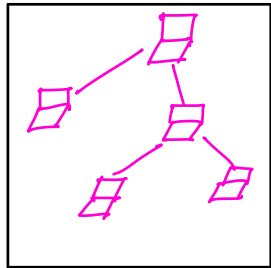


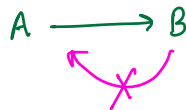
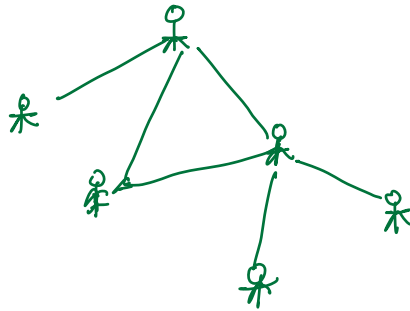
graphs :- network



collection of nodes and edges



social media



undirected / directed

Scaler.com

Home

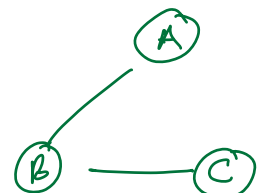
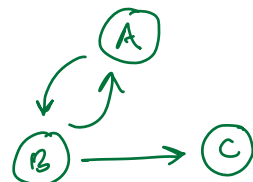
classroom

HR

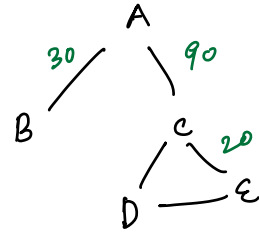
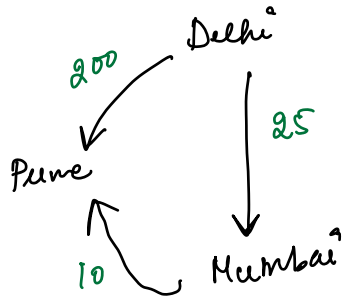
Mentor



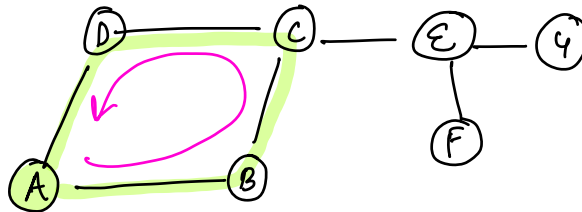
Tree - graph?



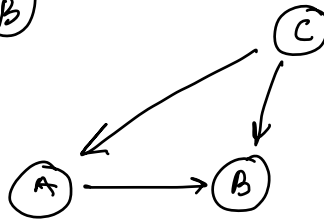
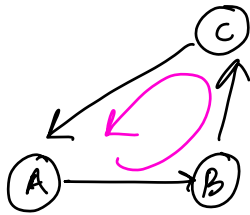
single graph  
unweighted / weighted



weighted / unweighted

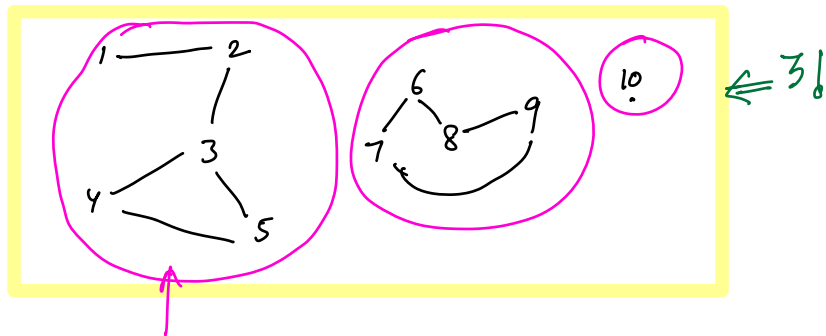


cyclic graph

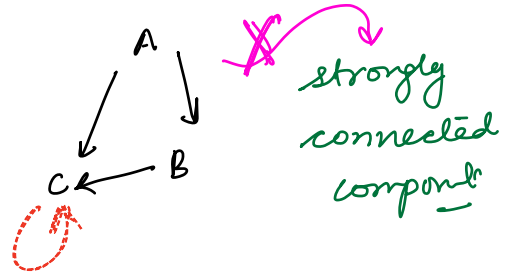


acyclic graph

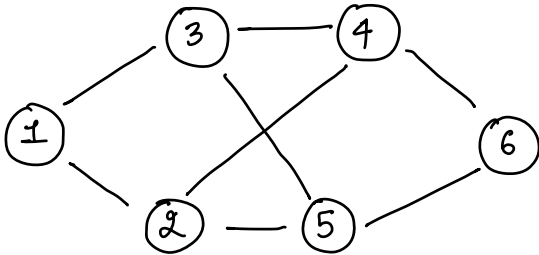
If you can start from a node and come back to the same node without covering any edge twice = cycle



