

Main page Contents Featured content Current events Random article Donate to Wkipedia Wkipedia store

Interaction

Help About Wikipedia Community portal Recent changes Contact page

Tools

What links here Related changes Upload file Special pages Permanent link Page information

Cite this page

Print/export

Wikidata item

Create a book
Download as PDF
Printable version

Languages Deutsch Français

Ædit links

Article Talk Read Edit View history Search Q

Maekawa's algorithm

From Wikipedia, the free encyclopedia (Redirected from Maekawa's Algorithm)



This article **does not cite any references or sources**. Please help improve this article by adding citations to reliable sources. Unsourced material may be challenged and removed. (*December 2009*)

Maekawa's algorithm is an algorithm for mutual exclusion on a distributed system. The basis of this algorithm is a quorum like approach where any one site needs only to seek permissions from a subset of other sites.

Contents [hide]

- 1 Algorithm
 - 1.1 Terminology
 - 1.2 Algorithm
 - 1.3 Performance
- 2 See also
- 3 References

Algorithm [edit]

Terminology [edit]

- A site is any computing device which is running the Maekawa's Algorithm
- For any one request of the critical section:
 - The requesting site is the site which is requesting entry into the critical section.
 - The receiving site is every other site which is receiving the request from the requesting site.
- ts refers to the local time stamp of the system according to its logical clock.

Algorithm [edit]

Requesting site:

- A requesting site P_i sends a message $ext{request}(ts,i)$ to all sites in its quorum set R_i

Receiving site:

- ullet Upon reception of a $\mathrm{request}(t_S,i)$ message, the receiving site P_i will:
 - If site P_j does not have an outstanding grant message (that is, a grant message that has not been released), then site P_j sends a grant (j) message to site P_i .
 - If site P_j has an outstanding grant message with a process with higher priority than the request, then site P_j sends a $\operatorname{failed}(j)$ message to site P_i and site P_j queues the request from site P_i .
 - If site P_j has an outstanding grant message with a process with lower priority than the request, then site P_j sends an $\operatorname{inquire}(j)$ message to the process which has currently been granted access to the critical section by site P_j . (That is, the site with the outstanding grant message.)
- ullet Upon reception of a $\mathrm{inquire}(j)$ message, the site P_k will:
 - Send a yield(k) message to site P_j if and only if site P_k has received a failed message from some
 other site or if P_k has sent a yield to some other site but have not received a new grant.
- Upon reception of a $\operatorname{vield}(k)$ message, site P_i will:
 - Send a $g_{rant}(j)$ message to the request on the top of its own request queue. Note that the requests at the top are the highest priority.
 - ullet Place P_{k} into its request queue.
- Upon reception of a release(i) message, site P_i will:
 - Delete P_i from its request queue.
 - Send a $\operatorname{grant}(i)$ message to the request on the top of its request queue.

Critical section:

- Site P_i enters the critical section on receiving a grant message from all sites in R_i
- ullet Upon exiting the critical section, P_i sends a $\mathrm{release}(i)$ message to all sites in R_i :

Quorum set (R_x):

A quorum set must abide by the following properties:

- 1. $\forall i \, \forall j \, [R_i \cap R_j \neq \emptyset]$
- 2. $\forall i [P_i \in R_i]$
- 3. $\forall i [|R_i| = K]$
- 4. Site P_i is contained in exactly K request sets

Therefore:

•
$$|R_i| \ge \sqrt{N-1}$$

Performance [edit]

- Number of network messages; $3\sqrt{N}$ to $6\sqrt{N}$
- Synchronization delay: 2 message propagation delays

See also [edit]

- Lamport's bakery algorithm
- Lamport's Distributed Mutual Exclusion Algorithm
- Ricart-Agrawala algorithm
- Raymond's algorithm

References [edit]

- Mamoru Maekawa, Arthur E. Oldehoeft, Rodney R. Oldehoeft (1987). Operating Systems: Advanced Concept. Benjamin/Cummings Publishing Company, Inc.
- B. Sanders (1987). The Information Structure of Distributed Mutual Exclusion Algorithms. ACM Transactions on Computer Systems, Vol. 3, No. 2, pp. 145–59.

Categories: Concurrency control algorithms

This page was last modified on 12 February 2014, at 18:24.

Text is available under the Oreative Commons Attribution-ShareAlike License; additional terms may apply. By using this site, you agree to the Terms of Use and Privacy Policy. Wikipedia® is a registered trademark of the Wikimedia Foundation, Inc., a non-profit organization.

Privacy policy About Wikipedia Disclaimers Contact Wikipedia Developers Mobile view



