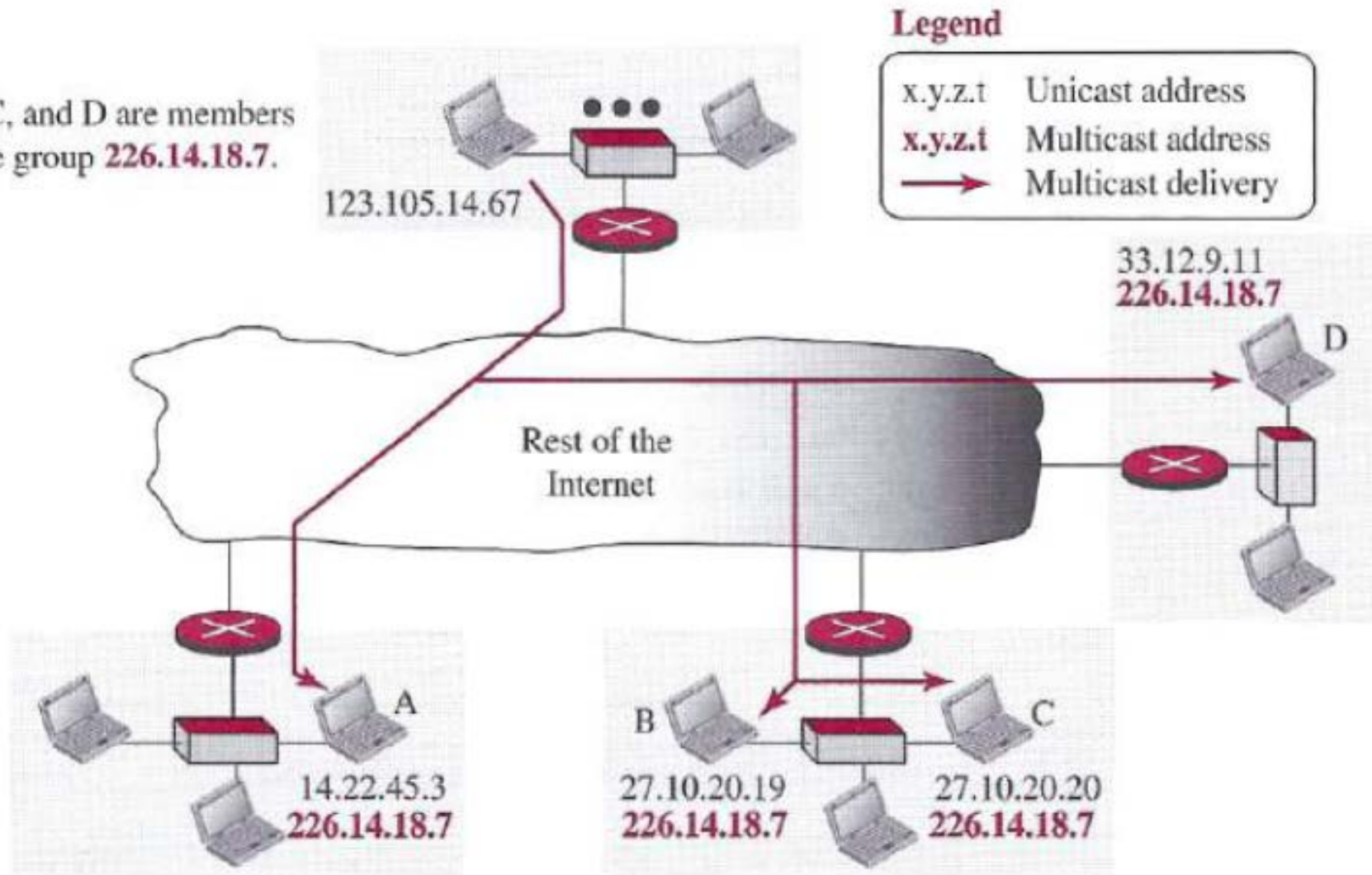
# Multicast Address in IPv4

- A router or a destination host needs to **distinguish between** a unicast and a multicast datagram.
- IPv4 assigns a block of addresses for this purpose
- In **classful addressing**, all of class D was composed of these addresses;
- **classless addressing** used the same block, but it was referred to as the block 224.0.0.0/4 (from 224.0.0.0 to 239.255.255.255).



# Example

A, B, C, and D are members of the group **226.14.18.7**.



# Multicast Sub-blocks

- Total number of multicast address blocks =  $2^{28}$
- The blocks are divided into multiple sub-blocks
  - Local Network Control Blocks: [224.0.0.0/24](#)
    - Multicast routing is used inside a network
    - Datagram cannot be forwarded by the router to outside
  - Internetwork Control Block: [224.0.1.0/24](#)
    - Routing protocol can used whole Internet
  - Source-specific Multicast Block: [232.0.0.0/8](#)
    - IGMP protocol use this
  - GLOP Block: [233.0.0.0/8](#)
    - To restrict inside an AS (autonomous system)
  - Administratively Scoped Block: [239.0.0.0/8](#)
    - To restrict inside an organization or an area

