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# Pollard's kangaroo algorithm

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In [computational number theory](#) and [computational algebra](#), **Pollard's kangaroo algorithm** (aka **Pollard's lambda algorithm**, see [Naming](#) below) is an [algorithm](#) for solving the [discrete logarithm](#) problem. The algorithm was introduced in 1978 by the number theorist [J. M. Pollard](#), in the same paper <sup>[1]</sup> as his better-known [p algorithm](#) for solving the same problem. Although Pollard described the application of his algorithm to the discrete logarithm problem in the multiplicative group of units modulo a prime *p*, it is in fact a generic discrete logarithm algorithm—it will work in any finite cyclic group.

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## The algorithm [\[edit\]](#)

Suppose *G* is a finite cyclic group of order *n* which is generated by the element *α*, and we seek to find the discrete logarithm *x* of the element *β* to the base *α*. In other words, we seek *x* ∈ *Z*<sub>*n*</sub> such that *α*<sup>*x*</sup> = *β*. The lambda algorithm allows us to search for *x* in some subset {*a*, . . . , *b*} ⊂ *Z*<sub>*n*</sub>. We may search the entire range of possible logarithms by setting *a* = 0 and *b* = *n* − 1, although in this case [Pollard's rho algorithm](#) is more efficient. We proceed as follows:

- Choose a set *S* of integers and define a [pseudorandom](#) map *f* : *G* → *S*.
- Choose an integer *N* and compute a sequence of group elements {*x*<sub>0</sub>, *x*<sub>1</sub>, . . . , *x*<sub>*N*</sub>} according to:

- x*<sub>0</sub> = *α*<sup>*b*</sup>
- x*<sub>*i*+1</sub> = *x*<sub>*i*</sub>*α*<sup>*f*(*x*<sub>*i*</sub>)</sup> for *i* = 0, 1, . . . , *N* − 1

- Compute

$$d = \sum_{i=0}^{N-1} f(x_i).$$

Observe that:

$$x_N = x_0 \alpha^d = \alpha^{b+d}.$$

- Begin computing a second sequence of group elements {*y*<sub>0</sub>, *y*<sub>1</sub>, . . . } according to:

- y*<sub>0</sub> = *β*
- y*<sub>*i*+1</sub> = *y*<sub>*i*</sub>*α*<sup>*f*(*y*<sub>*i*</sub>)</sup> for *i* = 0, 1, . . . , *N* − 1

and a corresponding sequence of integers {*d*<sub>0</sub>, *d*<sub>1</sub>, . . . } according to:

$$d_n = \sum_{i=0}^{n-1} f(y_i).$$

Observe that:

$$y_i = y_0 \alpha^{d_i} = \beta \alpha^{d_i} \text{ for } i = 0, 1, \dots, N - 1$$

- Stop computing terms of {*y*<sub>*i*</sub>} and {*d*<sub>*i*</sub>} when either of the following conditions are met:

A) *y*<sub>*j*</sub> = *x*<sub>*N*</sub> for some *j*. If the sequences {*x*<sub>*i*</sub>} and {*y*<sub>*j*</sub>} "collide" in this manner, then we have:

$$x_N = y_j \Rightarrow \alpha^{b+d} = \beta \alpha^{d_j} \Rightarrow \beta = \alpha^{b+d-d_j} \pmod{n} \Rightarrow x \equiv b+d-d_j \pmod{n}$$

and so we are done.

B) *d*<sub>*i*</sub> > *b* − *a* + *d*. If this occurs, then the algorithm has failed to find *x*. Subsequent attempts can be

made by changing the choice of  $S$  and/or  $f$ .

Complexity [\[edit\]](#)

Pollard gives the time complexity of the algorithm as  $O(\sqrt{b-a})$ , based on a probabilistic argument which follows from the assumption that  $f$  acts pseudorandomly. Note that when the size of the set  $\{a, \dots, b\}$  to be searched is measured in **bits**, as is normal in **complexity theory**, the set has size  $\log(b-a)$ , and so the algorithm's complexity is  $O(\sqrt{b-a}) = O(2^{\frac{1}{2} \log(b-a)})$ , which is exponential in the problem size. For this reason, Pollard's lambda algorithm is considered an **exponential time** algorithm. For an example of a **subexponential time** discrete logarithm algorithm, see the **index calculus algorithm**.

## Naming [\[edit\]](#)

The algorithm is well known by two names.

The first is "Pollard's kangaroo algorithm". This name is a reference to an analogy used in the paper presenting the algorithm, where the algorithm is explained in terms of using a *tame* [kangaroo](#) to trap a *wild* kangaroo. Pollard has explained<sup>[2]</sup> that this analogy was inspired by a "fascinating" article published in the same issue of *Scientific American* as an exposition of the [RSA public key cryptosystem](#). The article<sup>[3]</sup> described an experiment in which a kangaroo's "energetic cost of locomotion, measured in terms of oxygen consumption at various speeds, was determined by placing kangaroos on a [treadmill](#)".


The second is "Pollard's lambda algorithm". Much like the name of another of Pollard's discrete logarithm algorithms, [Pollard's rho algorithm](#), this name refers to the similarity between a visualisation of the algorithm and the [Greek letter lambda](#) ( $\lambda$ ). The shorter stroke of the letter lambda corresponds to the sequence  $\{x_i\}$ , since it starts from the position b to the right of x. Accordingly, the longer stroke corresponds to the sequence  $\{y_i\}$ , which "collides with" the first sequence (just like the strokes of a lambda intersect) and then follows it subsequently.

Pollard has expressed a preference for the name "kangaroo algorithm",<sup>[4]</sup> as this avoids confusion with some parallel versions of his rho algorithm, which have also been called "lambda algorithms".

See also [\[edit\]](#)

- Rainbow table

## References [\[edit\]](#)

1. <sup>^</sup> J. Pollard, *Monte Carlo methods for index computation mod p*, Mathematics of Computation, Volume 32, 1978
2. <sup>^</sup> J. M. Pollard, *Kangaroos, Monopoly and Discrete Logarithms*, Journal of Cryptology, Volume 13, pp 437-447, 2000
3. <sup>^</sup> T. J. Dawson, *Kangaroos*, Scientific American, August 1977, pp. 78-89
4. <sup>^</sup> <http://sites.google.com/site/jmptidcott2/> 

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