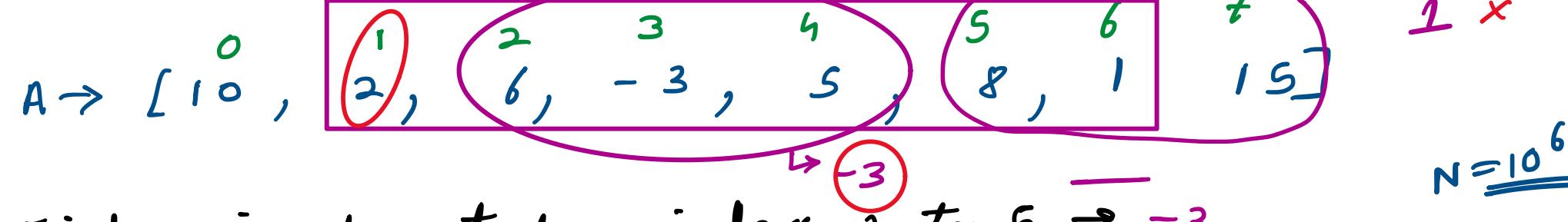


Basics of Segment Tree

Friday, 14 January 2022

9:06 PM



$$N = 10^6$$

Find min element from index 2 to 5 $\rightarrow -3$

Find max element from index 1 to 6 $\rightarrow 8$

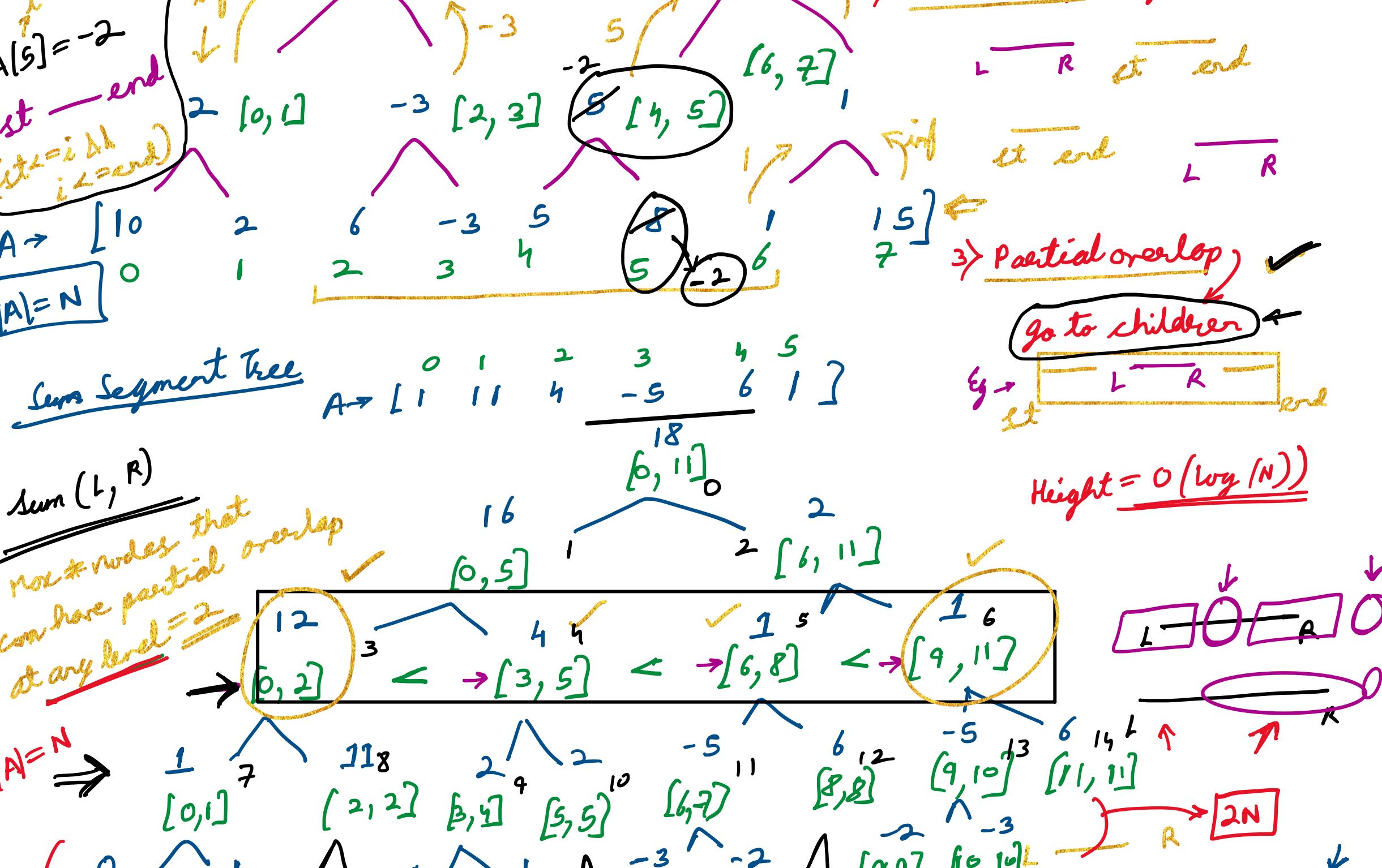
Find sum of elements from index 3 to 5 $\rightarrow 10$

$$T.C. = \frac{20}{\log N} \approx 10^6$$

Bruteforce \rightarrow linear traversal over the range $\Rightarrow T.C. = O(\text{Range}) = O(N)$