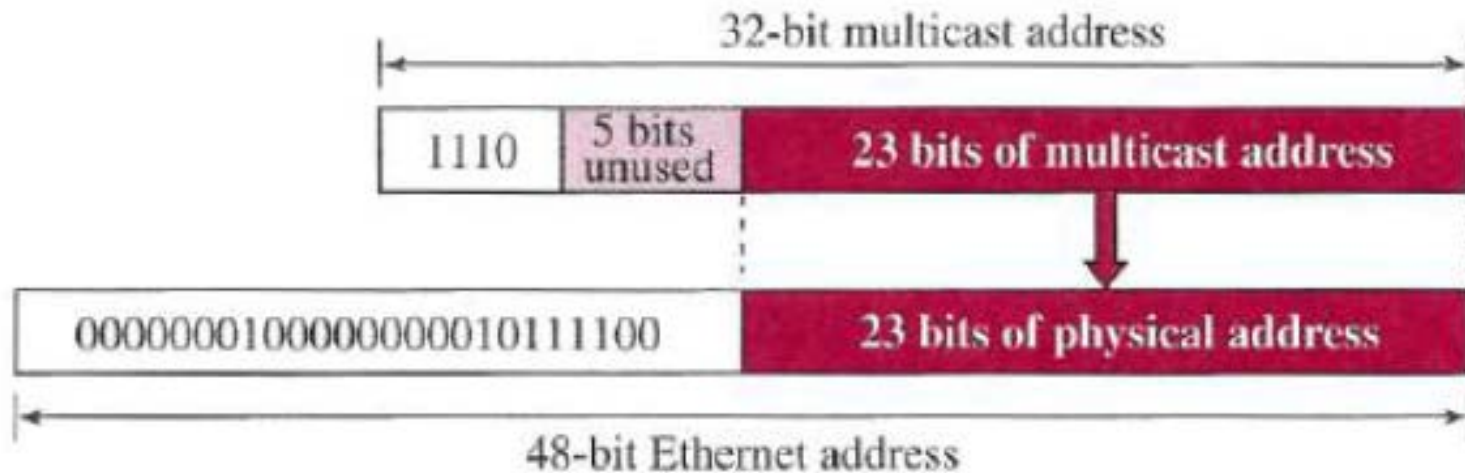
# Delivery at DLL

- In multicasting, the delivery at the **Internet level** is done using multicast IP addresses
- But, **data-link layer** multicast addresses are also needed to deliver a multicast packet encapsulated in a frame.
- **ARP protocol cannot help** in finding multicast MAC address
- Solution for two scenario:
  1. Network **with Multicast Support**

Most LANs (e.g. **Ethernet**) support physical multicast addressing.

If the **first 25 bits** in an Ethernet address are 00000001 00000000 01011110 0, this identifies a physical multicast address for the TCP/IP protocol.

# Cont...



- An Ethernet multicast physical address is in the range

01:00:5E:00:00:00 - 01:00:5E:7F:FF:FF

# Cont...

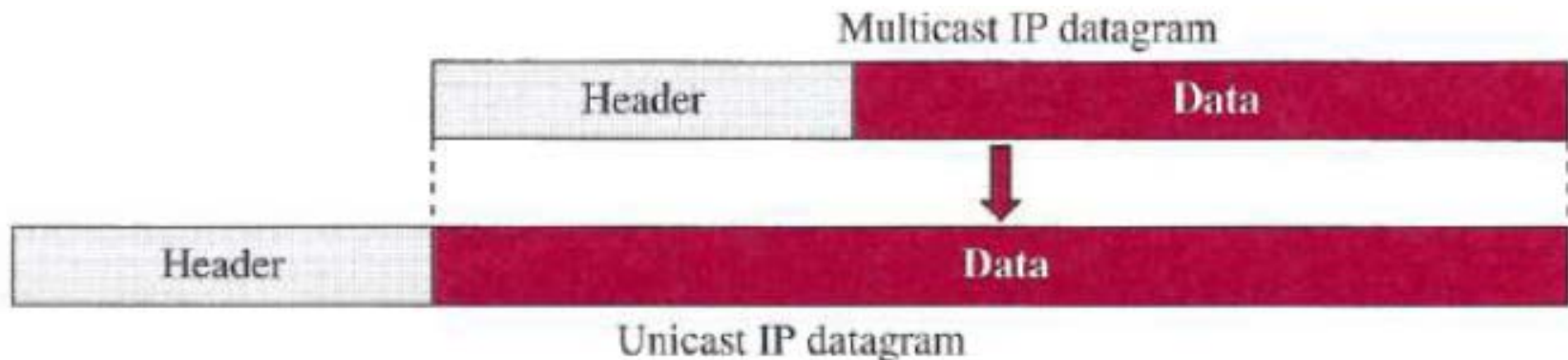


- *Example:* Change the multicast IP address **232.43.14.7** to an Ethernet multicast physical address.
- We can do this in **two steps**:
  - We write the rightmost 23 bits of the IP address in hexadecimal. Then subtracting 8 from the leftmost digit if it is greater than or equal to 8. In our example, the result is **2B:0E:07**
  - We add the result of part a to the starting Ethernet multicast address, which is **01:00:5E:00:00:00**. The result is **01:00:5E:2B:0E:07**

# Cont...

## 2. Network with **No Multicast Support**

- Most WANs do not support physical multicast addressing
- To send a multicast packet through these networks, a process called **tunneling** is used
- In tunneling, the multicast packet is encapsulated in a unicast packet and sent



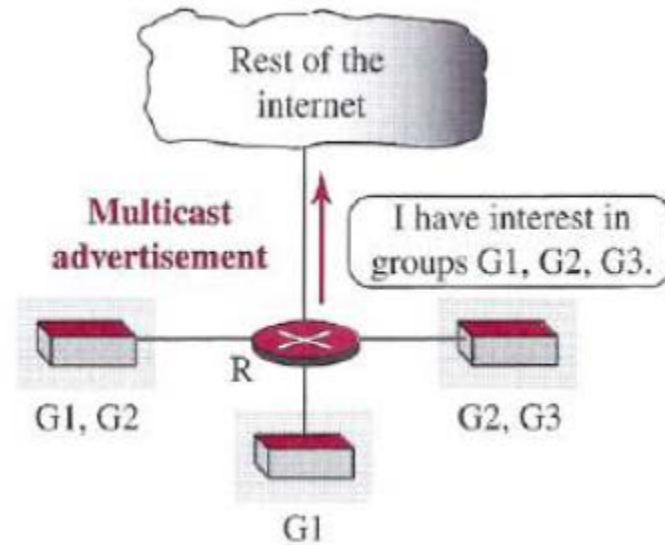
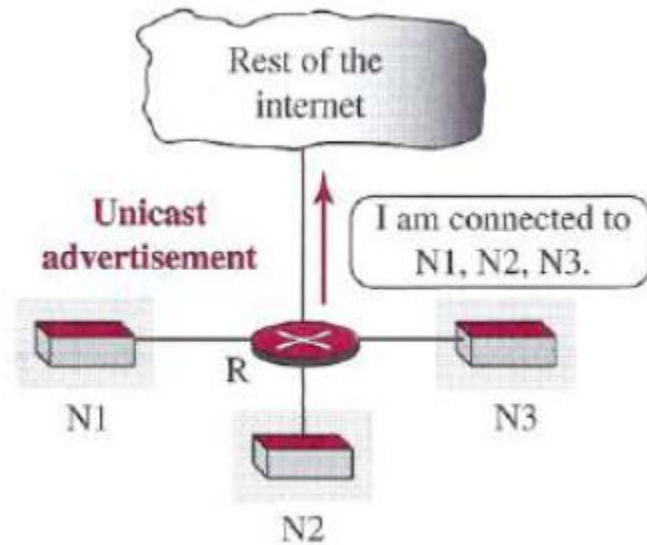
# Collecting Information about Groups



