

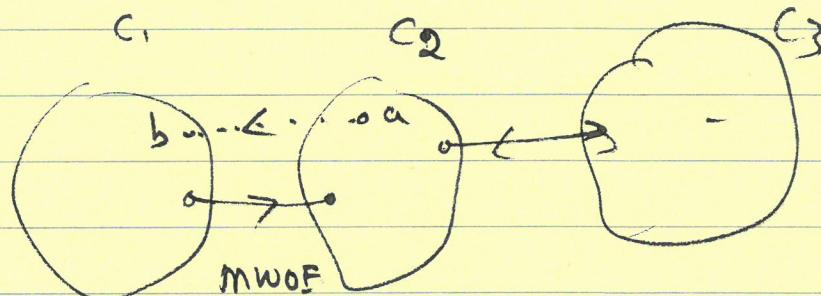
①

## Asynchronous MST or Asynch HTS

Complications because of asynchrony:

A. Some Components are "too fast"  
Compared to others.

All processes of a Combined Component  
need to know about new Component id.



C1 absorbed into C2

$C_2 \oplus$  combined with  $C_3$   
( $C_2 \& C_3$  merged)

Later a sends a test message  
or (a,b) to see if (a,b) is an  
outgoing edge

b may "erroneously" send an  
accept reply.

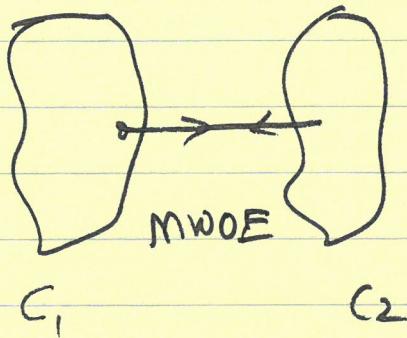
B. "Big" Component finds MWOE  
to a "small" Component  
This results in "big" message complexity

②

How to take care of these Complications?

We have levels for Components. (Initially level 0, ~~0~~ for each process).

(a)

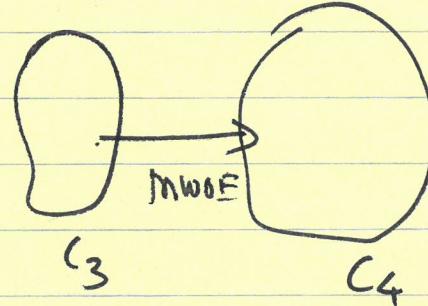


merge

C<sub>1</sub> at Level L & C<sub>2</sub> at Level L  
Combine (merge) using Common MWOF

C<sub>1</sub> ∪ C<sub>2</sub> ∪ {Common MWOF} is at  
level L+1

(b)



absorb

MWOF (C<sub>3</sub>) ~~takes~~ connects to C<sub>4</sub>  
and Level (C<sub>3</sub>) < Level (C<sub>4</sub>)

All processes in C<sub>3</sub> take on C<sub>4</sub>'s Component id & C<sub>4</sub>'s level.

③

Initially, each process is a component by itself and at level 0.

### Finding MWOE of a Component

Leader broadcasts an initiate ( $id$ , Level) message

$x$  needs to find ~~MWOE~~ MWOE candidate incident on  $x$ .

$x \xrightarrow{0} \text{test } y$   
 $\$(\text{Component id, Level})$   
wait for a reply

$y$  sends a reply if Level of  $y$

$\geq$  Level of  $x$ . and send an accept or ~~reject~~ reject

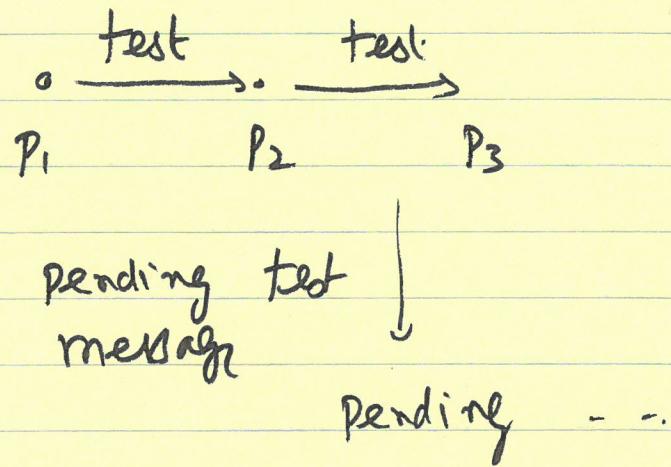
else { // Level( $x$ ) > Level( $y$ ) :

defer reply until  $y$ 's level is greater than or equal to  $x$

}

④

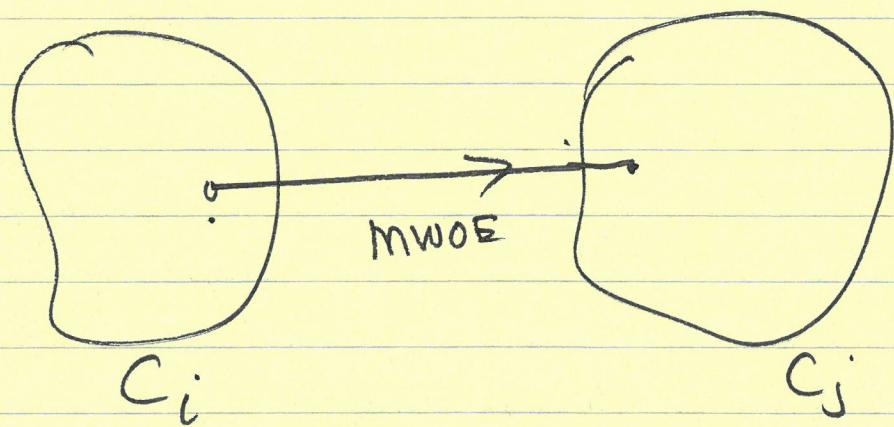
Deadlock possible because of  
these pending replies?



thus lead to successively smaller  
levels; no cycle; no deadlocks.

