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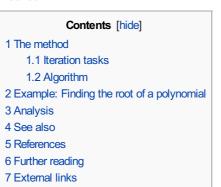
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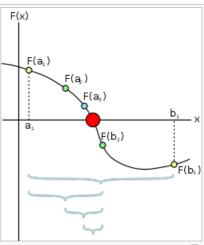
Bisection method

From Wikipedia, the free encyclopedia

This article is about searching continuous function values. For searching a finite sorted array, see binary search algorithm.

The **bisection method** in mathematics is a root-finding method that repeatedly bisects an interval and then selects a subinterval in which a root must lie for further processing. It is a very simple and robust method, but it is also relatively slow. Because of this, it is often used to obtain a rough approximation to a solution which is then used as a starting point for more rapidly converging methods. [1] The method is also called the **interval halving** method, [2] the **binary search method**, [3] or the **dichotomy method**, [4]





Afew steps of the bisection method applied over the starting range [a₁;b₁]. The bigger red dot is the root of the function.

The method [edit]

The method is applicable for numerically solving the equation f(x) = 0 for the real variable x, where f is a continuous function defined on an interval [a, b] and where f(a) and f(b) have opposite signs. In this case a and b are said to bracket a root since, by the intermediate value theorem, the continuous function f must have at least one root in the interval (a, b).

At each step the method divides the interval in two by computing the midpoint c = (a+b) / 2 of the interval and the value of the function f(c) at that point. Unless c is itself a root (which is very unlikely, but possible) there are now only two possibilities: either f(a) and f(c) have opposite signs and bracket a root, or f(c) and f(b) have opposite signs and bracket a root. [5] The method selects the subinterval that is guaranteed to be a bracket as the new interval to be used in the next step. In this way an interval that contains a zero of f is reduced in width by 50% at each step. The process is continued until the interval is sufficiently small.

Explicitly, if f(a) and f(c) have opposite signs, then the method sets c as the new value for b, and if f(b) and f(c) have opposite signs then the method sets c as the new a. (If f(c)=0 then c may be taken as the solution and the process stops.) In both cases, the new f(a) and f(b) have opposite signs, so the method is applicable to this smaller interval.^[6]

Iteration tasks [edit]

The input for the method is a continuous function f, an interval [a, b], and the function values f(a) and f(b). The function values are of opposite sign (there is at least one zero crossing within the interval). Each iteration performs these steps:

- 1. Calculate c, the midpoint of the interval, c = 0.5 * (a + b).
- 2. Calculate the function value at the midpoint, f(c).
- 3. If convergence is satisfactory (that is, a c is sufficiently small, or f(c) is sufficiently small), return c and stop iterating.
- 4. Examine the sign of f(c) and replace either (a, f(a)) or (b, f(b)) with (c, f(c)) so that there is a zero crossing within the new interval.

When implementing the method on a computer, there can be problems with finite precision, so there are often additional convergence tests or limits to the number of iterations. Although f is continuous, finite precision may preclude a function value ever being zero. For $f(x) = x - \pi$, there will never be a finite representation of x that gives zero. Floating point representations also have limited precision, so at some point the midpoint of [a, b] will be either a or b.

Algorithm [edit]

The method may be written in pseudocode as follows:[7]

```
INPUT: Function f, endpoint values a, b, tolerance TOL, maximum iterations NMAX CONDITIONS: a < b, either f(a) < 0 and f(b) > 0 or f(a) > 0 and f(b) < 0 OUTPUT: value which differs from a root of f(x)=0 by less than TOL

N \leftarrow 1

While N \le NMAX \# limit iterations to prevent infinite loop <math>c \leftarrow (a+b)/2 \# new \ midpoint

If f(c) = 0 or (b-a)/2 < TOL then \# solution found Output (c)

Stop

EndIf N \leftarrow N + 1 \# increment \ step \ counter

If sign(f(c)) = sign(f(a)) then a \leftarrow c else b \leftarrow c \# new \ interval

EndWhile Output ("Method failed.") \# max \ number \ of \ steps \ exceeded
```

Example: Finding the root of a polynomial [edit]

Suppose that the bisection method is used to find a root of the polynomial

$$f(x) = x^3 - x - 2$$
.

First, two numbers a and b have to be found such that f(a) and f(b) have opposite signs. For the above function, a=1 and b=2 satisfy this criterion, as

$$f(1) = (1)^3 - (1) - 2 = -2$$

and

$$f(2) = (2)^3 - (2) - 2 = +4$$
.