- Creation of **forwarding tables** in both **unicast** and **multicast** routing involves two steps:
  - A router needs to **know** to which **destinations** it is connected.
  - Each router needs to **propagate information** obtained in the first step to all other routers so that each router knows to which destination each other router is connected



# Cont...



- In **unicast routing**, the collection of the information in the first step is **automatic**;
- Each router knows to which network it is connected, and the prefix of the network (in CIDR) is what a router needs.
- In **multicast routing**, the collection of information in the first step is **not automatic**.

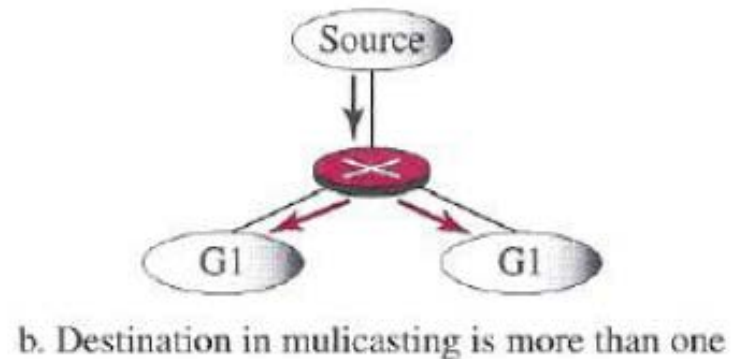
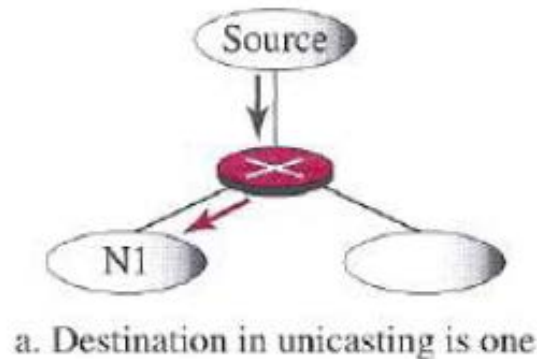
# Cont...



- Because,
  - a **router does not know** which host in the attached network is a member of a particular group; membership in the group does not have any relation to the prefix associated with the network.
  - the **membership is not a fixed attribute of a host**; a host may join some new groups and leave some others even in a short period of time.
- For unicasting, the router needs no help to collect;
- but for multicasting, it needs the help of **another protocol** namely Internet Group Management Protocol (IGMP)

# Multicast Forwarding

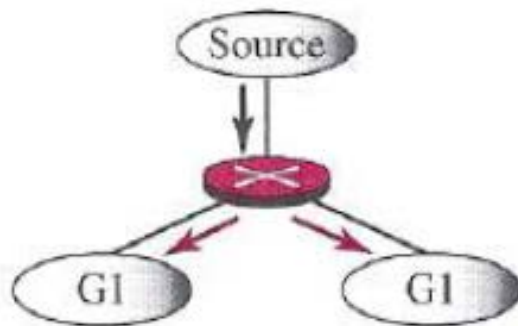
- a router needs to make a decision to forward a multicast packet



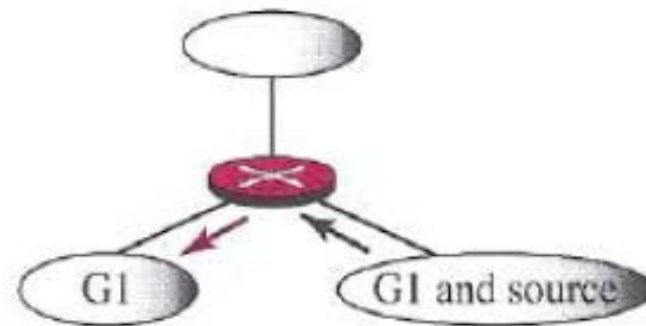
- In **unicast** communication, the **destination** address of the packet defines one single destination. So, forwarded through **one** interface.
- In **multicast** communication, the **destination** of the packet defines one group, but that group may have more than one member in the internet. So, forwarded through many interfaces.

# Cont...

- Forwarding decisions in **unicast** communication depend **only on the destination address** of the packet.
- Forwarding decisions in **multicast** communication depend on **both the destination and the source address** of the packet.



a. Packet sent out of two interfaces



b. Packet sent out of one interface

# Multicasting Approaches

- we need to create routing trees to optimally route the packets from their source to their destination.
  - **Source-Based Tree** Approach
    - each router needs to create a separate tree for each **source-group combination**.
    - In each tree, the corresponding source is the **root**, the members of the group are the **leaves**, and the router itself is somewhere on the tree.
  - **Group-Shared Tree** Approach
    - we designate a router to act as the **dummy source** for each group.
    - The designated router, which is called the **core router**, acts as the representative for the group.
    - Any source that has a packet to send to a member of that group **sends it to the core router** (unicast communication) and the **core router is responsible for multicasting**.

