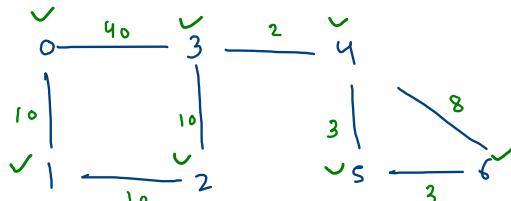


Shortest path in wt: single source to all dest

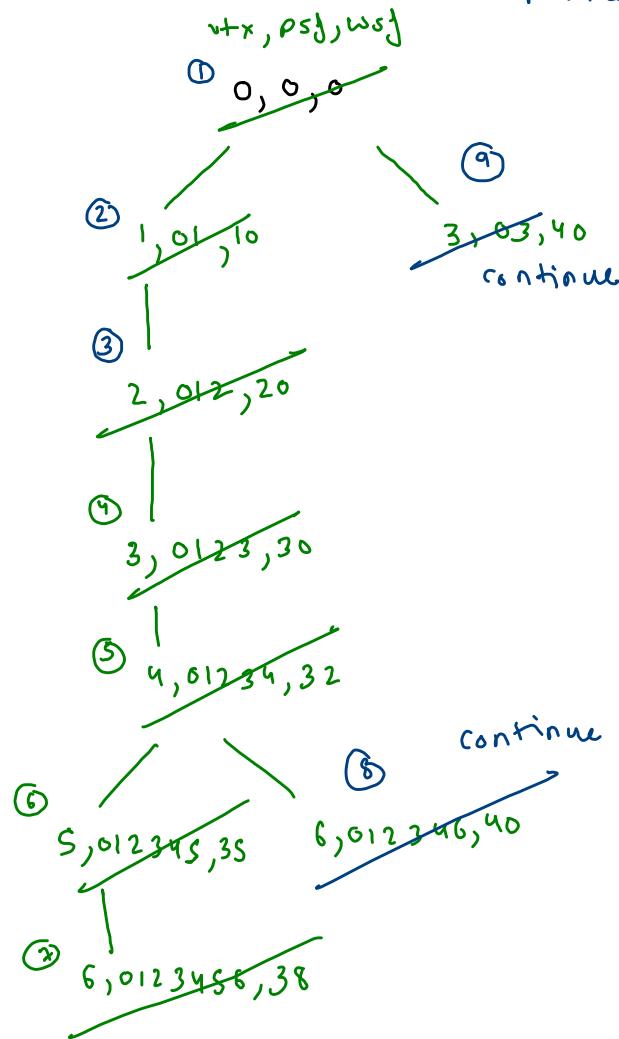


remove

mark*

work

add unvisited nbr



shortest in single traversal.

0 → 0 @ 0

1 → 01 @ 10 $src = 0$

2 → 012 @ 20

3 → 0123 @ 30

4 → 01234 @ 32

5 → 012345 @ 35

6 → 0123456 @ 38

BFS

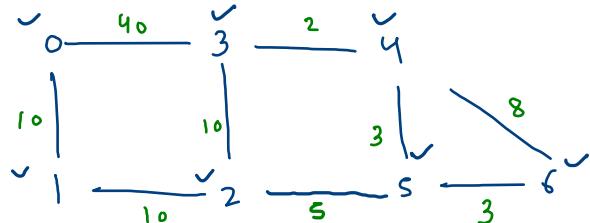
queue

path is shortest (edges)

Dijkstra

PQ

path is shortest (wt)

