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# Simple precedence parser

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In [computer science](#), a **simple precedence parser** is a type of [bottom-up parser](#) for [context-free grammars](#) that can be used only by [simple precedence grammars](#).

The implementation of the parser is quite similar to the generic [bottom-up parser](#). A stack is used to store a [viable prefix](#) of a [sentential form](#) from a [rightmost derivation](#). Symbols  $\lessdot$ ,  $\dot{\phantom{x}}$  and  $\gtrdot$  are used to identify the **pivot**, and to know when to **Shift** or when to **Reduce**.

## Implementation [\[edit\]](#)

- Compute the [Wirth–Weber precedence relationship](#) table.
- Start with a stack with only the **starting marker** \$.
- Start with the string being parsed (**Input**) ended with an **ending marker** \$.
- While not (Stack equals to \$\$ and Input equals to \$) (S = Initial symbol of the grammar)
  - Search in the table the relationship between Top(stack) and NextToken(Input)
  - if the relationship is  $\dot{\phantom{x}}$  or  $\lessdot$ 
    - **Shift**:
      - Push(Stack, relationship)
      - Push(Stack, NextToken(Input))
      - RemoveNextToken(Input)
    - if the relationship is  $\gtrdot$ 
      - **Reduce**:
        - SearchProductionToReduce(Stack)
        - RemovePivot(Stack)
        - Search in the table the relationship between the Non terminal from the production and first symbol in the stack (Starting from top)
        - Push(Stack, relationship)
        - Push(Stack, Non terminal)

SearchProductionToReduce (Stack)

- search the **Pivot** in the stack the nearest  $\lessdot$  from the top
- search in the productions of the grammar which one have the same right side than the **Pivot**

## Example [\[edit\]](#)

Given the language:

```
E --> E + T' | T'
T' --> T
T --> T * F | F
F --> ( E' ) | num
E' --> E
```

**num** is a terminal, and the [lexer](#) parse any integer as **num**.

and the Parsing table:


	E	E'	T	T'	F	+	*	(	)	num	\$
E						$\dot{\phantom{x}}$			$\gtrdot$		$\gtrdot$
E'									$\dot{\phantom{x}}$		
T						$\gtrdot$	$\dot{\phantom{x}}$		$\gtrdot$		$\gtrdot$
T'						$\gtrdot$			$\gtrdot$		$\gtrdot$
F						$\gtrdot$	$\gtrdot$		$\gtrdot$		$\gtrdot$


+										
*										
(										
)										
num										
\$										

STACK	PRECEDENCE	INPUT	ACTION
\$	<	2 * ( 1 + 3 ) \$	SHIFT
\$ < 2	>	* ( 1 + 3 ) \$	REDUCE (F -> num)
\$ < F	>	* ( 1 + 3 ) \$	REDUCE (T -> F)
\$ < T	=	* ( 1 + 3 ) \$	SHIFT
\$ < T = *	<	( 1 + 3 ) \$	SHIFT
\$ < T = * < (	<	1 + 3 ) \$	SHIFT
\$ < T = * < ( < 1	>	+ 3 ) \$	REDUCE 4 times (F -> num) (T -> F) (T' -> T) (E ->T ')
\$ < T = * < ( < E	=	+ 3 ) \$	SHIFT
\$ < T = * < ( < E = +	<	3 ) \$	SHIFT
\$ < T = * < ( < E = + < 3	>	) \$	REDUCE 3 times (F -> num) (T -> F) (T' -> T)
\$ < T = * < ( < E = + = T	>	) \$	REDUCE 2 times (E -> E + T) (E' -> E)
\$ < T = * < ( < E'	=	) \$	SHIFT
\$ < T = * < ( = E' = )	>	\$	REDUCE (F -> ( E' ))
\$ < T = * = F	>	\$	REDUCE (T -> T * F)
\$ < T	>	\$	REDUCE 2 times (T' -> T) (E -> T')
\$ < E	>	\$	ACCEPT

References [\[edit\]](#)

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