

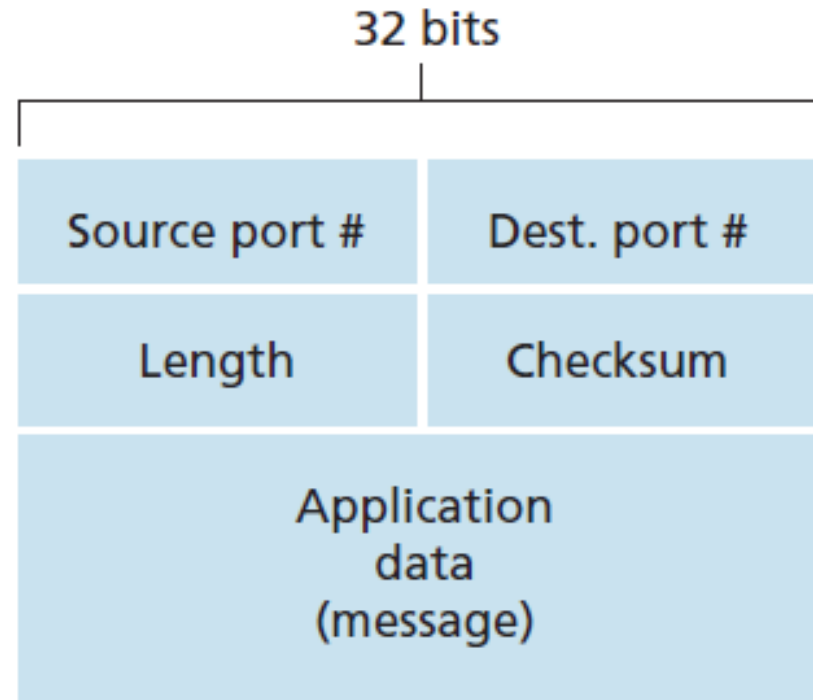
UDP and TCP

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Introduction of UDP

- UDP (User Datagram Protocol)
 - Transport-layer protocol
 - Connectionless
 - Simple
 - Efficient
 - Unreliable



UDP Services



- Process-to-process communication
 - Need socket address (IP + Port)
- Connectionless service
 - No sequence number
 - No relation between UDP datagrams
 - No connection establishment
 - Datagram can travel through different path
 - No segmentation (message size $< (65535-8)$)
- No flow control
 - No window mechanism
- No error control
 - Error detection through checksum but no control

Cont...



- No congestion control
 - Assumption is that congestion will not be created as UDP datagrams are small in size
- Encapsulation and decapsulation
 - Needs to send message from one process to another
- Multiplexing and demultiplexing
 - One UDP, but several process in application layer wants to use its services
- Queuing
 - Queues are associated with port

- Checksum
 - Consider three parts: Pseudoheader, UDP header, message from the upper layer; it is not mandatory;
 - Datalink layer has error detection mechanism. Why do we need checksum in transport layer?
 - neither link-by-link reliability nor in-memory error detection is guaranteed w.r.t. end-to-end service
 - Why do we need pseudoheader?
 - Socket-address need to be correct.
 - To ensure intended receiver (avoids in-memory error)

UDP Checksum



- What value is sent for the **checksum** in each one of the following hypothetical situations?
 - The sender decides **not to include** the checksum.
 - The sender decides to include the checksum, but the value of the **sum is all 1s**.
 - The sender decides to include the checksum, but the value of the **sum is all 0s**.
- Solution:
 - **All 0s**
 - When the sender complements the sum, the result is all 0s; the sender complements the result again before sending. The value sent for the checksum is all 1s. The **second complement** operation is needed to avoid confusion with the previous case. Note that this does not create confusion because the value of the checksum is never all 1s in a normal situation
 - This situation never happens because it implies that the value of every term included in the calculation of the sum is all 0s, which is impossible.

UDP Applications



- If the request and response can each fit in a single user datagram, a **connectionless** service may be preferable.
 - E.g. DNS request and response
 - But, not suitable in SMTP as e-mail size could be large
- **Lack of error control** is advantageous sometimes
 - E.g. real time communication through Skype
 - But, not suitable for file download
- **Lack of congestion control**
 - Advantageous in error-prone network