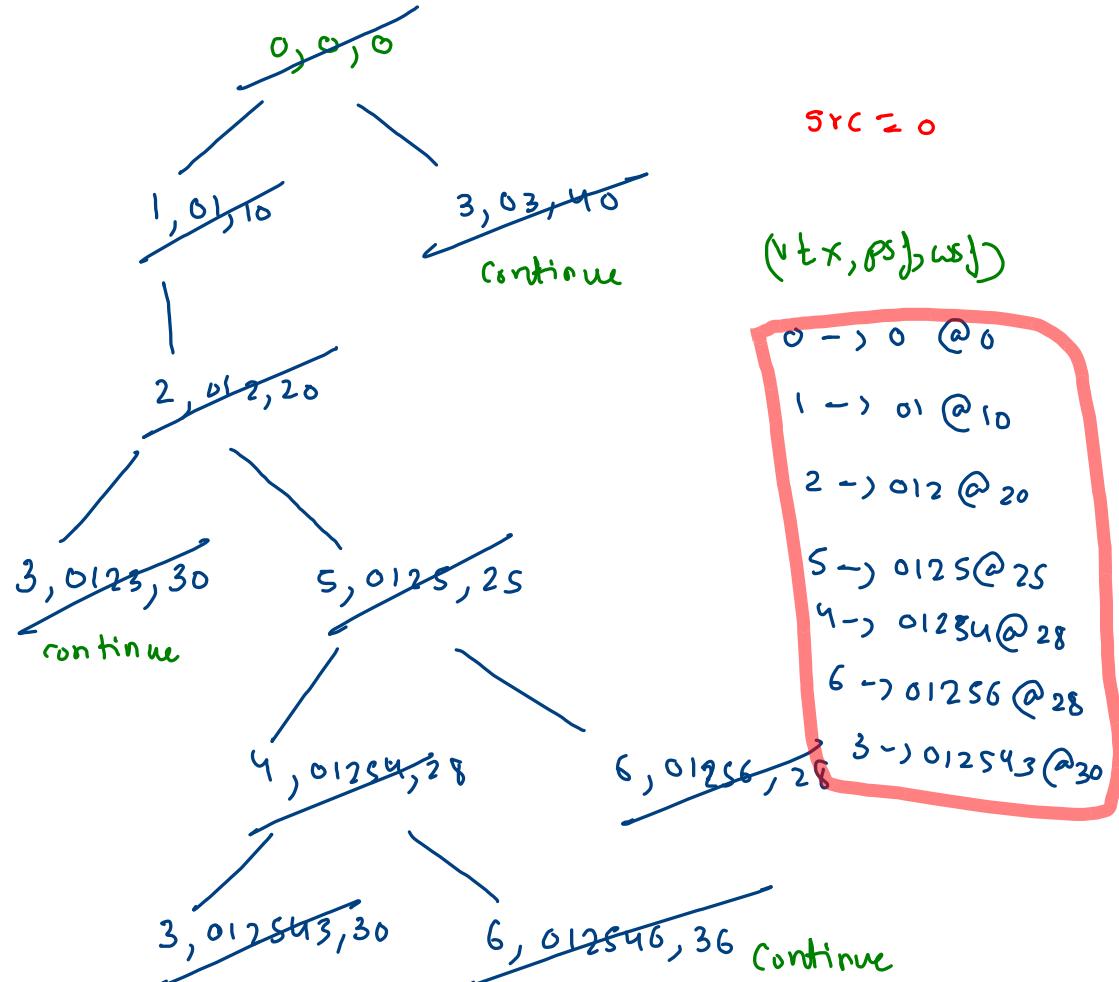


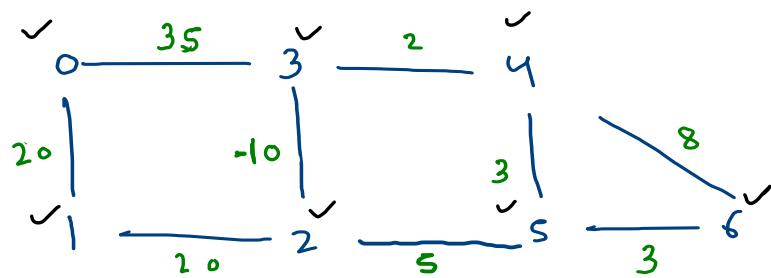
dijkstra → single src to all dest.

```

public static void dijkstra(ArrayList<Edge>[] graph, int src) {
    PriorityQueue<Pair> pq = new PriorityQueue<Pair>();
    boolean[] vis = new boolean[graph.length];
    pq.add(new Pair(src, src + "", 0));
    while(pq.size() > 0) {
        //remove
        Pair rem = pq.remove();
        //mark
        if(vis[rem.vtx] == true) {
            continue;
        }
        vis[rem.vtx] = true;
        //work
        System.out.println(rem.vtx + " via " + rem.psf + " @ " + rem.wsf);
        //add unvisited nbr
        for(Edge ne : graph[rem.vtx]) {
            if(vis[ne.nbr] == false) {
                pq.add(new Pair(ne.nbr, rem.psf + ne.nbr, rem.wsf + ne.wt));
            }
        }
    }
}

