

Week 2: Intelligent Agents and

<u>Course</u> > <u>Uninformed Search</u> Week 2 Project: Implementation

FAQs

> Week 2 Project: Search Algorithms >

Week 2 Project: Implementation FAQs

Q. Is there skeleton code available for the assignment?

A. Yes, there is **skeleton code available for use** (see the third tab of this assignment). This is **completely optional to use**, and you may alter it as much as you'd like. You may also complete the assignment without referring to the skeleton code at all.

Q. My search algorithm seems correct but is too slow. How can I reduce its running time?

A. Search algorithm is perhaps one of the best learning materials for computational complexity and Python's idiosyncrasies. There are four dos and don'ts:

1. Don't store possibly-large data member such as solution path in search tree node class. Instead, rethink what operation should be fast.

Explanation: Storing a path from the root node in each node class achieves O(1) lookup time at the expense of O(n) creation time. For example, if the current state is visited after 60,000 intermediate states, the current state has to allocate a list of 60,000 elements, and each of the children states have to allocate a list of 60,001 elements. This would soon use up physical memory, and typically your machine's Operating System kills the search process.

A key observation is that in the case of search algorithm, path lookup operation is executed just once after search finishes. Thus, the lookup operation is fine to be slow. You might consider another data structure having O(n) lookup time for solution path but requiring O(1) operations during search.

2. Don't be satisfied by just using list as frontier. Instead, design your

Frontier class which works faster.

Explanation: one major bottleneck of list, deque, or queue class in Python is that their **membership testing operation** is O(n). The membership testing speed is critical for search algorithm because that operation is executed for every child state. Coming up with using such list-like data structures is a good first step, but for using it with reasonable execution time, you might need one more trick.

Note that pseudocode in lecture slides does not necessarily reflect implementation details (i.e. time and space). Rather, it conceptually shows the algorithm's inputs, processing orders, and outputs. One of your missions in this assignment is to make the "frontier" thing into a reality by using Python's "low-level" primitives.

3. Don't use O(n) operation when you have another faster way to do the same thing.

Explanation: Roughly speaking, if you set one O(n) operation under your for neighbor in neighbors loop, your code will be highly likely to exceed grading time limit. In other words, it happens that your code executes drastically faster when you fix just one line of your code.

For example, merging two **set**s and checking an element is in the merged set is an expensive operation, while checking an element is in one of the two sets are O(1) operation.

4. Don't use copy.deepcopy() for list. Instead, use list1 = [5,6]; list2 = list(list1) or list2 = list1[:].

Explanation: copy.deepcopy() handles very rare recursive edge cases and is slow. When simply copying a list or other data structures, you can construct a new list by list() constructor or: operator. Some people avoid the second notation due to readability, but it's frequently used in the real world.

Q. How can I locate the slow code block?

A. The simplest way is to execute **python -m cProfile -s tottime driver.py <puzzle initial state>**.

• -m : specifies module name. cProfile is a built-in profiling Python module.

- https://courses.edx.org/courses/course-v1:Colu...
- -s: sorting by the tottime (total execution time) statistic.
- For meaning of each statistics, see this official documentation.
- The profiling result will be shown even if you stop your driver.py by Ctrl + C command.

If you would like to see the results not on stdout but on browser, <u>cprofilev</u> module might help. There is a well-written <u>tutorial</u> here.

Having said that, sometimes cProfile does not show which part of code makes search slow. Typical case is the first top entry when sorted by tottime is your entire search function (e.g. A-STAR-SEARCH, which is called only one time), and it takes 99% of your execution time.

One approach is to make a wrapper function for an operation which you think might cause the slowness. For example, if you're suspecting **set1** | **set2** (merging two sets) operation is slow, you can make the following wrapper function:

```
def merge_two_sets(set1, set2): # 0(1)? 0(n)? 0(len(set1) +
len(set2))?
    return set1 | set2
```

This simple trick triggers cProfile module to show how long this one-operation function takes. In my test script, 95% of execution time is actually caused by this operation. Recall that cProfile will return profiling results even if you pause your program by Ctrl + C command.

Q. Do I need to optimize my search algorithm as much as possible?

A. You don't need to squeeze your code's performance by fancy optimization techniques such as bit shifting or reducing the number of function calls (i.e. putting every operation in one function for reducing overhead of function calls). Except copy.deepcopy(), most of your design choices are about choosing best data structures in terms of time/space complexity.

Q. Is there any other test cases?

path to goals are truncated and running time/max ram usage are removed: python driver.py dfs 6,1,8,4,0,2,7,3,5 path_to_goal: ['Up', 'Left', 'Down', ... , 'Up', 'Left', 'Up', 'Left'] cost of path: 46142 nodes expanded: 51015 search_depth: 46142 max_search_depth: 46142 python driver.py bfs 6,1,8,4,0,2,7,3,5 path_to_goal: ['Down', 'Right', 'Up', 'Up', 'Left', 'Down', 'Right', 'Down', 'Left', 'Up', 'Left', 'Up', 'Right', 'Right', 'Down', 'Down', 'Left', 'Left', 'Up', 'Up'] cost_of_path: 20 nodes_expanded: 54094 search_depth: 20 max_search_depth: 21 python driver.py ast 6,1,8,4,0,2,7,3,5 path_to_goal: ['Down', 'Right', 'Up', 'Up', 'Left', 'Down', 'Right', 'Down', 'Left', 'Up', 'Left', 'Up', 'Right', 'Right', 'Down', 'Down', 'Left', 'Left', 'Up', 'Up'] cost_of_path: 20 nodes_expanded: 696 search_depth: 20 max search depth: 20 python driver.py dfs 8,6,4,2,1,3,5,7,0 path_to_goal: ['Up', 'Up', 'Left', ..., , 'Up', 'Up', 'Left'] cost_of_path: 9612 nodes_expanded: 9869 search depth: 9612 max_search_depth: 9612