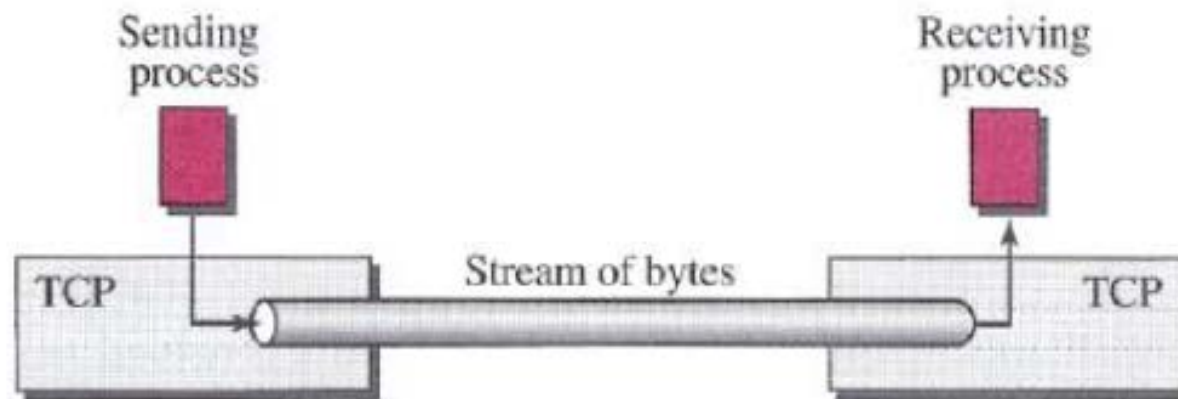
Introduction to TCP



- TCP (Transmission Control Protocol)
 - Connection-oriented (**but not virtual-circuit**)
 - Reliable
 - Create connection, do data transfer, tear down connection
 - Uses ARQ protocols (GBN and SR)
 - Checksum for error detection
 - Retransmission of lost / corrupted packets
 - Cumulative / selective ACK
 - Commonly used in Internet

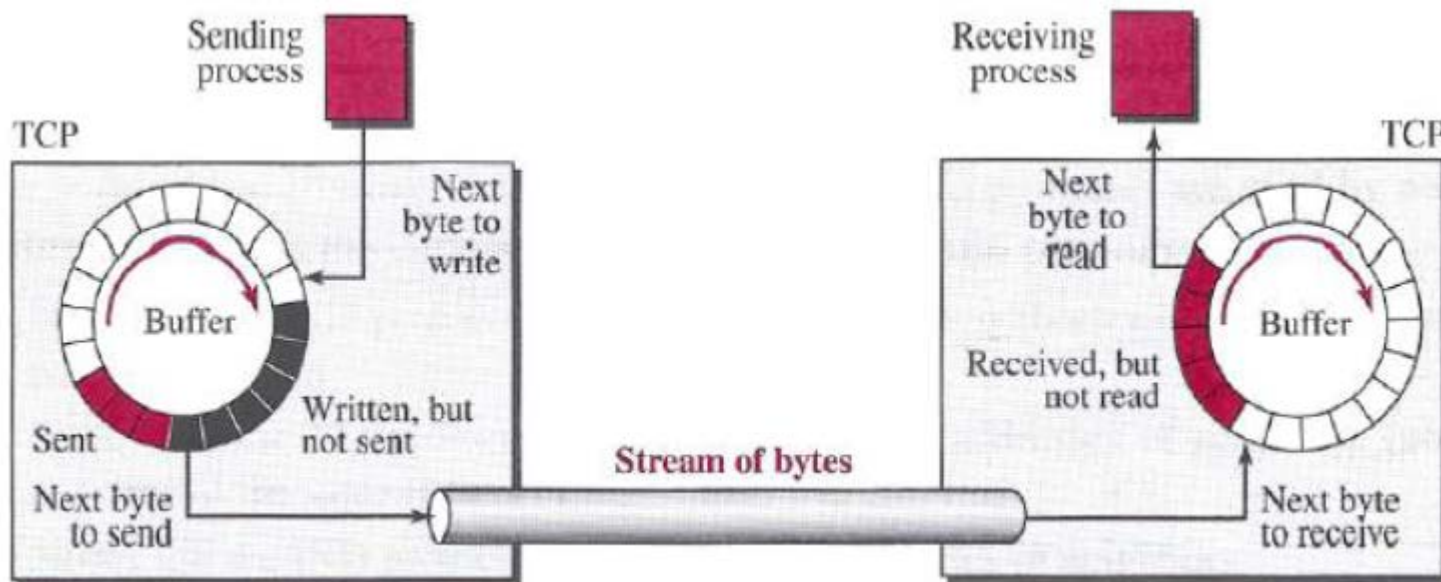
TCP Services

- Process-to-process communication
 - Needs socket address
- Stream Delivery Service



