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Wookyung Jeong, Junghyun Lee, Sungmoon Jung

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4	Tree	8	// #pragma GCC optimize("03")	
	4.1 Heavy-Light Decomposition	8	// #pragma GCC optimize("Ofast")	
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5	Network Flow	9	<pre>#include <bits stdc++.h=""></bits></pre>	
	5.1 Dinic's Algorithm	9	#include <cassert></cassert>	
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	5.3 MCMF	10	#define 11 long long	
		10	#define ull unsigned long long	
	String	10	#define ld long double	
	6.1 Rabin-Karp Algorithm		<pre>#define pii pair<int, int=""></int,></pre>	
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	6.4 Aho-Corasick		#define fr first	
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	6.7 Z Algorithm		#uerine 32(A) (int)(A).312e()	
			const double EPS = 1e-9;	
7	Geometry	15	const int INF = 1e9 + 7;	
	7.1 Convex Hull	15	const int MOD = 1e9 + 7;	
	7.2 Rotating Callipers		const int dy[] = { 0, 0, 1, -1, 1, 1, -1, -1 };	
	7.3 Ray Casting	16	const int $dx[] = \{ 1, -1, 0, 0, 1, -1, 1, -1 \};$	
	7.4 Sort by Angular			
	7.5 Bulldozer Trick		int main() {	
		18	<pre>// #ifndef ONLINE_JUDGE // freopen("Write your absolute/relative path of input.txt", "r", stdin);</pre>	
	Math 8.1 Sieve	18 18	<pre>// freopen("Write your absolute/relative path of output.txt", "w", stdout); // #endif</pre>	
	8.2 Euclidean Algorithms	19	// #endii	
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2 Data Structure

```
2.1 Segment Tree
// 2. Iterative Segment Tree
```

```
const int MAXN = 1010101; // limit for array size
struct Seg { // 0-indexed
 int n; // array size
 11 t[2 * MAXN];
  void build(int N) {
   n = N:
   for (int i = 0; i < n; i++) cin >> t[n + i];
    for (int i = n - 1; i >= 1; i--) t[i] = t[i << 1] + t[i << 1 | 1];
 void modify(int p, ll value) { // set value at position p
   for (t[p += n] = value; p > 1; p >>= 1) t[p >> 1] = t[p] + t[p ^ 1];
 11 query(int 1, int r) { // sum on interval [1, r)
   11 \text{ ret} = 0;
   for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
     if (1 & 1) ret += t[1++];
     if (r & 1) ret += t[--r]:
   }
    return ret;
 }
}seg;
// 5. Persistent Segment Tree
// TIME COMPLEXITY: O(n) for initialize PST, O(logn) for each query.
// SPACE COMPLEXITY: O(nlogm).
struct PST { // 1-indexed
    int flag; // array size
    struct Node { int 1, r; 11 val; };
    vector<Node> t;
    vector<int> root;
    void addNode() { t.push_back({ -1, -1, 0 }); }
    void build(int 1, int r, int n) {
        assert(0 \le n \&\& n \le sz(t)):
        if (l == r) { t[n].val = a[l]; return; }
        addNode():
        t[n].1 = sz(t) - 1;
        addNode();
        t[n].r = sz(t) - 1;
        int mid = (1 + r) >> 1:
        build(1, mid, t[n].1);
        build(mid + 1, r, t[n].r);
        t[n].val = t[t[n].1].val + t[t[n].r].val;
    void build(int Flag) {
        addNode():
        root.push_back(sz(t) - 1);
        flag = Flag;
        build(1, flag, root[0]);
    void modify(int p, ll val, int l, int r, int n1, int n2) {
        assert(0 \le n1 \&\& n1 \le sz(t));
        assert(0 \le n2 \&\& n2 \le sz(t));
        if (p < 1 || r < p) { t[n2] = t[n1]; return; }
```

```
if (1 == r) { t[n2].val = val; return; }
        int mid = (1 + r) >> 1;
        if (p <= mid) {</pre>
           t[n2].r = t[n1].r;
            addNode();
           t[n2].1 = sz(t) - 1:
            modify(p, val, 1, mid, t[n1].1, t[n2].1);
       }
        else {
            t[n2].1 = t[n1].1;
           addNode();
           t[n2].r = sz(t) - 1:
            modify(p, val, mid + 1, r, t[n1].r, t[n2].r);
        t[n2].val = t[t[n2].1].val + t[t[n2].r].val;
   }
   void modify(int p, ll val) {
        addNode();
       root.push back(sz(t) - 1):
        modify(p, val, 1, flag, root[sz(root) - 2], root[sz(root) - 1]);
   11 query(int 1, int r, int n, int nl, int nr) {
        assert(0 \le n \&\& n \le sz(t));
        if (r < nl || nr < 1) return 0;
        if (1 <= nl && nr <= r) return t[n].val:
        int mid = (nl + nr) >> 1;
        return query(1, r, t[n].1, n1, mid) + query(1, r, t[n].r, mid + 1, nr);
   11 query(int 1, int r, int n) {
        assert(n < sz(root));</pre>
        return query(1, r, root[n], 1, flag);
   }
}pst;
2.2 Fenwick Tree
// 1. Fenwick Tree
struct Fenwick { // 0-indexed
 int flag, cnt; // array size
 vector<ll> arr, t;
 void build(int n) {
   for (flag = 1; flag < n; flag <<= 1, cnt++);
   arr.resize(flag);
   t.resize(flag);
   for (int i = 0; i < n; i++) cin >> arr[i];
   for (int i = 0; i < n; i++) {
     t[i] += arr[i];
     if (i | (i + 1) < flag) t[i | (i + 1)] += t[i];
   }
 void add(int p, ll value) { // add value at position p
   arr[p] += value;
   while (p < flag) {
     t[p] += value;
     p | = p + 1;
   }
 void modify(int p, ll value) { // set value at position p
```

```
add(p, value - arr[p]);
 11 query(int x) {
   ll ret = 0:
    while (x \ge 0) ret += t[x], x = (x & (x + 1)) - 1;
    return ret;
 11 query(int 1, int r) {
    return query(r) - (1 ? query(1 - 1) : 0);
  int kth(int k) { // find the kth smallest number (1-indexed)
    assert(t.back() >= k):
    int l = 0, r = arr.size():
    for (int i = 0; i <= cnt; i++) {
     int mid = (1 + r) >> 1:
      11 val = mid ? t[mid - 1] : t.back();
      if (val >= k) r = mid;
      else 1 = mid, k -= val;
    return 1:
}fw;
// 3. 2D Fenwick Tree
// INPUT: Given an 2D array of integers of size N * M.
// Can modify the value of the (x, y)th element.
// Can find the sum of elements from (sx, sy) to (ex, ey).
// OUTPUT: Given the query (1 sx sy ex ey), output the sum of elements from the interval (sx,
sv) to (ex. ev)
// TIME COMPLEXITY: O(N * M) for initialize fenwick tree, O(\log N * \log M) for each query.
struct Fenwick2D { // O-indexed
    int n, m, real_n, real_m;
    vector<vector<ll>>> arr. t:
    void build(int N, int M) {
        real_n = N, real_m = M;
        n = m = 1:
        while (n < N) n <<= 1;
        while (m < M) m <<= 1:
        arr.resize(n, vector<11>(m));
        t.resize(n, vector<ll>(m)):
        for (int i = 0; i < real_n; i++) {</pre>
            for (int j = 0; j < real_m; j++) {
                cin >> arr[i][j];
            }
        }
        for (int i = 0; i < n; i++) {
            for (int j = 0; j < m; j++) {
                t[i][j] += arr[i][j];
                int ni = i | (i + 1), nj = j | (j + 1);
                if (ni < n) t[ni][j] += t[i][j];</pre>
                if (n; < m) t[i][n;] += t[i][;];</pre>
                if (ni < n && nj < m) t[ni][nj] -= t[i][j];
           }
```

```
void add(int x, int y, ll value) { // add value at position (x, y)
        assert(0 \le x \&\& x \le real_n \&\& 0 \le y \&\& y \le real_m);
        arr[x][y] += value;
        for (int i = x; i < n; i | = i + 1) {
            for (int j = y; j < m; j | = j + 1) {
                t[i][j] += value;
            }
        }
   }
    void modify(int x, int y, 11 value) { // set value at position (x, y)
        assert(0 \le x \&\& x \le real_n \&\& 0 \le y \&\& y \le real_m);
        add(x, y, value - arr[x][y]);
   }
   11 query(11 x, 11 y) {
        assert(0 <= x && x < real_n && 0 <= y && y < real_m);
        for (int i = x; i \ge 0; i = (i & (i + 1)) - 1) {
            for (int j = y; j >= 0; j = (j & (j + 1)) - 1) {
                ret += t[i][j];
            }
        }
        return ret;
   7
   11 query(ll sx, ll sy, ll ex, ll ey) {
        assert(0 <= sx && sx <= ex && ex < real_n);
        assert(0 <= sy && sy <= ey && ey < real_m);
        11 ret = query(ex, ey);
        if (sx) ret -= query(sx - 1, ey);
        if (sy) ret -= query(ex, sy - 1);
        if (sx \&\& sy) ret += query(sx - 1, sy - 1);
        return ret:
   }
}fw2d:
2.3 Li Chao Tree
// INPUT: Initially, a 2d plane in which no linear function exists is given.
// Two types of gueries are given.
// 1 a b : The linear function f(x) = ax + b is added.
// 2 x : Find the max(f(x)) among the linear functions given so far.
// OUTPUT: For each query 2 x, output the max(f(x)) among the linear functions given so far.
// TIME COMPLEXITY: O(glogg)
#define Line pair<11, 11>
const Line e = \{ 0, -1e18 \}:
struct LiChaoTree {
   11 f(Line 1, 11 x) { return 1.first * x + 1.second; }
    struct Node {
        ll xl, xr; int l, r;
        Line line;
   }:
    vector<Node> t;
    void build(ll xlb, ll xub) {
        t.push_back({ xlb, xub, -1, -1, e });
   }
   void insert(Line newLine, int n = 0) {
        11 x1 = t[n].x1, xr = t[n].xr;
        11 \text{ xmid} = (x1 + xr) >> 1;
```