# Intra-domain Multicast Protocol

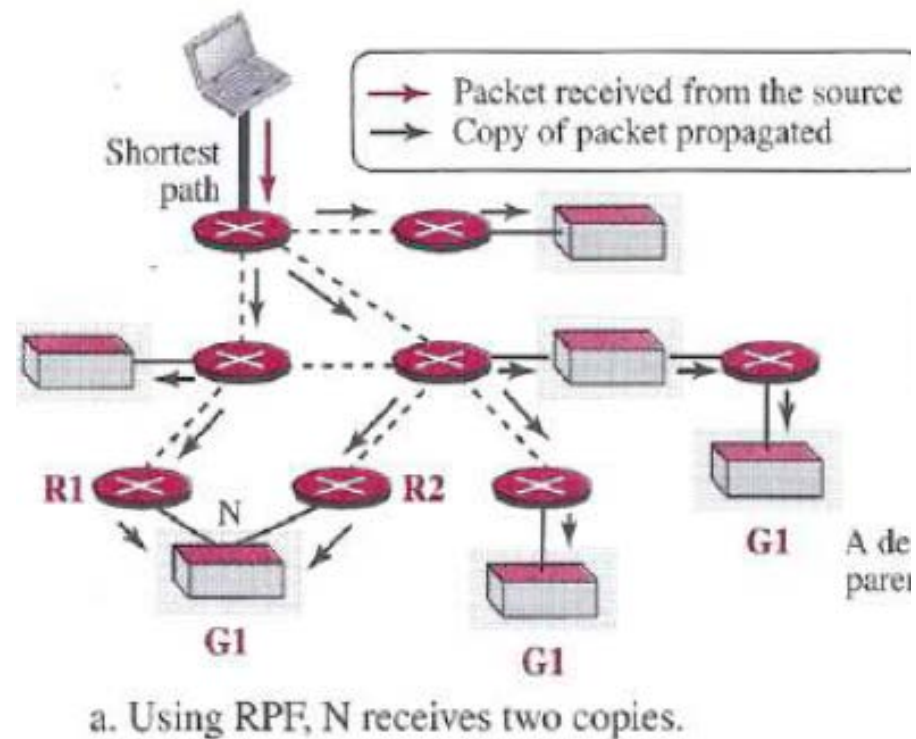


- Using **distance-vector + source-based tree** approach
  - Distance Vector Multicast Routing Protocol (**DVMRP**)
    - Extension of RIP and OSPF
- Using **link-state + source-based tree** approach
  - Multicast Open Shortest Path First (**MOSPF**)
- Using **distance-vector / link-state + source-based tree / group-shared tree** approach
  - Protocol Independent Multicast (**PIM**)

- Distance Vector Multicast Routing Protocol (DVMRP) is an **extension of RIP** for multicasting
- Router creates a multicast tree to forward multicast packet using the **3 steps**:
  - Router uses **reverse path forwarding** (RPF)
    - to create optimal source-based tree between source and itself
  - Router uses **reverse path broadcasting** (RPB)
    - to create a broadcast (spanning) tree whose root is router itself and whose leaves are all networks in the Internet
  - Router uses **reverse path multicasting** (RPM)
    - to create multicast tree by cutting some branches of the tree that end in network with no member in the group.

# Reverse Path Forwarding (RPF)

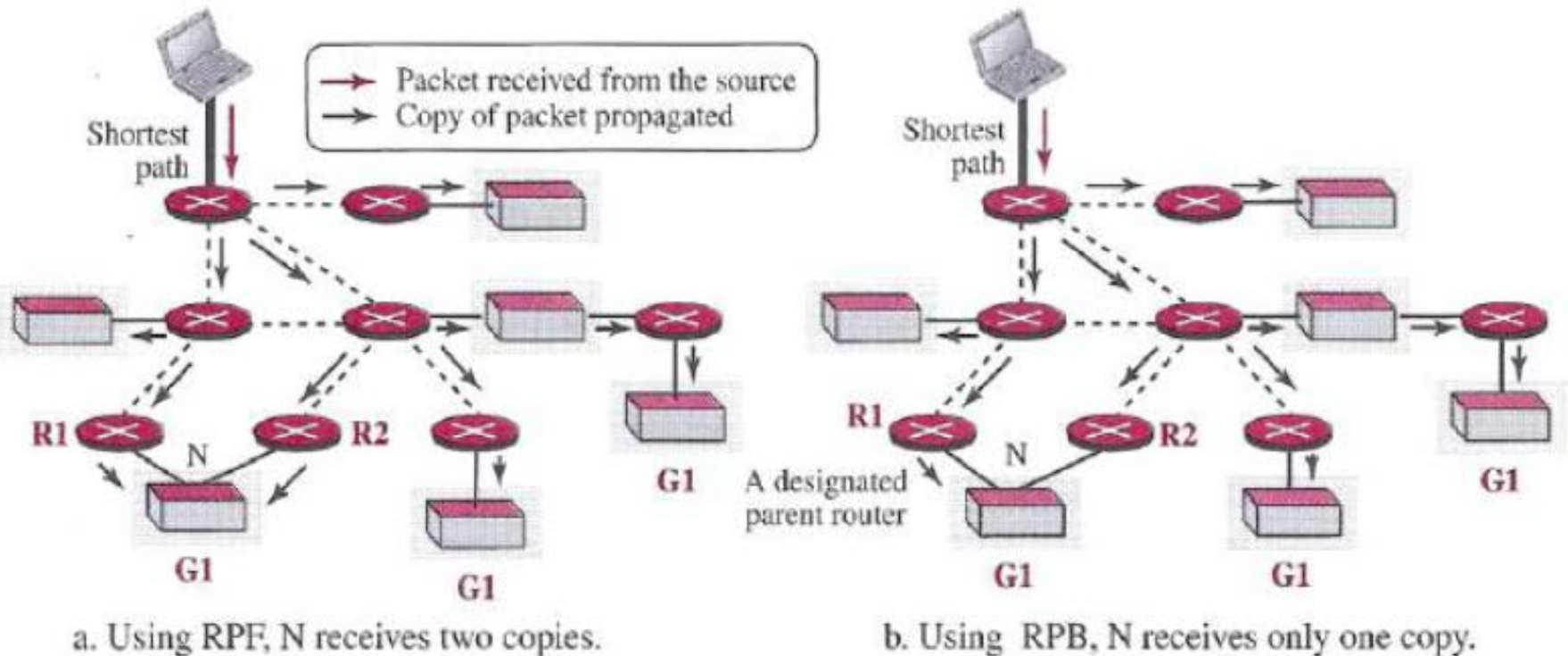
- Router forwards a multicast packet which has come through the **interface associated** with shortest path from source to the router
- Router does not know shortest path from source to itself; so, **consults with reverse path**
- prevents **looping & duplicate packet** receive





