

Dimensions in Pointer Analysis

CS6013: Modern Compilers - Theory and Practice

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April 8th, 2016

Pointer Analysis

- Establishes which pointers (or heap references) can point to which objects (or storage locations).
- Applications: Alias analysis, shape analysis, escape analysis, etc.



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- Applications: Alias analysis, shape analysis, escape analysis, etc.

```
class A {...}
class B extends A {...}
class C {
    public void foo() {
        A a1, a2, a3;
        a1 = new A();      //l1
        if(*)
            a2 = new A();  //l2
        else
            a2 = new B();  //l3
        a3 = a1;
    }
}
```



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- Establishes which pointers (or heap references) can point to which objects (or storage locations).
- Applications: Alias analysis, shape analysis, escape analysis, etc.

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            a2 = new B();  //l3
        a3 = a1;
    }
}
```

Points-to sets:

- $a1 \rightarrow \{l1\}$
- $a2 \rightarrow \{l2, l3\}$
- $a3 \rightarrow \{l1\}$



Overview

1 Pointer Analysis

2 Analysis Dimensions

- Flow-sensitivity
- Field-sensitivity
- Interprocedural analysis
- Context-sensitivity

3 Application

4 Conclusion



Flow-sensitivity

Flow-sensitive: Maintain information at each point of the program.

```
class C {  
    public void foo() {  
        A a1, a2, a3;  
1:      a1 = new A();      //l1  
2:      if(*)  
3:          a2 = new A();  //l2  
4:      else  
5:          a2 = new A();  //l3  
6:      a1 = a2;  
7:      a3 = a2;  
8:      ...  
    }  
}
```



Flow-sensitivity

Flow-sensitive: Maintain information at each point of the program.

```
class C {  
    public void foo() {  
        A a1, a2, a3;  
1:    a1 = new A();    //l1  
2:    if(*)  
3:        a2 = new A(); //l2  
4:    else  
5:        a2 = new A(); //l3  
6:    a1 = a2;  
7:    a3 = a2;  
8:    ...  
    }  
}
```

Flow-insensitive points-to sets:

- a1: {l1, l2, l3}
- a3: {l2, l3}

Flow-sensitivity

Flow-sensitive: Maintain information at each point of the program.

```
class C {  
    public void foo() {  
        A a1, a2, a3;  
1:      a1 = new A();      //l1  
2:      if(*)  
3:          a2 = new A();  //l2  
4:      else  
5:          a2 = new A();  //l3  
6:      a1 = a2;  
7:      a3 = a2;  
8:      ...  
    }  
}
```

Flow-insensitive points-to sets:

- a1: {11, 12, 13}
- a3: {12, 13}

Flow-sensitive points-to sets for:

- a1:
 - {} till line no. 1
 - {11} from line nos. 1 to 6
 - {12, 13} afterwards
- a3:
 - {} till line no. 7
 - {12, 13} afterwards



Field-sensitivity

Field-sensitive: Maintain information separately for fields of an object.

```
class A {A f1; A f2;}  
class C {  
    public void foo() {  
        A a1;  
        a1 = new A();    //l1  
        a1.f1 = new A(); //l2  
        a1.f2 = new A(); //l3  
    }  
}
```

Field-sensitivity

Field-sensitive: Maintain information separately for fields of an object.

```
class A {A f1; A f2;}  
class C {  
    public void foo() {  
        A a1;  
        a1 = new A();    //l1  
        a1.f1 = new A(); //l2  
        a1.f2 = new A(); //l3  
    }  
}
```

Field-insensitive points-to sets:

- $a1 \rightarrow \{l1, l2, l3\}$

Field-sensitivity

Field-sensitive: Maintain information separately for fields of an object.

```
class A {A f1; A f2;}  
class C {  
    public void foo() {  
        A a1;  
        a1 = new A();    //l1  
        a1.f1 = new A(); //l2  
        a1.f2 = new A(); //l3  
    }  
}
```

Field-insensitive points-to sets:

- $a1 \rightarrow \{l1, l2, l3\}$

Field-sensitive points-to sets:

- $a1 \rightarrow \{l1\}$
- $a1.f1 \rightarrow \{l2\}$
- $a1.f2 \rightarrow \{l3\}$

Intraprocedural vs Interprocedural analyses

Say, method `f0` calls method `f1`.

Intraprocedural analysis:

- Information computed for `f0` ignores the points-to results of `f1`.
- Conservative assumptions are made at call sites.



Intraprocedural vs Interprocedural analyses

Say, method `f0` calls method `f1`.

Intraprocedural analysis:

- Information computed for `f0` ignores the points-to results of `f1`.
- Conservative assumptions are made at call sites.

Interprocedural analysis:

- Information computed for `f0` considers the points-to results of `f1`.
- Requires a call-graph.



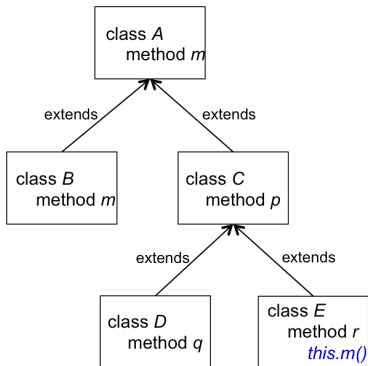
Call-Graph Construction

- A call-graph is needed to determine the possible callees at a call-site.
 - Offline as a pre-analysis.
 - On-the-fly using points-to results.



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- A call-graph is needed to determine the possible callees at a call-site.
 - Offline as a pre-analysis.
 - On-the-fly using points-to results.
- Class Hierarchy Analysis (CHA)



CHA helps in determining that only one implementation of *m* can be called.

Context-sensitivity

Context-sensitive: Maintain different results for different contexts from which a method is called.



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What is a context?



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Context-sensitive: Maintain different results for different contexts from which a method is called.

What is a context?

- Call-site-sensitivity
- Object-sensitivity



1-Call-site-sensitive