}

// TIME COMPLEXITY: O(VE)

```
Line llow = t[n].line, lhigh = newLine;
        if (f(llow, xl) >= f(lhigh, xl)) swap(llow, lhigh);
        if (f(llow, xr) <= f(lhigh, xr)) {</pre>
            t[n].line = lhigh;
            return;
        }
        else if (f(llow, xmid) < f(lhigh, xmid)) {
            t[n].line = lhigh;
            if (t[n].r == -1) {
                t[n].r = sz(t);
                 t.push_back({ xmid + 1, xr, -1, -1, e });
            }
            insert(llow, t[n].r);
        }
        else if (f(llow, xmid) >= f(lhigh, xmid)) {
            t[n].line = llow;
            if (t[n].1 == -1) {
                t[n].1 = sz(t);
                 t.push_back({ xl, xmid, -1, -1, e });
            }
            insert(lhigh, t[n].1);
    }
    11 \text{ query}(11 \text{ x, int } n = 0) 
        if (n == -1) return e.second:
        11 xl = t[n].xl, xr = t[n].xr;
        11 \text{ xmid} = (x1 + xr) >> 1:
        11 \text{ ret} = f(t[n].line, x);
        if (x \le xmid) ret = max(ret, query(x, t[n].1));
        else ret = max(ret, query(x, t[n].r));
        return ret;
}lct;
int main() {
    lct.build(-1'000'000'000'00011, 1'000'000'000'00011);
    int q; cin >> q;
    while (q--) {
        int op; ll a, b;
        cin >> op >> a;
        if (op == 1) {
            cin >> b;
            lct.insert({ a, b });
        if (op == 2) cout << lct.query(a) << '\n';
3 Graph
3.1 Bellman-Ford, Floyd-Warshall
// 2. Bellman-Ford Algorithm
// INPUT: Given a directed graph with weighted(possibly negative) edges and no negative cycles.
Given a starting vertex.
// OUTPUT: Outputs the shortest distance from the starting vertex to all vertices.
```

```
struct wv {
 ll w: int v:
int n, m;
vector<wv> adj[101010];
vector<ll> upper(101010, (11)1e18);
int bellmanFord() {
 upper[1] = 0;
 int update = 1;
 for (int i = 0; i <= n; i++) {
    update = 0;
   for (int now = 1; now <= n; now++) {
      if (upper[now] == INF) continue;
     for (wv e : adj[now]) {
        int next = e.v:
        if (upper[next] > upper[now] + e.w) {
          upper[next] = upper[now] + e.w;
          update = 1;
     }
   }
    if (!update) break;
 return !update; // Returns false <=> The graph has a negative cycle.
// 3. Floyd-Warshall Algorithm
// INPUT: Given a directed graph with weighted(possibly negative) edges and no negative cycles.
// OUTPUT: Outputs the shortest distance from all vertices to all vertices.
// TIME COMPLEXITY: O(V^3)
int n. m:
ll adj[1010][1010];
void floyd() {
 for (int i = 0; i < 1010; i++) {
   for (int j = 0; j < 1010; j++) {
      adj[i][j] = (11)1e18;
   }
 }
 for (int i = 1; i <= n; i++) adj[i][i] = 0;
 for (int k = 1; k \le n; k++) {
   for (int u = 1; u <= n; u++) {
      for (int v = 1; v \le n; v++) {
        adj[u][v] = min(adj[u][v], adj[u][k] + adj[k][v]);
   }
 }
3.2 SCC, 2-SAT
// 1. SCC (Kosaraju's Algorithm)
// INPUT: Given a directed graph.
// OUTPUT: Decompose this graph into SCCs and print them in lexicographical order.
// TIME COMPLEXITY: O(V + E)
// BOJ 2150 AC Code
// https://www.acmicpc.net/problem/2150
#include <bits/stdc++.h>
```

```
using namespace std;
#define sz(x) (int)(x).size()
const int MAXV = 10101;
int n, m;
vector<int> adj[MAXV], radj[MAXV];
int in[MAXV], out[MAXV], num, p[2 * MAXV];
int vi[MAXV], cnt;
vector<vector<int>> scc;
void input() {
  cin >> n >> m:
  for (int i = 0; i < m; i++) {
    int u. v:
    cin >> u >> v;
    adj[u].push_back(v);
    radj[v].push_back(u);
}
void dfs(int v) {
  in[v] = ++num:
  for (auto& i : radj[v]) {
    if (!in[i]) dfs(i);
  out[v] = ++num;
  p[num] = v;
void flood(int v) {
  scc[cnt].push_back(v);
  vi[v] = cnt;
  for (auto& i : adj[v]) {
    if (!vi[i]) flood(i);
 }
}
void kosaraju() {
  for (int v = 1; v <= n; v++) {
    if (!in[v]) dfs(v);
  for (int v = 2 * n; v >= 1; v--) {
    if (!p[v]) continue;
    if (vi[p[v]]) continue;
    cnt++:
    scc.resize(cnt + 1);
    flood(p[v]);
void print() {
  for (auto& i : scc)
    sort(i.begin(), i.end());
  sort(scc.begin(), scc.end());
  cout << sz(scc) - 1 << '\n';
  for (int i = 1; i < sz(scc); i++) {
    auto& arr = scc[i]:
```

```
for (auto& j : arr) cout << j << ' ';
    cout << -1 << '\n':
}
int main() {
 cin.tie(NULL); cout.tie(NULL);
  ios_base::sync_with_stdio(false);
  input();
  kosaraju();
 print();
// 3. 2-SAT
// INPUT: A 2-CNF is given. 2-CNF is a boolean expression in the form (x \lor y) \land (\_ y \lor z) \land
(x \lor \neg z) \land (z \lor y).
// OUTPUT: Determine whether there exists a case where a given 2-CNF expression can be true.
(2-Satisfiability Problem)
// TIME COMPLEXITY: O(n + m) = O(n) (m = 2n)
// BOJ 11281 AC Code
// https://www.acmicpc.net/problem/11281
#include <bits/stdc++.h>
using namespace std;
#define pii pair<int, int>
#define fr first
#define sc second
const int MAXV = 20202;
int n, m;
int dfsn[MAXV], dCnt, sNum[MAXV], sCnt;
int finished[MAXV];
vector<int> adj[MAXV];
stack<int> stk:
pii p[MAXV];
int ans[MAXV / 2]:
inline int inv(int x) {
 // negative number -a indicates -a.
 return (x > 0) ? 2 * (x - 1) : 2 * (-x - 1) + 1;
void twoCnf(int a, int b) {
 // (a \vee b) iff (\nega \rightarrow b) iff (\negb \rightarrow a)
  adj[inv(-a)].push_back(inv(b));
  adj[inv(-b)].push_back(inv(a));
void input() {
 cin >> n >> m;
 for (int i = 0; i < m; i++) {
   int a, b;
   cin >> a >> b;
    twoCnf(a, b);
```

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```

```
}
int dfs(int now) {
  int ret = dfsn[now] = ++dCnt:
  stk.push(now);
  for (int next : adj[now]) {
    if (dfsn[next] == -1) ret = min(ret, dfs(next));
    else if (!finished[next]) ret = min(ret, dfsn[next]);
  if (ret >= dfsn[now]) {
    while (1) {
      int t = stk.top();
      stk.pop();
      sNum[t] = sCnt;
      finished[t] = 1:
      if (t == now) break;
    sCnt++;
  }
  return ret:
int isSatisfiable() {
  // determining satisfiability
  int isS = 1;
  for (int v = 0: v < 2 * n: v += 2) {
    // if x and \neg x is in same scc, then the proposition is not satisfiable
    if (sNum[v] == sNum[v + 1]) {
      isS = 0:
      break;
    }
  return isS;
void findValueOfEachVariable() {
  // order of scc is the reverse of the topological sort
  for (int v = 0; v < 2 * n; v++) {
    p[v] = \{ sNum[v], v \};
  sort(p, p + 2 * n);
  // determining true/false of each variable
  for (int i = 2 * n - 1; i \ge 0; i--) {
    int v = p[i].sc;
    if (ans[v / 2 + 1] == -1)
      ans[v / 2 + 1] = (v & 1) ? 1 : 0;
  for (int v = 1; v \le n; v++)
    cout << ans[v] << ' ';
int main() {
  cin.tie(NULL); cout.tie(NULL);
  ios_base::sync_with_stdio(false);
  memset(dfsn, -1, sizeof(dfsn));
  memset(ans, -1, sizeof(ans));
```

```
input();
  // finding scc
  for (int \bar{v} = 0; v < 2 * n; v++)
    if (dfsn[v] == -1) dfs(v):
  if (isSatisfiable()) {
    cout << 1 << '\n';
    findValueOfEachVariable();
  else cout << 0;</pre>
 return 0:
3.3 BCC
// A Biconnected Component (BCC) is a subset of vertices in an undirected graph that satisfies
the following conditions:
// (1) If you delete any vertex from a subset, the remaining vertices are connected to each
// (2) Adding other vertices to this subset does not satisfy (1). (This is the largest set that
satisfies (1))
// TIME COMPLEXITY: O(V + E)
// A vertex at which the graph is divided into two or more components when the vertex is removed
is called a 'articulation point'.
// After decomposing the graph into BCCs, vertices belonging to two or more BCCs are
articulation point.
// Similarly, a edge at which the graph is divided into two or more components when the edge is
removed is called a 'articulation edge'.
// For all tree edges on a dfs spanning tree, if tmp > dfsn[now], the edge { now, next } is a
articulation edge.
#include <bits/stdc++.h>
using namespace std;
#define pii pair<int, int>
const int MAXV = 101010;
int n, m;
vector<int> adj[MAXV];
vector<vector<pii>>> bcc;
set<int> aPoint;
set<pii> aEdge;
void input() {
 cin >> n >> m;
 for (int i = 0; i < m; i++) {
   int u, v;
   cin >> u >> v:
    adj[u].push_back(v);
    adj[v].push_back(u);
 }
}
int dfsn[MAXV], dCnt;
stack<pii> stk;
int dfs(int now, int prv) {
```

```
int ret = dfsn[now] = ++dCnt;
  int childCnt = 0;
  for (int next : adj[now]) {
    if (next == prv) continue;
    // If an edge { now, next } has not yet been visited, puts an edge on the stack.
    if (dfsn[now] > dfsn[next]) stk.push({ now, next });
    // Back edge
    if (dfsn[next] != -1) ret = min(ret, dfsn[next]);
    // Tree edge
    else {
      childCnt++;
      int tmp = dfs(next, now);
      ret = min(ret, tmp);
      if (prv != -1 && tmp >= dfsn[now])
        aPoint.insert(now);
      if (tmp > dfsn[now])
        aEdge.insert({ min(now, next), max(now, next) }):
      // If next cannot go to ancestor node of now, find BCC
      if (tmp >= dfsn[now]) {
        vector<pii> nowBCC;
        while (true) {
          pii t = stk.top();
          stk.pop();
          nowBCC.push_back(t);
          if (t == make_pair(now, next)) break;
        bcc.push_back(nowBCC);
  if (prv == -1 && childCnt > 1)
    aPoint.insert(now);
  return ret;
}
void getBCC() {
 memset(dfsn, -1, sizeof(dfsn));
 for (int v = 1: v \le n: v++)
    if (dfsn[v] == -1) dfs(v, -1);
}
int main() {
  cin.tie(NULL): cout.tie(NULL):
 ios_base::sync_with_stdio(false);
 input();
  getBCC();
```

3.4 Euler Circuit

```
// Hierholzer's Algorithm
// INPUT: Given a undirected graph.
// OUTPUT: Print the path of the Euler circuit of the graph.
// Euler Path is a path in a finite graph that visits every edge exactly once.
// Similarly, an Euler Circuit is an Euler Path that starts and ends on the same vertex.
// TIME COMPLEXITY: O(VE)
// BOJ 1199 AC Code
// https://www.acmicpc.net/problem/1199
#include <bits/stdc++.h>
using namespace std;
const int MAXV = 1010:
int n, adj[MAXV][MAXV], nxt[MAXV];
vector<int> eulerCircult:
void input() {
 cin >> n:
 for (int i = 1; i <= n; i++) {
   for (int j = 1; j <= n; j++) {
      cin >> adj[i][j];
   }
 }
}
int doesEulerCircuitExist() {
 // If the degree of all nodes in the graph is even, then an euler circuit exists.
 // Otherwise, the euler circuit does not exist.
 // We can do similar way to determine the existence of euler path.
 // If only two vertices have odd degree, than an eular path exists. Otherwise, the euler path
 does not exist.
 for (int i = 1; i <= n; i++) {
   int deg = 0;
   for (int j = 1; j <= n; j++) {
      deg += adj[i][j];
    if (deg & 1) return 0;
 return 1;
void dfs(int now) {
 for (int& x = nxt[now]; x \le n; x++) {
   while (x \le n \&\& adj[now][x]) {
      adj[now][x]--;
      adj[x][now]--;
      dfs(x):
   }
  eulerCircult.push_back(now);
int main() {
 cin.tie(NULL); cout.tie(NULL);
 ios_base::sync_with_stdio(false);
```