# Reverse Path Broadcasting (RPB)

- helps a router to **forward only one copy** received from a source and drop the rest
- we need to **allow only one of the routers** attached to a network to pass the packet to the network



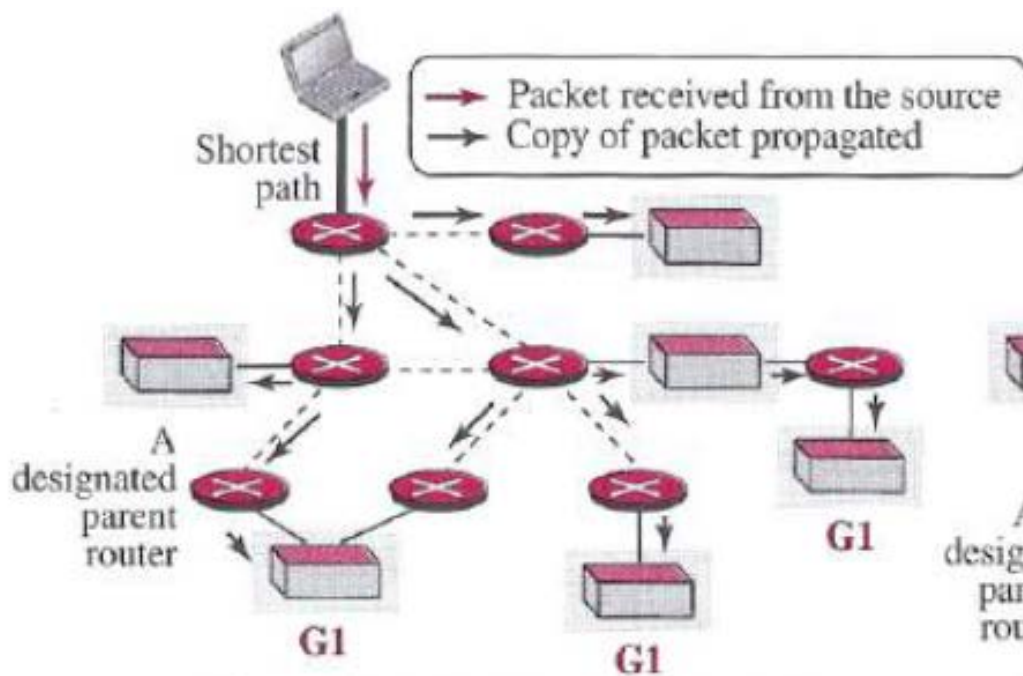
# Cont...



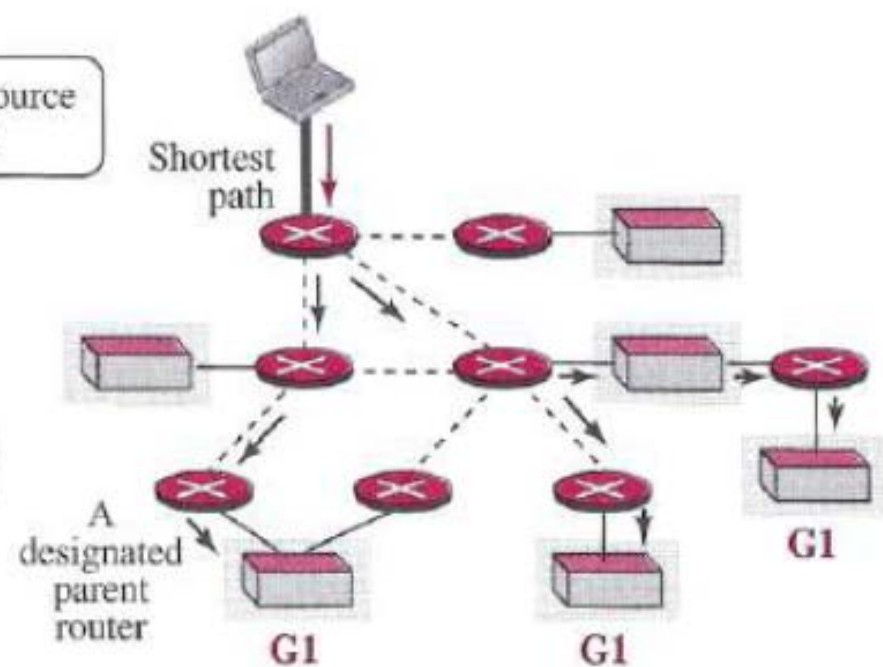
- designate only **one router as the *parent*** of a network related to a specific source.
  - Parent router forwards, others simply drop
- How to select the parent?
  - select the router that has **shortest path to the source**
  - If there is a tie in this case, the router with the **smaller IP address can be selected.**
- RPB creates broadcast **tree** from the **graph** created by RPF
- RPB **cuts** those **branches** of the tree that cause **cycles**
- Finally, we have a **shortest-path tree** with the **source** as the root and **all networks** (LANs) as the leaves.

# Reverse Path Multicasting (RPM)

- RPB does broadcast; so, not efficient
- To increase efficiency, we should do multicast



a. Using RPB, all networks receive a copy.



b. Using RPM, only members receive a copy.

