

Q-2.

To Prove - Backward Error analysis Results for bordered LU factorization algorithm.

Preparation: Before starting the proof lets define a few results.

From Corollary 6.3.3.2 of the textbook.

For the dot Product Backward and Forward Error.

$$R-LB \quad \tilde{x} = (x + \delta x)^T y \quad \text{where} \quad |\delta x| \leq \gamma_n |x| \quad - (1)$$

$$R-2B \quad \tilde{k} = x^T (y + \delta y) \quad \text{where} \quad |\delta y| \leq \gamma_n |y| \quad - (2)$$

$$R-LF \quad \tilde{k} = x^T y + \delta k \quad \text{where} \quad |\delta k| \leq \gamma_n |x^T y| \quad - (3)$$

From Theorem 6.6.2.14 of the textbook.

The Error analysis for the result of triangular system.

$$(L + \Delta L) \tilde{x} = y \quad \text{where} \quad |\Delta L| \leq \max(\gamma_2, \gamma_{n-1}) |L| \quad - (4)$$

From Corollary 6.6.2.15 of the textbook, the error analysis of the triangular system can also be written as.

$$R-LB \quad (L + \Delta L) \tilde{x} = y \quad \text{where} \quad |\Delta L| \leq \max(\gamma_2, \gamma_{n-1}) |L| \quad - (5)$$

$$R-LF \quad L \hat{x} = b + \delta b \quad \text{where} \quad |\delta b| \leq \max(\gamma_2, \gamma_{n-1}) |L| |x| \quad - (6)$$

Proof - we will impley the proof by induction.

* Base Case $n=1$

the base case of $n=1$, where $A = (a)$ and $L = (1)$,
 $U = (a)$

clearly we can see

$$A + \Delta A = LU \quad \text{holds for any } \Delta A = (\delta a)$$

$$\text{where } |\Delta A| \leq \gamma_1 |\delta a|.$$

* For the induction Step, we assume. that the statement holds. for matrix $A \in \mathbb{C}^{m \times n}$, with the nonsingular principle matrix. That is there exists ΔA such that

$$A + \Delta A = \tilde{L} \tilde{U} \quad \text{with } |\Delta A| \leq \gamma_n |\tilde{L}| |\tilde{U}|.$$

where \tilde{L} is unit lower triangular and \tilde{U} is upper triangular.

* Now we consider the case where $A \in \mathbb{C}^{m \times n+1}$ with a nonsingular leading principal submatrices. we partition A as.

for step n .

$$A \rightarrow \left(\begin{array}{c|c} A_{TL} & A_{TR} \\ \hline A_{BL} & A_{BR} \end{array} \right), \quad L \rightarrow \left(\begin{array}{c|c} L_{TL} & 0 \\ \hline L_{BL} & L_{BR} \end{array} \right), \quad U \rightarrow \left(\begin{array}{c|c} U_{TL} & U_{TR} \\ \hline 0 & U_{BR} \end{array} \right).$$

$$\text{with } \Delta A = \left(\begin{array}{c|c} \Delta A_{TL} & \Delta A_{TR} \\ \hline \Delta A_{BL} & \Delta A_{BR} \end{array} \right).$$

for step $n+1$, we partition as.

$$A \rightarrow \left(\begin{array}{c|c} A_{00} & a_{01} \\ \hline a_{10}^T & a_{11} \end{array} \right) \quad L \rightarrow \left(\begin{array}{c|c} \tilde{L}_{00} & 0 \\ \hline \tilde{l}_{10}^T & 1 \end{array} \right) \quad U \rightarrow \left(\begin{array}{c|c} \tilde{U}_{00} & \tilde{u}_{01} \\ \hline 0 & \tilde{u}_{11} \end{array} \right)$$

Where

$$\left(\begin{array}{c|c} A_{00} & q_{01} \\ \hline q_{10}^T & \alpha_{11} \end{array} \right) + \left(\begin{array}{c|c} \Delta A_{00} & \delta q_{01} \\ \hline \delta q_{10}^T & \delta \alpha_{11} \end{array} \right) = \left(\begin{array}{c|c} \tilde{L}_{00} & 0 \\ \hline \tilde{l}_{10}^T & 1 \end{array} \right) \left(\begin{array}{c|c} \tilde{U}_{00} & \tilde{u}_{01} \\ \hline 0 & \tilde{\alpha}_{11} \end{array} \right)$$