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```
input();
  if (!doesEulerCircuitExist()) {
    cout << -1:
    return 0;
  for (int i = 1; i <= n; i++) nxt[i] = 1;
  for (auto i : eulerCircult)
    cout << i << ' ';
}
4 Tree
4.1 Heavy-Light Decomposition
// 2. HLD with Segment Tree
#include <bits/stdc++.h>
using namespace std;
#define 11 long long
const int MAXV = 202020:
int flag; // array size
struct Seg { // 1-indexed
  vector<ll> t;
  void build(int n) {
    for (flag = 1; flag < n; flag <<= 1);
    t.resize(2 * flag);
  void modify(int p, ll value) { // set value at position p
    for (t[p += flag - 1] = value; p > 1; p >>= 1) t[p >> 1] = t[p] + t[p ^ 1];
  ll query(int l, int r, int n = 1, int nl = 1, int nr = flag) { // sum on interval [l, r]
    if (r < nl || nr < l) return 0;
    if (1 <= n1 && nr <= r) return t[n];
    int mid = (nl + nr) / 2;
    return query(1, r, n << 1, n1, mid) + query(1, r, n << 1 | 1, mid + 1, nr);
}seg;
vector<int> adj[MAXV], g[MAXV];
int siz[MAXV], dep[MAXV], par[MAXV];
int top[MAXV], in[MAXV], out[MAXV], pv;
void dfs(int v, int prv) {
  for (auto& i : adj[v]) {
    if (i == prv) continue;
    g[v].push_back(i);
    dfs(i, v);
int dfs1(int v) {
  siz[v] = 1;
  for (auto& i : g[v]) {
    dep[i] = dep[v] + 1, par[i] = v;
    siz[v] += dfs1(i);
    if (siz[i] > siz[g[v][0]]) swap(i, g[v][0]);
  return siz[v];
```

```
void dfs2(int v) {
 in[v] = ++pv:
 for (auto& i : g[v]) {
   top[i] = (i == g[v][0] ? top[v] : i);
   dfs2(i);
 out[v] = pv;
void modify(int v, ll value) {
 seg.modify(in[v], value);
11 query(int u, int v) {
 ll ret = 0:
 while (top[u] ^ top[v]) {
   if (dep[top[u]] < dep[top[v]]) swap(u, v);</pre>
   int st = top[u];
   ret += seg.query(in[st], in[u]);
   u = par[st];
 if (dep[u] > dep[v]) swap(u, v);
 ret += seg.query(in[u], in[v]);
 return ret;
int main() {
 cin.tie(NULL); cout.tie(NULL);
 ios_base::sync_with_stdio(false);
  int n, q;
  cin >> n >> a:
  for (int i = 0; i < n - 1; i++) {
   int u, v;
   cin >> u >> v;
    adj[u].push_back(v);
    adj[v].push_back(u);
 }
 dfs(1, 0);
 top[1] = 1;
  dfs1(1):
  dfs2(1);
  while (q--) {
   int op, a, b;
   cin >> op >> a >> b;
    if (op == 1) modify(a, b);
    else cout << query(a, b) << '\n';
 }
4.2 Centroid Decomposition
#include <bits/stdc++.h>
using namespace std;
const int MAXV = 202020;
vector<int> adj[MAXV];
int used[MAXV], siz[MAXV], dep[MAXV], cdtree[MAXV];
int getSize(int now, int prv) {
 siz[now] = 1;
 for (auto i : adj[now]) {
   if (used[i] || prv == i) continue;
    siz[now] += getSize(i, now);
```

```
return siz[now];
int getCent(int now, int prv, int cnt) {
  for (auto& i : adj[now]) {
    if (used[i] || i == prv) continue;
    if (siz[i] > cnt / 2) return getCent(i, now, cnt);
 return now;
}
void cd(int now, int prv) {
 int cnt = getSize(now, prv);
  int cent = getCent(now, prv, cnt);
  cdtree[now] = prv;
  used[cent] = 1;
  for (auto i : adj[cent])
    if (!used[i]) cd(i, cent);
}
int main() {
  cin.tie(NULL); cout.tie(NULL);
  ios_base::sync_with_stdio(false);
  cd(0, -1):
  return 0;
}
5 Network Flow
5.1 Dinic's Algorithm
// Dinic's Algorithm
// time complexity : O(V^2 * E)
#include <bits/stdc++.h>
using namespace std;
const int INF = 1e9 + 7;
const int MAXV = 505;
int N, st = 0, en = MAXV + 1;
vector<int> adj[MAXV + 5];
int c[MAXV + 5][MAXV + 5], f[MAXV + 5][MAXV + 5];
int level[MAXV + 5], work[MAXV + 5];
void input() {
 // TODO
void bfs() {
  memset(level, -1, sizeof(level));
 level[st] = 0;
  queue<int> q;
  q.push(st);
  while (!q.empty()) {
    int now = q.front();
    q.pop();
    for (int next : adj[now]) {
      if (level[next] == -1 && c[now][next] - f[now][next] > 0) {
        level[next] = level[now] + 1:
        q.push(next);
    }
int dfs(int now, int flow) {
  if (now == en) return flow;
 for (int& i = work[now]; i < adj[now].size(); i++) {</pre>
```

```
int next = adj[now][i];
    if (level[next] == level[now] + 1 && c[now][next] - f[now][next] > 0) {
      int df = dfs(next, min(c[now][next] - f[now][next], flow));
      if (df > 0) {
        f[now][next] += df;
        f[next][now] -= df;
        return df:
   }
 }
 return 0;
int dinic() {
 int ret = 0;
 while (true) {
   bfs();
   if (level[en] == -1) break;
   memset(work, 0, sizeof(work));
   while (true) {
     int flow = dfs(st, INF);
     if (flow == 0) break:
     ret += flow;
   }
 }
 return ret;
int main() {
 cin.tie(NULL): cout.tie(NULL):
 ios_base::sync_with_stdio(false);
 input();
 int total = dinic();
  cout << total << '\n';</pre>
5.2 Hopcroft-Karp Algorithm
// Bipartite Matching Algorithm
// time complexity : O(E * sqrt(V))
#include <bits/stdc++.h>
using namespace std;
const int INF = 1e9 + 7;
const int MAXV = 10101;
int n, A[MAXV], B[MAXV], dist[MAXV];
bool used[MAXV];
vector<int> adj[MAXV];
void input() {
 // TODO
}
void bfs() {
 queue<int> q;
 for (int i = 0; i < n; i++) {
   if (!used[i]) {
     dist[i] = 0;
      q.push(i);
    else dist[i] = INF;
 }
  while (!q.empty()) {
   int a = q.front();
```

```
q.pop();
    for (int b : adj[a]) {
      if (B[b] != -1 \&\& dist[B[b]] == INF) {
        dist[B[b]] = dist[a] + 1:
        q.push(B[b]);
    }
 }
bool dfs(int a) {
 for (int b : adj[a]) {
    if (B[b] == -1 \mid | (dist[B[b]] == dist[a] + 1 && dfs(B[b]))) {
      used[a] = true:
      A[a] = b;
      B[b] = a:
      return true;
  return false;
int hopcroft() {
  memset(A, -1, sizeof(A));
  memset(B, -1, sizeof(B));
  int ret = 0;
  while (true) {
    bfs():
    int flow = 0;
    for (int i = 0; i < n; i++)
      if (!used[i] && dfs(i)) flow++;
    if (flow == 0) break;
    ret += flow:
  return ret;
int main() {
  cin.tie(NULL); cout.tie(NULL);
  ios_base::sync_with_stdio(false);
  input();
 int total = hopcroft();
  cout << total << '\n';</pre>
}
5.3 MCMF
#include <bits/stdc++.h>
using namespace std;
const int INF = 1e9 + 7;
const int MAXV = 1010;
int N, M, st = 0, en = 1001;
int c[MAXV] [MAXV], f[MAXV] [MAXV];
int d[MAXV][MAXV], prv[MAXV];
vector<int> adj[MAXV];
int mFlow, mCost;
void input() {
 // TODO
void spfa() {
  memset(prv, -1, sizeof(prv));
  vector<int> dist(MAXV, INF);
```

```
vector<bool> inQ(MAXV);
  queue<int> q;
 q.push(st);
  dist[st] = 0, inQ[st] = true;
  while (!q.empty()) {
   int now = q.front();
   q.pop();
   inQ[now] = false;
   for (int next : adj[now]) {
      if (dist[now] + d[now] [next] < dist[next] && c[now] [next] - f[now] [next] > 0) {
        dist[next] = dist[now] + d[now][next];
        prv[next] = now;
        if (!inQ[next]) {
         inQ[next] = true;
         q.push(next);
   }
void flow() {
 int block = INF;
 for (int i = en; i != st; i = prv[i]) {
    block = min(block, c[prv[i]][i] - f[prv[i]][i]);
 for (int i = en; i != st; i = prv[i]) {
   mCost += d[prv[i]][i] * block;
   f[prv[i]][i] += block:
   f[i][prv[i]] -= block;
 mFlow += block;
void mcmf() {
 while (1) {
    spfa();
   if (prv[en] == -1) break;
   flow():
 }
}
int main() {
 cin.tie(NULL); cout.tie(NULL);
 ios_base::sync_with_stdio(false);
 input();
 mcmf():
 cout << mFlow << '\n' << mCost;</pre>
6 String
6.1 Rabin-Karp Algorithm
// BOJ 1786 AC Code
// https://www.acmicpc.net/problem/1786
#include <bits/stdc++.h>
using namespace std;
#define ll long long
#define sz(x) (int)(x).size()
const int MAX = 1010101;
const int MOD1 = 1e9 + 7, MOD2 = 1e9 + 9;
string T, P;
```

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```
11 d = 128, dexp1[MAX], dexp2[MAX];
vector<int> ans:
void rabinKarp() {
 int len = sz(P):
 11 p1 = 0, p2 = 0, t1 = 0, t2 = 0;
  for (int i = 0; i < len; i++) {
    p1 = (d * p1 + P[i]) \% MOD1;
    p2 = (d * p2 + P[i]) \% MOD2;
    t1 = (d * t1 + T[i]) \% MOD1;
    t2 = (d * t2 + T[i]) \% MOD2;
  if (p1 == t1 && p2 == t2) ans.push_back(0);
  for (int i = 1; i < sz(T) - len + 1; i++) {
    t1 = (d * (t1 - dexp1[len - 1] * T[i - 1]) + T[i + len - 1]) % MOD1;
    t1 = (t1 + MOD1) \% MOD1:
    t2 = (d * (t2 - dexp2[len - 1] * T[i - 1]) + T[i + len - 1]) % MOD2;
    t2 = (t2 + MOD2) \% MOD2;
    if (p1 == t1 && p2 == t2) ans.push_back(i);
}
int main() {
  cin.tie(NULL); cout.tie(NULL);
  ios_base::sync_with_stdio(false);
  dexp1[0] = dexp2[0] = 1;
  for (int i = 1; i < MAX; i++) {
    dexp1[i] = d * dexp1[i - 1] % MOD1;
    dexp2[i] = d * dexp2[i - 1] % MOD2;
  getline(cin, T);
  getline(cin, P);
  rabinKarp();
  cout << sz(ans) << '\n';
  for (int i : ans) cout << i + 1 << ' ';
6.2 KMP Algorithm
// BOJ 1786 AC Code
// https://www.acmicpc.net/problem/1786
#include <bits/stdc++.h>
using namespace std;
#define sz(x) (int)(x).size()
vector<int> getpi(const string& P) {
  vector<int> pi(sz(P));
  for (int i = 1, j = 0; i < sz(P); i++) {
    while (j > 0 \&\& P[i] != P[j]) j = pi[j - 1];
    if (P[i] == P[j]) pi[i] = ++j;
  return pi;
vector<int> kmp(const string& T, const string& P) {
  vector<int> ret;
  vector<int> pi = getpi(P);
  for (int i = 0, j = 0; i < sz(T); i++) {
    while (j > 0 \&\& T[i] != P[j]) j = pi[j-1];
    if (T[i] == P[j]) {
      if (i == sz(P) - 1) {
        ret.push_back(i - (sz(P) - 1));
        j = pi[j];
```

```
}
      else ++i:
 return ret;
int main() {
 cin.tie(NULL); cout.tie(NULL);
 ios_base::sync_with_stdio(false);
 string T, P;
 getline(cin, T);
 getline(cin, P);
 vector<int> ans = kmp(T, P);
 cout << sz(ans) << '\n';
 for (int i : ans)
   cout << i + 1 << '\n';
6.3 Trie
// 1. Trie (Array Index)
#include <bits/stdc++.h>
using namespace std;
const char st = '0';
const int MAXC = '9' - '0' + 1;
const int MAXN = 100 * 100 * MAXC + 1;
struct trie {
 int cnt, t[MAXN][MAXC];
 bool term[MAXN]:
 void clear() {
   memset(t, 0, sizeof(t));
   memset(term, 0, sizeof(term));
   cnt = 0;
 }
 void insert(string& s) {
   int here = 0:
   for (char& i : s) {
     if (!t[here][i - st]) t[here][i - st] = ++cnt;
     here = t[here][i - st];
   }
   term[here] = true;
 bool find(string& s) {
   int here = 0;
   for (int i = 0; i < s.size(); i++) {
     if (!t[here][s[i] - st]) return false;
     here = t[here][s[i] - st];
     if (i == s.size() - 1 && term[here]) return true;
   }
   return false;
 }
trie T;
int main() {
 cin.tie(NULL); cout.tie(NULL);
 ios_base::sync_with_stdio(false);
 int N; cin >> N;
 for (int i = 0; i < N; i++) {
   string s; cin >> s;
```

