

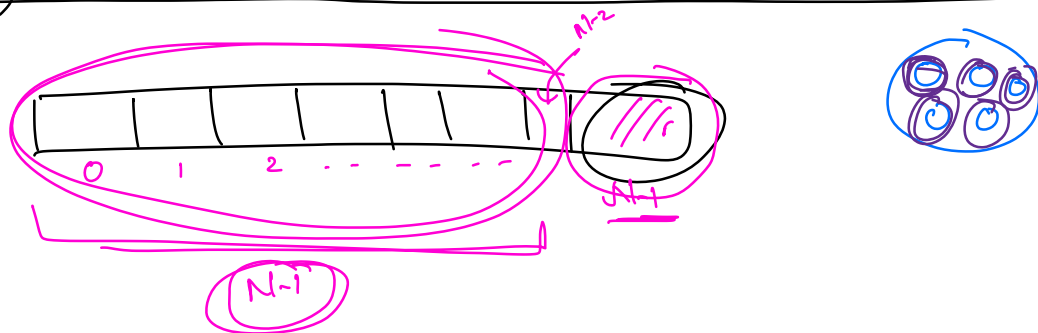
```

int SumCoins ( Coins[] , N ) {
    if (N == 0) return 0;

    return Coins[N-1] + SumCoins(Coins[], N-1);
}

```

0-1



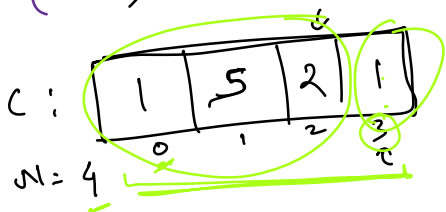
Recursion : A function calling itself.

```
int SumCoins ( Coins[] , N ) {
```

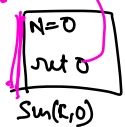
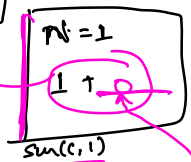
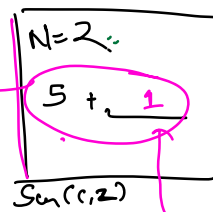
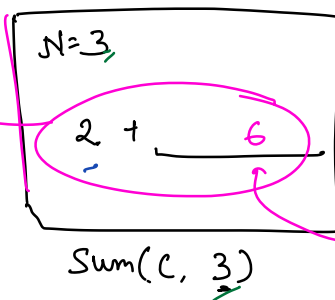
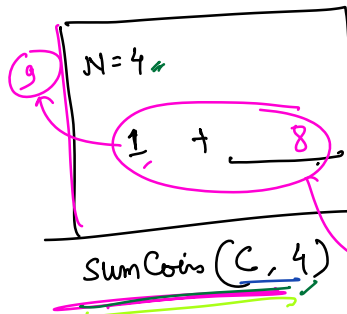
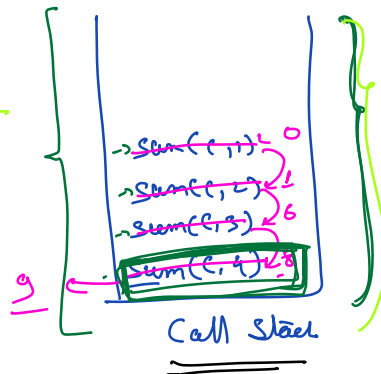
```
    if (N == 0) return 0;
```

```
    => return coins[N-1] + SumCoins (coins, N-1);
```

↪ Assembly/Trust



① + Sum()



LIFO

Stack

Call-Stack / Recursion Stack

$$TC = \frac{\text{Total}}{\text{No. of function calls}} \times TC \text{ of one function call} \left. \begin{array}{l} O(N) \\ \times \\ O(1) \end{array} \right\} = O(N)$$

$SC = \left[\begin{array}{l} \text{The max size of stack} \\ \text{SC of every function call} \end{array} \right] \Rightarrow O(N) \times O(1) = O(N)$
 Depth of rec tree.