Now for the $(n+1)^{th}$ step of Bordered algorithm we know,

$$\tilde{L}_{00} \tilde{u}_{01} = q_{01} \quad \leftarrow \text{triangular solve lower triangular mat}$$

$$\tilde{l}_{10}^T \tilde{U}_{00} = q_{10}^T \quad \leftarrow \text{triangular solve upper triangular mat}$$

$$\alpha_{11} := \tilde{\alpha}_{11} = \alpha_{11} - q_{10}^T q_{01} \quad \leftarrow \text{dot Product.}$$

with eqⁿ ⑤ defined above.

$$\tilde{L}_{00} \tilde{u}_{01} = q_{01} + \delta q_{01} \quad \text{where } |\delta q_{01}| \leq \max(\gamma_2, \gamma_{n+1}) \|\tilde{L}_{00}\| \|\tilde{u}_{01}\| \quad \leftarrow \text{⑦}$$

similarly with the eqⁿ ⑥ we can extrapolate the result for upper triangular solve.

$$\tilde{l}_{10}^T \tilde{U}_{00} = q_{10}^T$$

\leftarrow taking transpose on both sides.

$$\tilde{U}_{00}^T \tilde{l}_{10} = q_{10}.$$

this becomes a lower triangular solve. so.

$$\tilde{U}_{00}^T \tilde{l}_{10} = q_{10} + \delta q_{10} \quad \text{where } |\delta q_{10}| \leq \gamma_n \|\tilde{U}_{00}^T\| \|\tilde{l}_{10}\|. \quad \leftarrow \text{⑧}$$

$$\alpha_{11} := \tilde{\alpha}_{11} = \alpha_{11} - (q_{10} + \delta q_{10})^T (q_{01} + \delta q_{01})$$

$$= \alpha_{11} - q_{10}^T q_{01} - \delta \alpha_{11} \quad \leftarrow \text{⑨}$$

$$\text{where } \delta \alpha_{11} = q_{10}^T \delta q_{01} + \delta q_{10}^T q_{01} + \delta q_{10}^T \delta q_{01}$$

Putting eqⁿ ⑦, ⑧ and ⑨ together we get

$$\left(\begin{array}{c|c} A_{00} & q_{01} \\ \hline q_{10}^T & \alpha_{11} \end{array} \right) = \left(\begin{array}{c|c} \hat{L}_{00} \hat{U}_{00} = A_{00} + \Delta A_{00} & \check{L}_{00} \check{U}_{01} = q_{01} + \delta q_{01} \\ \hline (\check{U}_{00}^T \hat{L}_{10})^T = (q_{10} + \delta q_{10})^T & \check{U}_{11} = \alpha_{11} - q_{10}^T q_{01} - \delta \alpha_{11} \end{array} \right)$$

$$\Rightarrow \left(\begin{array}{c|c} \hat{L}_{00} \hat{U}_{00} & \check{L}_{00} \check{U}_{01} \\ \hline \hat{L}_{10}^T \hat{U}_{00} & \alpha_{11} - q_{10}^T q_{01} \end{array} \right) = \left(\begin{array}{c|c} A_{00} & q_{01} \\ \hline q_{10}^T & \alpha_{11} \end{array} \right) + \left(\begin{array}{c|c} \Delta A_{00} & \delta q_{01} \\ \hline \delta q_{10}^T & -\delta \alpha_{11} \end{array} \right)$$

So with this result we can see that after step $(n+1)$ the resultant matrix A can be written as.

$$A + \Delta A = \check{L} \check{U}$$

$$\begin{aligned} \text{for } m(A) &= n+1 \wedge \check{L} \check{U} \\ &= (A + \Delta A) \wedge |\Delta A| \\ &< \gamma_{n+1} |L| |U| \end{aligned}$$

Therefore we can see the backward error incurred with bordered LU factorization is

$$A + \Delta A = \check{L} \check{U} \quad \text{with } |\Delta A| \leq \gamma_n |\check{L}| |\check{U}|.$$

By The Principle of mathematical Induction this result holds.