

Election Algorithm

- ⇒ Many distributed algorithms need one process to act as a Coordinator
- Election alg : technique to pick a unique coordinator (leader election)
- Examples :- Take over the role of a failed process,
Pick a master in Berkeley clock synchronization Alg.
- Types of election algorithms :-
 - Bully algorithm
 - Ring algorithm

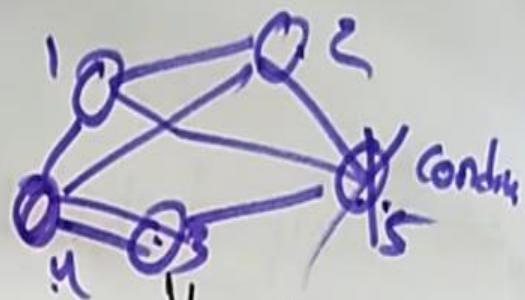
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Bully Algorithm

- Each process has a unique numerical ID
- Processes know the IDs and address of every other process
- Communication is assumed reliable.
- Key idea: Select process with highest ID
- Process initiates election if it just recovered from failure or if coordinator failed
- Several processes can initiate an election simultaneously
- $O(n^2)$ msg required with n processes

Bully Algorithm

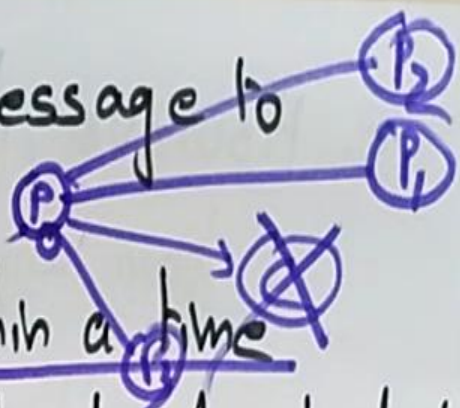


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Algorithm 1: Suppose Process P Sends a message to the coordinator.

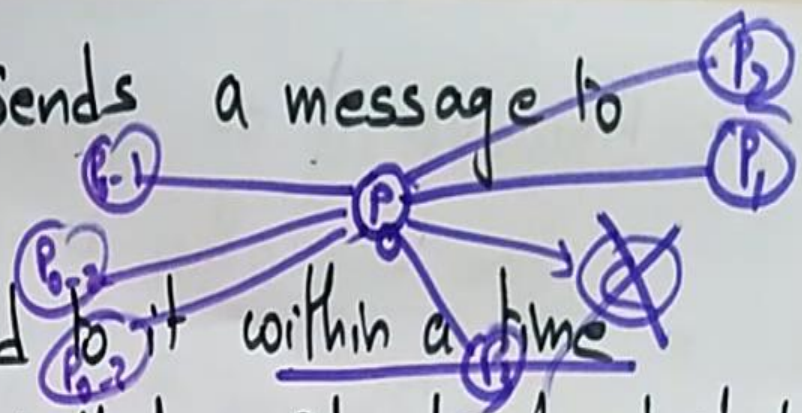
1. If coordinator does not respond to it within a time interval T , then it is assumed that coordinator has failed.
2. Now Process P Sends election message to every process with high priority number.
3. It waits for response, if no one responds for time interval T then process P elects itself as a coordinator.
4. Then it Sends a message to all lower priority numbers then it is elected as their new coordinator.

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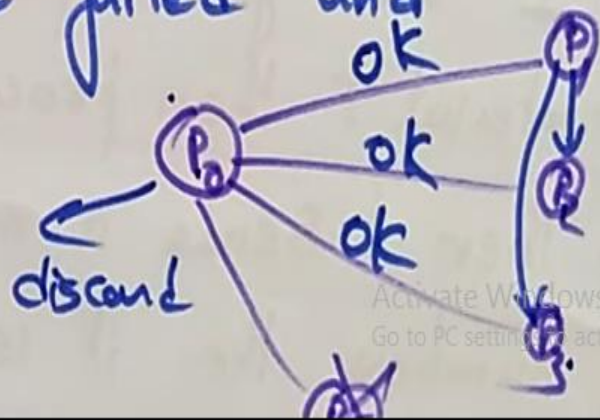