```
class A { fb()...}
class B extends A { fb()...}

class C {
    A a1;

    public void foo() {
        a1 = new A();    // l1
c1:   bar(a1);
        a1 = new B();    // l2
c2:   bar(a1);
    }

    public void bar(A p1) {
        p1.fb();
    }
}
```

1-Call-site-sensitive

```
class A { fb()...}  
class B extends A { fb()...}  
  
class C {  
    A a1;  
  
    public void foo() {  
        a1 = new A();    // 11  
c1:    bar(a1);  
        a1 = new B();    // 12  
c2:    bar(a1);  
    }  
  
    public void bar(A p1) {  
        p1.fb();  
    }  
}
```

Context-insensitive:

- $a1 \rightarrow \{11, 12\}$
- Both A's and B's fb can be called.

1-Call-site-sensitive

```

class A { fb()...}
class B extends A { fb()...}

class C {
    A a1;

    public void foo() {
        a1 = new A();    // 11
c1:   bar(a1);
        a1 = new B();    // 12
c2:   bar(a1);
    }

    public void bar(A p1) {
        p1.fb();
    }
}

```

Context-insensitive:

- $a1 \rightarrow \{11, 12\}$
- Both A's and B's fb can be called.

1-Call-site-sensitive:

- $a1 \rightarrow \{11\}$
A's fb will be called.
- $a1 \rightarrow \{12\}$
B's fb will be called.

2-Call-site-sensitive

```
class C {  
    public void foo1() {  
        bar();  
    }  
  
    public void foo2() {  
        bar();  
    }  
  
    public void bar() {  
        fb();  
    }  
}
```

2-Call-site-sensitive

```
class C {  
    public void foo1() {  
        bar();  
    }  
  
    public void foo2() {  
        bar();  
    }  
  
    public void bar() {  
        fb();  
    }  
}
```

1-Call-site-sensitive:

- 1 context for fb

2-Call-site-sensitive

```
class C {  
    public void foo1() {  
        bar();  
    }  
  
    public void foo2() {  
        bar();  
    }  
  
    public void bar() {  
        fb();  
    }  
}
```

1-Call-site-sensitive:

- 1 context for fb

2-Call-site-sensitive:

- 2 contexts for fb

1-Object-sensitive

Distinguish contexts based on the allocation site of the receiver.



1-Object-sensitive

Distinguish contexts based on the allocation site of the receiver.

```
main() {  
    o1 = new A();  
    o2 = new A();  
    o1.bar();  
    o2.foo();  
    o2.bar();  
    o2.foo();  
}  
  
foo() {  
    ...  
}  
  
bar() {  
    ...  
}
```

1 context for foo; 2 contexts for bar



2-Object-sensitive

Distinguish contexts based on:

- Allocation site of receiver
- Allocation site of allocator of receiver



2-Object-sensitive

Distinguish contexts based on:

- Allocation site of receiver
- Allocation site of allocator of receiver

```
main() {  
    o1 = new A();  
    o1.foo();  
    o1 = new A();  
    o1.foo();  
}  
  
foo() {  
    o2 = new A();  
    o2.bar();  
}  
  
bar() {...}
```

1-Object-sensitive:

- 1 context for bar

2-Object-sensitive:

- 2 contexts for bar

Which context-sensitivity is better?

```
main() {  
    o1 = new A();  
    o1.foo();  
    o1.foo();  
}
```

```
foo() {  
    o2 = new A();  
    o2.bar();  
}
```

```
bar() {...}
```

2-Object-sensitive:

- 1 context for bar

Which context-sensitivity is better?

```
main() {  
    o1 = new A();  
    o1.foo();  
    o1.foo();  
}
```

```
foo() {  
    o2 = new A();  
    o2.bar();  
}
```

```
bar() {...}
```

2-Object-sensitive:

- 1 context for bar

2-Call-site-sensitive:

- 2 contexts for bar
- No change in precision

Which context-sensitivity is better?