All processes of a component find  
MWOE

Combine with Component at other  
end of MWOE



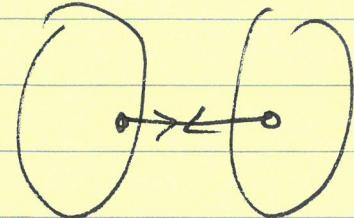
(3)

Two Cases

(a) ~~no edge~~ :  $MWOE(C_i) = MWOE(C_j)$

$$\text{Level}(C_i) = \text{Level}(C_j)$$

Common MWOE is the  
core edge.

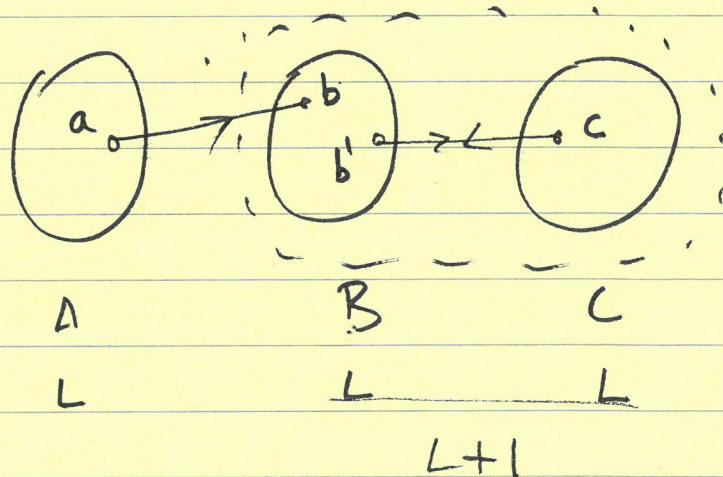


$$C_i \cup C_j \cup MWOE(C_i \cup C_j) c_i \quad c_j$$

is at level  $\text{Level}(C_i)$

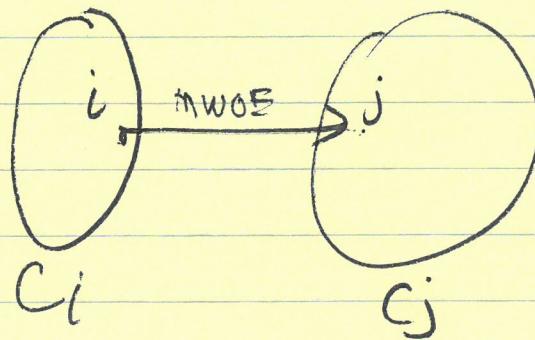
one end point of Common MWOE is the  
end of combined component

edge weight of Common MWOE is the  
new (combined) component's id.



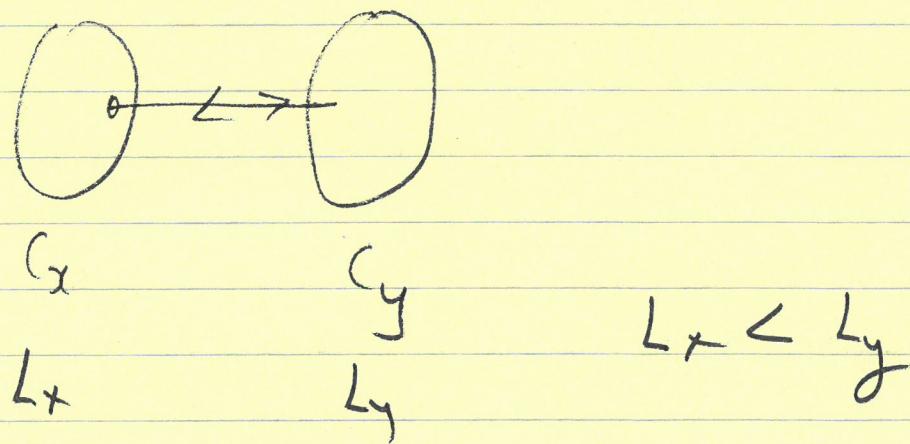
(b) ~~edge~~ Absorb

⑥ Absorb operation.



$$mwoE(C_j) \neq mwoE(C_i)$$

$\Rightarrow \text{Level}(C_i) < \text{Level}(C_j)$  at some  
some time in the future, and, at that  
time  $C_i$  will be absorbed into  $C_j$



When  $C_i$  at Level  $L_i$  is absorbed  
into  $C_j$  at Level  $L_j$  ( $L_j > L_i$ )  
processes in  $C_i$  take over  $C_j$ 's component id  
and Level  $L_j$

9

$C_i$ 's processes participate in finding

$MWSE$  of  $C_i \cup C_j \cup MWEC(C_i)$

o

- o Think about the 2 cases
- o  $(i) j$  already sent the Convergecast reply to  $C_j$ 's leader and

(ii)  $j$  has NOT yet sent the Convergecast reply to  ~~$C_j$~~   $j$ 's Parent

Max Level:  $\lceil \log_2 n \rceil$

min # of processes in any component  
of level  $L$  is  $2^L$  test/reject

Message Complexity:  $(\underbrace{IE}_1 + \underbrace{n \log n}_2)$

Time Complexity:  $O(n \log n)$

test/accept

Combine, new  
id.  
Convergecast  
loop