Undirected  $\Leftarrow$  Connected Components  
 every node is reachable from other nodes



## Representation of graph



$$n = 6$$

$$\text{edges} = 8$$

$$1-3 \quad 2-5$$

$$1-2 \quad 2-4$$

$$3-4 \quad 5-6$$

$$3-5 \quad 4-6$$

weighted  $\Rightarrow$  weight in  
replacement of 1

### Advantage

• access -  $O(1)$

• update -  $O(1)$   
of edge

• waste of space

```
int n; input(n)
int m; input(m)
for (int i=0; i<m; i++)
{
    int u, v; input(u, v);
    mat[u][v] = 1;
    mat[v][u] = 1;
}
```

$$\begin{cases} 1 \rightarrow n \\ 0 \rightarrow n-1 \end{cases}$$

### Adjacency Matrix

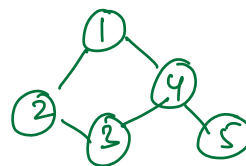
$$(n+1) \times (n+1)$$

$$1 \rightarrow n$$

	0	1	2	3	4	5	6
0							
1	0	0	1	1		0	0
2		1				1	
3		1			1	1	
4				1			1
5			1	1			1
6					1	1	

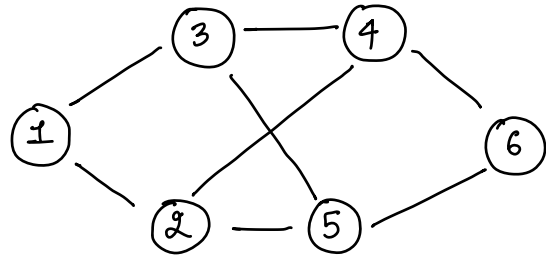
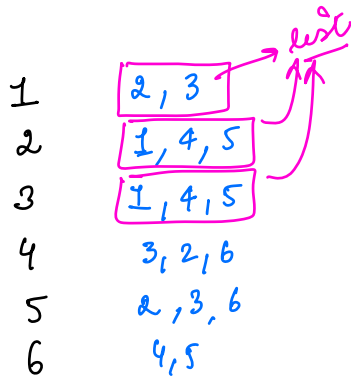
mat[i][j]  $\rightarrow$  0  $i \neq j$   
 $\rightarrow$  1  $i = j$

dense sparse



5  
5  
1 2  
2 3  
3 4  
4 5  
1 4

## # Adjacency list



```

hashmap < int, list > graph;
list < int > graph(n+1);
    
```

```

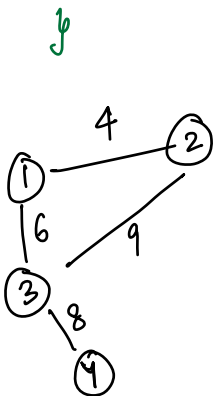
int n;
int m;
list < int > graph(n+1);
    
```

```

for (i=0; i<m; i++)
{
    int u, v;
    graph[u].insert(v);
    graph[v].insert(u);
}
    
```

HashSet < int > graph(n+1);  
 when you want to have  
 info about the  
 edge.

```
graph[u].insert(v);
```



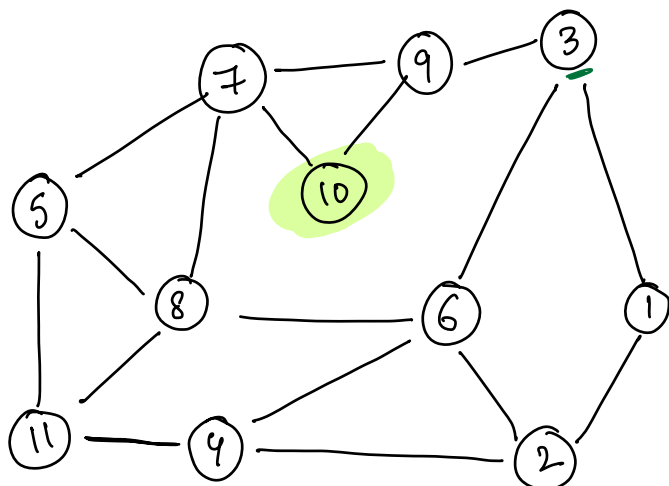
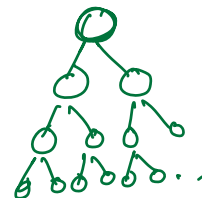
```

1  {2, 4} {3, 6}
2  {1, 4} {3, 9}
3  {1, 6} {4, 8}
4  {3, 8}
    
```

```
list < pair > graph(n+1);
```

Traversals  $\begin{cases} \rightarrow \text{BFS} \\ \rightarrow \text{DFS} \end{cases}$

