

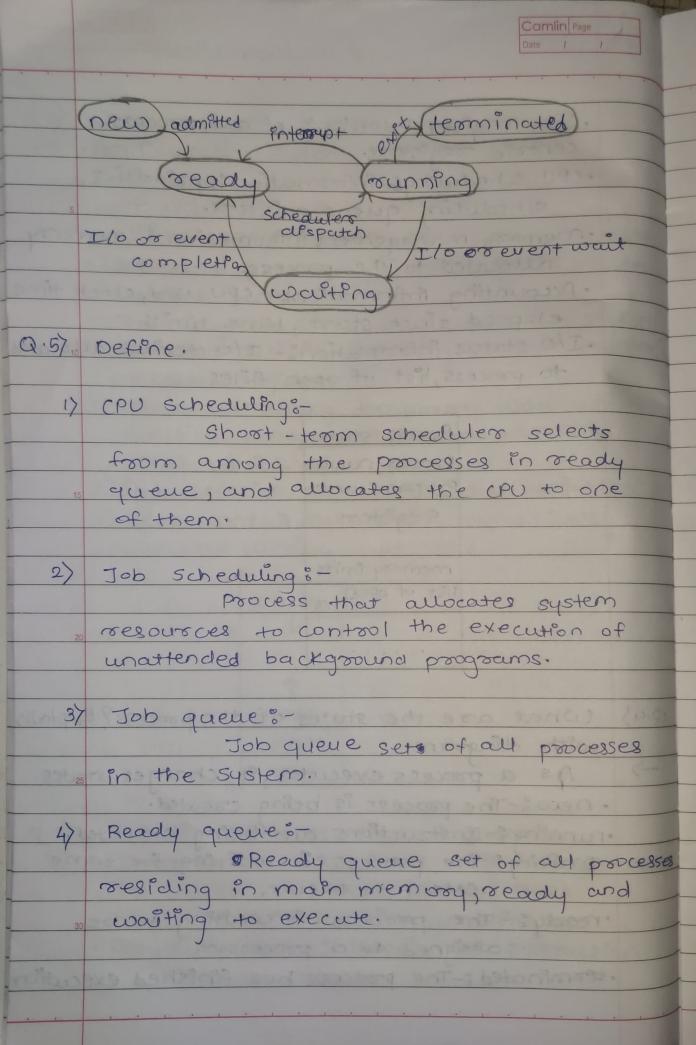
- .CPU registers:- contents of all process.
- . CPU scheduling s'information à provosities, scheduling queue pointers.
- memory management information: memory allocated to the process.
- · Accounting information: CPU used, clock time elapsed since start, time limits.
- . Ilo status information: Ilo devices allocated to process, list of open files.

Process state
Process numbers
Program Counters
Registers
memory limits
Ist of open files

with diagram.

As a process executes, 9+ changes stutes.

- · new :- The process is being created.
- ·running :- Instructions are being executed.
- event to occur.
- · ready: The process is waiting to be assigned to a processor.
- · terminated: The process has finished execution



5) Device queue:

Device queue set of processes waiting for an I/o device.

6) Scheduler:-

Scheduler:
A special type of system saftware

that handles process scheduling in numerous

The selects which process should be executed next and allocates CPU.

8) Long-term scheduler:-It selects which process should be

brought into the ready queue.

9) Degree of multiprogramming:

It describes the maximum number

of processes that a single processor

system can accompdate efficiently.

10) Context Switch :-

when CPU switches to another process, the system must save the state of the old processes and load the sowed steete for the new processes via context switch.

(2.6) Explain medium term scheduling with diagram.

· medium-term scheduler can be added if degree of multiple programming needs to decrease.

· BIt removes process from memory, Store on disk, bring back in from disk to continue execution: swapping.

swapped-out processes

ready queue cpu jend

I/O I/O waiting
queues

Q.7) What is interprocess communication?

-> Processes within a system may be independent or cooperating.

- Cooperating process can affect or be affected by other processes, including sharing data.
- · Reasons for cooperating process.
 - 6) Compt Computation Speedup.
- · Cooperating system processes need interproocess communication.
- or threads communicate among each others.

(a.8) What are the advantages of cooperative process? > Independent process cannot affect or be affected by the execution of another process. · Cooperating process can affect or be affected by the execution of another process. · Advantages of cooperative process are as follows: a) Information sharing. b) Computation speedup. c) modularity. d) convenience. (9) Explain models used in IPCE. models used in the are as follows: 1) shared memory: . An arrea of memory shared among the processes that wish to communicate. · The communication is under control of user processes not os. · Sya majors issues is to provide mechanism that will allow user processes to synchronize their actions when they access shared memory. process A shared memory process B Keronel

2) message passing:-

end to synchronize their actions.

· message system- processes communicate with each other without resorting to shared vorsicubles.

. If processes Pand Q wish to communicate, they need to:

· Establish a communication link between them.