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```
T.insert(s);
  int Q; cin >> Q;
  while (Q--) {
    string s; cin >> s;
    if (T.find(s)) cout << "Is exist.\n";</pre>
    else cout << "Is not exist.\n";</pre>
}
// 2. Trie (Pointer)
#include <bits/stdc++.h>
using namespace std;
const char st = 'a';
const int MAXC = 'z' - 'a' + 1:
struct trie {
  trie* child[MAXC];
  bool term:
  trie() {
    fill(child, child + MAXC, nullptr);
    term = false:
  }
  ~trie() {
    for (int i = 0; i < MAXC; i++)
      if (child[i]) delete child[i];
  void insert(const string& s, int key = 0) {
    if (s.size() == kev) term = true:
    else {
      int next = s[key] - st;
      if (!child[next]) child[next] = new trie;
      child[next]->insert(s, key + 1);
  }
  bool find(const string& s, int key = 0) {
    if (s.size() == key) return term;
    else {
      int next = s[key] - st;
      if (!child[next]) return false;
      else return child[next]->find(s, key + 1);
    }
  }
};
int main() {
  cin.tie(NULL); cout.tie(NULL);
  ios_base::sync_with_stdio(false);
  trie* root = new trie;
  int N; cin >> N;
  for (int i = 0; i < N; i++) {
    string s: cin >> s:
    root->insert(s);
  int Q; cin >> Q;
  while (Q--) {
    string s; cin >> s;
    if (root->find(s)) cout << "Is exist.\n";</pre>
    else cout << "Is not exist.\n";</pre>
```

```
delete root;
6.4 Aho-Corasick
// BOJ 9250 AC Code
// https://www.acmicpc.net/problem/9250
#include <bits/stdc++.h>
using namespace std;
const char st = 'a';
const int MAXC = 'z' - 'a' + 1;
struct trie {
 trie* child[MAXC];
 trie* fail:
 bool term;
 trie() {
   fill(child, child + MAXC, nullptr);
   fail = nullptr;
   term = false:
 }
  ~trie() {
   for (int i = 0: i < MAXC: i++)
      if (child[i]) delete child[i];
  void insert(const string& s, int key = 0) {
    if (s.size() == key) term = true;
    else {
      int next = s[key] - st;
     if (!child[next]) child[next] = new trie;
      child[next] -> insert(s, key + 1);
   }
 }
trie* root = new trie;
void getFail() {
 queue<trie*> q;
 q.push(root);
 root->fail = root:
  while (!q.empty()) {
   trie* now = q.front();
   q.pop();
    for (int i = 0; i < MAXC; i++) {
     trie* next = now->child[i];
      if (!next) continue;
      if (now == root) next->fail = root;
      else {
        trie* t = now->fail;
        while (t != root && !t->child[i])
         t = t->fail;
        if (t->child[i]) t = t->child[i];
        next->fail = t:
      if (next->fail->term) next->term = true;
      q.push(next);
 }
bool isMatch(const string& s) {
 trie* now = root:
```

```
bool ret = false;
  for (int c = 0: c < s.size(): c++) {
    int next = s[c] - st;
    while (now != root && !now->child[next])
      now = now->fail:
    if (now->child[next])
      now = now->child[next]:
    if (now->term) {
      ret = true;
      break;
  return ret:
int main() {
  cin.tie(NULL); cout.tie(NULL);
  ios_base::sync_with_stdio(false);
  int N; cin >> N;
  for (int i = 0; i < N; i++) {</pre>
    string s: cin >> s:
    root->insert(s):
  getFail();
  int M; cin >> M;
  for (int i = 0; i < M; i++) {
    string s: cin >> s:
    if (isMatch(s)) cout << "YES\n";</pre>
    else cout << "NO\n":
  }
  delete root;
}
6.5 Suffix Array
// Manber-Myers Algorithm for Suffix Array
// Time Complexity: O(mlog^2n)
// Kasai's Algorithm for LCP(Longest Common Prefix)
// Time Complexity: O(n)
// BOJ 9248 AC Code
// https://www.acmicpc.net/problem/9248
#include <bits/stdc++.h>
using namespace std;
#define sz(x) (int)(x).size()
vector<int> buildsa(const string& s) {
    int n = sz(s);
    vector\langle int \rangle sa(n), r(n + 1), nr(n + 1);
    for (int i = 0; i < n; i++) sa[i] = i, r[i] = s[i];
    for (int d = 1; d < n; d <<= 1) {
        auto cmp = [&](int i, int j) {
            if (r[i] ^ r[j]) return r[i] < r[j];</pre>
            return r[i + d] < r[i + d]:
        sort(sa.begin(), sa.end(), cmp);
        nr[sa[0]] = 1:
        for (int i = 1; i < n; i++)
            nr[sa[i]] = nr[sa[i-1]] + cmp(sa[i-1], sa[i]);
    }
    return sa:
```

```
vector<int> buildlcp(const string& s. const vector<int>& sa) {
    int n = sz(s);
   vector<int> lcp(n), isa(n):
   for (int i = 0; i < n; i++) isa[sa[i]] = i;
    for (int k = 0, i = 0; i < n; i++) if (isa[i]) {
        for (int j = sa[isa[i] - 1]; s[i + k] == s[j + k]; k++);
        lcp[isa[i]] = (k ? k-- : 0);
   }
   return lcp;
int main() {
    cin.tie(NULL): cout.tie(NULL):
   ios_base::sync_with_stdio(false);
    string s; cin >> s;
    vector<int> sa = buildsa(s);
    vector<int> lcp = buildlcp(s, sa);
   for (auto& i : sa) cout << i + 1 << ' ';
   cout << '\n';
   cout << "x ":
   for (int i = 1: i < sz(lcp): i++) cout << lcp[i] << ' ':
// Manber-Myers Algorithm for Suffix Array
// Time Conplexity: O(nlogn)
// Kasai's Algorithm for LCP(Longest Common Prefix)
// Time Complexity: O(n)
// BOJ 9248 AC Code
// https://www.acmicpc.net/problem/9248
#include <bits/stdc++.h>
using namespace std;
#define sz(x) (int)(x).size()
vector<int> buildsa(const string& s) {
    int n = sz(s), m = max(256, n) + 1;
    vector\langle int \rangle sa(n), r(2 * n), nr(2 * n), cnt(m), idx(n);
    for (int i = 0; i < n; i++) sa[i] = i, r[i] = s[i];
    for (int d = 1; d < n; d <<= 1) {
        auto cmp = [&](int i, int j) {
            if (r[i] ^ r[j]) return r[i] < r[j];</pre>
            return r[i + d] < r[j + d];
        }:
        for (int i = 0; i < m; i++) cnt[i] = 0;
        for (int i = 0; i < n; i++) cnt[r[i + d]]++:
        for (int i = 1; i < m; i++) cnt[i] += cnt[i - 1];
        for (int i = n - 1; ~i; i--) idx[--cnt[r[i + d]]] = i;
        for (int i = 0; i < m; i++) cnt[i] = 0;
        for (int i = 0; i < n; i++) cnt[r[i]]++;
        for (int i = 1; i < m; i++) cnt[i] += cnt[i - 1];
        for (int i = n - 1; ~i; i--) sa[--cnt[r[idx[i]]]] = idx[i];
        nr[sa[0]] = 1;
        for (int i = 1; i < n; i++) nr[sa[i]] = nr[sa[i - 1]] + cmp(sa[i - 1], sa[i]);
        for (int i = 0; i < n; i++) r[i] = nr[i];
        if (r[sa[n-1]] == n) break;
   }
   return sa;
vector<int> buildlcp(const string& s, const vector<int>& sa) {
    int n = sz(s);
    vector<int> lcp(n), isa(n);
```

```
for (int i = 0; i < n; i++) isa[sa[i]] = i;
    for (int k = 0, i = 0; i < n; i++) if (isa[i]) {
        for (int j = sa[isa[i] - 1]; s[i + k] == s[j + k]; k++);
        lcp[isa[i]] = (k ? k-- : 0):
    }
    return lcp;
}
int main() {
    cin.tie(NULL); cout.tie(NULL);
    ios_base::sync_with_stdio(false);
    string s; cin >> s;
    vector<int> sa = buildsa(s);
    vector<int> lcp = buildlcp(s, sa);
    for (auto& i : sa) cout << i + 1 << ' ';
    cout << '\n':
    cout << "x ";
    for (int i = 1; i < sz(lcp); i++) cout << lcp[i] << ' ';
}
6.6 Manacher's Algorithm
// Manacher's Algorithm
// Find all palindromes in string in O(N)
// BOJ 14444 AC Code
// https://www.acmicpc.net/problem/14444
#include <bits/stdc++.h>
using namespace std;
#define sz(x) (x).size()
int n: // n: length of string
vector<int> p; // p[i]: the radius of the palindrome at the current position i
void manacher() {
    // Preprocessing for determining even-length pelindromes
    n = sz(s);
    s.resize(n << 1 | 1);
    p.resize(n << 1 | 1);
    for (int i = n - 1; i \ge 0; i--) {
        s[i \ll 1 \mid 1] = s[i]:
        s[i << 1] = '#';
    }
    n <<= 1;
    s[n++] = '#';
    // Processing
    int r = -1, c = -1;
    // r: end of palindrome
    // c: center of palindrome
    for (int i = 0; i < n; i++) {
        if (i \le r) p[i] = min(r - i, p[c * 2 - i]);
        else p[i] = 0;
        while (1) {
            if (i - p[i] - 1 < 0 \mid | i + p[i] + 1 >= n) break;
            if (s[i + p[i] + 1] != s[i - p[i] - 1]) break;
            p[i]++;
        }
        if (i + p[i] > r) {
            r = i + p[i], c = i;
    }
}
```

