Dynamic Programming Scoring matrices BLAST

Section 3

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Outline

- Recursion
- More Dynamic Programming
- Scoring Matrices
- BLAST

Recursion

General Idea: Solve a problem by solving related subproblems and combining the results.

Advantage: Sub-problems are easier to solve and their total run time can be much less than original problem.

Example: Merge Sort – To sort a large list, split the list in half, sort each half, and merge the two sorted lists.

Example: Factorial(x) $x! = x^*(x-1)^*(x-2)^*...^*1$ $5! = 5^*4^*3^*2^*1 = 120$

Recursion - Factorial Iterative Algorithm: $\begin{cases} \$x = 5; \\ \text{for } (\$n = \$x - 1; \$n > 0; \$n -) \{ \\ \$x * = \$n; \} \end{cases}$ Recursive Algorithm: $\begin{cases} \$x = \text{factorial}(5); \\ \text{sub factorial } \{ \\ \text{my}(\$n) = \$ _[0]; \\ \text{if } (\$n <= 1) \{ \\ \text{return } 1; \\ \text{else } \{ \\ \text{return } \$n \text{ factorial}(\$n - 1); \\ \} \end{cases}$

Dynamic Programming

General Idea: Solve a problem by solving related sub-problems and combining the results.

Key Point: In recursive algorithms (e.g., Merge Sort), the sub-problems are independent. In dynamic programming the sub-problems share sub-subproblems.

Implementation: Solve each sub-sub-problem only once and store the answers in an array.

Dynamic Programming - Step 1

Characterize the structure of the optimal solution.

Example: Sequence Alignment

Three choices at each step:

- 1. Align a base in X with a base in Y
- 2. Align a base in X with a gap
- 3. Align a base in Y with a gap

Goal:

Maximize the total score

Dynamic Programming - Step 2

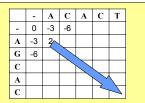
Recursively define the value of an optimal solution.

Example: Needleman-Wunsch (Global Alignment)

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F(i,j) = \max \{ \\ F(i-1, j-1) + s(x_i, y_j), \\ F(i-1, j) - d, \\ F(i, j-1) - d \\ \}
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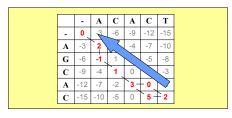
Dynamic Programming – Step 3

Compute the optimal solution, starting from the easiest sub-sub-problems and saving the results.



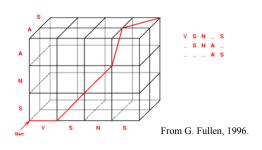
Dynamic Programming - Step 4

Construct an optimal solution from the computed information. (Perform a trace-back.)



Multiple Sequence Alignment

Dynamic programming in many dimensions!!!



Dynamic Programming Summary

Advantage: Guaranteed to find the mathematically optimal solution (highest scoring alignment). Much faster than trying all possible alignments.

Disadvantage: For large problems (long sequences or many sequences), dynamic programming can be very slow and require lots of computer memory.

Scoring Matrices

Definition: A table (matrix) of values that describe the probability of a residue (amino acid or base) pair occurring in an alignment of related sequences.

Why a matrix? Evolution tends to favor different sets of substitutions. Thus, a single probability does not accurately describe all residue pairs

Example: PAM

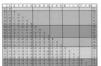
PAM = Percent Accepted Mutation

- Ranks a.a. substitutions on how well they're "accepted" by natural selection, without changing the function of the protein.
- Derived from global alignments of closely related sequences (85% identical).
- PAM1 models the mutation rate in 1 "PAM Unit" of evolutionary time.
- PAMn = (PAM1)ⁿ = mutation rate in n PAM Units
- Examples: PAM40, PAM100

Example: BLOSUM

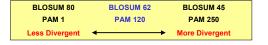
BLOSUM = Blocks Amino Acid Substitution Matrices

- Derived from large data sets of local, ungapped alignments of biochemically related sequences.
- · Each matrix is derived from different data sets.
- BLOSUMn means no pair of sequences in the data set used to derive the matrix had more than n% similarity.
- BLOSUM62 indicates the substitution frequencies for sequences that are up to 62% similar



PAM vs BLOSUM

- · BLOSUM better than PAM for local similarity
- BLOSUM better in similarity searches in databases
- PAM incorporates an evolutionary model representing species divergence
- BLOSUM tolerant of hydrophobic changes and of cystein and tryptophan mismatches
- PAM tolerant of change to/from hydrophilic a.a.



BLAST

BLAST = Basic Local Alignment Search Tool

- At least 50 times faster than dynamic programming
- Not guaranteed to find the optimal solution, but...
- Individual alignments need not be perfect, you can always fine-tune later.
 - Most sequences will be completely unrelated to the query.
- Will always find related sequences in a database that meet some similarity criteria.
- BLAST is a heuristic algorithm (usually works well)
- Always translate DNA into protein sequences (if possible) before running BLAST.

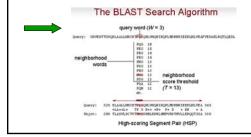
BLAST

Assumption: True alignments are very likely to contain a short stretch of identities, or very high scoring matches.



BLAST: K-mer Words

- Create a list of all "words" of length K (default is 3 amino acids or 11 nucleotides)
- Example: TPQGQR → TPQ, PQG, QGQ, GQR



BLAST: K-mer Hashing

- For each k-mer, list all locations where it occurs in the database.
- Match the query to k-mers (or "neighborhood words" above a threshold) in the list. This forms seeds.



BLAST: Extending Seeds

- · Connect "close" seeds into single long sequences.
- · Extend seeds until the cumulative score drops.
- High-scoring Segment Pair = locally maximal segment.



BLAST: Evaluating HSPs

- · Bit Score: independent of the scoring matrix used.
- E-value: "Expected number of High-Scoring Segment Pairs". The smaller the number, the better the alignment.
- E-values of 0.1 or 0.05 are typically used as cutoffs.



Next Week

- Microarrays
- Sequencing

Acknowledgement / References

This handout includes material written by Suzanne Komili, Yonatan Grad, Doug Selinger, and Zhou Zhu.

Mount, Bioinformatics – Sequence and Genome Analysis; Durbin et al., Biological Sequence Analysis; Cormen et al., Introduction to Algorithms

http://www.ncbi.nlm.nih.gov/BLAST