A. The following two test cases might help for your stat validation. Note that long

```
python driver.py bfs 8,6,4,2,1,3,5,7,0
path_to_goal: ['Left', 'Up', 'Up', 'Left', 'Down', 'Right',
'Down', 'Left', 'Up', 'Right', 'Right', 'Up', 'Left', 'Left',
'Down', 'Right', 'Right', 'Up', 'Left', 'Down', 'Down',
'Right', 'Up', 'Left', 'Up', 'Left']
cost_of_path: 26
nodes_expanded: 166786
search depth: 26
max search depth: 27
python driver.py ast 8,6,4,2,1,3,5,7,0
path_to_goal: ['Left', 'Up', 'Up', 'Left', 'Down', 'Right',
'Down', 'Left', 'Up', 'Right', 'Right', 'Up', 'Left', 'Left',
'Down', 'Right', 'Right', 'Up', 'Left', 'Down', 'Down',
'Right', 'Up', 'Left', 'Up', 'Left']
cost_of_path: 26
nodes_expanded: 1585
search_depth: 26
max_search_depth: 26
```

Note that these two test cases are different from ones used for grading. Coming up with tricky test cases also help you understand search algorithm behaviours deeply.

Q. (Windows Users) Is there "resource" module in Windows?

If you use Python in Cygwin, resource module is available. If otherwise, one possible workaround is to use third-party module only if the machine is Windows:

```
import sys
if sys.platform == "win32":
    import psutil
    print("psutil", psutil.Process().memory_info().rss)
else:
    # Note: if you execute Python from cygwin,
    # the sys.platform is "cygwin"
    # the grading system's sys.platform is "linux2"
```

```
import resource
print("resource",
resource.getrusage(resource.RUSAGE_SELF).ru_maxrss)
```

Note that **the values of max_ram_usage and running_time are not graded** as stated in project instruction page. These stats are only for helping your study on search algorithm's time/space metrics.

Q. Why does the example of python driver.py dfs 1,2,5,3,4,0,6,7,8 return ['Up', 'Left', 'Left'] instead of ['Up', 'Left', 'Down', ...] (a solution path with 31 moves)? We are using UDLR (Up, Down, Left, Right) order, and Down move should be executed before Left move. Isn't the ['Up', 'Left', 'Left'] solution resulted from optimization forbidden in project instruction?

No, ['Up', 'Left', 'Left'] solution does not use the forbidden optimization and is a correct answer. Compared to the simpleness of the 3-move solution path, you would notice that "nodes_expanded" statistic is extremely large (181437 states). In fact, the total reachable states in (solvable) 8-puzzle is 9!/2 = 181440 states, so this statistic suggests depth first search constantly overlooks the goal state and expands more than 99.9% of possible states in the search space.

Why is this happening?
Please think about the reason for one minute before reading the following explanations.

There are two facts we can read from dfs pseudocode in class slides:

Fact 1. goalTest() function is only called for states which are just popped out from frontier. In other words, goalTest() is not applied to **neighbor** states **even if the neighbor is actually the solution state** (because we prohibited such an optimization). Fact 2. A child state (neighbor) is only pushed into frontier when it's not already in frontier (and explored). More specifically, **the goal state** is only pushed into frontier when it's not already in frontier.

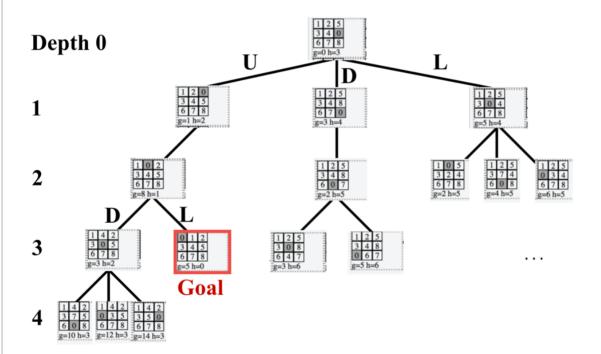
Combining those two facts reaches to the conclusion: since the goal state is already pushed into frontier (i.e. the state corresponding to ['Up', 'Left', 'Left'] move), any

subsequent encounter to the solution state cannot execute frontier.push() or goalTest() and is effectively meaningless. Thus, the dfs algorithm extensively searches through numerous states in 8-puzzle (without putting the solution state into frontier again), gradually goes back to the previously pushed states, and finally find the solution state which is pushed at the very beginning of the search.

This question would arise when you remove/forget neighbor in frontier checking in your pseudocode. And if you add the checking, you will face the slowness of membership checking. In that case, please see the first question of this page ("Q. My search algorithm seems correct but is too slow. How can I reduce its running time?").

Q. Why the max_search_depth of python driver.py bfs 1,2,5,3,4,0,6,7,8 is 4 even though the goal state is at depth 3?

The following figure would be useful (please ignore g and h values):



Q. Why does my driver_3.py return nothing on Vocareum?

If you have default driver.py in your submission, please remove that file. Grading script does not work correctly if both of driver.py and driver_3.py exist.

Q. My python driver.py dfs 1,2,5,3,4,0,6,7,8 returns max_search_depth as 66126 not 66125. Where does this off-by-one difference come from?

One probable reason is you are updating your max_search_depth too early. If you increase max_search_depth after generate_successor() but before checking membership, the generated child state could be actually already in frontier or explored (thus the state cannot be appended to the search tree).

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