```
main() {  
    o1 = new A();  
    o1.foo();  
    o1.foo();  
}
```

```
foo() {  
    o2 = new A();  
    o2.bar();  
}
```

```
bar() {...}
```

2-Object-sensitive:

- 1 context for bar

2-Call-site-sensitive:

- 2 contexts for bar
- No change in precision

There is no thumb-rule for choosing the type of context-sensitivity; it depends on the application and the desired precision.



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Escape Analysis

Definition

An object is said to *escape* from a method/thread if it can be accessed in another method/thread.



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An object is said to *escape* from a method/thread if it can be accessed in another method/thread.

In Java, an object may escape the allocating method when:

- Passed as an argument to another method.
- Returned from the method.
- Accessible by a static (global) variable. (thread-escape)



Escape Analysis

Definition

An object is said to *escape* from a method/thread if it can be accessed in another method/thread.

In Java, an object may escape the allocating method when:

- Passed as an argument to another method.
- Returned from the method.
- Accessible by a static (global) variable. (thread-escape)

Escape analysis helps in:

- Stack allocation
- Synchronization elimination



Example

```
class A {...}
class C {
    static A global;
    public void foo() {
        A a1, a2, a3;
        a1 = new A();    //l1
        a2 = new A();    //l2
        a3 = new A();    //l3
        global = a2;
        a1.m2();
    }
}
```

Example

```
class A {...}
class C {
    static A global;
    public void foo() {
        A a1, a2, a3;
        a1 = new A();    //11
        a2 = new A();    //12
        a3 = new A();    //13
        global = a2;
        a1.m2();
    }
}
```

Points-to sets:

- a1 → {11}
- a2 → {12}
- a3 → {13}
- global → {12}

Example

```
class A {...}
class C {
    static A global;
    public void foo() {
        A a1, a2, a3;
        a1 = new A();    //l1
        a2 = new A();    //l2
        a3 = new A();    //l3
        global = a2;
        a1.m2();
    }
}
```

Points-to sets:

- a1 → {l1}
- a2 → {l2}
- a3 → {l3}
- global → {l2}

Escape analysis results:

- l1 escapes foo()
- l2 escapes foo() as well as the thread
- l3 does not escape



Exercise: Flow-sensitivity

```
class C {  
    static A global;  
  
    public void foo() {  
p1:    A a1 = new A(); //l1  
p2:    A a2 = new A(); //l2  
p3:    global = a1;  
        ...  
    }  
}
```

Exercise: Flow-sensitivity

```
class C {  
    static A global;  
  
    public void foo() {  
p1:      A a1 = new A(); //l1  
p2:      A a2 = new A(); //l2  
p3:      global = a1;  
        ...  
    }  
}
```

Flow-insensitive:

- l1 escapes the thread.

Flow-sensitive:

- l1 escapes the thread after the point p3.



Exercise: Field-sensitivity

```
class C {  
    static A global;  
  
    public void foo() {  
        A a1 = new A(); //l1  
        a1.f = new A(); //l2  
        A a2 = new A(); //l3  
        global = a1.f;  
    }  
}
```

Exercise: Field-sensitivity

```
class C {  
    static A global;  
  
    public void foo() {  
        A a1 = new A(); //11  
        a1.f = new A(); //12  
        A a2 = new A(); //13  
        global = a1.f;  
    }  
}
```

Field-insensitive:

- 11 and 12 escape the thread.

Field-sensitive:

- 12 escapes the thread.



Exercise: Context-sensitivity

```
class A {  
    A f;  
    public void bar() {  
        A b3 = new A(); //l4  
        this.f = b3;  
    }  
}  
  
class C {  
    static A global;  
    public void foo() {  
        A a1 = new A(); //l1  
        a1.bar();        //c1  
        global = a1;  
        a1.bar();        //c2  
    }  
}
```



Exercise: Context-sensitivity

```
class A {  
    A f;  
    public void bar() {  
        A b3 = new A(); //l4  
        this.f = b3;  
    }  
}
```

```
class C {  
    static A global;  
    public void foo() {  
        A a1 = new A(); //l1  
        a1.bar(); //c1  
        global = a1;  
        a1.bar(); //c2  
    }  
}
```

Context-insensitive (bar):

- 14 escapes the thread.

Context-sensitive (bar):

- 14 does not escape the thread from the call at c1.
- 14 escapes the thread from the call at c2.



Conclusion

- There are various dimensions along which the precision of a pointer analysis can be improved.
- Usually there is a tradeoff between the precision and the efficiency of an analysis.
- The dimensions that we discussed can be applied to improve the precision of other program analyses as well.



Conclusion

- There are various dimensions along which the precision of a pointer analysis can be improved.
- Usually there is a tradeoff between the precision and the efficiency of an analysis.
- The dimensions that we discussed can be applied to improve the precision of other program analyses as well.

Thank You.



Pointers for the enthusiast

- Vivien F. and Rinard M., *Incrementalized Pointer and Escape Analysis*, PLDI 2001.
- Hardekopf B. and Lin C., *The Ant and the Grasshopper: Fast and Accurate Pointer Analysis for Millions of Lines of Code*, PLDI 2007.
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- Slides: www.cse.iitm.ac.in/~manas/docs/cs6013-dpa.pdf