

# Time Complexity

1.

```
for i in range(n):  
    i = i*2
```

Since, the loop iterates for  $n$  number of times so the time complexity is  $O(n)$ .

2.

```
def fact_iter(n):  
    "assumes n an int >= 0"  
    answer = 1  
    while n > 1:  
        answer *= n  
        n -= 1  
    return answer
```

The loop iterates  $n$  times because it starts with  $n$  and decrements  $n$  by 1 in each iteration until it reaches 1. Hence, the time complexity is  $O(n)$ .

3.

```
A = [1, 2, 3, 4]  
B = [2, 3, 4, 5, 6]  
for i in A:  
    for j in B:  
        if i < j:  
            print('{}, {}'.format(i, j))
```

Here, the outer loop is iterated till the length of  $A$  and the inner loop is iterated till the length of  $B$ .

If we consider the length of  $A$  and  $B$  as  $m$  and  $n$  respectively, then the time complexity will be  $O(m*n)$ .

## 4.

```
L = [1, 2, 3, 4, 5, 6, 7, 8]
for i in range(len(L)//2):
    other = len(L) - i - 1
    temp = L[i]
    L[i] = L[other]
    L[other] = temp
```

The loop iterates  $\text{len}(L)//2$  times where  $\text{len}(L)$  denotes the length of the list. The approximated time complexity of the above code is  $O(n/2)$  which ultimately simplifies to  $O(n)$  where  $n$  is the length of the list  $L$ .

## 5.

```
def fib(n):
    if n==1 or n==0:
        return 1
    else:
        return fib(n-1) + fib(n-2)
```

Here, the number of recursive function calls doubles for each increment of  $n$ . In the worst case, when  $n$  is large, the recursive calls will create a binary tree-like structure, where each node has two child nodes. Therefore, the number of nodes (recursive calls) in the tree will be approximately  $2^n$ , leading to an exponential number of function calls. As a result, the time complexity for this recursive function is  $O(2^n)$ .