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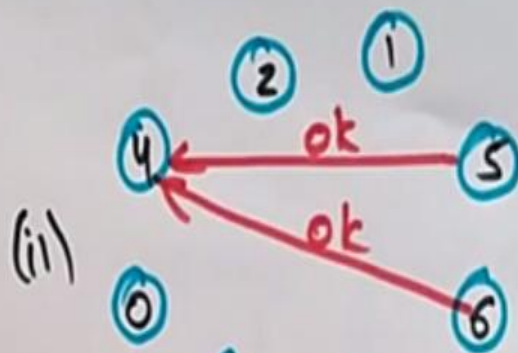
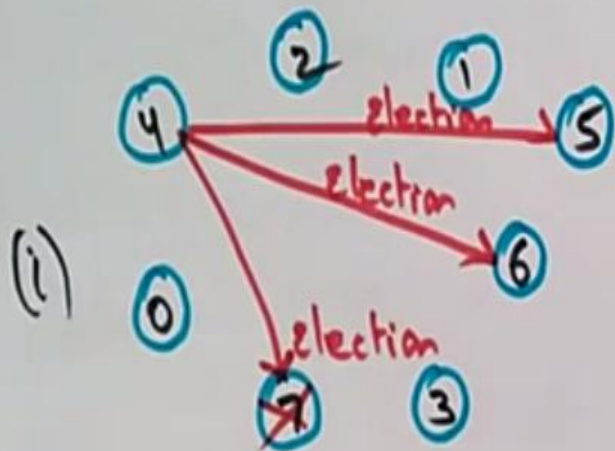
5. However, if an answer is received within time T from any other process Q ,

→ Process P again waits for time interval T to receive another message from Q that it has been elected as coordinator.

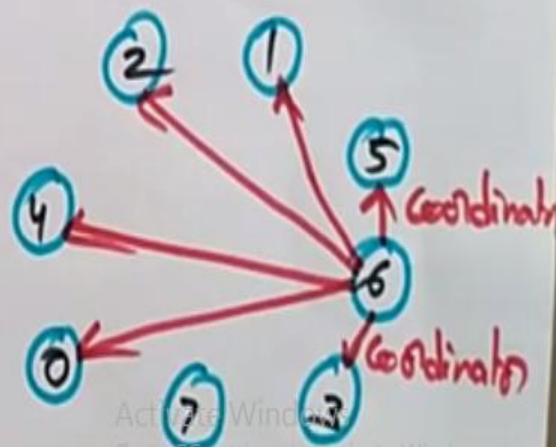
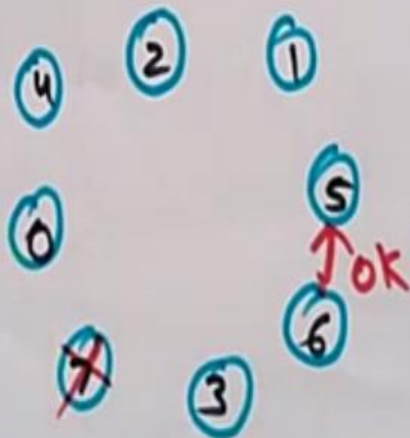
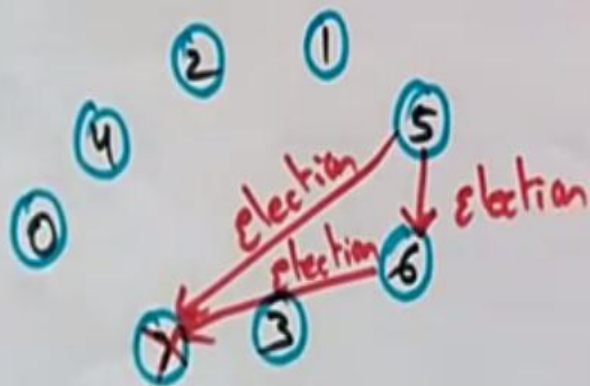
→ If Q doesn't respond within time interval T , then it is assumed to have failed and algorithm is restarted.



Example: 'Bully algorithm'