```





$1 \rightarrow 01 @ 20$

$3 \rightarrow 03 @ 35$

dijkstra : $2 \rightarrow 032 @ 25$

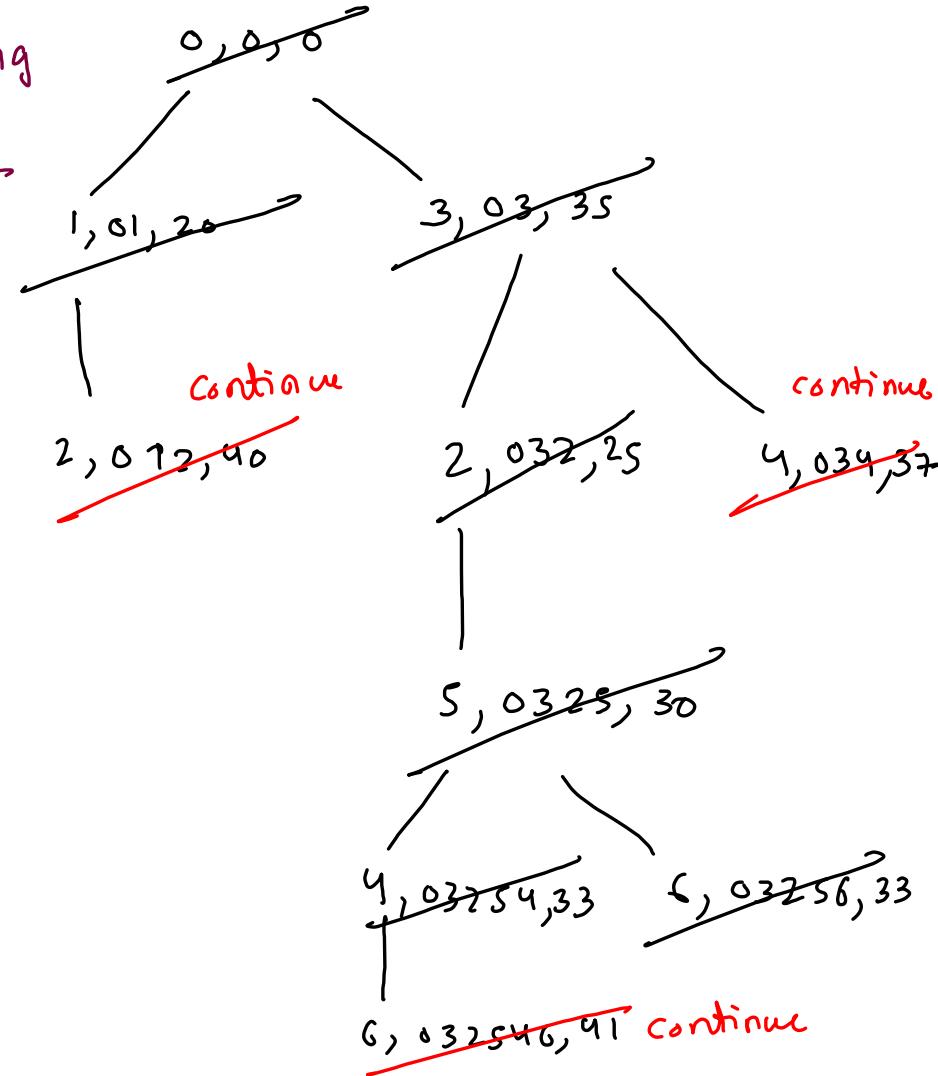
jails or -ve $5 \rightarrow 0325 @ 30$

edge wt $4 \rightarrow 03254 @ 33$

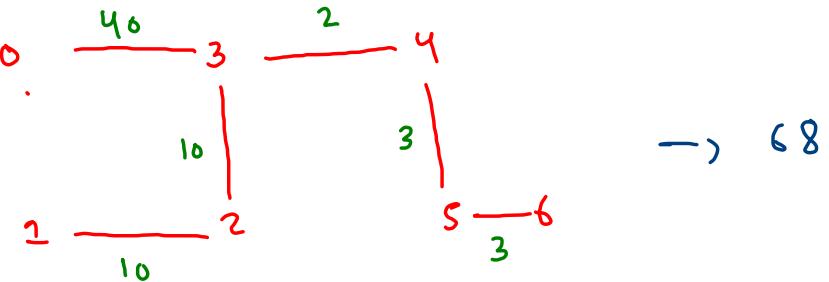
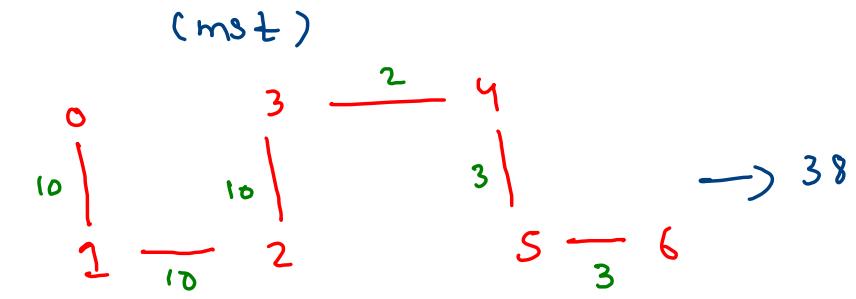
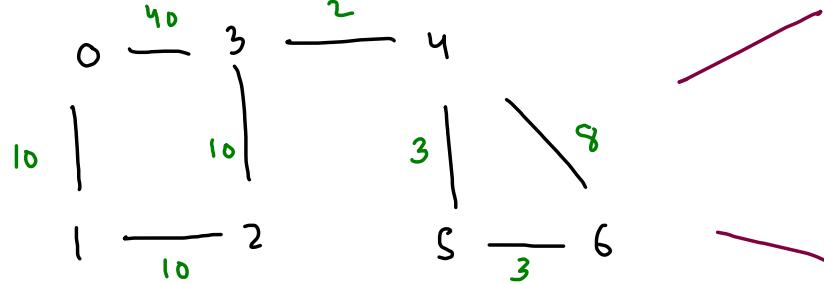
$6 \rightarrow 03256 @ 33$

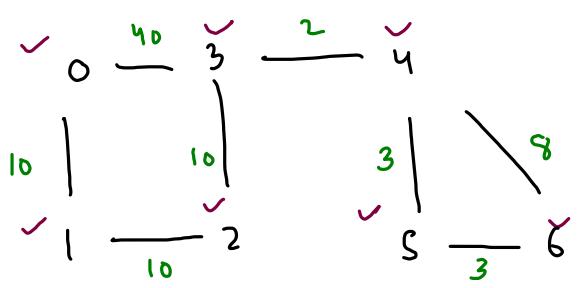
Failing

case



Prim's \rightarrow MST (min spanning tree)



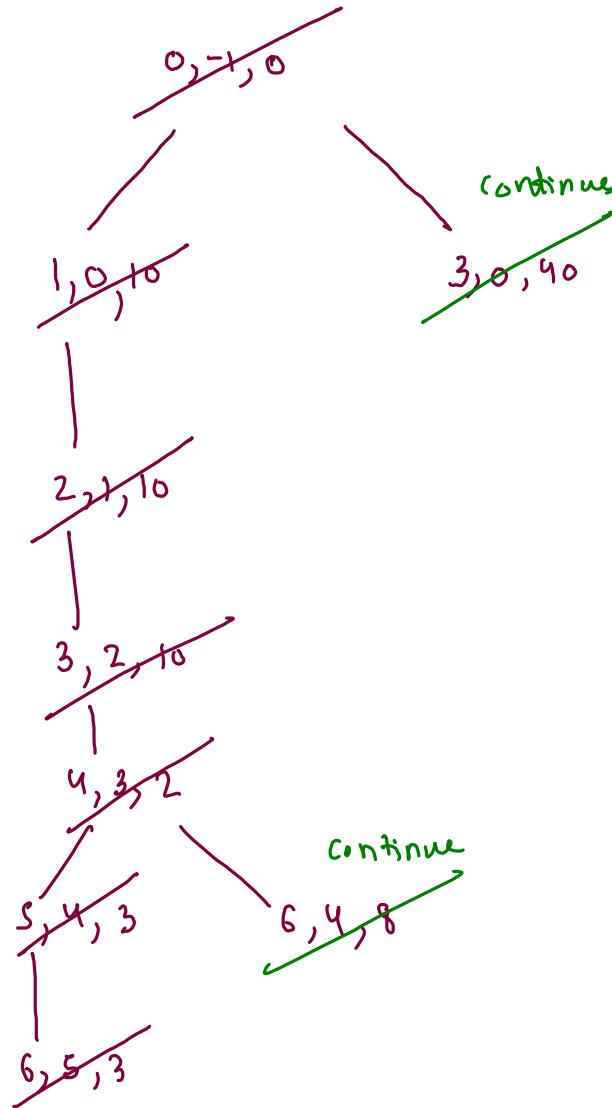
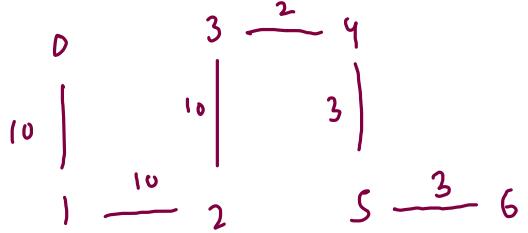


remove

mark*

work \rightarrow add edge

add nbr



pair:

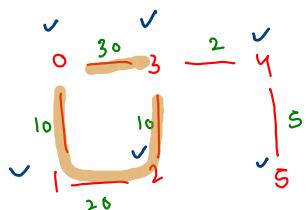
$v \times a_v$

wt.

Dijkstra

(i) shortest path (in terms of w_b)

(ii) $\text{src} \rightarrow$ single src to all dest
shortest path.



$0 \rightarrow 0 @ 0$

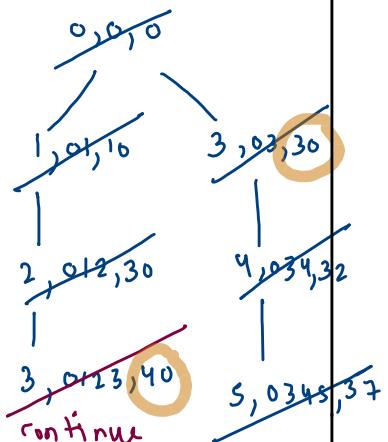
$1 \rightarrow 01 @ 10$

$2 \rightarrow 012 @ 30$

$3 \rightarrow 03 @ 30$

$4 \rightarrow 034 @ 32$

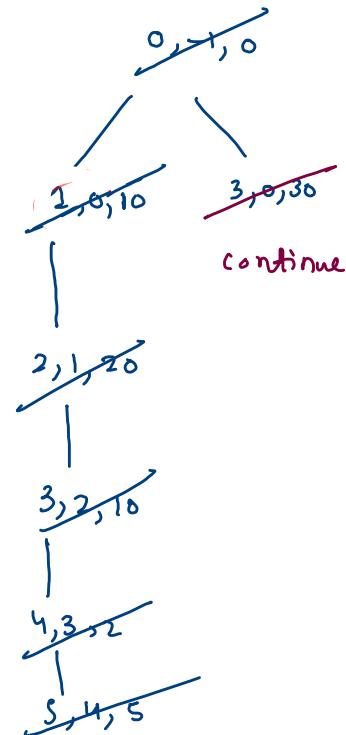
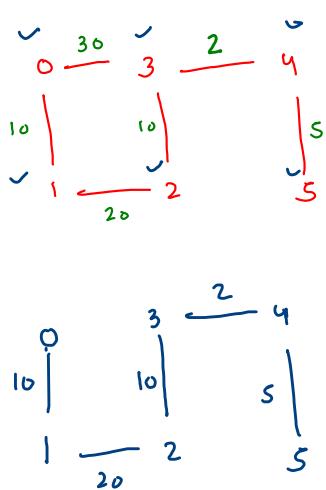
$5 \rightarrow 0345 @ 37$

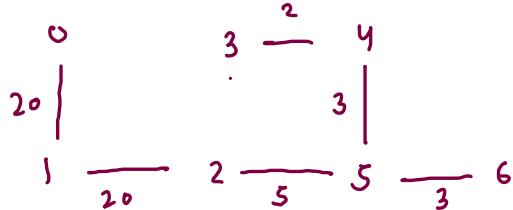
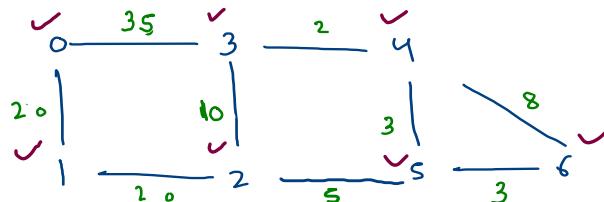


Prim

(i) MST (connect all vertices in min cost).

