

# Cover tree

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The **cover tree** is a type of [data structure](#) in [computer science](#) that is specifically designed to facilitate the speed-up of a [nearest neighbor search](#). It is a refinement of the Navigating Net data structure, and related to a variety of other data structures developed for indexing intrinsically low-dimensional data.<sup>[1]</sup>

The tree can be thought of as a hierarchy of levels with the top level containing the root [point](#) and the bottom level containing every point in the metric space. Each level  $C$  is associated with an integer value  $i$  that decrements by one as the tree is descended. Each level  $C$  in the cover tree has three important properties:

- **Nesting:**
- **Covering:** For every point , there exists a point  such that the distance from  to  is less than or equal to  and exactly one such  is a parent of .
- **Separation:** For all points , the distance from  to  is greater than .

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## Complexity[\[edit\]](#)

### Find[\[edit\]](#)

Like other [metric trees](#) the cover tree allows for nearest neighbor searches in  where  is a constant associated with the dimensionality of the dataset and  $n$  is the cardinality. To compare, a basic linear search requires , which is a much worse dependence on . However, in high-dimensional [metric spaces](#) the  constant is non-trivial, which means it cannot be ignored in complexity analysis. Unlike other metric trees, the cover tree has a theoretical bound on its constant that is based on the dataset's [expansion constant](#) or doubling constant (in the case of approximate NN retrieval). The bound on search time is  where  is the expansion constant of the dataset.

### Insert[\[edit\]](#)

Although cover trees provide faster searches than the naive approach, this advantage must be weighed with the additional cost of maintaining the data structure. In a naive approach adding a new point to the dataset is trivial because order does not need to be preserved, but in a cover tree it can take  time. However, this is an upper-bound, and some techniques have been implemented that seem to improve the performance in practice.<sup>[2]</sup>

### Space[\[edit\]](#)

The cover tree uses implicit representation to keep track of repeated points. Thus, it only requires  $O(n)$  space.

## See also[\[edit\]](#)

- [Nearest neighbor search](#)
- [kd-tree](#)

## References[\[edit\]](#)

Notes

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