Multiple Queries \rightarrow Find for a range $\Rightarrow T.C. = O(N \times Q) = O(Q \times \log N)$

Min Segment Tree



Sum Segment Tree

$$A \rightarrow [10, 2, 6, -3, 5, 8, 1, 15]$$

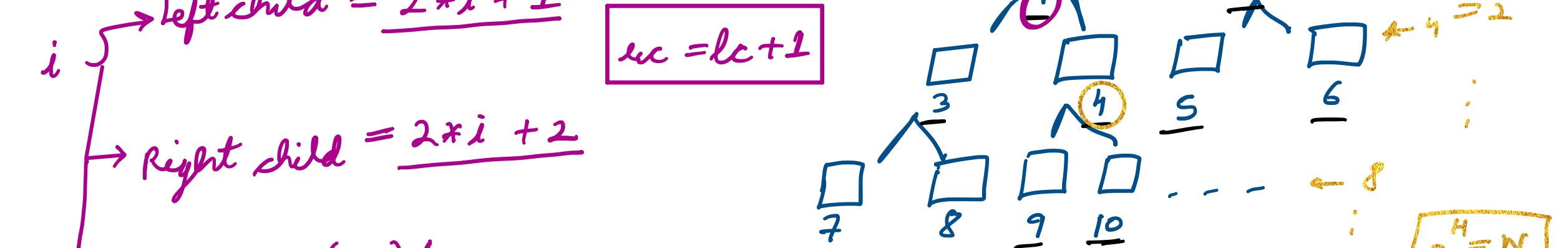
$$T.C. = \frac{20}{\log N}$$

sum(L, R)

Non # nodes that have partial overlap at any level = 2

$$T.C. \text{ per query} = O(\log N)$$

$$\text{Height} = O(\log N)$$



$$T.C. \text{ per query} = O(\log N)$$

$$(0-1)(2-3)(4-5)\dots$$

Build Segment Tree

0 1 2 3 4 5 6 7 ...

$$i \rightarrow \text{left child} = 2 \times i + 1$$

$$ec = lc + 1$$

$$i = 1 \rightarrow 2^0$$

$$i \rightarrow \text{right child} = 2 \times i + 2$$

$$i = 2 \rightarrow 2^1$$

$$i \rightarrow \text{parent} = \frac{i-1}{2}$$

$$i = 4 \rightarrow 2^2$$



$$i = 1 \rightarrow 2^0 \quad lc = 2 \times 1 + 1 = 3$$

$$parent = \frac{1-1}{2} = 0 \quad ec = 2 \times 1 + 2 = 4$$

$$i = 4 \rightarrow 2^3 \quad lc = 2 \times 4 + 1 = 9$$

$$ec = 2 \times 4 + 2 = 10$$

$$parent = \frac{4-1}{2} = \frac{3}{2} = 2$$

$$|A| = N$$

$$|\text{tree}[1]| = 4N$$

$$SC = O(N)$$

$$\Rightarrow 1 + 2 + 4 + 8 + \dots = 2^H$$

$$\Rightarrow 2^0 + 2^1 + 2^2 + \dots = \frac{2^0(2^{H+1}-1)}{2-1} = 2^{H+1} = 2 \cdot 2^H - 1 = 2N-1$$

Build $\text{void build}(idx, st, end)$ // idx \rightarrow index in tree array

if(st == end) { $\text{tree}[idx] = A[st]$; } // leaf node

else { mid = (st + end)/2; $lc = 2 \times idx + 1$; $ec = 2 \times idx + 2$; build(lc, st, mid); build(ec, mid+1, end); }

$\text{tree}[idx] = \min(\text{tree}[lc], \text{tree}[ec])$; // max/min/avg/gcd

$T.C. = O(N)$

Query $\text{int query}(idx, st, end, L, R)$ // complete overlap

$L \leq st \leq end \leq R$ return $\text{tree}[idx]$; // non-overlap

$mid = (st + end)/2$; $lc = 2 \times idx + 1$; $ec = 2 \times idx + 2$;

return $\min(\text{query}(lc, st, mid, L, R), \text{query}(ec, mid+1, end, L, R))$; // partial overlap