(ii) src : does not matter.





```

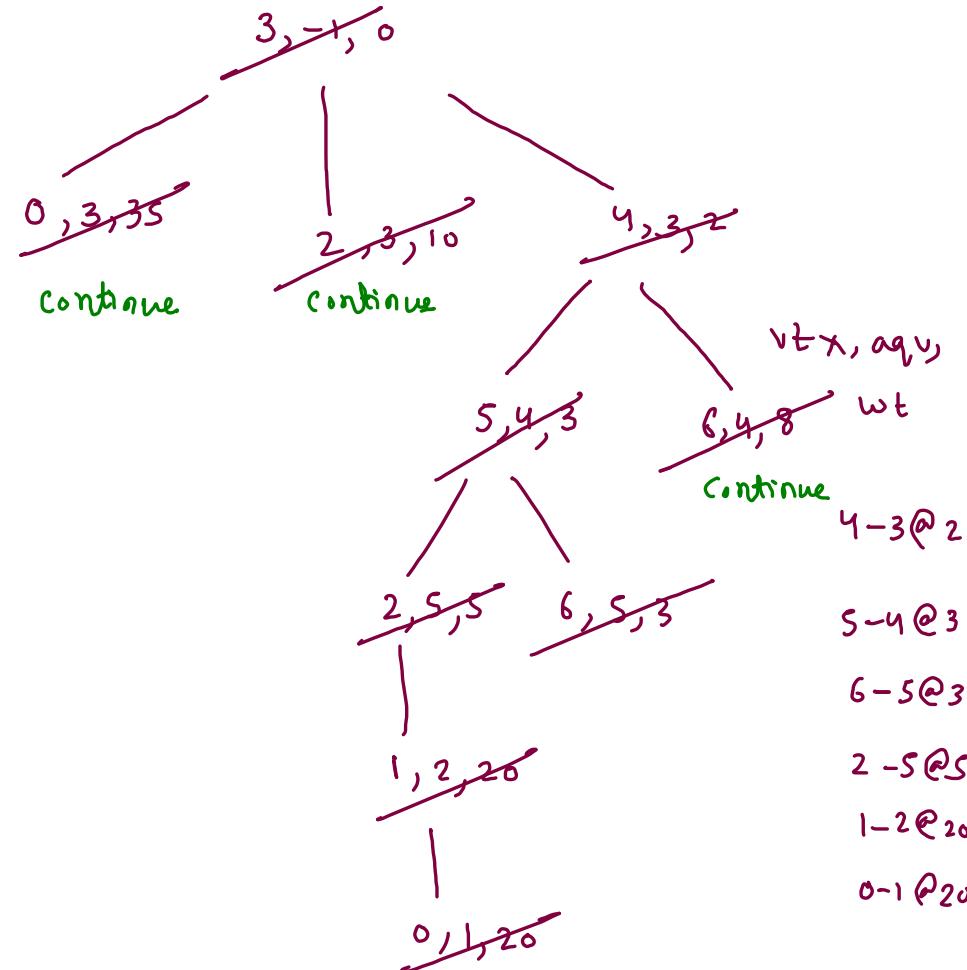
while(pq.size() > 0) {
    //remove
    Pair rem = pq.remove();

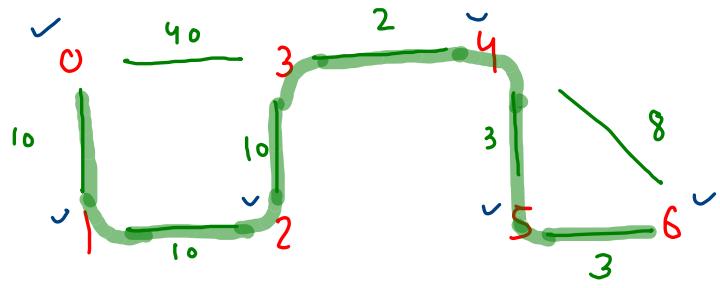
    //mark*
    if(vis[rem.vtx] == true) {
        continue;
    }
    vis[rem.vtx] = true;

    //work
    if(rem.aqv != -1) {
        System.out.println("[ " + rem.vtx + " - " + rem.aqv + "@" + rem.wt + "]");
    }

    //add nbr
    for(Edge ne : graph[rem.vtx]) {
        if(vis[ne.nbr] == false) {
            pq.add(new Pair(ne.nbr, rem.vtx, ne.wt));
        }
    }
}

```





```

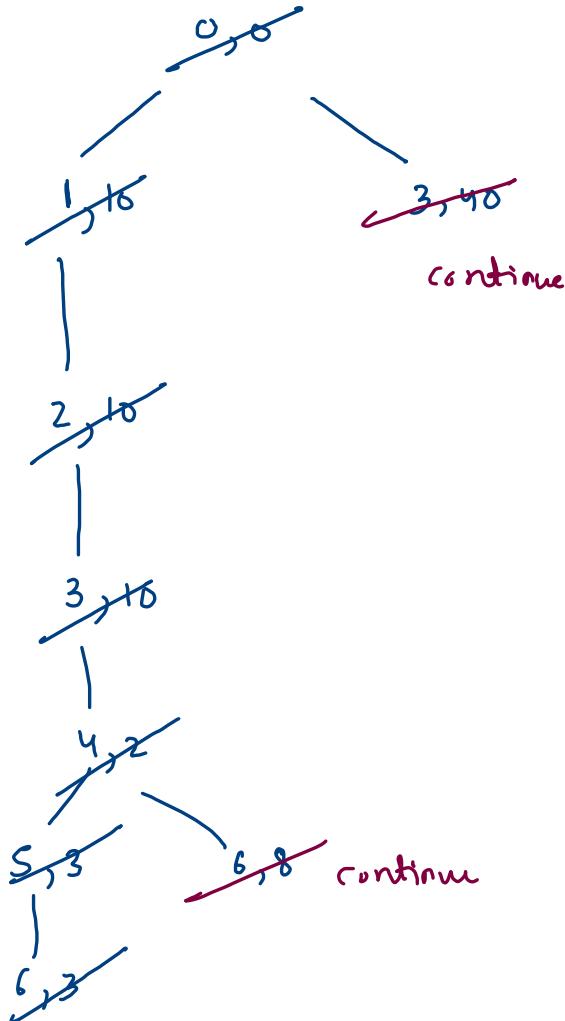
while(pq.size() > 0) {
    //remove
    Edge rem = pq.remove();

    //mark*
    if(vis[rem.v] == true) {
        continue;
    }
    vis[rem.v] = true;

    //work
    cost += rem.wt;

    //add nbr
    for(Edge ne : graph.get(rem.v)) {
        int nbr = ne.v;
        if(vis[nbr] == false) {
            pq.add(new Edge(nbr,ne.wt));
        }
    }
}