```
int main() {
    cin.tie(NULL): cout.tie(NULL):
    ios_base::sync_with_stdio(false);
    cin >> s:
   manacher():
   // Get answer
   int ans = 0:
   for (int i = 0; i < n; i++) {
        ans = max(ans, p[i]);
   }
    cout << ans;</pre>
6.7 Z Algorithm
// Z Algorithm
// Given a string S of length n, the Z Algorithm produces an array Z
// where Z[i] is the length of the longest substring starting from S[i] which is also a prefix
// BOJ 13713 AC Code
// https://www.acmicpc.net/problem/13713
#include <bits/stdc++.h>
using namespace std;
#define sz(x) (int)(x).size()
const int MAXS = 1010101;
string s;
int z[MAXS];
void input() {
   string du; cin >> du;
   for (int i = sz(du) - 1; i \ge 0; i--)
        s.push_back(du[i]);
void zfunction() {
   z[0] = sz(s);
   int 1 = 0, r = 0;
   for (int i = 1; i < sz(s); i++) {
        if (i > r) {
           1 = r = i:
            while (r < sz(s) \&\& s[r - 1] == s[r]) r++;
            z[i] = r - 1; r--:
        }
        else {
            int k = i - 1;
            if (z[k] < r - i + 1) z[i] = z[k];
            else {
                while (r < sz(s) \&\& s[r - 1] == s[r]) r++;
                z[i] = r - 1; r--;
           }
        }
   }
int main() {
    cin.tie(NULL); cout.tie(NULL);
   ios_base::sync_with_stdio(false);
   input();
   zfunction();
   int q; cin >> q;
    while (q--) {
```

int x; cin >> x;

cout $\langle\langle z[sz(s) - x] \langle\langle ' n';$

```
}
           Geometry
                Convex Hull
 // 1. Graham Scan
 // BOJ 1708 AC Code
 // https://www.acmicpc.net/problem/1708
 #include <bits/stdc++.h>
 using namespace std;
 #define ll long long
 struct point {
     11 x, y;
     bool operator<(const point& rhs) const {</pre>
           if (y != rhs.y) return y < rhs.y;</pre>
          return x < rhs.x:
     }
};
int N:
vector<point> p;
 vector<int> st;
 11 ccw(const point& a, const point& b, const point& c) {
      // res > 0 -> ccw, res < 0 -> cw, res = 0 -> colinear
     ll res = (b.x - a.x) * (c.y - a.y) - (c.x - a.x) * (b.y - a.y);
     return (res > 0 ? 1 : (res < 0 ? -1 : 0));
void input() {
      cin >> N;
     for (int i = 0; i < N; i++) {
          int x, y;
           cin >> x >> y;
           p.push_back({ x, y });
11 dist(const point& p1, const point& p2) {
           return (p1.x - p2.x) * (p1.x - p2.x) + (p1.y - p2.y) * (p1.y - p2.y);
 bool cmp(const point& p1, const point& p2) {
           return (ccw(p[0], p1, p2) > 0 \mid | (ccw(p[0], p1, p2) == 0 && dist(p[0], p1) < dist(p[0], p1) | (ccw(p[0], p1, p2) | (ccw(p[0], p1, p2
           p2)));
void grahamScan() {
      sort(p.begin(), p.end());
      sort(p.begin() + 1, p.end(), cmp);
      st.push_back(0);
      st.push_back(1);
      for (int next = 2; next < N; next++) {</pre>
           while (st.size() >= 2) {
                int first = st.back();
                 st.pop_back();
                int second = st.back();
                 if (ccw(p[second], p[first], p[next]) > 0) {
                      st.push_back(first);
                     break;
                }
          }
```

```
st.push_back(next);
 }
}
// 2. Monotone Chain
// BOJ 1708 AC Code
// https://www.acmicpc.net/problem/1708
#include <bits/stdc++.h>
using namespace std;
#define 11 long long
struct point {
 11 x, y;
 bool operator<(const point& rhs) const {</pre>
   if (x != rhs.x) return x < rhs.x;
    else return y < rhs.y;</pre>
 }
};
int N;
vector<point> p;
vector<int> dh, uh;
11 ccw(const point& a, const point& b, const point& c) {
 // res > 0 -> ccw, res < 0 -> cw, res = 0 -> colinear
 11 res = (b.x - a.x) * (c.y - a.y) - (c.x - a.x) * (b.y - a.y);
 return (res > 0 ? 1 : (res < 0 ? -1 : 0));
void input() {
 cin >> N;
 for (int i = 0; i < N; i++) {
   int x, y;
   cin >> x >> y;
    p.push_back({ x, y });
void monotoneChain() {
  sort(p.begin(), p.end());
 // calculate lower hull
  dh.push_back(0);
  dh.push_back(1);
  for (int next = 2; next < N; next++) {</pre>
    while (dh.size() >= 2) {
      int first = dh.back();
      dh.pop_back();
      int second = dh.back();
      if (ccw(p[second], p[first], p[next]) > 0) {
        dh.push_back(first);
        break;
   }
    dh.push_back(next);
 // calculate upper hull
  uh.push_back(N - 1);
  uh.push_back(N - 2);
 for (int next = N - 3; next >= 0; next--) {
    while (uh.size() >= 2) {
      int first = uh.back();
      uh.pop_back();
      int second = uh.back();
```

```
if (ccw(p[second], p[first], p[next]) > 0) {
        uh.push_back(first);
        break;
      }
    }
    uh.push_back(next);
}
int main() {
  cin.tie(NULL); cout.tie(NULL);
  ios_base::sync_with_stdio(false);
  input();
  monotoneChain();
  cout << (int)dh.size() + (int)uh.size() - 2;</pre>
7.2 Rotating Callipers
// BOJ 10254 AC Code
// https://www.acmicpc.net/problem/10254
#include <bits/stdc++.h>
using namespace std;
#define ll long long
struct point {
    11 x, y;
    bool operator<(const point& rhs) const {
        if (x != rhs.x) return x < rhs.x;</pre>
        return y < rhs.y;</pre>
    }
};
int n;
vector<point> p, ch;
point ans1, ans2;
void init() {
    n = 0;
    p.clear();
    ch.clear();
    ans1 = ans2 = \{0, 0\}:
void input() {
    cin >> n;
    p.resize(n);
    ch.resize(n);
    for (auto& i : p) cin >> i.x >> i.y;
}
11 ccw(const point& a, const point& b, const point& c) {
  // res > 0 -> ccw, res < 0 -> cw, res = 0 -> colinear
  ll res = (b.x - a.x) * (c.y - a.y) - (c.x - a.x) * (b.y - a.y);
  return (res > 0 ? 1 : (res < 0 ? -1 : 0));
11 cccw(point a, point b, point c, point d) {
    d.x = c.x - b.x;
    d.y = c.y - b.y;
    return ccw(a, b, d);
11 dist(const point& p1, const point& p2) {
    return (p1.x - p2.x) * (p1.x - p2.x) + (p1.y - p2.y) * (p1.y - p2.y);
bool cmp(const point& p1, const point& p2) {
```

```
p2)));
void rotatingCallipers() {
    sort(p.begin(), p.end());
    sort(p.begin() + 1, p.end(), cmp);
    ch[0] = p[0];
    ch[1] = p[1];
    11 f1 = 2, cnt = 2;
    for (int i = 2; i < n; i++) {
        while (f1 \ge 2 \&\& ccw(ch[f1 - 2], ch[f1 - 1], p[i]) \le 0) f1--;
        ch[fl] = p[i], fl++;
    }
    11 f12 = 1, mx = 0;
    for (int i = 0; i < fl; i++) {
        while ((fl2 + 1) != i \&\& cccw(ch[i], ch[i + 1], ch[fl2 % fl], ch[(fl2 + 1) % fl]) > 0) {
            if (mx < dist(ch[i], ch[f12 % f1])) {</pre>
                ans1 = ch[i], ans2 = ch[fl2 % fl];
                mx = dist(ch[i], ch[fl2 % fl]);
           }
           f12++:
        }
        if (mx < dist(ch[i], ch[f12 % f1])) {</pre>
            ans1 = ch[i], ans2 = ch[fl2 \% fl];
            mx = dist(ch[i], ch[f12 % f1]);
        }
    }
}
int main() {
    cin.tie(NULL); cout.tie(NULL);
    ios_base::sync_with_stdio(false);
    int tc; cin >> tc;
    while (tc--) {
        init();
        input();
        rotatingCallipers();
        cout << ans1.x << ' ' << ans1.y << ' ';
        cout << ans2.x << ' ' << ans2.y << '\n';
    }
7.3 Ray Casting
// BOJ 1688 AC Code
// https://www.acmicpc.net/problem/1688
#include <bits/stdc++.h>
using namespace std;
#define 11 long long
#define pii pair<int, int>
struct point {
    bool operator==(const point& rhs) const {
        return x == rhs.x && y == rhs.y;
    bool operator<=(const point& rhs) const {</pre>
        if (x < rhs.x || (x == rhs.x && y <= rhs.y)) return 1;
        else return 0;
    }
∣};
```

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```
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```

```
int n;
vector<point> p;
point a, b, c;
void input() {
    cin >> n:
    p.resize(n);
    for (auto& i : p) {
        cin >> i.x >> i.y;
    p.push_back(p[0]);
    cin >> a.x >> a.v;
    cin >> b.x >> b.y;
    cin >> c.x >> c.y;
11 ccw(const point& a, const point& b, const point& c) {
    // res > 0 -> ccw, res < 0 -> cw, res = 0 -> colinear
    ll res = (b.x - a.x) * (c.y - a.y) - (c.x - a.x) * (b.y - a.y);
    return (res > 0 ? 1 : (res < 0 ? -1 : 0));
// Does the line segment ab and cd intersect?
bool isCross(point a, point b, point c, point d) {
    ll ab = ccw(a, b, c) * ccw(a, b, d);
    11 cd = ccw(c, d, a) * ccw(c, d, b);
    if (ab == 0 && cd == 0) {
        pii A = \{a.x, a.y\}, B = \{b.x, b.y\}, C = \{c.x, c.y\}, D = \{d.x, d.y\};
        if (A > B) swap(A, B):
        if (C > D) swap(C, D);
        return (A <= D && C <= B):
    }
    else return (ab <= 0 && cd <= 0);
bool insidePolygon(point v) {
    point u = \{ 101010101011, v.v + 1 \};
    for (int i = 0; i < n; i++) {
        if (p[i] == v) return 1;
    for (int i = 0: i < n: i++) {
        if (!ccw(p[i], p[i + 1], v) \&\& (p[i] \le v \hat{p}[i + 1] \le v)) return 1;
    }
    int cnt = 0;
    for (int i = 0; i < n; i++) {
        cnt += isCross(p[i], p[i + 1], u, v);
    }
    return cnt & 1;
int main() {
    cin.tie(NULL); cout.tie(NULL);
    ios_base::sync_with_stdio(false);
    input();
    cout << (insidePolygon(a) ? 1 : 0) << '\n';</pre>
    cout << (insidePolygon(b) ? 1 : 0) << '\n';</pre>
    cout << (insidePolygon(c) ? 1 : 0) << '\n';</pre>
}
7.4 Sort by Angular
// Sort by Angular (Relative to Origin)
struct Point {
    11 x, y;
```