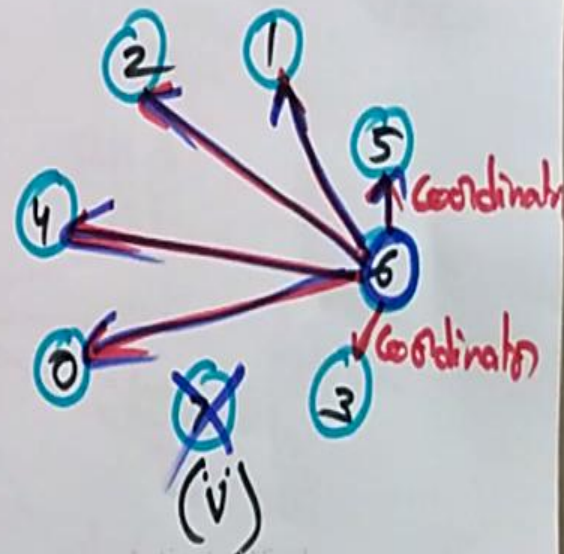
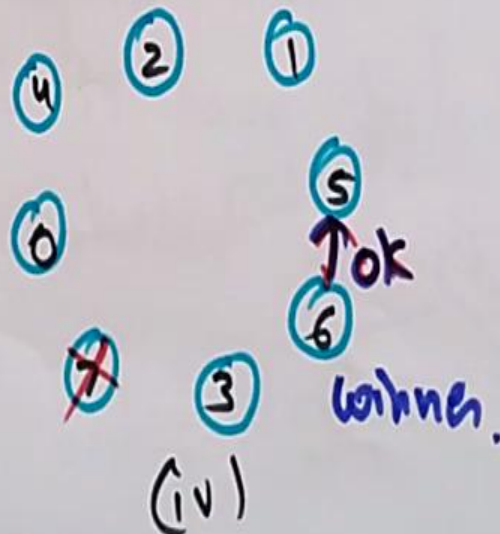
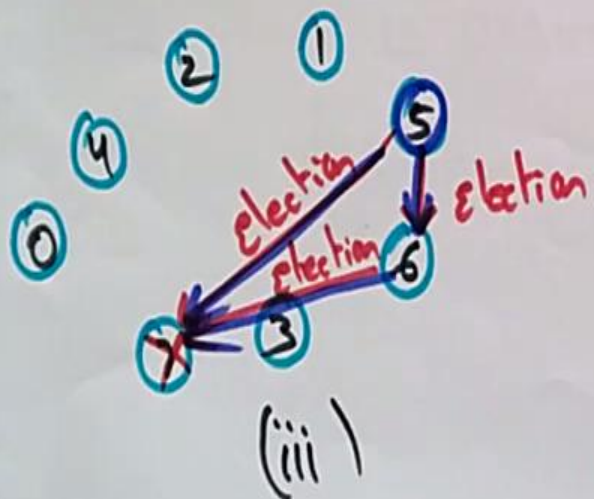
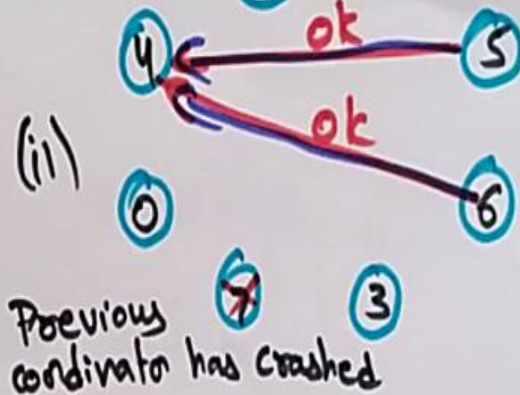
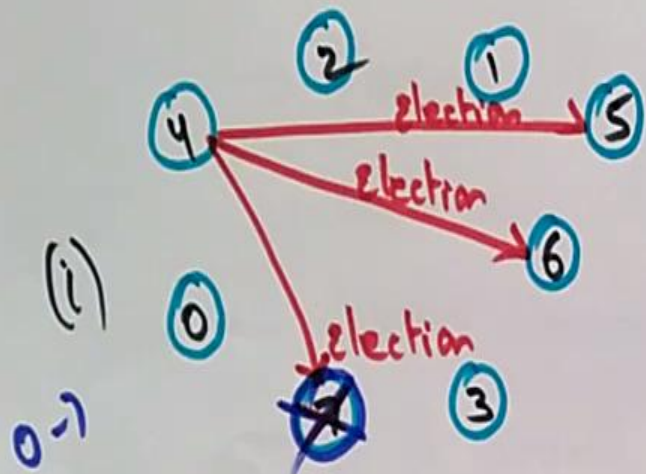
Previous coordinator has crashed



Activate Windows
Go to Settings to activate Windows.

Example: Bully algorithm

job is over

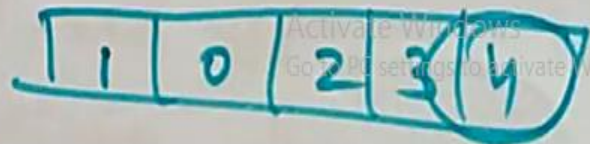


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Ring Algorithm

- This Alg applies to system organized as a ring (logically or physically).
- In this alg we assume that the link between the process are unidirectional and every process can manage to the process on its right only.
- Data structure that this algorithm uses is active list, a list that has priority number of all active processes in the system.