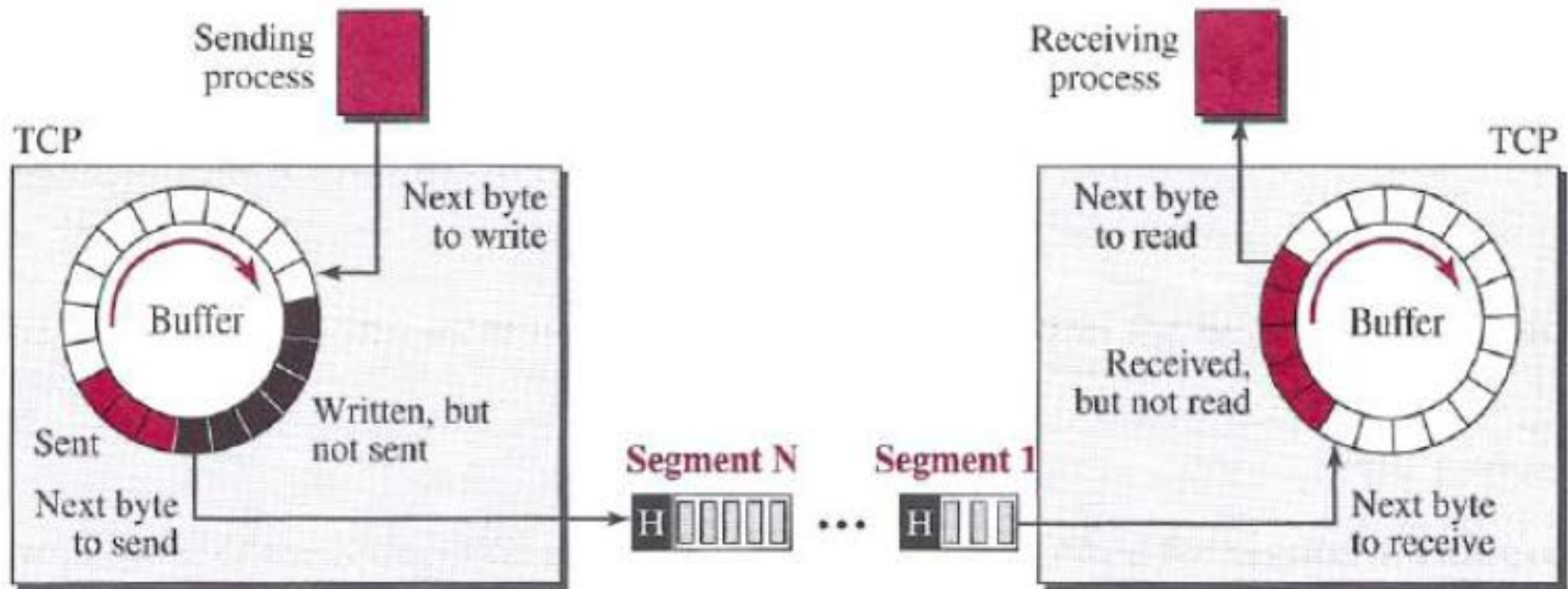
Cont...

- Sending and receiving buffer
 - may not necessarily write or read data at the same rate
 - Buffer allows to have flow control
- Full-Duplex Communication
 - Each TCP endpoint has its own sending and receiving buffer



Cont...

- Segmentation



- Reliable service
 - Using ACK

Cont...



- Multiplexing and Demultiplexing
 - But need connection between processes
- Connection-oriented service
 - two TCP's establish a logical connection
 - This is not a virtual-circuit switching
 - Connections are at two end systems only
 - Data are exchanged in both directions
 - The connection is terminated at the end of transmission

Numbering in TCP

- No segment number
- Instead, we have **sequence** number and **acknowledgement** number
 - These are **byte numbers**, but not segment number
- Number is independent in each direction
- TCP **numbers all data bytes** transmitted in a connection
- TCP chooses an arbitrary number between 0 and $2^{32} - 1$ for numbering the first byte

