



HOME CONTESTS GYM PROBLEMSET GROUPS RATING API RCC W VK CUP CROC COC

INDY256 BLOG TEAMS SUBMISSIONS GROUPS CONTESTS

indy256's blog

Dynamic Programming Optimizations

By indy256, 3 years ago, 3, 2

Several recent problems on Codeforces concerned dynamic programming optimization techniques.

The following table summarizes methods known to me.

→ Pay attention

Before contest Codeforces Round #356 (Div. 1) 6 days

Before contest
Codeforces Round #356 (Div. 2)
6 days

Like 64 people like this. Sign Up to see what your friends like.

Name	Original Recurrence	Sufficient Condition of Applicability	Ori Co	→ Top rated		
				#	User	Rating
Convex Hull Optimization1	$dp[i] = min_{j < i} \{dp[j] + b[j] \star a[i]\}$	$b[j] \ge b[j+1]$ optionally $a[i] \le a[i+1]$	0(21	tourist	3554
				2	Petr	3392
'		, , [] _ []		3	TooDifficult	3240
Convex Hull Optimization2	$dp[i][j] = min_{k < j} \{ dp[i-1][k] + b[k] * a[j] \}$	$b[k] \ge b[k+1]$ optionally $a[j] \le a[j+1]$	O(1	24	rng_58	3102
				5	vepifanov	3077
Divide and	$dp[i][j] = min_{k < j} \{dp[i-1][k] + C[k][j]\}$	$A[i][j] \le A[i][j+1]$	O(1	6	jcvb	3008
				7	izrak	3007
Conquer Optimization				8	Egor	2972
Opumzauom				9	ikatanic	2964
Knuth	$dp[i][j] = min_{i < k < j} \{dp[i][k] + dp[k][j]\} + C[i][j]$	$A[i,j-1] \le A[i,j] \le A[i+1,j]$	0(³ 10	² anta	2963
Optimization				Countr	ies <u>Cities</u> <u>Organizations</u>	<u>View all →</u>

Notes:

- A[i][j] the smallest k that gives optimal answer, for example in dp[i][j] = dp[i-1][k] + C[k][j]
- C[i][j] some given cost function
- We can generalize a bit in the following way: $dp[i] = min_{j < i} \{F[j] + b[j] * a[i]\}$, where F[j] is computed from dp[j] in constant time.
- It looks like Convex Hull Optimization2 is a special case of Divide and Conquer Optimization.
- It is claimed (in the references) that **Knuth Optimization** is applicable if C[i][j] satisfies the following 2 conditions:
- quadrangle inequality:

$$C[a][c] + C[b][d] \le C[a][d] + C[b][c], a \le b \le c \le d$$

- monotonicity: $C[b][c] \le C[a][d], \ a \le b \le c \le d$
- It is claimed (in the references) that the recurrence $dp[j] = min_{i < j} \{dp[i] + C[i][j]\}$ can be solved in O(nlogn) (and even O(n)) if C[i][j] satisfies **quadrangle inequality**. **WJMZBMR** described how to solve some case of this problem.

Open questions:

- 1. Are there any other optimization techniques?
- 2. What is the sufficient condition of applying **Divide and Conquer Optimization** in terms of function C[i][j]? Answered

→ Top contributors

#	User	Contrib.
1	Errichto	174
2	Petr	162
2	Zlobober	162
4	Edvard	160
5	gKseni	155
6	Swistakk	153
7	chrome	151
8	Xellos	142
9	I_love_Hoang_Yen	139
10	amd	137
		View all →

→ Find user					
Handle:					
	Find				

References:

- "Efficient dynamic programming using quadrangle inequalities" by F. Frances Yao. find
- "Speed-Up in Dynamic Programming" by F. Frances Yao. find
- "The Least Weight Subsequence Problem" by D. S. Hirschberg, L. L. Larmore. find
- "Dynamic programming with convexity, concavity and sparsity" by Zvi Galil, Kunsoo Park. find
- "A Linear-Time Algorithm for Concave One-Dimensional Dynamic Programming" by Zvi Galil, Kunsoo Park. find

Please, share your knowledge and links on the topic.

dynamic programming, knuth optimization, convex hull optimization



△ +388 ▽



indy256







Comments (32)

Write comment?

3 years ago, # |



A +27

Here is another way to optimize some 1D1D dynamic programming problem that I know.

Suppose that the old choice will only be worse compare to the new choice(it is quite common in such kind of problems).

Then suppose at current time we are deal with dp_i , and we have some choice $a_0 \le a_1 \le a_2, ..., a_{k-1} \le a_k$ then we know at current time a_i should be better than a_{i+1} . Otherwise it will never be better than a_{i+1} , so it is useless.



we can use a deque to store all the a_i .