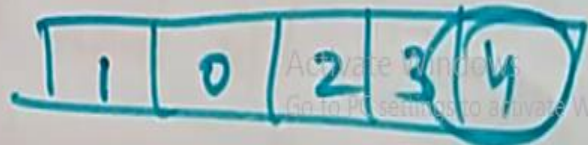
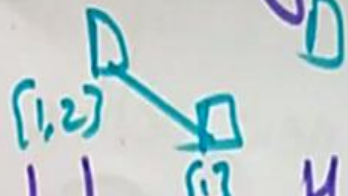


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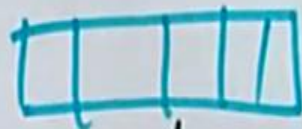


Algorithm :-

① If process P_1 detects a coordinator failure, it creates a new active list which is empty initially. It sends election message to its neighbour on right and adds number 1 to its active list.

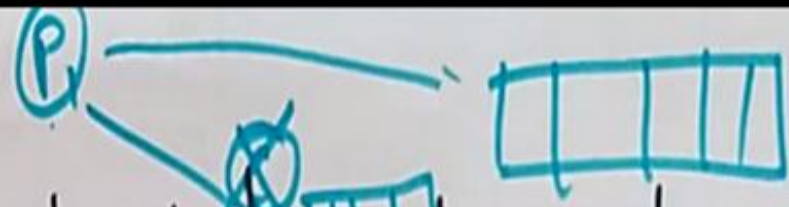
② If process P_2 receives message elect from processes on left, it responds in 3 ways.

(i) If msg received does not contain 1 in active list then P_1 adds 2 to its active list & forwards the message.

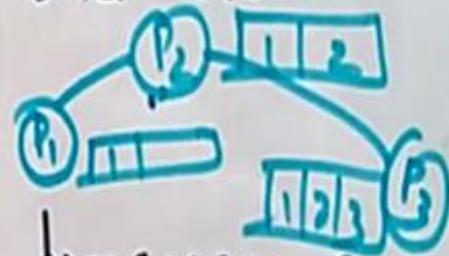


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(iv) If this is the first election message it has received or sent P_i creates new active list with numbers 1 & 2. It then sends election message 1 followed by 2.

(ii) If process P_i receives its own election message 1 then active list for P_i now contains numbers of all the active processes in the system. Now process P_i detects highest priority number from list & elects it as the new coordinator.