BFS: Breadth first search } level order



~~10~~ ~~7~~ ~~9~~ 5 8 10 9 7 10 3 6

bool visited [n+1];

0	1	2	3	4	5	6	7	8	9	10	11
F	F	F	F	F	F	F	F	F	F	F	F
T	T	T	T	T	T	T	T	T	T	T	T

10 9 7 3 5 8 9 new  
6 1 11 4 2

~~10~~ ~~9~~ ~~7~~ ~~3~~ ~~5~~ ~~8~~ ~~6~~ ~~1~~ ~~11~~ ~~4~~ ~~2~~

```
queue <int> q;
bool visited (n+1); // initially false
```

$1 \rightarrow n$

```
for(int source=1; source<=n; source++)
```

```
{ if(visited[source]) continue;
```

```
q.push(source);
visited[source] = true;
```

```
while(!q.empty())
```

```
{
```

```
int u = q.front();
```

```
q.pop();
```

```
print(u);
```

```
for(int i=0; i<graph[u].size(); i++)
```

```
{
```

```
int v = graph[u][i];
```

```
if(!visited[v])
```

```
{
```

```
visited[v] = true;
```

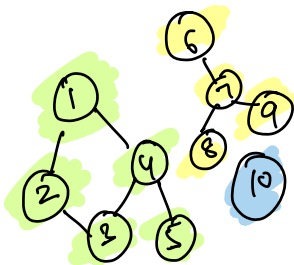
```
q.push(v);
```

```
}
```

```
}
```

```
}
```

```
}
```



T.C:  $O(n+m)$

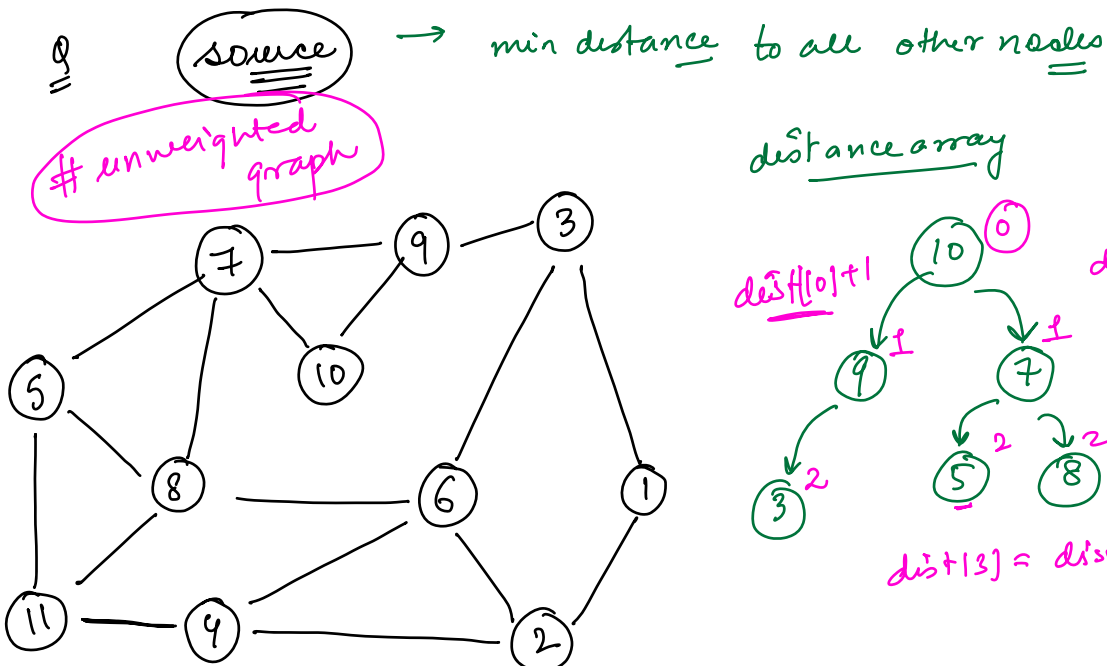
no. of edges

S.C:  $O(n)$

source  $\longrightarrow$  destination  $\xrightarrow{\text{valid}}$

{ go from source to destination }

start doing BFS from source  
if destination is already visited  $\Rightarrow$  True





```
queue <int> q;  
bool visited [n+1]; // initially false  
int dist [n+1];
```

```
dist[source] = 0;  
q.push(source);  
visited[source] = true;
```

```
while (! q.empty())  
{  
    int u = q.front();  
    q.pop();  
    print(u);  
    for (int i=0; i < graph[u].size(); i++)  
    {  
        int v = graph[u][i];  
        if (! visited[v])  
        {  
            dist[v] = dist[u] + 1;  
            visited[v] = true;  
            q.push(v);  
        }  
    }  
}
```