```
bool operator<(const Point& rhs) const {</pre>
        return x ^ rhs.x ? x < rhs.x : y < rhs.y;
}:
const Point o = \{ 0, 0 \};
vector<Point> p;
11 ccw(const Point& a, const Point& b, const Point& c) {
    // res > 0 -> ccw, res < 0 -> cw, res = 0 -> colinear
    ll res = (b.x - a.x) * (c.y - a.y) - (c.x - a.x) * (b.y - a.y);
    return (res > 0 ? 1 : (res < 0 ? -1 : 0));
inline 11 dist(const Point& v) { return v.x * v.x + v.y * v.y; }
// If the angle between any two points and the origin is less than 180 degrees,
// they can be sorted through the cross product of the vectors.
// Therefore, the points were divided into 1st and 4th quadrants and 2nd and 3rd quadrants,
// and the points in the 1st and 4th quadrants were arranged in front.
void sortBvAngular() {
    sort(p.begin(), p.end(), [&](const Point& lhs, const Point& rhs) {
        if ((lhs < o) ^ (rhs < o)) return (lhs < o) < (rhs < o);
        if (ccw(o, lhs, rhs)) return ccw(o, lhs, rhs) > 0;
        return dist(lhs) < dist(rhs);</pre>
   });
7.5 Bulldozer Trick
// Bulldozer Trick
// There are a total of O(N^2) results of sorting points on a two-dimensional plane based on an
arbitrary axis.
// The Bulldozer Trick traverses all O(N^2) results in O(N^2logN) time.
struct Point {
    11 x, y;
    bool operator<(const Point& rhs) const {</pre>
        return tie(x, y) < tie(rhs.x, rhs.y);</pre>
    }
}:
struct Line {
    int u, v; 11 dx, dy; // u < v, dx >= 0;
    bool operator<(const Line& rhs) const {</pre>
        if (dy * rhs.dx != rhs.dy * dx) return dy * rhs.dx < rhs.dy * dx;
        return tie(u, v) < tie(rhs.u, rhs.v);</pre>
    bool operator==(const Line& rhs) const {
        return dy * rhs.dx == rhs.dy * dx;
}:
int n, pos[2020];
Point p[2020];
void bulldozerTrick() {
    sort(p + 1, p + 1 + n);
    for (int i = 1; i <= n; i++) pos[i] = i;
```

}

}

```
// find the slope between every two points.
    vector<Line> arr:
    for (int i = 1; i <= n; i++) {
        for (int j = i + 1; j \le n; j++) {
            arr.push_back({ i, j, p[j].x - p[i].x, p[j].y - p[i].y });
    }
    sort(arr.begin(), arr.end());
    // can check one of the results of sorting points at here.
    for (int i = 0, j = 0; i < arr.size(); i = j) {</pre>
        while (j < arr.size() && arr[j] == arr[i]) j++; // all lines in [i, j) are same
        for (int k = i; k < j; k++) {
            int u = arr[k].u, v = arr[k].v:
            swap(p[pos[u]], p[pos[v]]);
            swap(pos[u], pos[v]);
        }
        // can check one of the results of sorting points at here.
7.6 Minimum Enclosing Circle
// INPUT: Given N points in a 2D plane with integer coordinates.
// OUTPUT: Find the center and the radius of the minimum enclosing circle.
// A minimum enclosing circle is a circle in which all the points lie either inside the circle
or on its boundaries.
// TIME COMPLEXITY: O(N) (using random)
// BOJ 2626 AC Code
// https://www.acmicpc.net/problem/2626
#include <bits/stdc++.h>
using namespace std;
struct Point { long double x, y; };
struct Circle { Point c: long double r: }:
long double dist(const Point& a, const Point& b) {
    return sqrt(pow(b.x - a.x, 2) + pow(b.y - a.y, 2));
Point getCircleCenter(const Point& a, const Point& b) {
    long double A = a.x * a.x + a.y * a.y;
    long double B = b.x * b.x + b.y * b.y;
    long double C = a.x * b.y - a.y * b.x;
    return { (b.y * A - a.y * B) / (2 * C), (a.x * B - b.x * A) / (2 * C) };
Circle circleFrom(const Point& a, const Point& b, const Point& c) {
    Point i = getCircleCenter({ b.x - a.x, b.y - a.y }, { c.x - a.x, c.y - a.y });
    i.x += a.x;
    i.y += a.y;
    return { i, dist(a, i) };
Circle circleFrom(const Point& a, const Point& b) {
    Point c = \{ (a.x + b.x) / 2.0, (a.y + b.y) / 2.0 \};
    return { c, dist(a, b) / 2.0 };
```

```
Circle minimumEnclosingCircle(int n. const vector<Point>& p) {
    Circle ret = { { 0, 0 }, 0 };
    for (int i = 0: i < n: i++) {
        if (dist(ret.c, p[i]) <= ret.r) continue;</pre>
        ret.c = p[i], ret.r = 0;
        for (int j = 0; j < i; j++) {
            if (dist(ret.c, p[j]) <= ret.r) continue;</pre>
            ret = circleFrom(p[i], p[j]);
            for (int k = 0; k < j; k++) {
                if (dist(ret.c, p[k]) <= ret.r) continue;</pre>
                 ret = circleFrom(p[i], p[j], p[k]);
            }
        }
    }
    return ret;
int main() {
    cin.tie(NULL): cout.tie(NULL):
    ios base::svnc with stdio(false):
    int n; cin >> n;
    vector<Point> p(n);
    for (auto& i : p)
        cin >> i.x >> i.y;
    random_shuffle(p.begin(), p.end());
    Circle ans = minimumEnclosingCircle(n, p);
    cout << fixed;</pre>
    cout.precision(3);
    cout << ans.c.x << ' ' << ans.c.y << '\n' << ans.r;
8 Math
8.1 Sieve
// Sieve of Eratosthenes
// TIME COMPLEXITY: O(Nlog(log(N)))
const int MAX = 1e6;
bool isPrime[MAX + 1]:
vector<int> prime(1, 2);
void getPrime() {
 fill(isPrime + 2, isPrime + MAX + 1, 1);
  for (11 i = 4; i \le MAX; i += 2)
    isPrime[i] = 0;
  for (11 i = 3; i <= MAX; i += 2) {
    if (!isPrime[i]) continue;
    prime.push back(i):
    for (ll j = i * i; j \le MAX; j += i * 2)
      isPrime[j] = 0;
  }
}
// Linear Sieve
// BOJ 16563 AC Code
// https://www.acmicpc.net/problem/16563
```

```
#include <bits/stdc++.h>
using namespace std:
#define ll long long
const int MAXN = 5000000:
vector<int> sp(MAXN + 1);
vector<ll> prime;
// Determine prime numbers between 1 and MAXN in O(MAXN)
void linearSieve() {
  for (int i = 2;i <= MAXN; i++) {</pre>
    if (!sp[i]) {
      prime.push_back(i);
      sp[i] = i;
    for (auto j : prime) {
      if (i * j > MAXN) break;
      sp[i * j] = j;
      if (i % j == 0) break;
 }
// factorization in O(log x)
void factorization(int x) {
  while (x > 1) {
    cout << sp[x] << ' ';
    x /= sp[x];
  cout << '\n';
int main() {
  cin.tie(NULL); cout.tie(NULL);
  ios_base::sync_with_stdio(false);
  linearSieve();
  int n: cin >> n:
  while (n--) {
    int x; cin >> x;
    factorization(x);
}
8.2 Euclidean Algorithms
// GCD, LCM
11 gcd(ll a, ll b) {
 if (b == 0) return a;
  else return gcd(b, a % b);
11 lcm(ll a, ll b) {
 return a * b / gcd(a, b);
// Extended GCD
// Bezout's Lemma
// Let a and b be integers with d := gcd(a, b).
// Then there exist integers x and y such that ax + by = d.
// Moreover, the integers of the form ax + by are exactly the multiples of d.
// If the integers x1 and y1 satisfy a * x1 + b * y1 = d,
// x2 := x1 + k * b / gcd(a, b) and
// y2 := y1 - k * a / gcd(a, b) also satisfy a * x2 + b * y2 = d for some integer k.
```

```
// Results from Bezout's Lemma
// (1) : gcd(a,b) = d \Rightarrow gcd(a/d, b/d) = 1.
// (2) : gcd(a,b) = 1, a|c, b|c => ab|c.
// (3) : a|bc, gcd(a,b) = 1 => a|c.
// (4) : d|a, d|b <=> d|gcd(a,b)
// (5) : gcd(ab,ac) = |a|gcd(b,c)
// (6) : a|bc <=> (a/gcd(a,b))|c
// Linear Congreunce
// Finding set of x satisfies ax = b \pmod{n}.
// We can think ax = b \pmod{n} as ax + ny = b.
// Then we need to find all (x, y) that satisfy ax + ny = b.
// ax + ny = b has a solution iff gcd(a,n) \mid b
// Modular Inverse
// In particular, if gcd(a,n) = 1, then we can find a unique x in (mod n) such that ax = 1 (mod
// This is called the modular inverse of a over (mod n), sometimes written a^{-1}.
// Multiplying the modular inverse is equivalent to dividing by a.
// TIME COMPLEXITY: O(log(AB))
// BOJ 14565 AC Code
// https://www.acmicpc.net/problem/14565
#include <bits/stdc++.h>
using namespace std;
#define 11 long long
#define pll pair<11, 11>
pair<pll, 11> egcd(11 a, 11 b) {
 ll s = 0, olds = 1;
 11 t = 1, oldt = 0;
 11 r = b, oldr = a;
  while (r != 0) {
   11 q = oldr / r;
   11 \text{ tmp} = \text{oldr} - \text{q} * \text{r};
    oldr = r, r = tmp:
    tmp = olds - q * s;
    olds = s, s = tmp:
   tmp = oldt - q * t;
    oldt = t, t = tmp;
 // a * olds + b * oldt = d
 // oldr = gcd(a, b)
 return { { olds, oldt }, oldr };
11 linearCongruence(11 a, 11 b, 11 n) { // Find x such that ax = b (mod n).
 pair<pll, 11> res = egcd(a, n);
 // ax + ny = b has a solution iff gcd(a,n) | b.
  if (b % res.second) return -1:
 return (res.first.first + n) % n;
11 modInv(ll a, ll p) { // Find x such that ax = 1 (mod p).
 pair<pll, 11> res = egcd(a, p);
  // Modular inverse exists iff gcd(a, p) = 1.
  if (res.second == 1) return (res.first.first + p) % p;
  else return -1;
```

```
int main() {
 11 N. A:
  cin >> N >> A;
  cout << N - A << ' ' << modInv(A, N):
8.3 Fermat's Little Theorem
// BOJ 11401 AC Code
// https://www.acmicpc.net/problem/11401
#include <bits/stdc++.h>
using namespace std;
#define 11 long long
const int MOD = 1e9 + 7:
// Fermat's little theorem
// A / B = A * B^{p - 2} \pmod{p}
11 powxy(11 x, 11 y) {
    if (y == 0) return 1;
    if (v == 1) return x:
    ll res = powxy(x, y / 2);
    return res * res % MOD * (y & 1 ? x : 1) % MOD;
}
int main() {
    cin.tie(NULL); cout.tie(NULL);
    ios_base::sync_with_stdio(false);
    11 \text{ fac}[4040404] = \{ 1, \};
    for (int i = 1; i < 4040404; i++)
        fac[i] = i * fac[i - 1] % MOD;
    int n, r;
    cin >> n >> r:
    // print nCr (mod 1e9+7)
    cout << fac[n] * powxy(fac[r], MOD - 2) % MOD * powxy(fac[n - r], MOD - 2) % MOD;
}
8.4 Euler's Phi Function
// INPUT: Given a natural number n.
// OUTPUT: Find the number of natural numbers 1 \le k \le n such that GCD(n, k) = 1.
// TIME COMPLEXITY: O(sart(n))
// BOJ 11689 AC Code
// https://www.acmicpc.net/problem/11689
#include <bits/stdc++.h>
using namespace std;
#define 11 long long
#define pll pair<11, 11>
#define fr first
#define sc second
ll phi(ll x) { // Find phi(x) in O(sqrt(x)).
    vector<pll> p;
    // Factorization in O(sqrt(x)).
    for (11 i = 2; i \le sqrt(x); i++) {
        11 \text{ res} = 1:
        while (x \% i == 0) {
            x /= i, res *= i;
        if (res > 1) p.push_back({ res, i });
    if (x > 1) p.push_back({ x, x });
    // Find phi(x).
    // phi(p^k) = p^{k-1} * (p-1) for any prime number p.
```