```

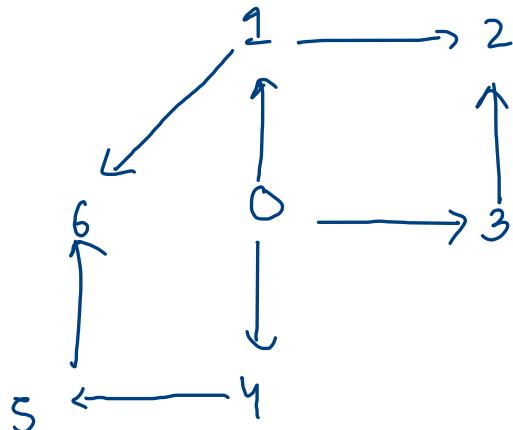


$$\begin{aligned}
 \text{Cost} &= 10 + 10 \\
 &\quad + 10 + 2 \\
 &\quad + 3 + 3 \\
 &= 38
 \end{aligned}$$

Topological sort -> A permutation of vertices for a directed acyclic graph is called topological sort if
for all directed edges uv , u appears before v in the graph

DA67

Directed acyclic graph



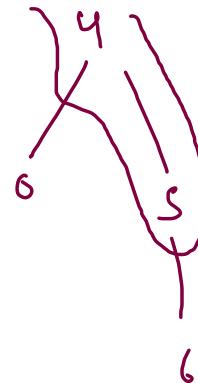
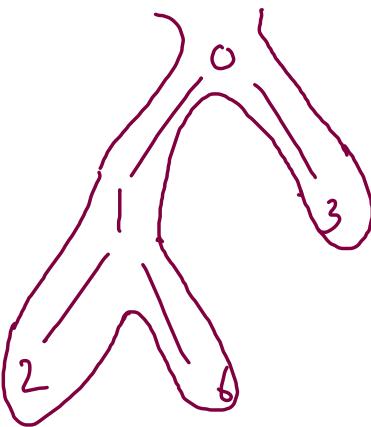
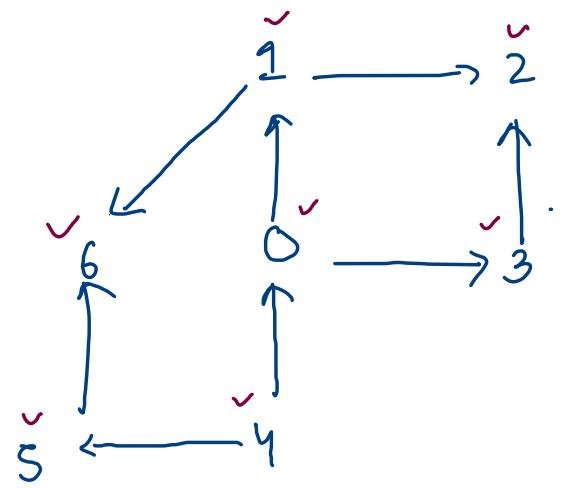
$u \rightarrow v$

u should come before v

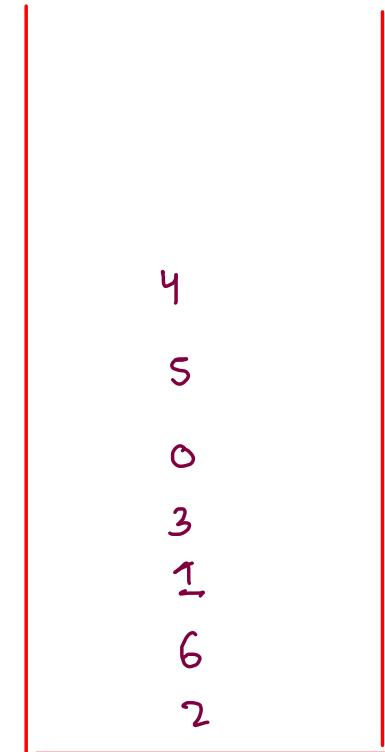
u is dependent on v

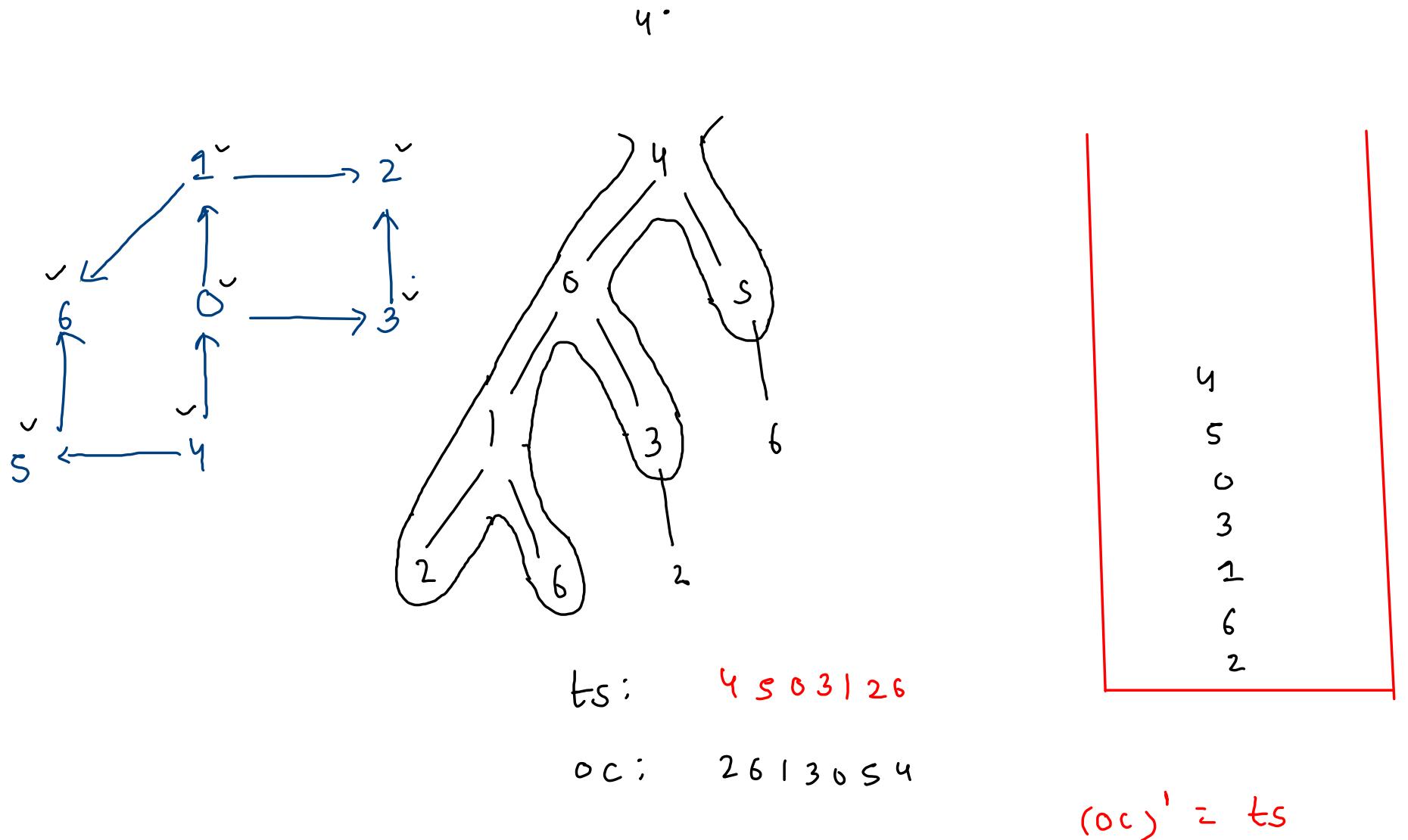
0 1 3 2 4 5 6 (topological sort)

