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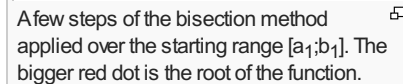
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The **bisection method** in **mathematics** is a **root-finding method** that repeatedly **bisects** an **interval** and then selects a subinterval in which a **root** must lie for further processing. It is a very simple and robust method, but it is also relatively slow. Because of this, it is often used to obtain a rough approximation to a solution which is then used as a starting point for more rapidly converging methods.<sup>[1]</sup> The method is also called the **interval halving** method,<sup>[2]</sup> the **binary search method**,<sup>[3]</sup> or the **dichotomy method**.<sup>[4]</sup>

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1. Calculate  $c$ , the midpoint of the interval,  $c = 0.5 * (a + b)$ .
2. Calculate the function value at the midpoint,  $f(c)$ .
3. If convergence is satisfactory (that is,  $a - c$  is sufficiently small, or  $f(c)$  is sufficiently small), return  $c$  and stop iterating.
4. Examine the sign of  $f(c)$  and replace either  $(a, f(a))$  or  $(b, f(b))$  with  $(c, f(c))$  so that there is a zero crossing within the new interval.

When implementing the method on a computer, there can be problems with finite precision, so there are often additional convergence tests or limits to the number of iterations. Although  $f$  is continuous, finite precision may preclude a function value ever being zero. For  $f(x) = x - \pi$ , there will never be a finite representation of  $x$  that gives zero. Floating point representations also have limited precision, so at some point the midpoint of  $[a, b]$  will be either  $a$  or  $b$ .

### Algorithm [\[edit\]](#)

The method may be written in [pseudocode](#) as follows:<sup>[7]</sup>

```

INPUT: Function  $f$ , endpoint values  $a, b$ , tolerance  $TOL$ , maximum iterations  $NMAX$ 
CONDITIONS:  $a < b$ , either  $f(a) < 0$  and  $f(b) > 0$  or  $f(a) > 0$  and  $f(b) < 0$ 
OUTPUT: value which differs from a root of  $f(x)=0$  by less than  $TOL$ 

 $N \leftarrow 1$ 
While  $N \leq NMAX$  # limit iterations to prevent infinite loop
     $c \leftarrow (a + b)/2$  # new midpoint
    If  $f(c) = 0$  or  $(b - a)/2 < TOL$  then # solution found
        Output( $c$ )
        Stop
    EndIf
     $N \leftarrow N + 1$  # increment step counter
    If  $\text{sign}(f(c)) = \text{sign}(f(a))$  then  $a \leftarrow c$  else  $b \leftarrow c$  # new interval
EndWhile
Output("Method failed.") # max number of steps exceeded

```

### Example: Finding the root of a polynomial [\[edit\]](#)

Suppose that the bisection method is used to find a root of the polynomial

$$f(x) = x^3 - x - 2.$$

First, two numbers  $a$  and  $b$  have to be found such that  $f(a)$  and  $f(b)$  have opposite signs. For the above function,  $a = 1$  and  $b = 2$  satisfy this criterion, as

$$f(1) = (1)^3 - (1) - 2 = -2$$

and

$$f(2) = (2)^3 - (2) - 2 = +4.$$

Because the function is continuous, there must be a root within the interval  $[1, 2]$ .

In the first iteration, the end points of the interval which brackets the root are  $a_1 = 1$  and  $b_1 = 2$ , so the midpoint is

$$c_1 = \frac{2 + 1}{2} = 1.5$$

The function value at the midpoint is  $f(c_1) = (1.5)^3 - (1.5) - 2 = -0.125$ . Because  $f(c_1)$  is negative,  $a = 1$  is replaced with  $a = 1.5$  for the next iteration to ensure that  $f(a)$  and  $f(b)$  have opposite signs. As this continues, the interval between  $a$  and  $b$  will become increasingly smaller, converging on the root of the function. See this happen in the table below.

