

Given a set of non-negative integers, and a value *sum*, determine if there is a subset of the given set with sum equal to given *sum*.

Examples: set[] = {3, 34, 4, 12, 5, 2}, sum = 9

Output: True //There is a subset (4, 5) with sum 9.

Let isSubSetSum(int set[], int n, int sum) be the function to find whether there is a subset of set[] with sum equal to *sum*. n is the number of elements in set[].

The isSubSetSum problem can be divided into two subproblems

...a) Include the last element, recur for n = n-1, sum = sum – set[n-1]

...b) Exclude the last element, recur for n = n-1.

If any of the above the above subproblems return true, then return true.

Following is the recursive formula for isSubSetSum() problem.

```
isSubSetSum(set, n, sum) = isSubSetSum(set, n-1, sum) ||
                           isSubSetSum(arr, n-1, sum-set[n-1])
```

Base Cases:

isSubSetSum(set, n, sum) = false, if sum > 0 and n == 0

isSubSetSum(set, n, sum) = true, if sum == 0

Following is naive recursive implementation that simply follows the recursive structure mentioned above.

```
// A recursive solution for subset sum problem
#include <stdio.h>
```

```
// Returns true if there is a subset of set[] with sun equal to given sum
bool isSubSetSum(int set[], int n, int sum)
{
    // Base Cases
    if (sum == 0)
        return true;
    if (n == 0 && sum != 0)
        return false;

    // If last element is greater than sum, then ignore it
    if (set[n-1] > sum)
        return isSubSetSum(set, n-1, sum);

    /* else, check if sum can be obtained by any of the following
       (a) including the last element
       (b) excluding the last element */
    return isSubSetSum(set, n-1, sum) || isSubSetSum(set, n-1, sum-set[n-1]);
}
```

```
// Driver program to test above function
int main()
{
    int set[] = {3, 34, 4, 12, 5, 2};
    int sum = 9;
    int n = sizeof(set)/sizeof(set[0]);
    if (isSubSetSum(set, n, sum) == true)
        printf("Found a subset with given sum");
    else
        printf("No subset with given sum");
    return 0;
}
```

Output:

Found a subset with given sum

The above solution may try all subsets of given set in worst case. Therefore time complexity of the above solution is exponential. The problem is in-fact **NP-Complete** (There is no known polynomial time solution for this problem).

We can solve the problem in Pseudo-polynomial time using Dynamic programming. We create a boolean 2D table subset[][] and fill it in bottom up manner. The value of subset[i][j] will be true if there is a subset of set[0..j-1] with sum equal to i., otherwise false. Finally, we return subset[sum][n]

```
// A Dynamic Programming solution for subset sum problem
#include <stdio.h>
```

```
// Returns true if there is a subset of set[] with sun equal to given sum
bool isSubsetSum(int set[], int n, int sum)
{
    // The value of subset[i][j] will be true if there is a subset of set[0..j-1]
    // with sum equal to i
    bool subset[sum+1][n+1];

    // If sum is 0, then answer is true
    for (int i = 0; i <= n; i++)
        subset[0][i] = true;

    // If sum is not 0 and set is empty, then answer is false
    for (int i = 1; i <= sum; i++)
        subset[i][0] = false;

    // Fill the subset table in bottom up manner
    for (int i = 1; i <= sum; i++)
    {
        for (int j = 1; j <= n; j++)
        {
            subset[i][j] = subset[i][j-1];
            if (i >= set[j-1])
                subset[i][j] = subset[i][j] || subset[i - set[j-1]][j-1];
        }
    }

    /* // uncomment this code to print table
    for (int i = 0; i <= sum; i++)
    {
        for (int j = 0; j <= n; j++)
            printf ("%4d", subset[i][j]);
        printf("\n");
    } */

    return subset[sum][n];
}
```

```
// Driver program to test above function
int main()
{
    int set[] = {3, 34, 4, 12, 5, 2};
    int sum = 9;
    int n = sizeof(set)/sizeof(set[0]);
    if (isSubsetSum(set, n, sum) == true)
        printf("Found a subset with given sum");
    else
        printf("No subset with given sum");
    return 0;
}
```

Output:

Found a subset with given sum

Time complexity of the above solution is $O(\text{sum} \times n)$.