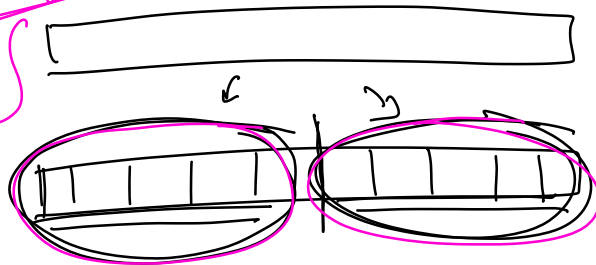
Scope

$$= \left[\text{Coins}[N-1] + \frac{\text{SumCoin}(\text{Coins}, N-1)}{\text{Sub problem.}} \right]$$



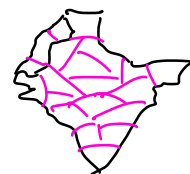
SubProblem > 1

Merge Sort



No overlap

Divide & Conquer



Fibonacci

$\left(\begin{smallmatrix} 0 \\ 0 \end{smallmatrix} \right), \left(\begin{smallmatrix} 1 \\ 1 \end{smallmatrix} \right), \frac{1}{2}, 2, 3, 5, 8, 13, \dots$

Print N^{th} Fibonacci No.

$$\text{fib}(N) = \text{fib}(N-1) + \text{fib}(N-2);$$

```

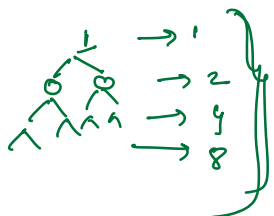
int fib(N) {
  if (N <= 1) return N;
  return fib(N-1) + fib(N-2);
}
  
```

L -

Dynamic Program

$$TC = \text{No of fn calls} \times TC \text{ of 1 fn call}$$

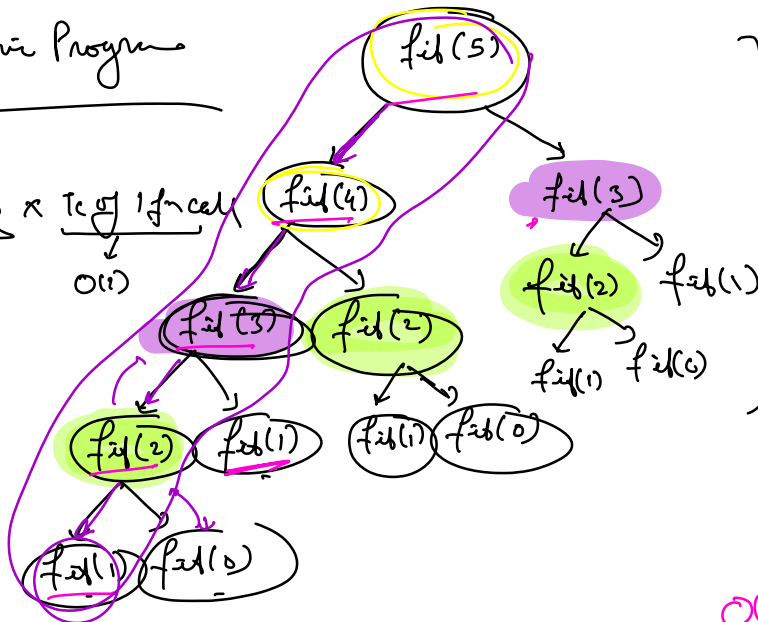
$O(1)$



$$1 + 2 + 4 + 8 + 16 + 32 \dots$$

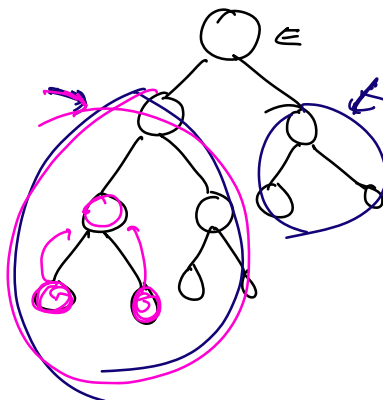
$$\sum_{i=0}^{n-1} 2^i = \frac{2^n - 1}{2 - 1} = \frac{2^n - 1}{1} = 2^n - 1$$