Iteration	$a_n$	$b_n$	$c_n$	$f(c_n)$
1	1	2	1.5	-0.125
2	1.5	2	1.75	1.6093750
3	1.5	1.75	1.625	0.6660156
4	1.5	1.625	1.5625	0.2521973
5	1.5	1.5625	1.5312500	0.0591125
6	1.5	1.5312500	1.5156250	-0.0340538
7	1.5156250	1.5312500	1.5234375	0.0122504
8	1.5156250	1.5234375	1.5195313	-0.0109712
9	1.5195313	1.5234375	1.5214844	0.0006222
10	1.5195313	1.5214844	1.5205078	-0.0051789

11	1.5205078	1.5214844	1.5209961	−0.0022794
12	1.5209961	1.5214844	1.5212402	−0.0008289
13	1.5212402	1.5214844	1.5213623	−0.0001034
14	1.5213623	1.5214844	1.5214233	0.0002594
15	1.5213623	1.5214233	1.5213928	0.0000780

After 13 iterations, it becomes apparent that there is a convergence to about 1.521: a root for the polynomial.

## Analysis [\[edit\]](#)

The method is guaranteed to converge to a root of  $f$  if  $f$  is a [continuous function](#) on the interval  $[a, b]$  and  $f(a)$  and  $f(b)$  have opposite signs. The [absolute error](#) is halved at each step so the method [converges linearly](#), which is comparatively slow.

Specifically, if  $c_1 = (a+b)/2$  is the midpoint of the initial interval, and  $c_n$  is the midpoint of the interval in the  $n$ th step, then the difference between  $c_n$  and a solution  $c$  is bounded by<sup>[8]</sup>

$$|c_n - c| \leq \frac{|b - a|}{2^n}.$$

This formula can be used to determine in advance the number of iterations that the bisection method would need to converge to a root to within a certain tolerance. The number of iterations needed,  $n$ , to achieve a given error (or tolerance),  $\epsilon$ , is given by:  $n = \log_2 \left( \frac{\epsilon_0}{\epsilon} \right) = \frac{\log \epsilon_0 - \log \epsilon}{\log 2}$ ,

where  $\epsilon_0 = \text{initial bracket size} = b - a$ .

Therefore, the linear convergence is expressed by  $\epsilon_{n+1} = \text{constant} \times \epsilon_n^m$ ,  $m = 1$ .

## See also [\[edit\]](#)

- [Secant method](#)
- [Newton's method](#)
- [Root-finding algorithm](#)
- [Binary search algorithm](#)
- [Lehmer–Schur algorithm](#), generalization of the bisection method in the complex plane
- [Nested intervals](#)
- [Brent's method](#)

## References [\[edit\]](#)

- ↑ Burden & Faires 1985, p. 31
- ↑ <http://siber.cankaya.edu.tr/NumericalComputations/ceng375/node32.html>
- ↑ Burden & Faires 1985, p. 28
- ↑ *Encyclopedia of Mathematics*
- ↑ If the function has the same sign at the endpoints of an interval, the endpoints may or may not bracket roots of the function.
- ↑ Burden & Faires 1985, p. 28 for section
- ↑ Burden & Faires 1985, p. 29. This version recomputes the function values at each iteration rather than carrying them to the next iterations.
- ↑ Burden & Faires 1985, p. 31, Theorem 2.1

- Burden, Richard L.; Faires, J. Douglas (1985), "2.1 The Bisection Algorithm", *Numerical Analysis* (3rd ed.), PWS Publishers, [ISBN 0-87150-857-5](#)

### Further reading [\[edit\]](#)

- Corliss, George (1977), "Which root does the bisection algorithm find?", *SIAM Review* **19** (2): 325–327, [doi:10.1137/1019044](#) [↗](#), [ISSN 1095-7200](#) [↗](#)
- Kaw, Autar; Kalu, Egwu (2008), *Numerical Methods with Applications* [↗](#) (1st ed.)

### External links [\[edit\]](#)

- [Weisstein, Eric W.](#), "Bisection" [↗](#), *MathWorld*.
- [Bisection Method](#) [↗](#) Notes, PPT, Mathcad, Maple, Matlab, Mathematica from [Holistic Numerical Methods Institute](#) [↗](#)



Wikiversity has learning materials about *[The bisection method](#)*



The Wikibook *[Numerical Methods](#)* has a page on the topic of: *[Equation Solving](#)*

Categories: [Root-finding algorithms](#)

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