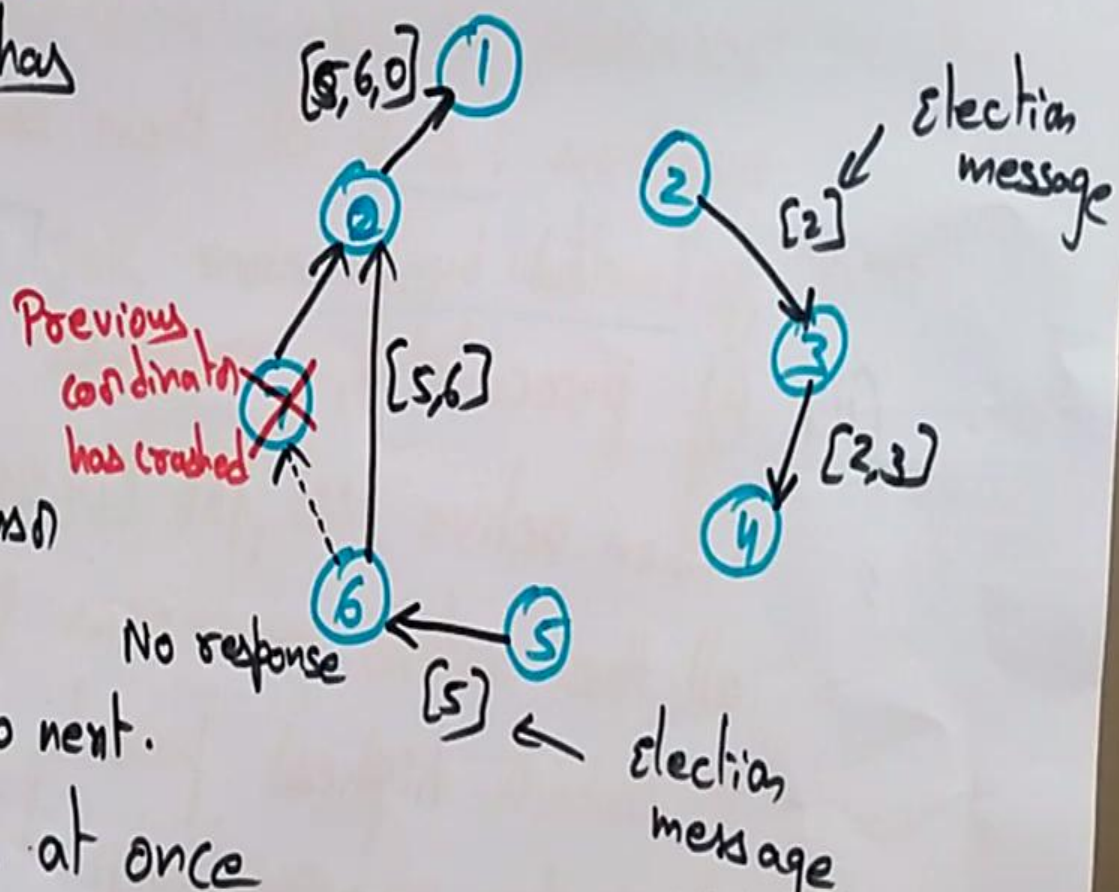
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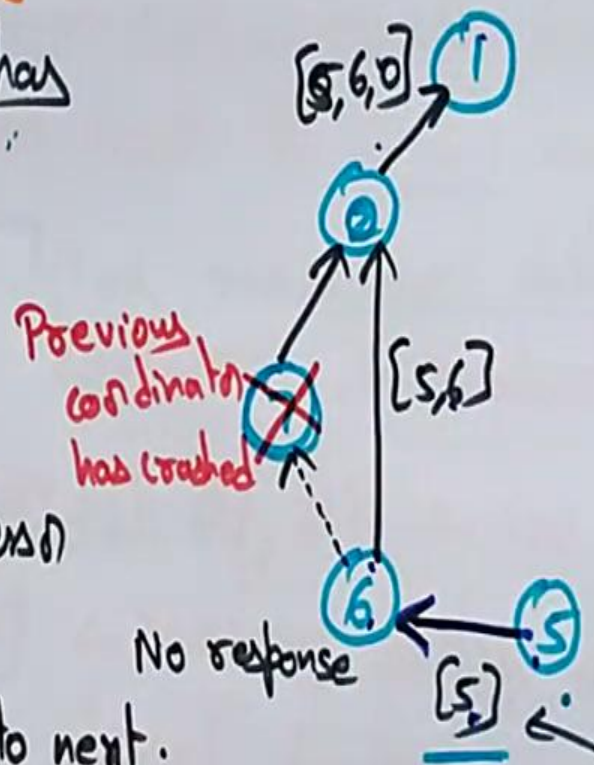
Example 1 "Ring Algorithm"

- P thinks the Coordinator has crashed; builds an election message which contains its own ID no.
- Sends to first live successor
- Each process adds its own number and forwards to next.
- Ok to have two elections at once

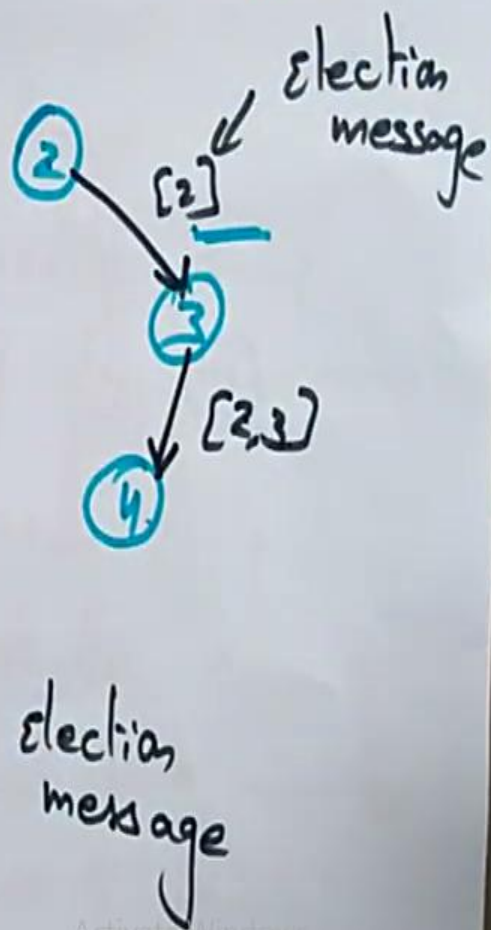


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0-7 - participating n/w.



→ when the message returns to P, it sees its own process ID in the list & knows that the circuit is complete.

→ P circulates a COORDINATOR message with the new high number.

→ Here, both 2 and 5 elected 6:

[5, 6, 0, 1, 2, 3, 4]

[2, 3, 4, 5, 6, 0, 1]

6-coordinator