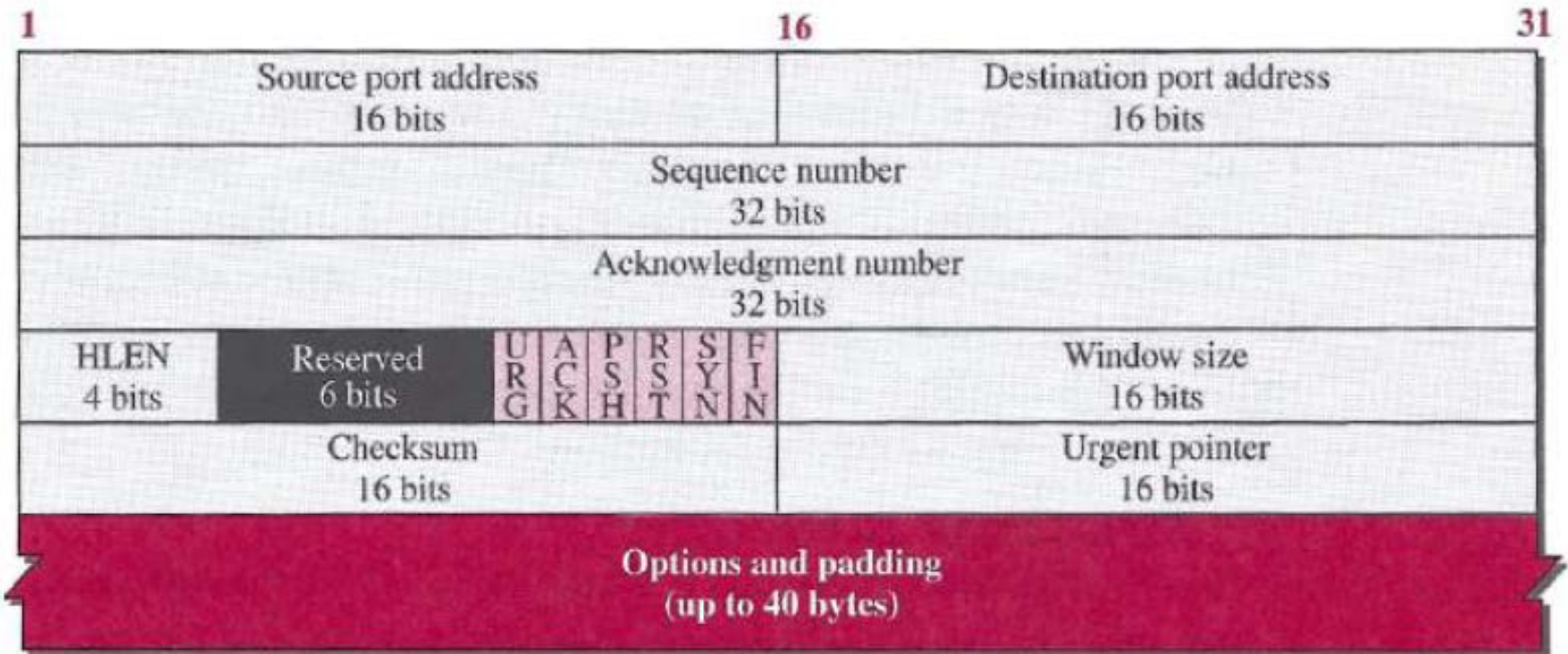
SEQ and ACK numbers

- After the bytes have been numbered, TCP assigns a **SEQ number to each segment** that is being sent
 - For **1st segment**: SEQ number is random
 - For any **other segment**: SEQ number of previous segment + number of byte in the previous segment
- Each party also uses an **ACK number** to confirm the bytes it has received.
- The ACK number defines the **number of the next byte** that the party expects to receive.
- The ACK number is **cumulative**

TCP Segment Format



a. Segment



b. Header

Fields in a Segment

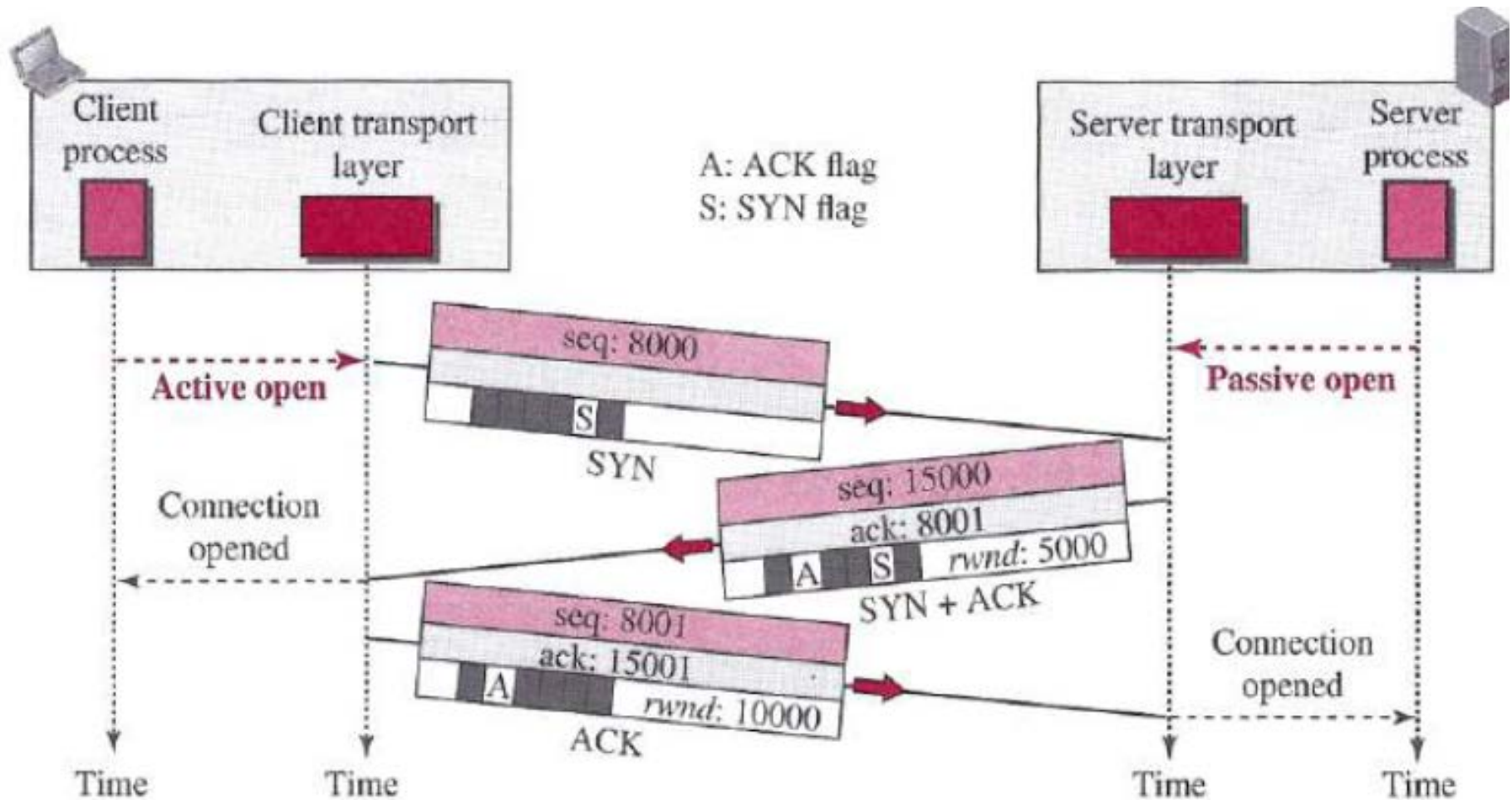
- Source & Destination port address.
- Sequence number & Acknowledgement number.
- Header length.
- Control. (6 control flags)
- Window size (window size of the sending TCP in bytes)
- Checksum. (pseudo header + TCP header + data)
- Urgent pointer (valid only if urgent flag is set)
- Options (optional information in the TCP header)