```
// phi(mn) = phi(m) * phi(n) if gcd(m, n) = 1
   ll ret = 1:
   for (auto& i : p) {
       ret *= (i.fr / i.sc) * (i.sc - 1);
   }
   return ret;
int main() {
   ll n; cin >> n;
    cout << phi(n);</pre>
8.5 Chinese Remainder Theorem
// INPUT: Given 4 integers, M, N, p, q. (1 <= M, 1 <= N, 0 <= p < M, 0 <= q < N)
// OUTPUT: Solve a system of linear congruence, x = p \pmod{M}, x = q \pmod{N}.
// TIME COMPLEXITY: O(log(max(M, N)))
// BOJ 6064 AC Code
// https://www.acmicpc.net/problem/6064
#include <bits/stdc++.h>
using namespace std;
#define ll long long
#define pll pair<11, 11>
#define fr first
#define sc second
11 gcd(11 x, 11 y) {
 if (!v) return x;
 return gcd(y, x % y);
11 minv(11 x, 11 y) {
 if (x == 0 && y == 1) return 0;
 if (x == 1) return 1;
 return y - minv(y \% x, x) * y / x;
// x = U.fr (mod U.sc)
// x = V.fr \pmod{V.sc}
// returns solutions as x = ret.fr (mod ret.sc)
// if no solution, returns { -1, -1 }
pll crt(pll U, pll V) {
 if (U.sc == -1 || V.sc == -1) return { -1, -1 };
 if (U.sc == 1) return V:
 if (V.sc == 1) return U;
 ll g = gcd(U.sc, V.sc);
 11 1 = U.sc * V.sc / g;
 // (U and V have a solution) iff (U.fr = U.sc (mod gcd(U.sc,V.sc)))
 // also the solution is unique in the range [0, lcm(U.sc,V.sc)).
 if ((V.fr - U.fr) % g) return { -1, -1 };
 11 u = U.sc / g, v = V.sc / g;
 11 mul = (V.fr - U.fr) / g;
 mul = mul * minv(u % v, v) % v;
 pll ret = { mul * U.sc + U.fr, 1 };
 ret.fr %= ret.sc, ret.fr = (ret.fr + ret.sc) % ret.sc;
 return ret;
pll solvingSystemOfLinearCongruence(const vector<pll>& a) {
```

```
if (a.size() == 1) return a[0];
  pll ret = crt(a[0], a[1]):
  for (int i = 2; i < a.size(); i++) ret = crt(ret, a[i]);</pre>
  return ret:
}
int main() {
  cin.tie(NULL); cout.tie(NULL);
  ios_base::sync_with_stdio(false);
  int tc; cin >> tc;
  while (tc--) {
    11 M, N, p, q;
    cin >> M >> N >> p >> q;
    p--, q--; // from the given input, 1 \le p \le M, 1 \le q \le N
    vector<pll> cg;
    cg.push_back({ p, M });
    cg.push_back({ q, N });
    pll ans = solvingSystemOfLinearCongruence(cg);
    cout << (ans.fr == -1 ? ans.fr : ans.fr + 1) << '\n':
}
8.6 Binomial Coefficient
// nCr in O(r)
// Beware of integer overflow
11 binom(int n, int r) {
  if (r < 0 \mid | n < r) return 0:
  r = min(r, n - r);
  11 \text{ ret} = 1;
  for (ll i = 1; i <= r; i++) {
    ret *= n + 1 - i;
    ret /= i;
  return ret;
// nCr (Pascal's Rule)
ll binomDP[1010][1010];
void init() {
  for (int i = 0; i < 1010; i++) {
    for (int j = 0; j < 1010; j++) {
      binomDP[i][j] = -1;
    }
 }
11 binom(int n, int r) {
  if (r < 0 \mid \mid n < r) return 0;
  11& ret = binomDP[n][r];
  if (ret != -1) return ret:
  if (n == 1) return ret = 1;
  return binom(n-1, r-1) + binom(n-1, r);
}
// nCr mod p in O(1)
// BOJ 13977 AC Code
// https://www.acmicpc.net/problem/13977
#include <bits/stdc++.h>
```

```
using namespace std;
#define 11 long long
const int MOD = 1e9 + 7;
const int MAXN = 4040404:
11 fac[MAXN], inv[MAXN], facInv[MAXN];
11 binom(int n, int r) {
 return fac[n] * facInv[r] % MOD * facInv[n - r] % MOD:
int main() {
 cin.tie(NULL); cout.tie(NULL);
 ios_base::sync_with_stdio(false);
 // Preprocessing in O(N)
 fac[0] = fac[1] = inv[1] = 1;
 facInv[0] = facInv[1] = 1;
 for (int i = 2; i < MAXN; i++) {
   fac[i] = i * fac[i - 1] % MOD;
   inv[i] = -(MOD / i) * inv[MOD % i] % MOD;
   if (inv[i] < 0) inv[i] += MOD;</pre>
   facInv[i] = facInv[i - 1] * inv[i] % MOD;
 // Answer each query in O(1)
  int q; cin >> q;
 while (q--) {
   int n, r;
   cin >> n >> r;
    cout << binom(n, r) << '\n';
 }
}
8.7 Matrix
// BOJ 11444 AC Code
// https://www.acmicpc.net/problem/11444
#include <bits/stdc++.h>
using namespace std;
#define 11 long long
#define sz(x) (int)(x).size()
const int MOD = 1e9 + 7:
struct Matrix {
   vector<vector<ll>> a:
   Matrix operator*(const Matrix& rhs) const {
        Matrix ret;
        ret.a.resize(sz(a), vector<ll>(sz(rhs.a[0])));
        for (int y = 0; y < sz(ret.a); y++) {
            for (int x = 0; x < sz(ret.a[y]); x++) {
                11 \text{ sum} = 0:
                for (int i = 0; i < sz(a[y]); i++) {
                    sum = (sum + a[y][i] * rhs.a[i][x]) % MOD;
                ret.a[y][x] = sum;
            }
        }
        return ret;
   7
Matrix matrixPower(const Matrix& val, 11 exp) {
    if (exp == 1) return val;
    Matrix res = matrixPower(val, exp / 2);
   Matrix ret = res * res:
```

```
if (exp & 1) ret = ret * val;
    return ret:
int main() {
    ll n; cin >> n;
    if (n == 1) {
        cout << 1:
        return 0;
    // Base Matrix
    Matrix base:
    base.a.resize(2, vector<11>(2));
    base.a[0][0] = base.a[0][1] = base.a[1][0] = 1;
    // Matrix Exponentiation
    Matrix ans = matrixPower(base, n - 1):
    cout << ans.a[0][0];
}
8.8 Catalan Number, Derangement Number
// Catalan Number
// BOJ 9343 AC Code
// https://www.acmicpc.net/problem/9343
#include <bits/stdc++.h>
using namespace std;
#define ll long long
const int MOD = 1e9 + 7:
const int MAXN = 2020202;
ll fac[MAXN]. inv[MAXN]. facInv[MAXN]:
11 catalanNumber(int n) { // Cn = 2nCn / (n + 1) = (2n)! / (n!(n + 1)!)
    return fac[2 * n] * facInv[n] % MOD * facInv[n + 1] % MOD;
int main() {
    cin.tie(NULL); cout.tie(NULL);
    ios_base::sync_with_stdio(false);
    // Preprocessing in O(N)
    fac[0] = fac[1] = inv[1] = 1;
    facInv[0] = facInv[1] = 1:
    for (int i = 2; i < MAXN; i++) {</pre>
        fac[i] = i * fac[i - 1] % MOD;
        inv[i] = -(MOD / i) * inv[MOD % i] % MOD;
        if (inv[i] < 0) inv[i] += MOD;</pre>
        facInv[i] = facInv[i - 1] * inv[i] % MOD;
    // Answer each query in O(1)
    int q; cin >> q;
    while (q--) {
        int n; cin >> n;
        cout << catalanNumber(n) << '\n';</pre>
    }
// Derangement Number
// Counting derangements of a set amounts to the hat-check problem,
// in which one considers the number of ways in which n hats
// can be returned to n people such that no hat makes it back to its owner.
// Recurrence relation:
// f1 = 0, f2 = 1.
// fi = (i - 1) * (f{i-1} + f{i-2}) (i >= 3)
```

```
#include <bits/stdc++.h>
using namespace std:
#define ll long long
const int MOD = 1e9 + 7:
const int MAX = 101010;
11 dp[MAX];
int main() {
    cin.tie(NULL); cout.tie(NULL);
   ios_base::sync_with_stdio(false);
   dp[1] = 0, dp[2] = 1;
   for (int i = 3; i < MAX; i++) {
        dp[i] = (i - 1) * (dp[i - 1] + dp[i - 2]) % MOD;
   }
   int n; cin >> n;
    cout << dp[n]:
8.9 FFT
typedef complex<double> base;
void fft(vector<base> &a, bool inv) {
 int n = a.size(), j = 0;
 vector<base> roots(n / 2):
 for (int i = 1; i < n; i++){
   int bit = (n >> 1);
   while (j >= bit) {
     j -= bit;
     bit >>= 1:
   j += bit;
    if (i < j) swap(a[i], a[j]);</pre>
 double ang = 2 * acos(-1) / n * (inv ? -1 : 1);
 for (int i = 0; i < n / 2; i++){
   roots[i] = base(cos(ang * i), sin(ang * i));
 for (int i = 2: i \le n: i \le 1)
    int step = n / i;
   for (int j = 0; j < n; j += i){
     for (int k = 0; k < i / 2; k++) {
        base u = a[j + k], v = a[j + k + i / 2] * roots[step * k];
        a[i + k] = u + v:
        a[i + k + i / 2] = u - v;
   }
 }
 if (inv) for (int i = 0; i < n; i++) a[i] /= n;
void multiply(const vector<11> &v, const vector<11> &w, vector<11>& res) {
 vector<base> fv(v.begin(), v.end()), fw(w.begin(), w.end());
 int n = 2; while (n < v.size() + w.size()) n <<= 1;</pre>
 fv.resize(n); fw.resize(n);
 fft(fv, 0); fft(fw, 0);
 for (int i = 0; i < n; i++) fv[i] *= fw[i];
 fft(fv. 1):
 res.resize(n);
 for (int i = 0; i < n; i++) res[i] = (ll)round(fv[i].real());</pre>
```

8.10 Gauss-Jordan Elimination