# Cont...



- change the **broadcast shortest-path tree** to a **multicast shortest-path tree**
- How?
  - each router needs to **prune the interfaces** that do not reach a network with **active members** corresponding to a particular **source-group combination**.
  - Follow bottom-up approach
  - At the **leaf level**, the routers connected to the network **collect the membership information** using the **IGMP**
  - The parent router of the network can then **disseminate this information upward** using the reverse shortest-path tree from the router to the source
  - disseminated periodically; so **joining** and **leaving** is updated dynamically

# Multicast Link-State (MOSPF)

- MOSPF (Multicast Open Shortest Path First) is the extension of link-state unicast protocol OSPF
- Uses **source-based tree** approach
- In Link-state, router uses **LSDB** (link-state database) to create shortest-path tree
- For multicasting, router needs another database
  - to show which interface has active member in a particular group
- Let a router has received a packet from source S and to be sent to group G

# Cont...

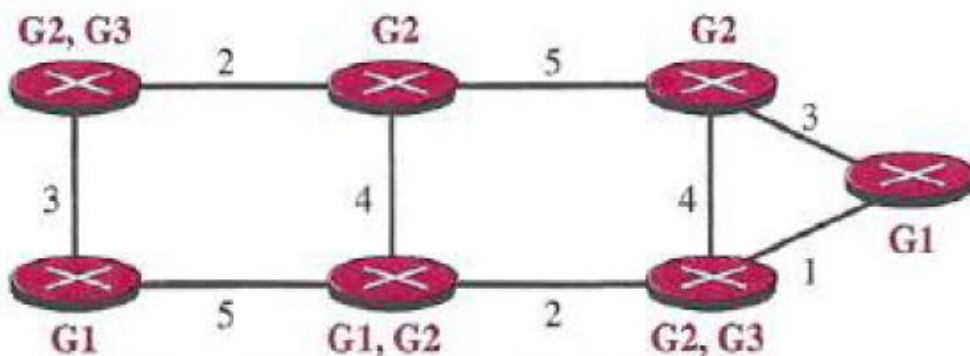


- Let a router has received a packet from source S and to be sent to group G
  - Router uses Dijkstra algorithm to create a shortest-path tree with **S as the root** (unlike unicasting in which **router itself is the root**) and all destinations in the internet as the leaves.
  - the router creates a shortest-path **subtree** with **itself as the root** of the subtree from the above tree.
  - The router prunes the shortest-path (broadcast) **subtree** to change it to a multicast tree.
  - How to get the membership information?
    - Using IGMP at the leaf level
    - Update the link state by flooding
  - Router then forwards through appropriate interface

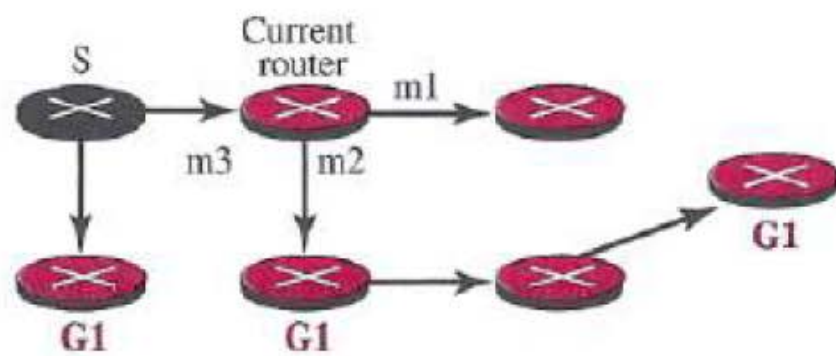


# Cont...

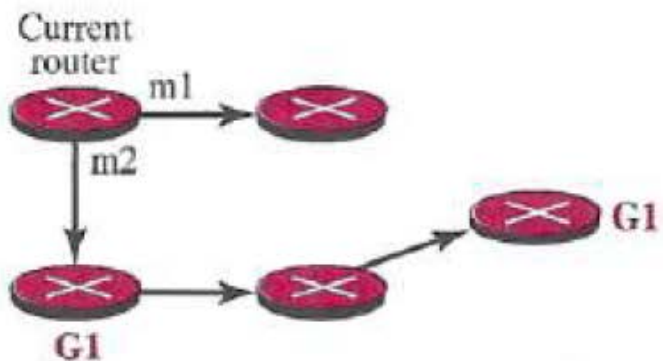
- Source S is attached with the top-left router
- Destination is G1



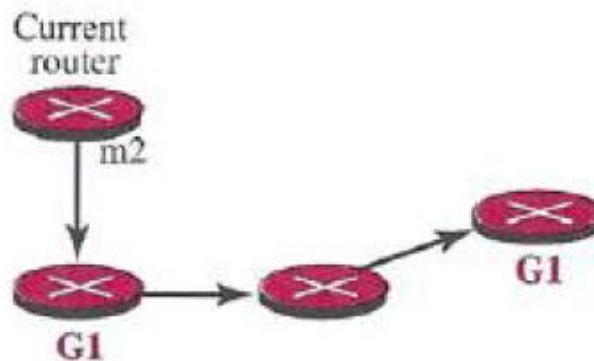
a. An internet with some active groups



b. S-G1 shortest-path tree



c. S-G1 subtree seen by current router



d. S-G1 pruned subtree

Forwarding table  
for current router

Group-Source	Interface
S, G1	m2
...	...

