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Gosper's algorithm

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In mathematics, **Gosper's algorithm** is a procedure for finding sums of hypergeometric terms that are themselves hypergeometric terms. That is: suppose we have a(1) + ... + a(n) = S(n) - S(0), where S(n) is a hypergeometric term (i.e., S(n + 1)/S(n) is a rational function of n); then necessarily a(n) is itself a hypergeometric term, and given the formula for a(n) Gosper's algorithm finds that for S(n).

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Outline of the algorithm [edit]

Step 1: Find a polynomial p such that, writing b(n) = a(n)/p(n), the ratio b(n)/b(n-1) has the form q(n)/r(n) where q and r are polynomials and no q(n) has a nontrivial factor with r(n+j) for j=0,1,2,... (This is always possible, whether or not the series is summable in closed form.)

Step 2: Find a polynomial f such that S(n) = q(n+1)/p(n) f(n) a(n). If the series is summable in closed form then clearly a rational function f with this property exists; in fact it must always be a polynomial, and an upper bound on its degree can be found. Determining f (or finding that there is no such f) is then a matter of solving a system of linear equations.

Relationship to Wilf-Zeilberger pairs [edit]

Gosper's algorithm can be used to discover Wilf–Zeilberger pairs, where they exist. Suppose that F(n+1,k) - F(n,k) = G(n,k+1) - G(n,k) where F is known but G is not. Then feed a(k) := F(n+1,k) - F(n,k) into Gosper's algorithm. (Treat this as a function of k whose coefficients happen to be functions of n rather than numbers; everything in the algorithm works in this setting.) If it successfully finds S(k) with S(k) - S(k-1) = a(k), then we are done: this is the required G. If not, there is no such G.

Definite versus indefinite summation [edit]

Gosper's algorithm finds (where possible) a hypergeometric closed form for the *indefinite* sum of hypergeometric terms. It can happen that there is no such closed form, but that the sum over *all* n, or some particular set of values of n, has a closed form. This question is only meaningful when the coefficients are themselves functions of some other variable. So, suppose a(n,k) is a hypergeometric term in both n and k: that is, a(n, k)/a(n - 1, k) and a(n, k)/a(n, k - 1) are rational functions of n and k. Then Zeilberger's algorithm and Petkovšek's algorithm may be used to find closed forms for the sum over k of a(n, k).

History [edit]

Bill Gosper discovered this algorithm in the 1970s while working on the Macsyma computer algebra system at SAIL and MIT.

Further reading [edit]

- Marko Petkovšek, Herbert Wilf and Doron Zeilberger, A = B, AK Peters 1996, ISBN 1-56881-063-6. Full text online.[1] 🖓
- Gosper's 1977 article in PNAS is reporting on the algorithm.

Categories: Computer algebra | Hypergeometric functions

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