```
// Inverse Matrix
void inverse_matrix(vector<vector<double>> &a){
    int n = a.size():
    int m = n + n:
   for(int i = 0; i < n; ++i)
        for(int j = 0; j < n; ++j)
            a[i].push_back(i==j);
   for(int c = 0, r = 0; c < m && r < n; ++c){
        int p = r; // pivot row
        for(int i = r; i < n; ++i)
            if(a[p][c] < a[i][c])
                p = i:
        if(a[p][c] == 0){ puts("no inverse"); return; };
        for(int i = 0: i < m: ++i)
            swap(a[p][j], a[r][j]);
        double t = a[r][c];
        for(int j = 0; j < m; ++j)
           a[r][i] /= t;
        for(int i = 0; i < n; ++i) if(i != r){
           double t = a[i][c]:
           for(int j = c; j < m; ++j)
                a[i][j] -= a[r][j] * t;
       }
        ++r;
   }
   for(int i=0;i<n;++i,puts(""))</pre>
        for(int i=0:i<n:++i)</pre>
           printf("%lf ",a[i][n+j]);
// Gauss-Jordan Elimination modulo p
vector<int> gauss_mod(vector<vector<int>> &a,int mod){
   vector<int> inv(mod); // modulo inverse 전처리
   inv[1] = 1:
   for(int i = 2; i < mod; ++i)
        inv[i] = mod - (mod/i) * inv[mod%i] % mod:
    int n = a.size():
    int m = a[0].size();
    vector<int> w(m, -1): // i번째 열에 있는 pivot이 몇 번째 행에 있는지 저장
   for(int c = 0, r = 0; c < m && r < n; ++c){
        int p = r; // pivot row
        for(int i = r: i < n: ++i)
            if(a[p][c] < a[i][c])
                p = i:
        if(a[p][c] == 0) continue; // free variable
        for(int j = 0; j < m; ++j)
            swap(a[p][j], a[r][j]);
        w[c] = r;
        int t = a[r][c];
        for(int i = 0: i < m: ++i)
           a[r][j] = a[r][j] * inv[t] % mod;
        for(int i = 0; i < n; ++i) if(i != r){
           int t = a[i][c]:
           for(int j = c; j < m; ++j)
                a[i][j] = (a[i][j] - a[r][j] * t % mod + mod) % mod;
       }
        ++r;
   }
```

```
for(int i = 0; i < n; ++i) // existence of solution</pre>
        if(count(a[i],begin(), --a[i],end(), 0) == m-1 && a[i][m-1])
            return vector<int>(); // no solution
    vector<int> ans(m):
   for(int i = 0; i < m; ++i)
        if(w[i]) ans[i] = a[w[i]][m-1];
   return ans: // solution exist
// Gauss-Jordan Elimination modulo 2
const int sz = 500:
bitset<sz> gauss bit(vector<bitset<sz>> &a){
   int n = a.size():
   int m = a[0].size():
   vector<int> w(m, -1);
   for(int c = 0, r = 0; c < m && r < n; ++c){
        for(int i = r; i < n; ++i)
            if(a[i][c]){
                swap(a[i],a[r]);
                break;
        if(a[r][c] == 0) continue:
        w[c] = r;
        for(int i = 0; i < n; ++i) if(i != r)
            if(a[i][c]) a[i] ^= a[r];
        ++r;
   }
    // .. same
9 Misc
9.1 DP Opt
// 1. Convex Hull Trick
// Recurrence: dp[i] = min(dp[j] + a[i] * b[j]) (j < i)
// Condition: b[i] >= b[i + 1]
// Naive Complexity: O(n^2)
// Optimized Complexity: O(nlogn) (if a[i] <= a[i + 1], it can also be done in O(n))
// BOJ 13263 AC Code
// https://www.acmicpc.net/problem/13263
#include <bits/stdc++.h>
using namespace std;
#define 11 long long
struct Line \{ // f(x) = px + q, x >= s \}
   11 p, q;
   double s:
   Line(): Line(1, 0) {}
   Line(ll sp, ll sq): p(sp), q(sq), s(0) {}
double cross(const Line& u, const Line& v) {
   return (double)(v.q - u.q) / (u.p - v.p);
int n:
ll a[101010], b[101010];
ll dp[101010];
Line ch[101010]:
void input() {
    cin >> n:
    for (int i = 1: i <= n: i++) cin >> a[i]:
```

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```
for (int i = 1; i \le n; i++) cin >> b[i];
}
void convexHullTrick() {
    int top = 1:
    for (int i = 2; i <= n; i++) {
        Line g(b[i-1], dp[i-1]);
        while (top > 1) {
            g.s = cross(ch[top - 1], g);
            if (ch[top - 1].s < g.s) break;</pre>
            --top;
        }
        ch[top++] = g;
        int 1 = 1, r = top - 1;
        while (1 < r) {
            int mid = (1 + r + 1) >> 1;
            if (a[i] < ch[mid].s) r = mid - 1;
            else 1 = mid;
        }
        int fpos = 1;
        dp[i] = ch[fpos].p * a[i] + ch[fpos].q;
    }
}
int main() {
    cin.tie(NULL); cout.tie(NULL);
    ios_base::sync_with_stdio(false);
    input():
    convexHullTrick();
    cout << dp[n]:
// 2. Knuth Optimization
// Recurrence: DP[i][j] = min(DP[i][k] + DP[k + 1][j]) + C[i][j] (i <= k < j)
// Condition: C[i][j] is a monge array (satisfies C[a][c] + C[b][d] \leftarrow C[a][d] + C[b][c]),
              and satisfies C[a][d] >= C[b][c] for a <= b <= c <= d
// Naive Complexity: O(n^3)
// Optimized Complexity: O(n^2)
// Letopt[i][j] be the value of k that minimizes DP[i][j]
// The following holds: opt[i][j-1] \le opt[i][j] \le opt[i+1][j]
// BOJ 13974 AC Code
// https://www.acmicpc.net/problem/13974
#include <bits/stdc++.h>
using namespace std;
#define ll long long
const 11 INF = 1e18;
int n, opt[5050][5050];
ll a[5050], DP[5050][5050], psum[5050];
int main() {
    cin.tie(NULL): cout.tie(NULL):
    ios_base::sync_with_stdio(false);
    int tc; cin >> tc;
    while (tc--) {
        cin >> n:
        for (int i = 1; i <= n; i++) {
            cin >> a[i];
            psum[i] = a[i] + psum[i - 1];
```

```
for (int i = 1; i <= n; i++) {
           DP[i][i] = 0:
            opt[i][i] = i;
       for (int i = n - 1; i >= 1; i --) {
            for (int j = i + 1; j \le n; j++) {
               11 mn = INF, mnk = -1;
                for (int k = opt[i][j-1]; k \le opt[i+1][j]; k++) {
                    ll res = DP[i][k] + DP[k + 1][j] + (psum[j] - psum[i - 1]);
                    if (res < mn) {
                        mn = res:
                        mnk = k;
                    }
                }
                DP[i][j] = mn;
                opt[i][j] = mnk;
           }
       }
        cout << DP[1][n] << '\n';
   }
}
// 3. Divide and Conquer Optimization
// Recurrence: dp[t][i] = min(dp[t - 1][j] + c[j][i]) (j < i)
// Condition: Let opt[t][i] be j with the smallest value of dp[t-1][j]+c[j][i]. It must
satisfy opt[t][i] \le opt[t][i + 1].
// Naive Complexity: O(m * n^2)
// Optimized Complexity: O(m * nlogn)
// BOJ 13261 AC Code
// https://www.acmicpc.net/problem/13261
#include <bits/stdc++.h>
using namespace std;
#define ll long long
const ll INF = 1e18;
int n, m;
ll a[8080], psum[8080];
11 dp[808][8080];
void f(int gr, int l, int r, int nl, int nr) {
    int mid = (1 + r) >> 1, idx = -1;
   11& res = dp[gr][mid];
   res = INF:
   for (int i = nl; i <= min(mid, nr); i++) {</pre>
        assert(i <= mid):
       ll val = dp[gr - 1][i] + (mid - i) * (psum[mid] - psum[i]);
        if (res > val) {
           res = val, idx = i;
   }
   if (1 < r) {
        f(gr, 1, mid, nl, idx);
        f(gr, mid + 1, r, idx, nr);
   }
int main() {
   cin.tie(NULL); cout.tie(NULL);
   ios_base::sync_with_stdio(false);
   // input
```

}

}

```
cin >> n >> m;
    for (int i = 1; i \le n; i++) cin >> a[i]:
    // build prefix sum
    for (int i = 1; i \le n; i++)
        psum[i] = a[i] + psum[i - 1];
    // dp (dnc opt)
   for (int i = 1: i \le n: i++)
        dp[1][i] = i * psum[i];
    for (int i = 2; i <= m; i++)
        f(i, 0, n, 0, n);
    // output
    cout << dp[m][n];
9.2 Sqrt Decomposition, Mo's Algorithm
#include <bits/stdc++.h>
using namespace std;
int sq;
struct se {
 int s, e, idx;
 bool operator<(const se& rhs) const {
   if (s / sq != rhs.s / sq) return s / sq < rhs.s / sq;
    return e < rhs.e:
// Zigzag Mo's (faster than basic Mo's Algorithm)
/*struct se {
 int s. e. idx:
 bool operator<(const se &rhs) const {
   if(s / sq != rhs.s / sq) return s / sq < rhs.s / sq;
   else return (s / sq) & 1 ? e < rhs.e : e > rhs.e:
 }
};*/
vector<se> q;
vector<int> ans;
void input() {
 // TODO: 1. receive input 2. resize q, ans 3. calculate sq
void add(int idx) {
 // TODO: add value at idx from data structure
void del(int idx) {
 // TODO: remove value at idx from data structure
 // TODO: extract the current answer of the data structure
void f() {
 int s = q[0].s, e = q[0].e;
  // TODO: initialize data structure
  ans[q[0].idx] = query();
  for (int i = 1; i < q.size(); i++) {</pre>
    while (q[i].s < s) add(--s);
    while (e < q[i].e) add(++e);
    while (s < q[i].s) del(s++);
    while (q[i].e < e) del(e--);
    ans[q[i].idx] = query();
```

```
int main() {
 cin.tie(NULL); cout.tie(NULL);
 ios_base::sync_with_stdio(false);
 input();
  sort(q.begin(), q.end());
 f():
 for (auto& i : ans)
    cout << i << '\n';
9.3 Rotation Matrix, Manhattan Distance, Chebyshev Distance
Rotation Matrix (In Two Dimensions)
x' = x * cos\theta - v * sin\theta
y' = x * sin\theta + y * cos\theta
Manhattan Distance, Chebyshev Distance
For any two points (x1, y1), (x2, y2) on two-dimensional coordinates,
- the Manhattan distance is |x1 - x2| + |y1 - y2|.
- the Chebyshev distance is max(|x1 - x2|, |y1 - y2|).
Relationship between Manhattan distance and Chebyshev distance
Let x1' = x1 - y1, y1' = x1 + y1,
x2' = x2 - y2, y2' = x2 + y2. (using Rotation matrix)
Then max(|x1 - x2|, |y1 - y2|)
= (|x1' - x2'| + |y1' - y2'|) / 2.
Thus, Manhattan distance can be solved by replacing it with Chebyshev distance, and vice versa.
9.4 Random
#include <random>
using namespace std;
const int MAX = 1'000;
srand(time(NULL));
int x = rand() % MAX;
9.5 Ternary Search
void ternarySearch() {
   int l = 1, r = 10'000;
    while (r - 1 >= 3) {
        int mid1 = (2 * 1 + r) / 3;
        int mid2 = (1 + 2 * r) / 3:
        int res1 = f(mid1), res2 = f(mid2);
        if (res1 <= res2) l = mid1;</pre>
        else r = mid2:
   }
   int res = 0;
   for (int i = 1; i <= r; i++)
        res = max(res, f(i));
```