TCP Connection

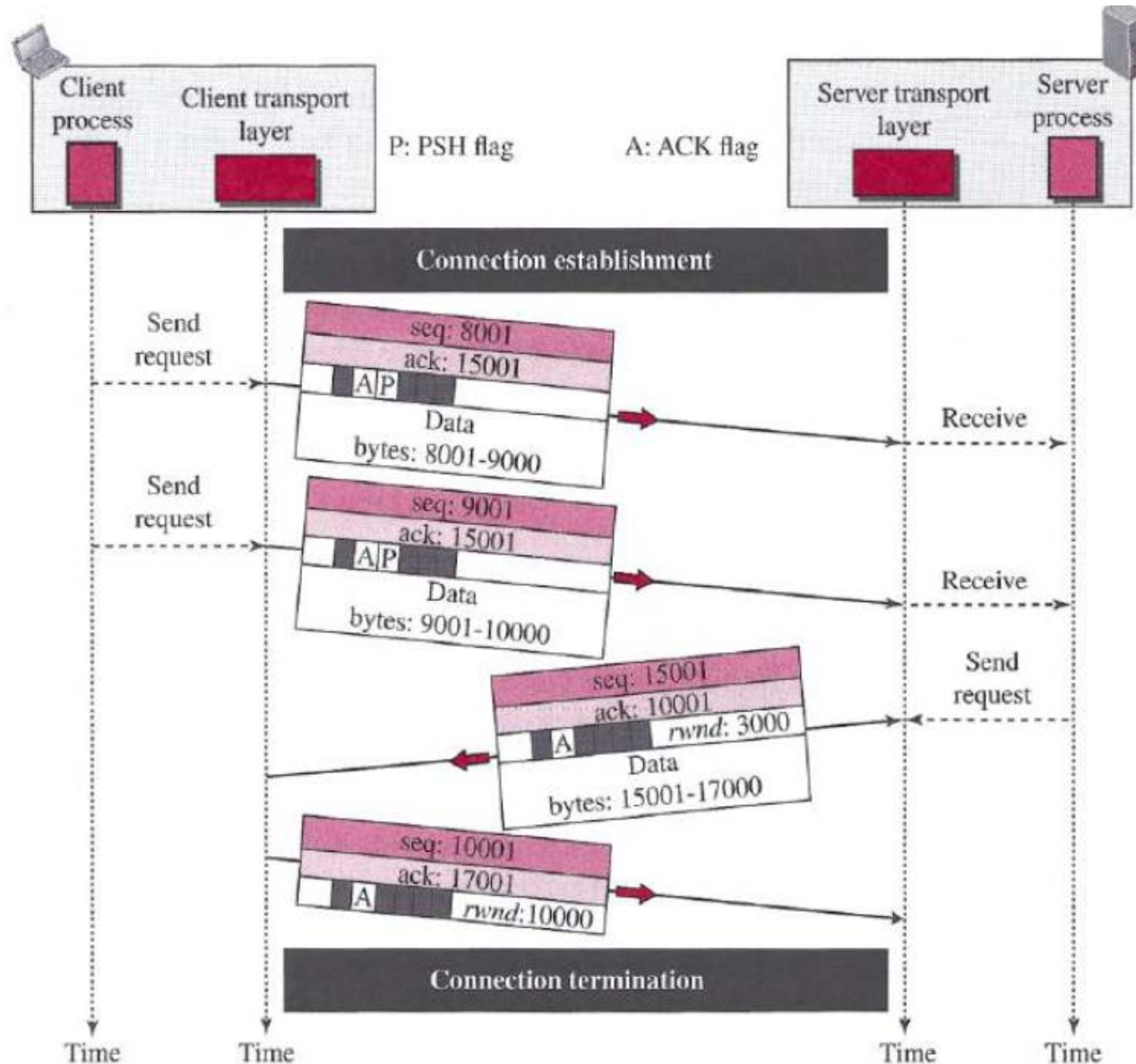
- TCP connection is **logical**, not physical.
- TCP is connection-oriented although IP is connectionless
- TCP operates in full-duplex mode
- TCP uses **three-way-handshaking**
- Let, an application program, called the **client**, wants to make a connection with another application program, called the **server**, using TCP
- The process **starts with** the server.
 - **Passive open** (server process informs transport layer of server that it is ready)
 - **Active open** (client issues request to client transport layer)
 - Now Client transport layer starts **three-way-handshaking**