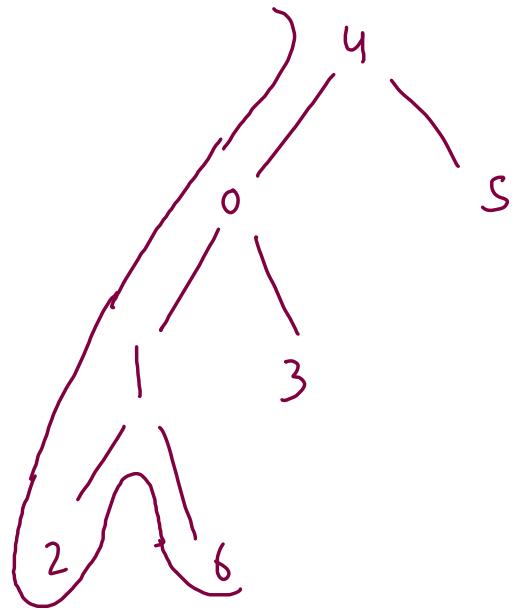
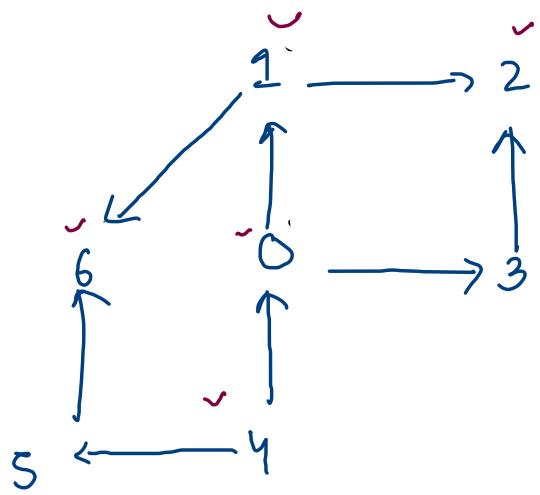
order of compilation -> (topological sort)



ts: 4 5 0 3 1 6 2







$\begin{array}{r} 4 & 0 & 1 & 2 & 6 \\ \hline \end{array}$
 Pre X

```

public static void dfs(ArrayList<Edge>[] graph, int src, boolean[] vis, Stack<Integer> st) {
    vis[src] = true;

    for(Edge ne : graph[src]) {
        if(vis[ne.nbr] == false) {
            dfs(graph, ne.nbr, vis, st);
        }
    }

    st.push(src);
}

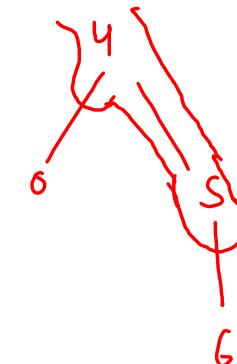
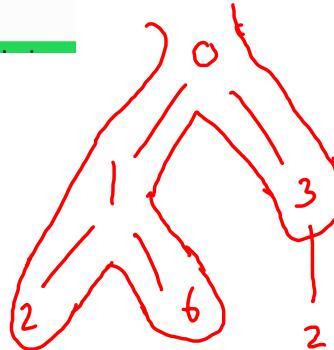
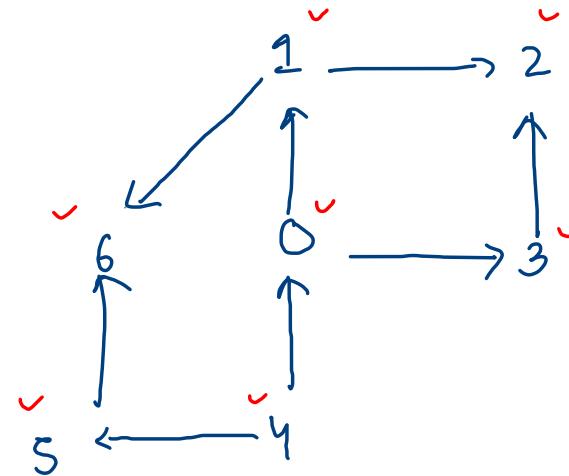
public static void topological_sort(ArrayList<Edge>[] graph) {
    boolean[] vis = new boolean[graph.length];
    Stack<Integer> st = new Stack<>();

    for(int i=0; i < graph.length; i++) {
        if(vis[i] == false) {
            dfs(graph, i, vis, st);
        }
    }

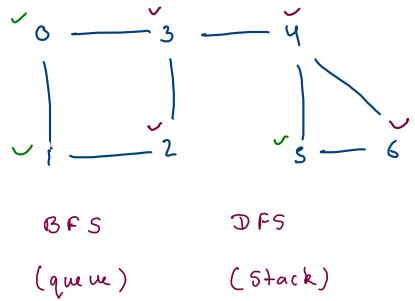
    //stack has ts
    while(st.size() > 0) {
        System.out.println(st.pop());
    }
}