Overlapping sub problem

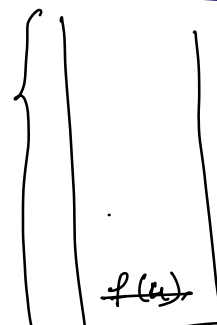
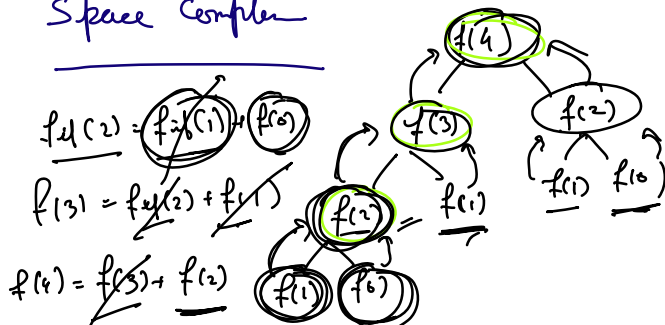


~~fib(1)~~
~~fib(2)~~
~~fib(3)~~
~~fib(4)~~
~~fib(5)~~

$O(2^n)$



Space Complex



Size: 4

ManSig: 4

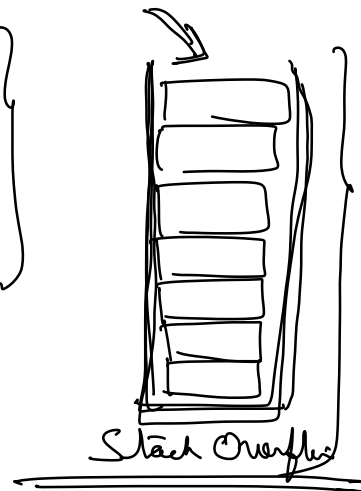
3 Steps

- o Assume & trust: That recursion will work for smaller sub-problem.
- o Recursive eq: Ret the ans of bigger problem using the ans of sub problem.
(logic)
- o Base/Termination Condition

```
while (true) {  
    print(0);  
    ,  
}
```

0
6
0
6
0
6
1
;

```
void printZero() {  
    print(0);  
    printZero();  
}
```



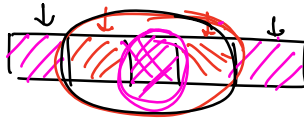
Q Given a string. Check if it is palindrome.

* Recursively

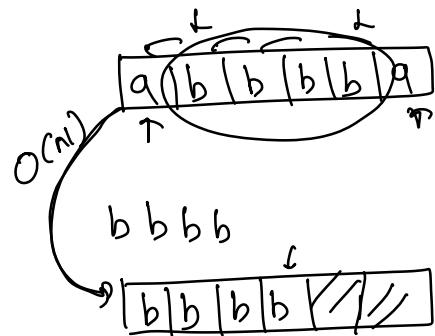
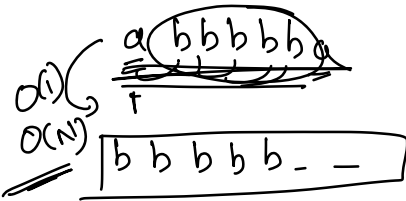
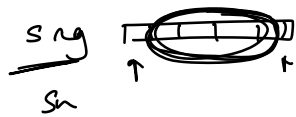
↳ ARORA
↳ MADAM
↳ NITIN
↳ CTC

IsPalin (str, s, e)

Assume ; IsPalin (smaller str) will give correct output.

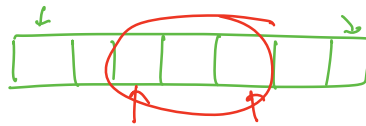


Logic/Rec eq :



if (first & last char are same)
 ret IsPal (str b/w them)

else
 ret false;



boolean isPal (str, s, e) {

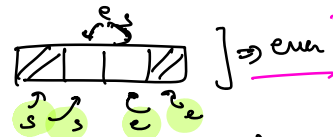
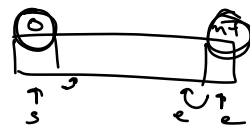
→ if (s >= e) ret true ;

if (str[s] == str[e])

ret isPal (str, s+1, e-1) ;

else

ret false ;



isPal (str, 0, n-1)

Print all elements of an array

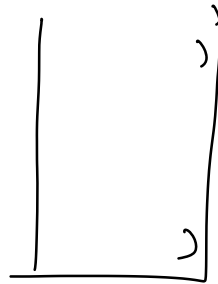
```
void PrintArr ( A[], N, s ) {
    if ( s == N ) return;