· Enchange messages via send/recève

· Implementation of communication link

· Physical:

a) shared memory.

b> Hardware bus.

c) Network,

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· Logical;

a) Direct or indirect.

b) synchronous or asynchronous.

c) Automatic or explicit buffering.

process A

process B

message queue

mo mi m2...mn

Kernel

Keaus

- Q-11) Explain communication methods used in Client-server system.
 - -) 1) Sockets:-
 - · A socket is defined as an enapoint for communication.
 - · Concatenation of IP address and post a number included at start of message packet.
 · Communication consists between a pair of sockets.
 - · Three types of sockers:
 - Connection-ordented (TCP).
 - T Connection less (UDP).
 - Multisocket
 - 2) Remote Procedure Calls:
 - · Remote procedure calls abstracts procedure calls between processes on networked system
 - · The client-side stub locates server and marshalls the parameters.
 - Remote communication has more failure scenarios than local messages can be delievered exactly once rather than at most once.
 - 3) Pipes:
 - · Acts as conduit allowing two processes to communicate.
 - · Orodfnary Pipes: cannot be accessed from outside process that created it. Typically, a parent process creates a pipe and uses it to communicate with child processe.

- · Named Pipes: con be accessed without pourent child relationship.
- Q.12 Define.
 - This process execution consists a of a cycle of CPU execution and I/o wait.
 - Dispatchers
 Dispatchers module gives control of

 the CPU to the process selected by short
 term schedulers.
 - 3) Disperten Lateray:
 Time it takes for the dispertenters

 to stop one process and start another

 running.
 - 4) CPU Utilization:
 CPU utilization is the process which

 keeps the CPU to as busy as possible.
 - throughput:
 # of processes that complete

 their execution per time unit.
 - amount of time to execute a paroticular process.

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	1		

2) Wouting time:amount of time a process has been wouting in the ready queue.

8) Response time:-

amount of time it toures from when a request was submitted until the first response is produced, not output.

Q.13) Calculate augerage waiting time & average turnaround time.

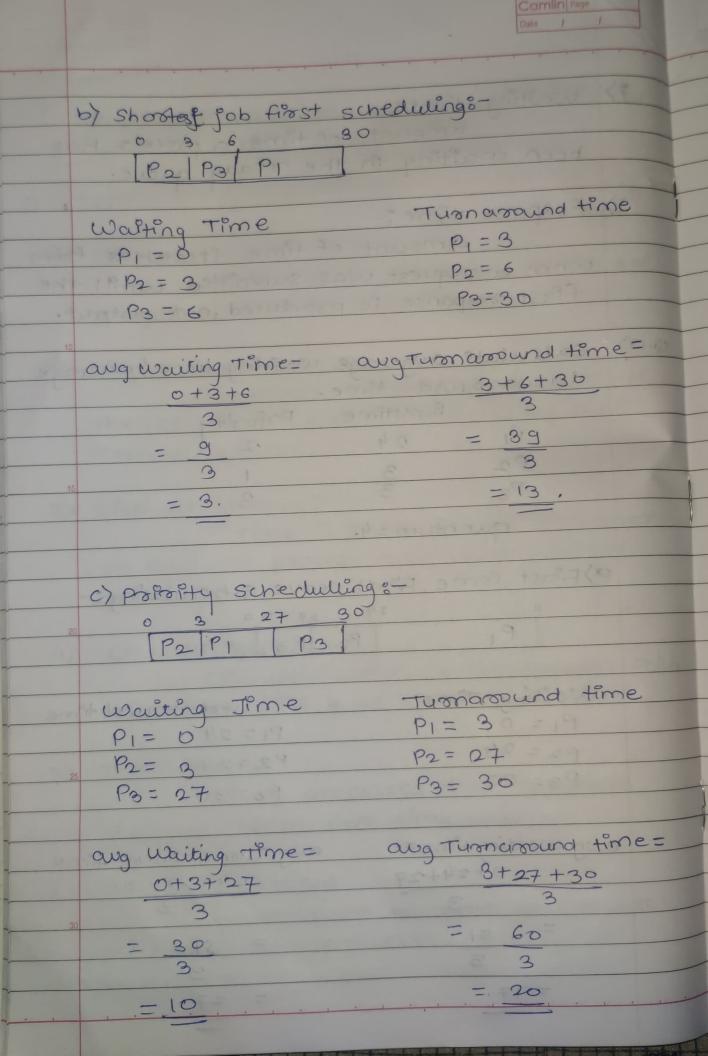
Busstime Parosity P₁ 24 2 P₂ 3 1 P₃ 3

quantum = 4.

a) First come, First serve scheduling :-P₁ P₂ P₃

Turnaround time Waiting Time P1 = 24 P1=0 p2 = 24 P2=27 P3=300 P3= 27

and Thonacound time and waiting time = 0+24+27 24+27+30



d> Round-Roben algorithms-9 4 7 10 14 18 22 26 30 P1 P2 P3 P1 P1 P1 P1 P1

Waiting time Turn around time P1 = 30 - 24 = 6 P2 = 7 - 3 = 4 P3 = 10 - 3 = 7 P3 = 10

aug waiting $t^{\circ}me = \frac{30 + 7 + 10}{30 + 7 + 10}$ 3 = 17 = 47 3 = 5.6 = 15.6

short job first scheduling