# Protocol Independent Multicast

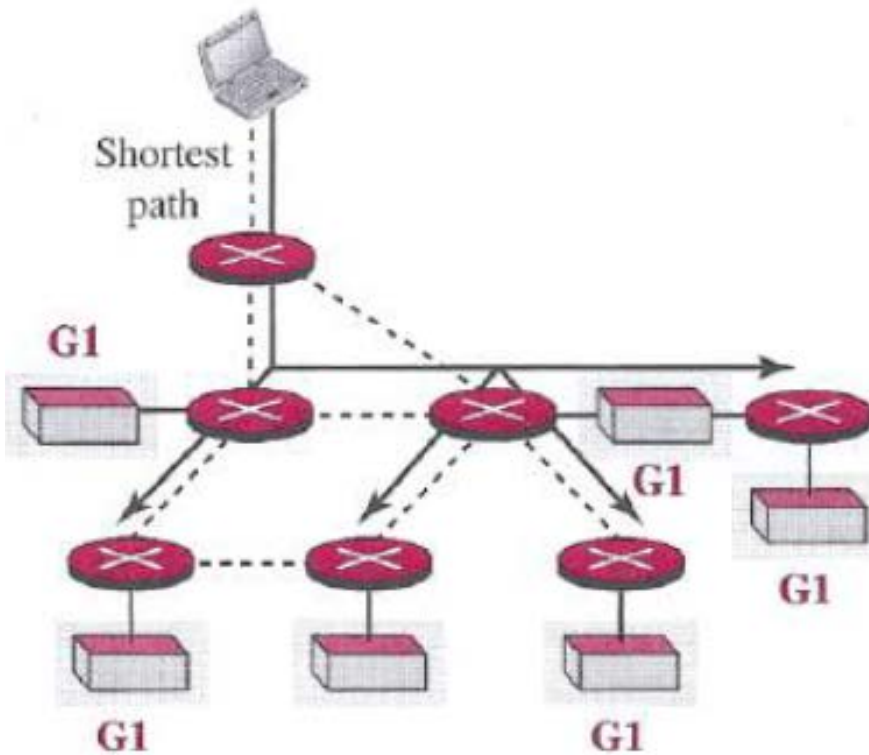


- PIM takes help of
  - Any type of unicast algo: distance vector / link state
  - Any type of multicast tree: source-based / group-shared
- PIM works in two modes:
  - PIM-DM for dense mode; uses source-based tree
  - PIM-SM for sparse mode; uses group-shared tree
- The term *dense* here means that the number of active members of a group in the internet is large
  - e.g. DM: popular teleconference that has a lot of members
  - e.g. SM: technical teleconference where a number of members are spread somewhere in the internet

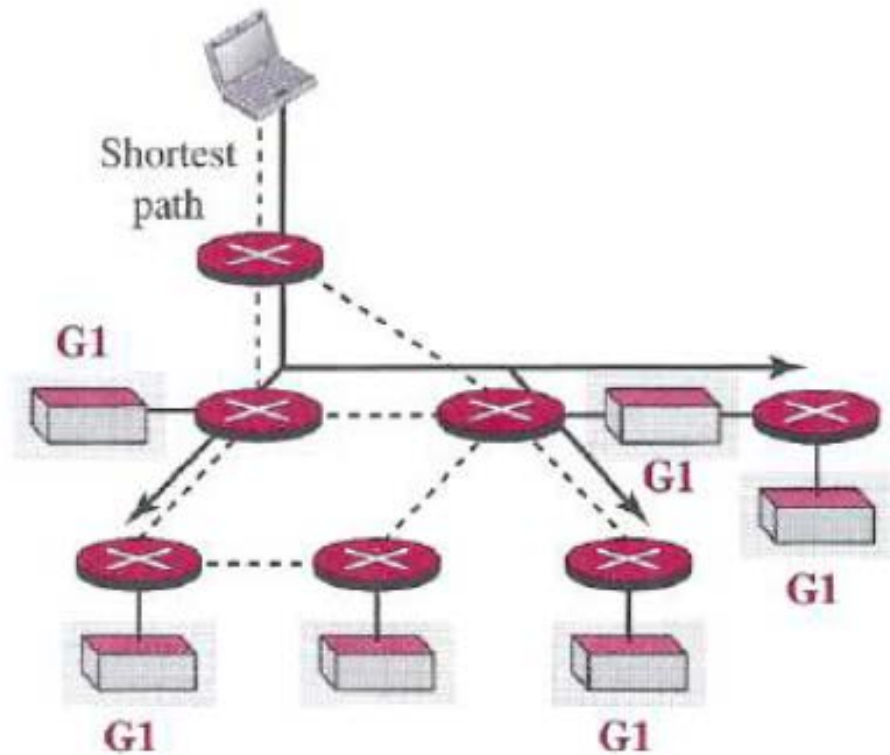


# PIM-DM

- PIM-DM uses only two strategies:
  - RPF & RPM. (no need of RPB)



a. First packet is broadcast.



b. Second packet is multicast.

# Cont...

- RPF is used to avoid receiving a **duplicate packet**
- If the packet has not arrived **from the next router in the reverse direction**, it drops the packet and sends a **prune message** in that direction to prevent receiving future packets related to pair (S, G).
- Else, router **forwards the packet** to all interfaces except
  - the receiving interface &
  - the interface from which it has already received a prune message related to (S, G).
- Initially, it is broadcast; but **over the time it switches to multicast** as pruned message arrives to the routers

# PIM-SM

- In this environment, the use of a protocol that **broadcasts** the packets **until the tree is pruned** is not justified
- PIM-SM uses a **group-shared tree approach** to multicasting
- In PIM-SM, we designate a router to act as **core router / rendezvous point** (RP)
- Multicast communication is achieved in **two steps**:
  - From **source** to **core router**: **Unicast**
  - From **core router** to **group members** : **Multicast**

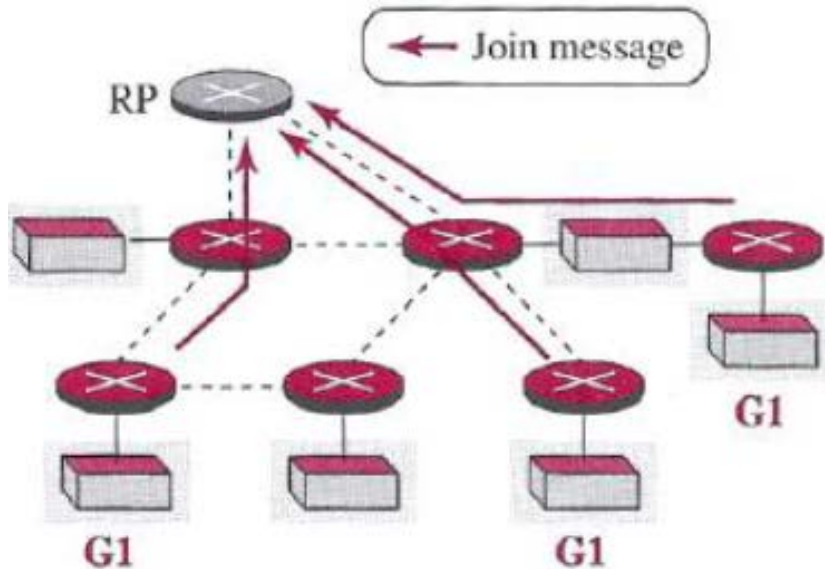
# Cont...



