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# Cristian's algorithm

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**Cristian's Algorithm** (introduced by Flaviu Cristian in 1989)<sup>[1]</sup> is a method for [clock synchronization](#) which can be used in many fields of distributive computer science but is primarily used in low-latency [intranets](#). Cristian observed that this simple algorithm is probabilistic, in that it only achieves synchronization if the [round-trip time](#) (RTT) of the request is short compared to required accuracy. It also suffers in implementations using a single server, making it unsuitable for many distributive applications where redundancy may be crucial.

## The algorithm [\[edit\]](#)

**Cristian's Algorithm** works between a process P, and a time server S — connected to a source of [UTC](#) (Coordinated Universal Time). Put simply:

1. P requests the time from S
2. After receiving the request from P, S prepares a response and appends the time T from its own clock.
3. P then sets its time to be  $T + \text{RTT}/2$

This method assumes that the RTT is split equally between request and response, which may not always be the case but is a reasonable assumption on a LAN connection.

Further accuracy can be gained by making multiple requests to S and using the response with the shortest RTT.

We can estimate the accuracy of the system as follows. Let *min* be the minimum time to transmit a message one-way. The earliest point at which S could have placed the time T, was *min* after P sent its request.

Therefore, the time at S, when the message is received by P, is in the range  $(T + \text{min})$  to  $(T + \text{RTT} - \text{min})$ . The width of this range is  $(\text{RTT} - 2*\text{min})$ . This gives an accuracy of  $(\text{RTT}/2 - \text{min})$ .

## References [\[edit\]](#)

- ↑ Cristian, F. (1989), "Probabilistic clock synchronization" [↗](#), *Distributed Computing* (Springer) **3** (3): 146–158, doi:10.1007/BF01784024 [↗](#)

## See also [\[edit\]](#)

- [Allan variance](#)
- [Clock synchronization](#)
- [International Atomic Time](#)
- [ntpd](#), [OpenNTPD](#) and [Ntupdate](#)
- [NTP pool](#), a collection of worldwide computers that provide a highly accurate time via the Network Time Protocol
- [NTP server misuse and abuse](#)
- [Synchronization](#)
- [Time server](#)

Other time synchronization protocols:

- [Berkeley algorithm](#)
- [DAYTIME protocol](#), older time synchronization protocol using TCP or UDP port 13
- [ICMP Timestamp](#) and [ICMP Timestamp Reply](#), older time synchronization protocol using ICMP
- [Precision Time Protocol](#)
- [TIME protocol](#), older time synchronization protocol using TCP or UDP port 37

Categories: [Distributed algorithms](#) | [Synchronization](#)



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