Connection Creation : step 1

- Using three-way-handshaking

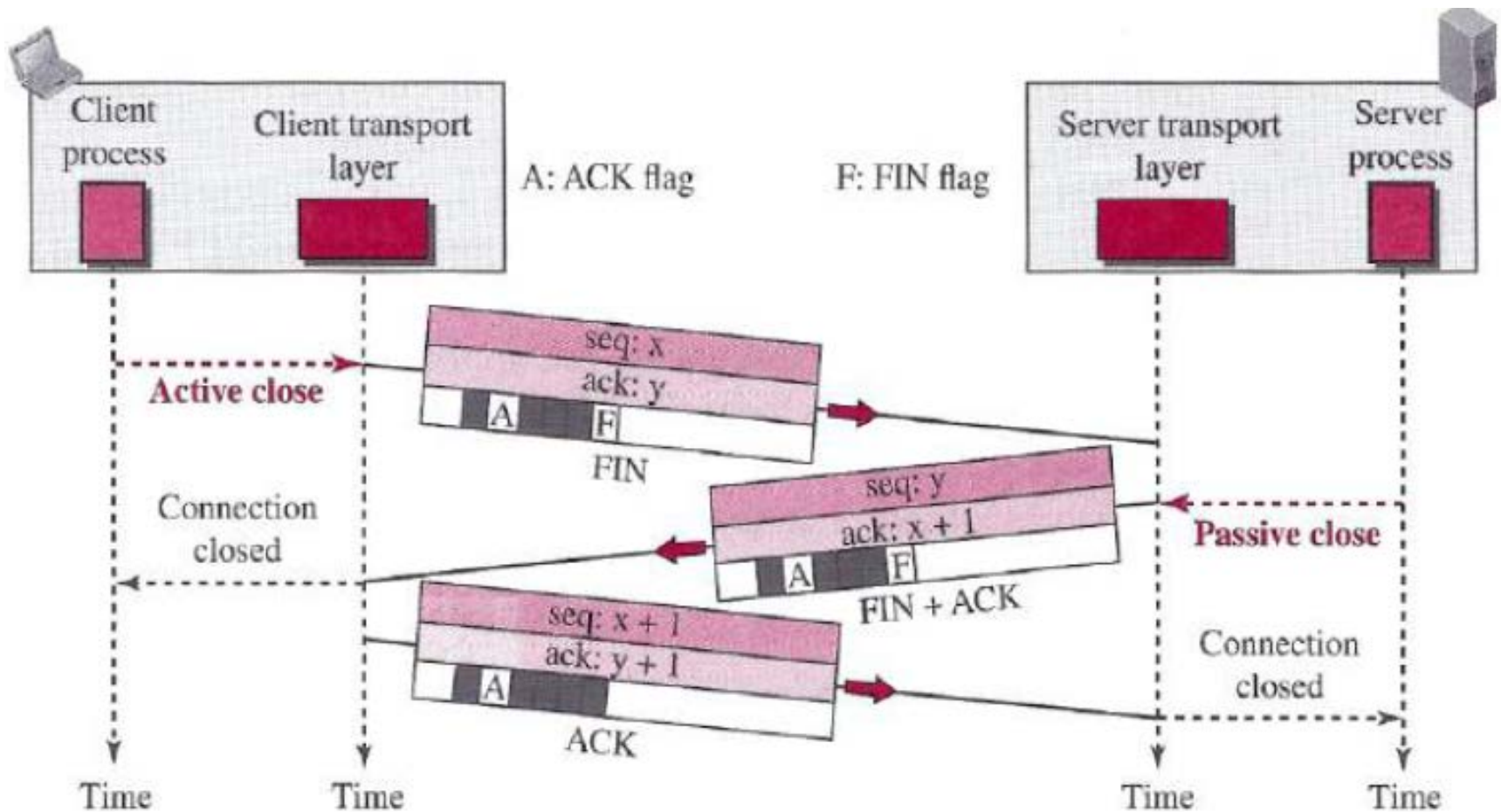


Data Transfer : step 2



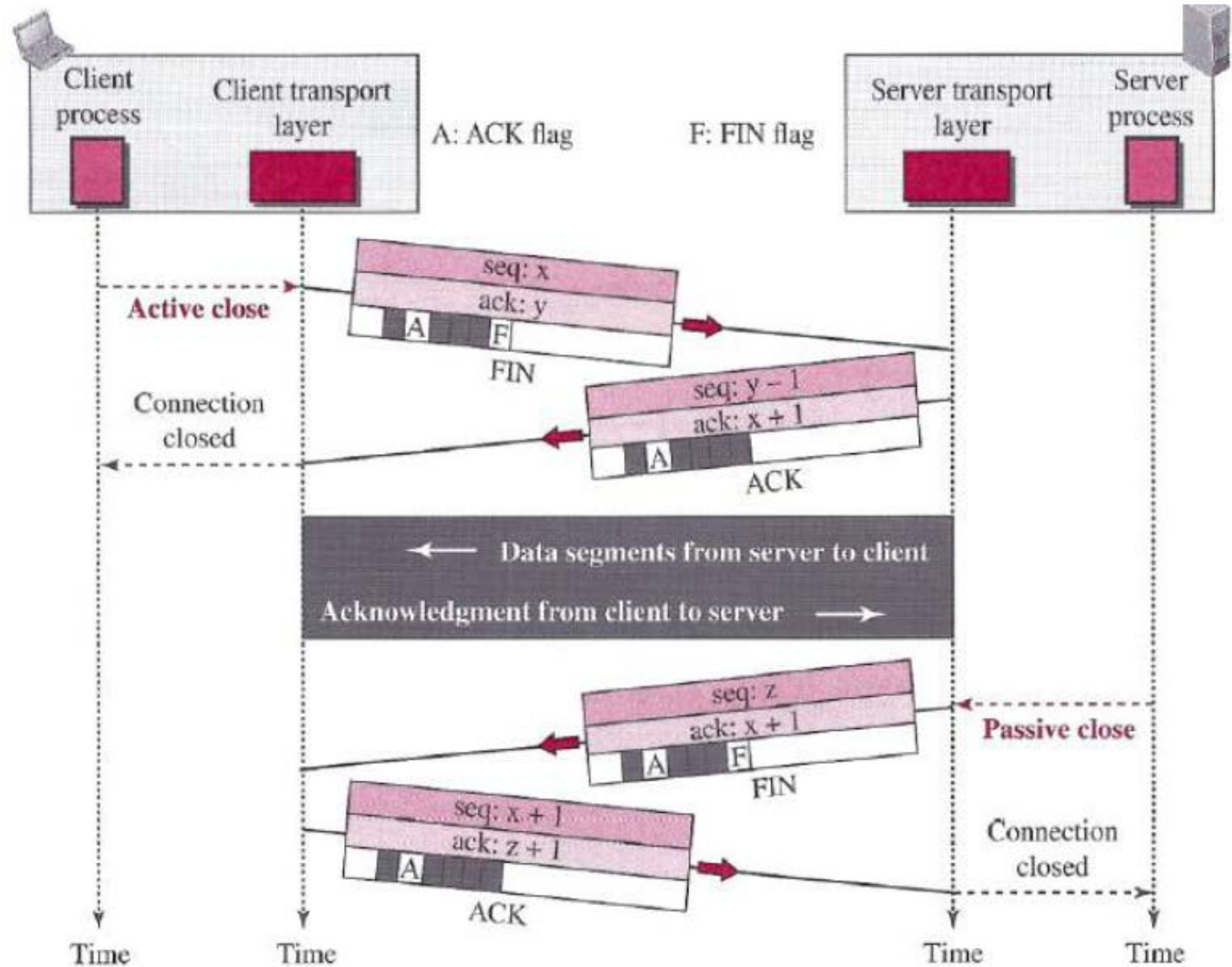
Connection Termination : step 3

- Using three-way-handshaking