And Also Let us denote D(a, b) as the smallest i such that choice b will be better than a

If $D(a_i, a_{i+1}) > D(a_{i+1}, a_{i+2})$, we can find a_{i+1} is also useless because when it overpass a_i , it is already overpass by a_{i+2} .

So we also let $D(a_i, a_{i+1}) \le D(a_{i+1}, a_{i+2})$. then we can find the overpass will only happen at the front of the deque.

So we can maintain this deque quickly, and if we can solve D(a, b) in O(1), it can run in O(n).

→ Reply



3 years ago, # ^ |

▲ +3 ▼

→ +5 ▼

could you please give some example problems?

→ <u>Reply</u>



3 years ago, <u>#</u> |

For question 2: The sufficient condition is:

 $C[a][d] + C[b][c] \ge C[a][c] + C[b][d]$ where a < b < c < d.

→ Reply

3 years ago, # $^{\wedge}$ | A 0 Is it quadrangle inequalities? ∀i≤ j,w[i, il+w[i+1 i+1]<w[i+1 il+w[i i+1] and are these two

→ Recent actions

sajalhsn13 → How can I optimize my code?

Wild_Hamster → Codeforces Round #355 (Div. 2) 💭

Whistle → what the story behind your ACM team name?

kursatbakis0 → I have been coding for 5 years and I am giving up 🦃

Yousef_Salama → Time to revive the Arab and Africa Region 📡

Wild_Hamster → Codeforces Round #355 (Div. 2) Editorial 💭

iscsi → Invitation to participate in CodeChef

komendart → Codeforces Round #353 (Div. 2) Editorial 💭

gsmcoder97 → B.Little Elephant And Array

belowthebelt → Invitation to June-Easy '16 on HackerEarth 💭

Bredor → Unbelieveable story of success at ACM ICPC final ©

zerocooL_tryHard → Contest Frequency Reduced.

csacademy → Beta Round #6 csacademy.com 📡

zholnin → <u>IPSC</u> — <u>looking for teammates</u>

MikeMirzayanov \rightarrow g++: beware of the signed overflow ©

 $\textbf{MerelyMHP} \rightarrow \underline{\text{Why most programmer prefer}}$ C++ than C? ©

muratt → Codeforces Round #352 Editorial

Bobur → How to find minimim path graph with bees colony algorithm?

Sawako → Complete and Utter Contest Devastation, and why Failure is Golden 📡

el_falso → Help! 💭

rachitjain → Unsuccessful Hacking Attempt

MikeMirzayanov → Codeforces Rating

Detailed →

+18



ון מווע מוב באר מווע פון און הארן הארבן הארב inequalities equivalent except the >= & <=?</pre> → Reply



There is one more optimization of dimanic progamming: 101E - Candies and Stones (editoral)

→ Reply

3 years ago, <u>#</u> |



▲ +13 ▼ 3 years ago, # |

you have put problem "B. Cats Transport" in "Convex Hull Optimization1",

kingofnumbers

actually it belongs to "Convex Hull Optimization2" → Reply



3 years ago, # $^{\wedge}$ | ▲ +5 ▼ fixed → Reply



3 years ago, # | ← Rev. 2 ▲ +55 V

For this moment it's the most useful topic of this year. Exactly in the middle: June 30th, 2013.

→ Reply



A +8 V 3 years ago, <u>#</u> |

this one seemed a nice dp with optimization to me:https://www.hackerrank.com/contests/monthly/challenges/alienlanguages

→ Reply

3 years ago, <u>#</u> |





hogloid

The problem mentioned in the article (Breaking Strings) is "Optimal Binary Search Tree Problem", traditional one.

It can be solved by simple DP in O(N^3), by using Knuth's optimization , in O(N^2) . But it still can be solved in O(NlogN) — http://poj.org/problem? id=1738 (same problem but bigger testcases) (I don't know how to solve it. I hear the algorithm uses meld-able heap)

 \rightarrow Reply

A +20 V 3 years ago, <u>#</u> |

Convex Hull Optimization 1 Problems:

- APIO 2010 task Commando
- TRAKA
- ACQUIRE
- SkyScrapers (+Data Structures)



Convex Hull Optimization 2 Problems:

BAABO

Giorgos Christoglou

Convex Hull Optimization 3 Problems (No conditions for a[] array and b[] array):

- GOODG
- BOI 2012 Day 2 Balls
- · Cow School

△ 0 ▼

A 0 V

- COM SCHOOL
- Solution-Video
- \rightarrow Reply



5 months ago, $\mbox{$\frac{\#}{\ }$}$ \ \(\triangle \) GOODG can be solved with Type 1

3 weeks ago, <u>#</u> <u>^</u> ∣

→ <u>Reply</u>





How? I noticed that, in this problem, b[j] follows no order and a[i] can be either decreasing or increasing, depending on how the equation is modeled. I was able to solve it using the fully dynamic variant, but I can't see how to apply the "type 1" optimization.

