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Ricart-Agrawala algorithm

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The **Ricart-Agrawala Algorithm** is an algorithm for mutual exclusion on a distributed system. This algorithm is an extension and optimization of Lamport's Distributed Mutual Exclusion Algorithm, by removing the need for **Telegise** messages. It was developed by Glenn Ricart and Ashok Agrawala.

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Algorithm [edit]

Terminology [edit]

- A site is any computing device which is running the Ricart-Agrawala Algorithm
- The requesting site is the site which is requesting entry into the critical section.
- The receiving site is every other site which is receiving the request from the requesting site.

Algorithm [edit]

Requesting Site

 Sends a message to all sites. This message includes the site's name, and the current timestamp of the system according to its logical clock (which is assumed to be synchronized with the other sites)

Receiving Site

- Upon reception of a request message, immediately send a timestamped reply message if and only if:
 - the receiving process is not currently interested in the critical section OR
 - the receiving process has a lower priority (usually this means having a later timestamp)
- Otherwise, the receiving process will defer the reply message. This means that a reply will be sent only after the receiving process has finished using the critical section itself.

Critical Section:

- · Requesting site enters its critical section only after receiving all reply messages.
- Upon exiting the critical section, the site sends all deferred reply messages.

Performance [edit]

- Number of network messages; 2*(N-1)
- Synchronization Delays: One message propagation delay

Common Optimizations [edit]

Once site P_i has received a reply message from site P_j , site P_i may enter the critical section multiple times without receiving permission from P_j on subsequent attempts up to the moment when P_i has sent a reply message to P_i . This is called Roucairol-Carvalho optimization or Roucairol-Carvalho algorithm.

Problems [edit]

One of the problems in this algorithm is failure of a node. In such a situation a process may starve forever. This problem can be solved by detecting failure of nodes after some timeout.

See also [edit]

- Lamport's Bakery Algorithm
- Lamport's Distributed Mutual Exclusion Algorithm
- Maekawa's Algorithm
- Suzuki-Kasami algorithm
- Raymond's Algorithm
- Naimi-Trehel's Algorithm

References [edit]

 Maekawa, M.,Oldehoeft, A.,Oldehoeft, R.(1987). Operating Systems: Advanced Concept.Benjamin/Cummings Publishing Company, Inc.

Categories: Distributed algorithms

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