Half-close Connection

Example:
Sorting
at server



Scenario

Create
Connection

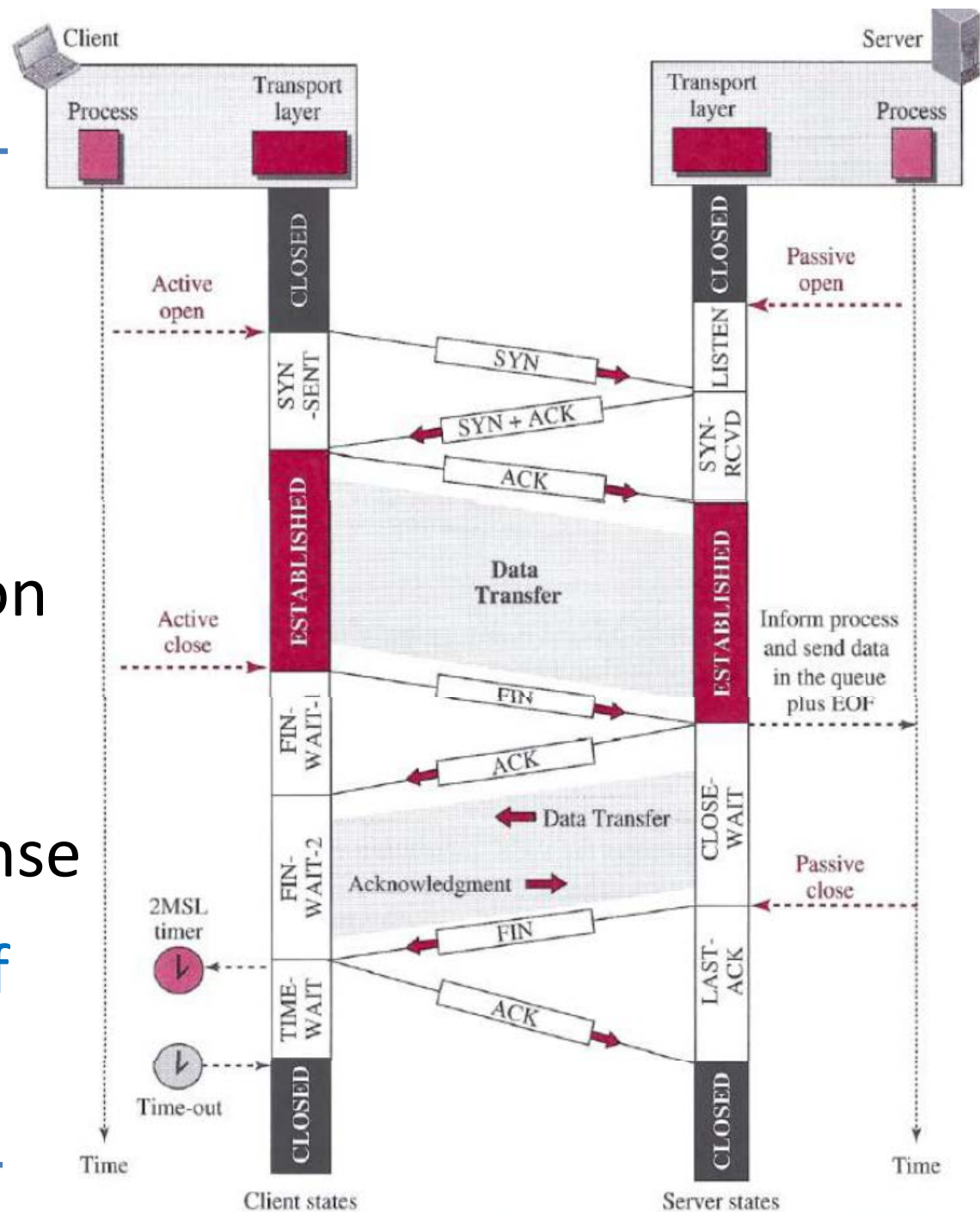
Data Transfer
in both direction

Half-close

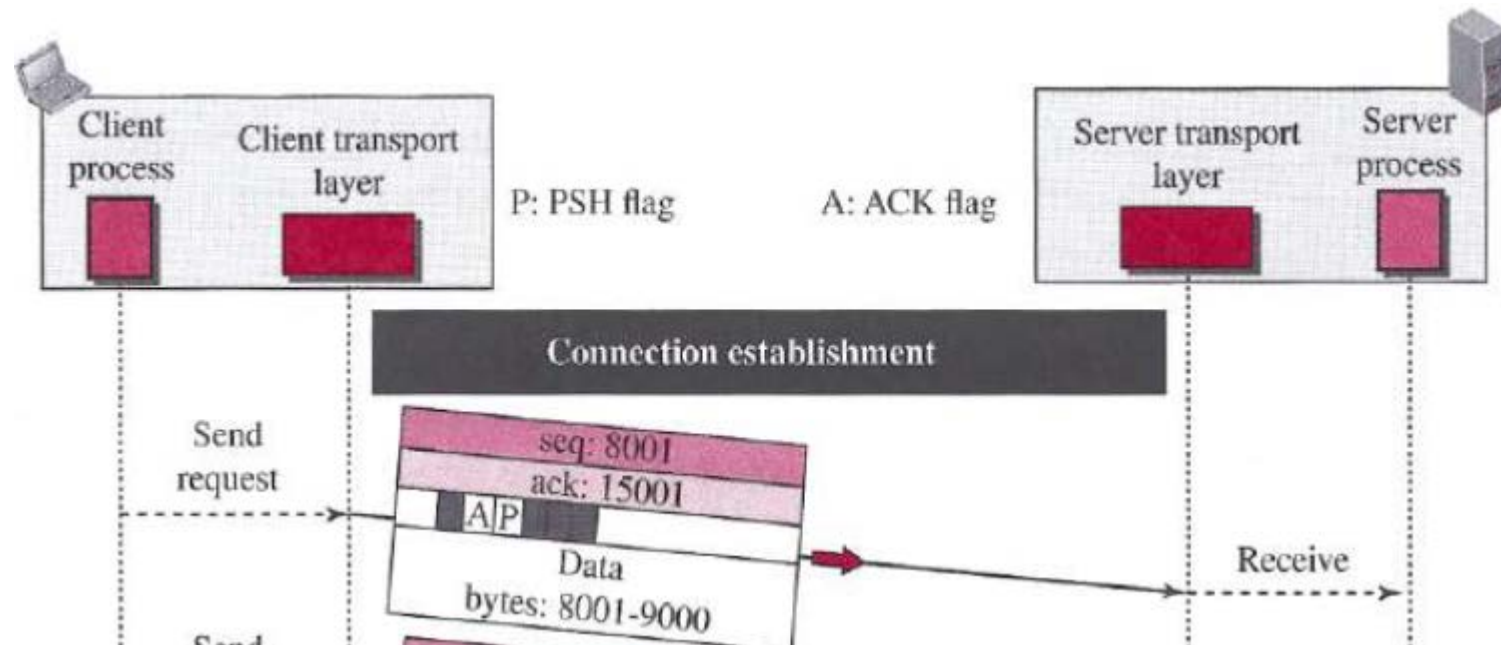
Receive Response

Close other half

Closed status



PUSH Flag



- PUSH means sending TCP must not wait for the window to be filled, and then send the segment
- PUSH flag informs the receiving TCP to deliver the received segment immediately to application program

Thanks!