Because the function is continuous, there must be a root within the interval [1, 2].

In the first iteration, the end points of the interval which brackets the root are $a_1=1$ and $b_1=2$, so the midpoint is

$$c_1 = \frac{2+1}{2} = 1.5$$

The function value at the midpoint is $f(c_1)=(1.5)^3-(1.5)-2=-0.125$. Because $f(c_1)$ is negative, a=1 is replaced with a=1.5 for the next iteration to ensure that f(a) and f(b) have opposite signs. As this continues, the interval between a and b will become increasingly smaller, converging on the root of the function. See this happen in the table below.

Iteration	a_n	b_n	c_n	$f(c_n)$
1	1	2	1.5	-0.125
2	1.5	2	1.75	1.6093750
3	1.5	1.75	1.625	0.6660156
4	1.5	1.625	1.5625	0.2521973
5	1.5	1.5625	1.5312500	0.0591125
6	1.5	1.5312500	1.5156250	-0.0340538
7	1.5156250	1.5312500	1.5234375	0.0122504
8	1.5156250	1.5234375	1.5195313	-0.0109712
9	1.5195313	1.5234375	1.5214844	0.0006222
10	1.5195313	1.5214844	1.5205078	-0.0051789

11	1.5205078	1.5214844	1.5209961	-0.0022794
12	1.5209961	1.5214844	1.5212402	-0.0008289
13	1.5212402	1.5214844	1.5213623	-0.0001034
14	1.5213623	1.5214844	1.5214233	0.0002594
15	1.5213623	1.5214233	1.5213928	0.0000780

After 13 iterations, it becomes apparent that there is a convergence to about 1.521: a root for the polynomial.

Analysis [edit]

The method is guaranteed to converge to a root of f if f is a continuous function on the interval [a, b] and f(a) and f(b) have opposite signs. The absolute error is halved at each step so the method converges linearly, which is comparatively slow.

Specifically, if $c_1 = (a+b)/2$ is the midpoint of the initial interval, and c_n is the midpoint of the interval in the nth step, then the difference between c_n and a solution c is bounded by^[8]

$$|c_n - c| \le \frac{|b - a|}{2^n}.$$

This formula can be used to determine in advance the number of iterations that the bisection method would need to converge to a root to within a certain tolerance. The number of iterations needed, n, to achieve a given

need to converge to a root to within a certain tolerance. The number of iterater error (or tolerance),
$$\epsilon$$
, is given by: $n = \log_2\left(\frac{\epsilon_0}{\epsilon}\right) = \frac{\log \epsilon_0 - \log \epsilon}{\log 2},$

where $\epsilon_0 = \text{initial bracket size} = b - a$.

Therefore, the linear convergence is expressed by $\epsilon_{n+1} = {\rm constant} \times \epsilon_n^m, \ m=1.$

See also [edit]

- Secant method
- Newton's method
- · Root-finding algorithm
- Binary search algorithm
- Lehmer-Schur algorithm, generalization of the bisection method in the complex plane
- Nested intervals
- Brent's method

References [edit]

- 1. ^ Burden & Faires 1985, p. 31
- 2. ^ http://siber.cankaya.edu.tr/NumericalComputations/ceng375/node32.html ☑
- 3. ^ Burden & Faires 1985, p. 28
- 4. ^ Encyclopedia of Mathematics ₺
- If the function has the same sign at the endpoints of an interval, the endpoints may or may not bracket roots of the function.
- 6. A Burden & Faires 1985, p. 28 for section
- Burden & Faires 1985, p. 29. This version recomputes the function values at each iteration rather than carrying them to the next iterations.
- 8. A Burden & Faires 1985, p. 31, Theorem 2.1

• Burden, Richard L.; Faires, J. Douglas (1985), "2.1 The Bisection Algorithm", *Numerical Analysis* (3rd ed.), PWS Publishers, ISBN 0-87150-857-5

Further reading [edit]

- Corliss, George (1977), "Which root does the bisection algorithm find?", *SIAM Review* **19** (2): 325–327, doi:10.1137/1019044 & ISSN 1095-7200 &
- Kaw, Autar; Kalu, Egwu (2008), Numerical Methods with Applications № (1st ed.)

External links [edit]

- Weisstein, Eric W., "Bisection" ₺, MathWorld.
- Bisection Method ₺ Notes, PPT, Mathcad, Maple, Matlab, Mathematica from Holistic Numerical Methods Institute ₺



Wikiversity has learning materials about *The* bisection method



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