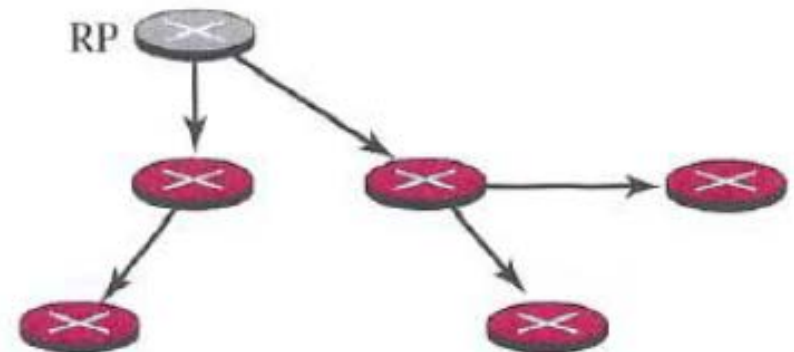
- How to select the **core router** for a group?
  - Following any suitable method
- PIM-SM uses a **spanning multicast tree** rooted at the **core router** with leaves pointing to designated routers connected to each network with an active member.
- How to form the **multicast tree** for a group?
  - The router should **know the unique interface** from which it should accept a multicast packet destined for a group. (likewise RPF)
  - It needs to **avoid delivering more than one copy** of the same packet to a network through several routers. (likewise RPB)
  - The router should **know the interface** or **interfaces** from which it should send out a multicast packet destined for a group. (like RPM)

# Cont...

- PIM-SM uses *join* and *prune* messages to create a multicast tree rooted at the RP (core router)
- Router maintains a *join counter*. It increases for each interface after receiving a join message through that interface.

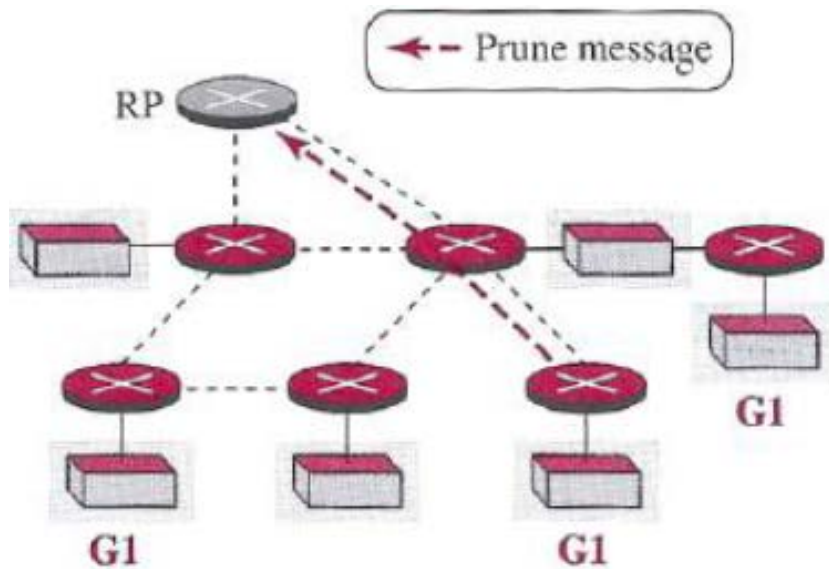


a. Three networks join group G1

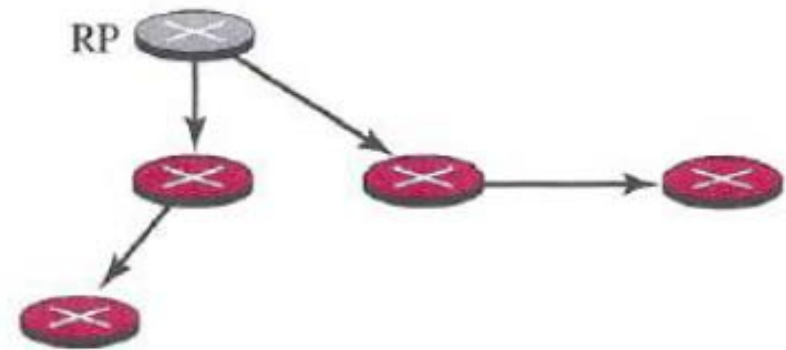


b. Multicast tree after joins

# Cont...



c. One network leaves group G1



d. Multicast tree after pruning

- When a router receives a prune message, it decrements the **join count** for the interface through which the message has arrived and forwards it to the next router.
- When the **join count for an interface reaches zero**, that interface is not part of the multicast tree anymore.

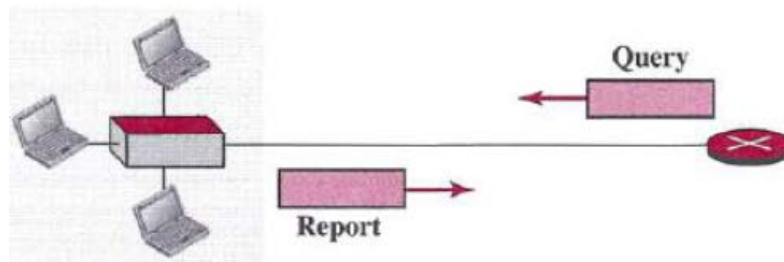
# Interdomain Multicast Protocol



- When the members of the groups are spread among different autonomous domains (ASs), we need an **interdomain multicast routing** protocol.
  - Multicast BGP
- MBGP provides **two paths** between ASs:
  - one for unicasting
  - one for multicasting
- Information about multicasting is exchanged **between border routers** in different ASs.
- MBGP is a **group-shared multicast** routing protocol in which one router in **each AS** is chosen as the core router.

# IGMP

- IGMP: Internet Group Management Protocol
- IGMP messages, like ICMP messages, are encapsulated in an IP datagram.
- IGMP uses two messages: **Query** and **Report**
- A **query message** is periodically sent by a router to all hosts attached to it to ask them to report their interests about membership in groups.
- A **report message** is sent by a host as a response to a query message.
- After a router has collected **membership information** from the hosts and other routers at its own level in the tree, it can **propagate it** to the router located in a higher level of the tree.





# Thanks!