```

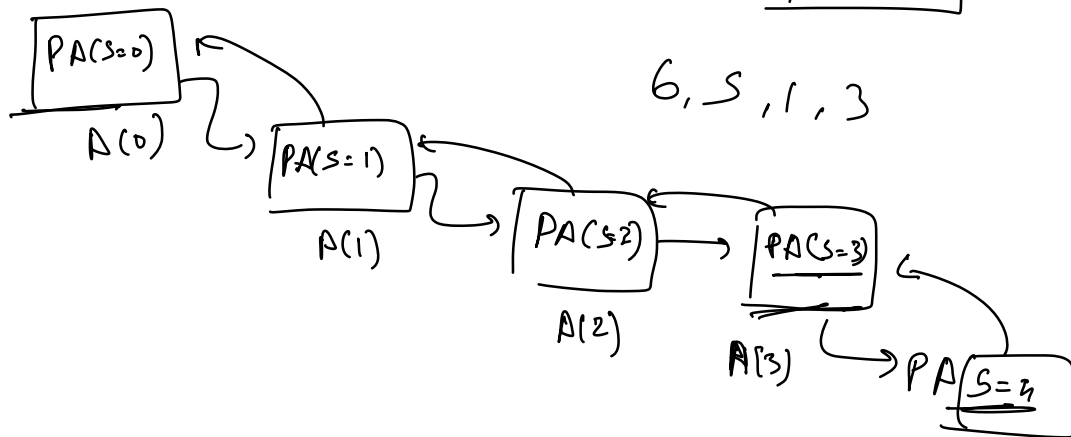
```
    Print ( A[s] );
    PrintArr ( A, N, s+1 );
}

```

3, 4, 5, 6
0 1 2 3
N=4



6, 5, 1, 3

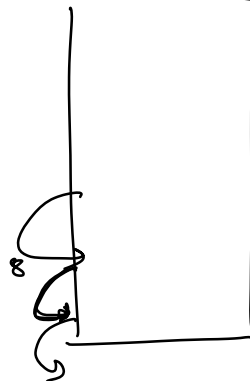


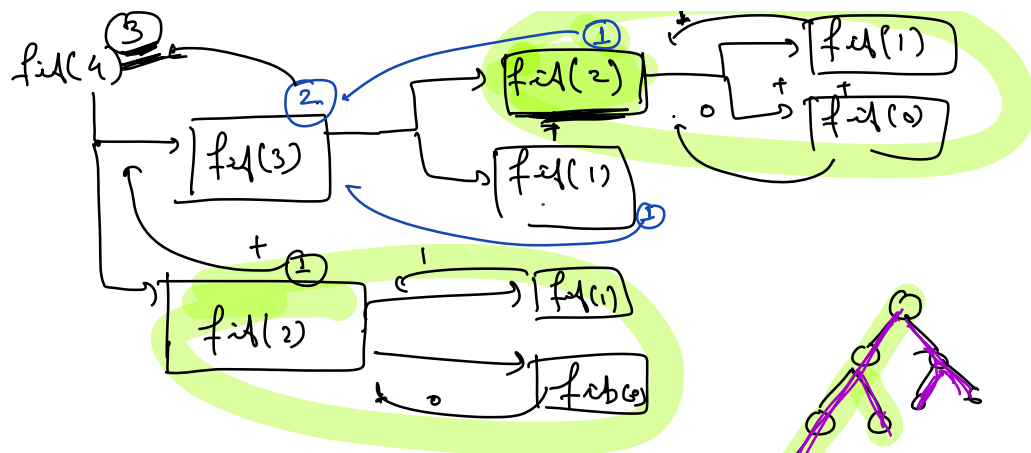
```
int add ( a, b ) {
    return a+b;
}

```

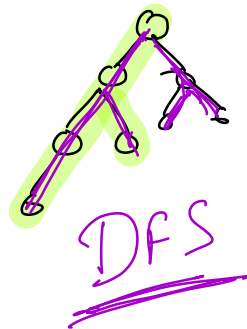
```
main ( ) {
    Print ( add ( 5, 3 ) );
}

```





0, 1, 1, 2, 3
 1, 2, 3, 4



$$\text{fib}(4) = \text{fib}(3) + \text{fib}(2)$$

$\text{fib}(3)$ $\text{fib}(2)$

$\text{fib}(2) + 1$
 $\text{fib}(1)$