3 weeks ago, # \land | \leftarrow Rev. 2

→ Reply





Can you add a link to your code I tried to implement the dynamic variant few weeks ago but there were so many bugs in my code :(.Maybe yours can help:/.

→ Reply



A +3 🔻 5 months ago, # $^{\wedge}$ |

New link for Commando:



→ Reply



zscefn

3 years ago, # |

△ 0 ▼

A +4 V

▲ +1 ▼

For some reason I cannot open the links with firefox because they go over the Top Rated table.

→ <u>Reply</u>



3 years ago, # <u>^</u> | Try to zoom out, pressing Ctrl + -→ Reply

indy256



← Rev. 2 **▲ +8** ▼ 3 years ago, # |



One more problem where Knuth Optimization is used: Andrew Stankevich Contest 10, Problem C.

BTW, does anybody know how to insert a direct link to a problem from gyms?

→ Reply





22 months ago, # |

I need some problems to solve on Divide and Conquer Optimization. Where can I find them? An online judge / testdata available would be helpful.

→ Reply



22 months ago, # ^ |

Check this one: Guardians of the Lunatics

 \rightarrow Reply



22 months ago, # ^ |

△ 0 ▼

Learnt Divide and Conquer Optimization just from there. :P That is why I'm asking for more problems to practice.

:D → Reply



8 months ago, # ^ |

Is this the best complexity for this problem? Can't we do any better? Can't we somehow turn the logL needed into a constant?

→ Reply



micklepru

8 months ago, # ^ | **△** 0 ▼ We can, using that opt[i-1][j] <= opt[i] $[j] \leftarrow opt[i][j+1]$.

Key thing is to see that opt function is monotone for both arguments. With that observation, we don't need to use binary

Check out my submission.

→ Reply



92anuraq

22 months ago, # |

▲ +3 ▼

can anyone provide me good editorial for dp with bitmask .

 \rightarrow Reply



22 months ago, # |

△ 0 ▼

Has matrix-exponent optimizations been included here?

→ Reply



20 months ago, # |

A +2 V

Can matrix chain multiplication problem b also optimized by knuth optimization? If not, dn why?

 \rightarrow Reply

19 months ago, # ^ |

▲ +3 ▼

Quote from the first of the references above:

The monotonicity property for the division points does not hold for the matrix multiplication chain problem...



indy256

Consider the matrices M1,M2,M3,M4 with dimensions 2x3, 3x2, 2x10, and 10x1, respectively. As can be easily verified, the proper order to compute M1M2M3 is to parenthesize it as (M1M2)M3, while the optimal computation of M1M2M3M4 corresponds to M1(M2(M3M4)).

The second reference gives $O(n^2)$ dynamic programming solution, based on some properties of the matrix chain multiplication problem.

There is also an $O(n * \log n)$ algorithm by Hu and Shing.

→ Reply



13 months ago, # ^ |

Link to the Hu and Shing algorithm?

→ <u>Reply</u>

mayankp

15 months ago, # | ▲ +1 ▼

What are some recent USACO questions that use this technique or variations of it?

→ Reply

11 months ago, # |

← Rev. 6 **△** 0 ▼

Can this problem be solved using convex hull optimization?

You are given a sequence A of N positive integers. Let's define "value of a splitting" the sequence to K blocks as a sum of maximums in each of Kblocks. For given K find the minimal possible value of splittings.



 $N <= 10^5$

$$K \le 100$$

Input: 5 2

1 2 3 4 5

→ Reply



11 months ago, # ^ |

Output:

△ 0 ▼

I don't think so, but I guess it can be solved by Divide And Conquer optimization.

→ <u>Reply</u>

10 months ago, $\ \underline{\#}$ |



Could you elaborate a little me more in the "Convex Hull Optimization2" and other sections for the clearer notations.



For example, You have "k" — a constant in O(kn^2). So the first dimension is of the length K and the second dimension is of the length N?

I think it would be clearer if you can write dp[n], dp[k][n] ... instead of dp[i], dp[i][j] .

Best regards,

→ Reply

0

Codeforces (c) Copyright 2010-2016 Mike Mirzayanov The only programming contests Web 2.0 platform Server time: Jun/02/2016 20:14:18^{UTC+5.5} (p1). Desktop version, switch to mobile version.