```

4 5 0 3 1 6 2



4
5
0
3
1
6
2



```

public static void iterative_dfs(ArrayList<Edge>[] graph,int src) {
    boolean[] vis = new boolean[graph.length];
    Stack<Pair> st = new Stack<>();
    st.push(new Pair(src,src+""));

    while(st.size() > 0) {
        //remove
        Pair top = st.pop();

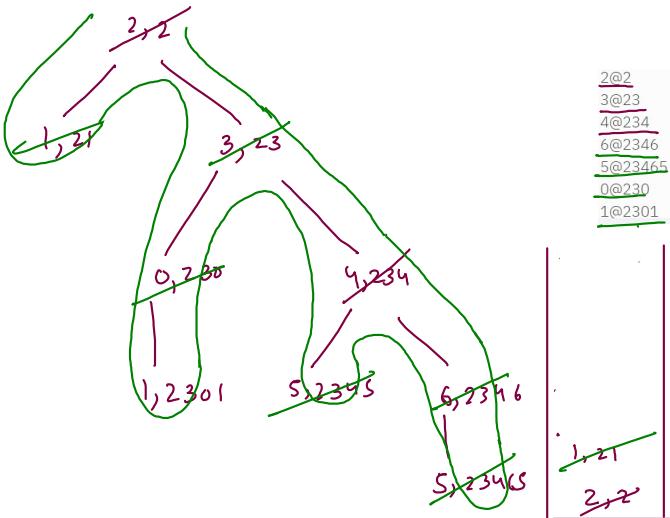
        //mark*
        if(vis[top.vtx] == true) {
            continue;
        }

        vis[top.vtx] = true;
        //work
        System.out.println(top.vtx + "@" + top.psf);

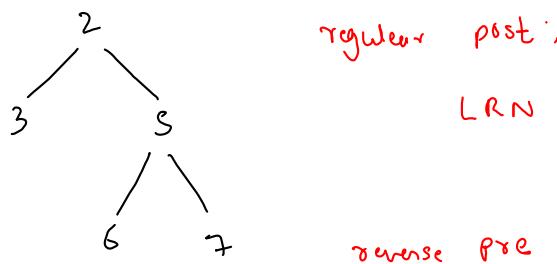
        //add unvisited nbr
        for(Edge ne : graph[top.vtx]) {
            int nbr = ne.nbr;
            if(vis[nbr] == false) {
                st.push(new Pair(nbr,top.psf + nbr));
            }
        }
    }
}

```

iterative DFS
src = 2

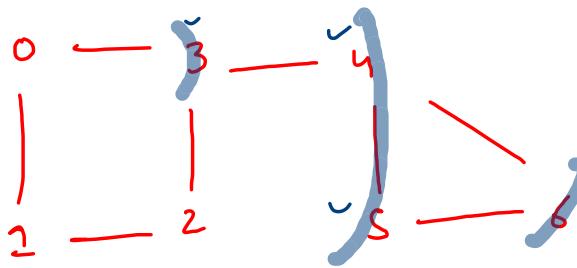


2@2
3@23
4@234
6@2346
5@23465
0@230
1@2301



NRL

Spread of Infection



```

while(q.size() > 0) {
    int count = q.size();

    if(lev > time) {
        break;
    }

    while(count-- > 0) {
        //remove
        int rem = q.remove();

        //mark*
        if(vis[rem] == true) {
            continue;
        }
        vis[rem] = true;

        //work
        ans += 1;

        //add nbr
        for(Edge ne : graph[rem]) {
            if(vis[ne.nbr] == false) {
                q.add(ne.nbr);
            }
        }
    }

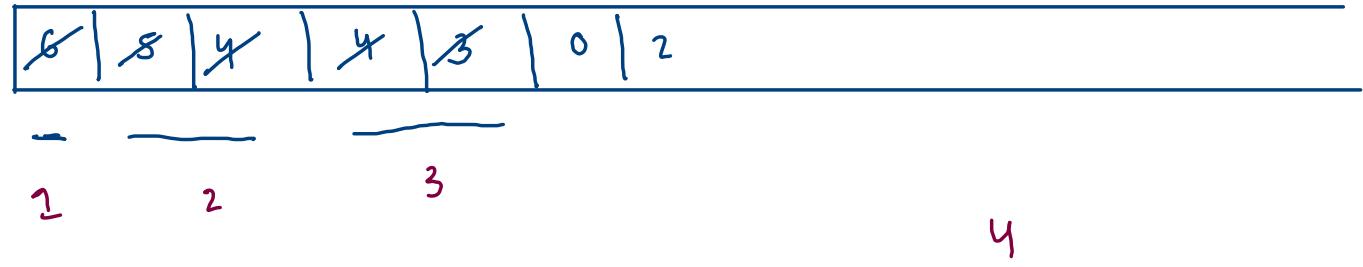
    lev++;
}

```

$$\text{ans} = \underline{1 + 1 + 1 + 1}$$

$$\text{count} = 2$$

$$\text{lev} = \cancel{2} \cancel{3} \cancel{4}$$



y

$$\text{inf} = 6$$

$$t = 3$$