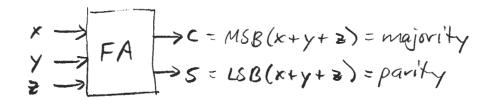
6,896 2/9/04 L2,1

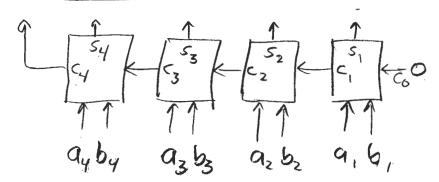
((Sort with Oligh) procs in linear array)) (Firing squad - Oli) size state) Addition

Basic component: Full adder - combinational

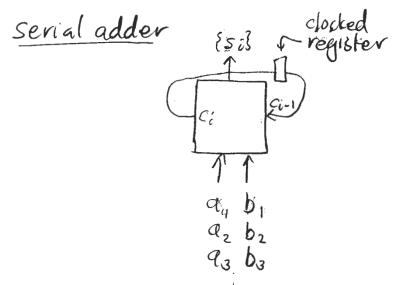


Problem: Add Z N-bit numbers

Ripple-carry adder



N-bit #'s => O(N) time, O(N) HW, combinational



O(N) time, O(1) hardware, sequential (clocked)

j

Fast addition

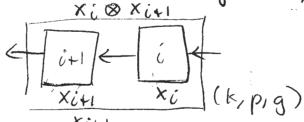
Idea: carries are the hard part.

Know carries => compute sum in O(1) time
How? Away of full adders. «Show on ripple-carry adder»)

Classify stages:

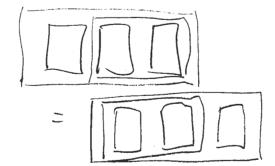
Carry into stage = { o otherwise

When do 2 consecutive stages kill, prop, gen?



	\otimes	K	P	9	
	k	k	k	9	
׿	Р	k	P	9	
	9	k	9	9	

Associative!



Theorem. Let $x_i = carry status of stage i$, where $x_0 = k$. Define $y_i = x_0 \otimes x_1 \otimes \cdots \otimes x_i$.

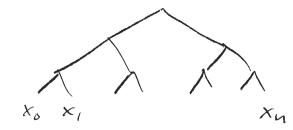
Then $y_i = k \Rightarrow c_i = 0$ $y_i = g \Rightarrow c_i = 1$ $y_i = p$ does not occur.

Proof. Induction on i. \

Log-time circuit:

 $y_0 = x_0$ $y_1 = x_0 \otimes x_1$ $y_2 = x_0 \otimes x_1 \otimes x_2$ $y_N = x_0 \otimes x_1 \otimes x_2$ $y_N = x_0 \otimes x_1 \otimes \dots \otimes x_N$

Use tree for each calculation:



Use tree to broadcast inputs (bounded-degree network):



Time = O(lgN), HW = O(N2).

Carry-lookahead addition

6,896 2/9/04 LZ, 4

OllgN) time, O(N) HW.

"Parallel prefix"

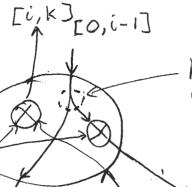
Let [i, j] denote xi & xi+1 & ... & xj

Lemma. [i,j] & [j+1, k] = [i,k] \

xi=[i,i]

Yi = [0, i]

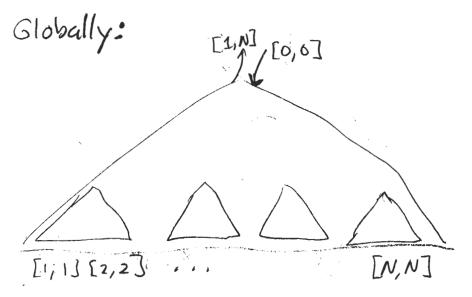
Build tree:



insert identity elem, or not bounded degree.

[i,j] [0,i-1]

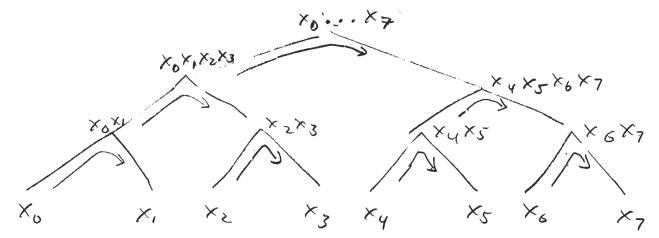
[jt]k] [0,j]



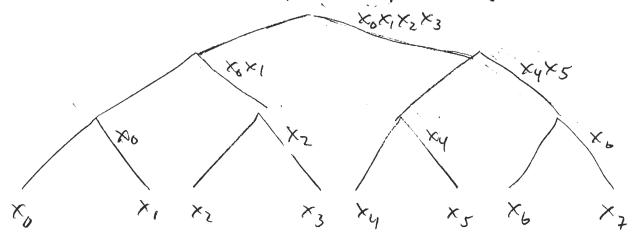
Left child values are passed up.

% .

Similar method:



Left child values are passed up and right



Postscript Kill, propagate, generate first used in Harrand relay calculator circa mid-1940's.

O(1)-time addition (in their model).