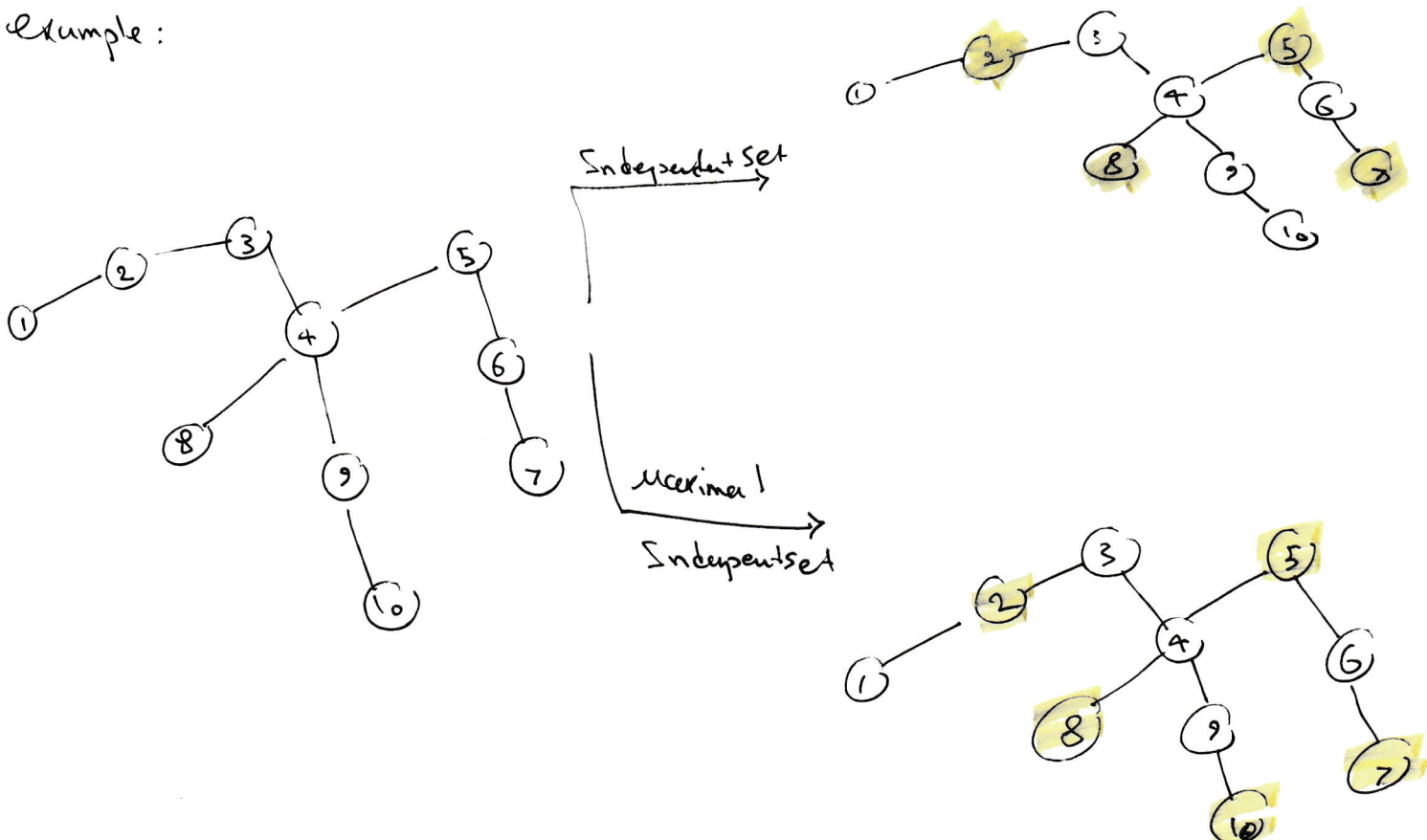


# MSS8

- + set of nodes is independent if no two nodes are adjacent.  
(one-hop away)
  - + Independent Set (IS) is maximal if it cannot be increased to form a larger set.
  - + An independent set of maximum cardinality is called maximum.
  - + We are interested in Maximal not Maximum.
  - + Motivation: Broadcast difficult problem.
- 

Example:



↳ Sequential Algorithm

$$G = (V, E)$$

$$I \leftarrow \emptyset ; \text{Nodes} \leftarrow V ; \text{Edges} \leftarrow E ;$$

while  $\text{Nodes} \neq \emptyset$

{ choose a non-empty set  $I' \subseteq \text{Nodes}$   
 $I'$  is an Independent Set;  
 $I \leftarrow I \cup I'$ ;  
 $\text{Nodes} \leftarrow \text{Nodes} - \{I'\} - \{\text{neighbors of } I'\}$   
 $\text{Edges} \leftarrow \text{Edges} - \{\text{edges incident of } I'\} -$   
 $\{\text{edges incident of neighbors of } I'\}$

↳ How to compute in distributed way?

↳ How to choose  $I'$ ?

↳ Randomization is used.

each process  $i$  chooses  $val_i$  randomly in the range  $\{1, \dots, n^4\}$

←  
sufficiently large to  
reduce possibility where 2 process  
choose same value.

$j \in I'$  if  $val_j$  is max among all its neighbors.

we define  $I'$  to consist of all nodes  $j$  that are local winners.

↪ It means, those nodes  $j$  s.t.  $val_j > val_i$  for all neighbors

$I'$  might be empty set at some stage  $\rightarrow$  (wasted stage)

The algorithm works in Stages, each consisting of 3 rounds.

Round 1: choose Vales locally and send to neighbors.

Round 2: winners notify neighbors.

By the end of this round, loser are known.

$I'$  is known.

loser: the nodes that have neighbors in  $I'$ .

Round 3: each loser notifies its neighbors.

winner & loser remove themselves.

↳ it means that, loser and winners

discontinue participate after this stage.

loser's neighbors remove cell edges.

↳ How many stages?

Depend on  $|I|$ .

In the randomized algorithm, Result is correct ; # of Stage may vary.

Lemma 4.7: Expected # of edges removed from  $G$  in single stage is at least  $\frac{|E|}{8}$ .

↳ the algorithm ensures that every edge with at least one endpoint in  $I$  is removed.

Lemma 4.8: with prob at least  $\frac{1}{16}$ , # number of edges removed from  $G$  in single stage is at least  $\frac{|E|}{16}$

Th 4.9 # of Stages is  $O(\log n)$  with high probability.