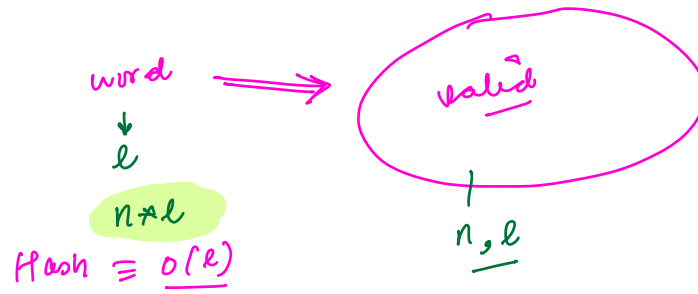


spelling

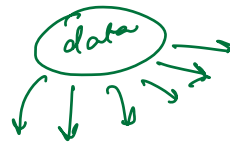
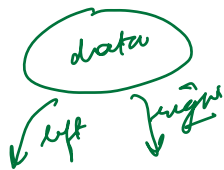


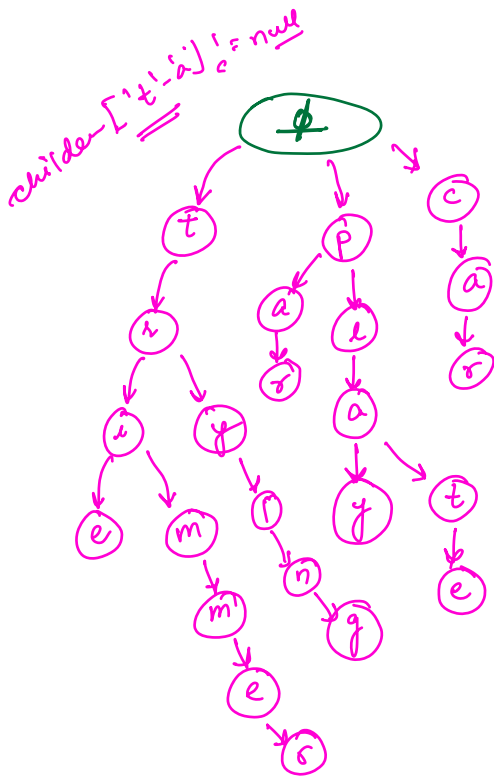
autocomplete

play      playground  
            play  
            player  
            ⋮

Trie is tree based DS which stores the info from top to bottom  
    ↳ prefix-trie

Trie	Try	Trim	play
plate	car	pal	trimmer
trying	pla		





```

class Node {
    char data;
    Node children [26];
    bool isEnd;
}

```

'a' → 0  
 'b' → 1  
 'c' → 2  
 'd' → 3  
 ⋮  
 's'

ch - 'a'

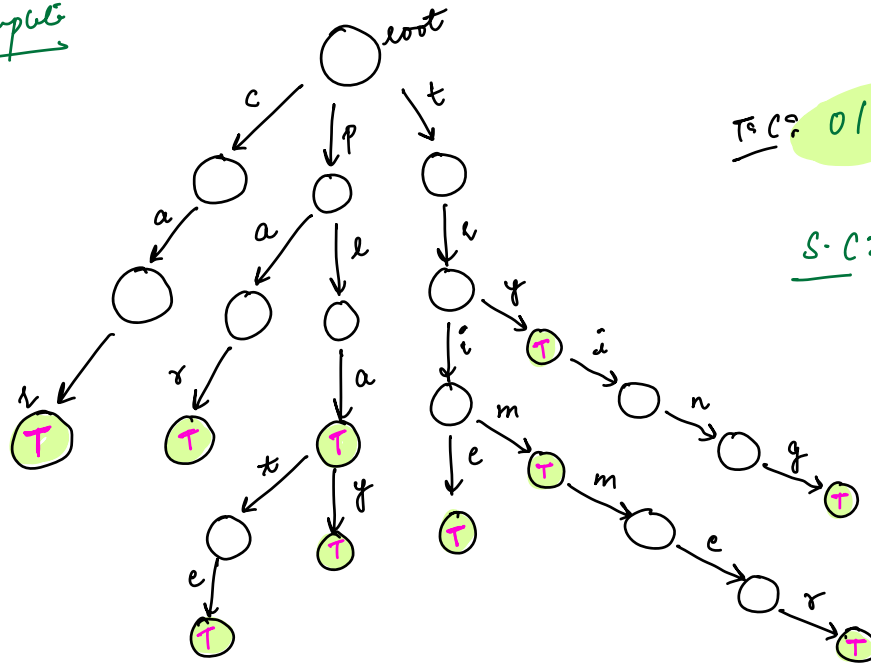
search(trap)  
(tri)  
pla

Traverse the string & search char by char in the trie

if a character's node doesn't exist  
↓  
word is invalid

find all characters  
 ↓  
 last chara node's isEnd = true → valid  
 ↓  
 isEnd = false → invalid

auto complete  
tr



$T.C: O(l)$

$S.C: (n) \times l \rightarrow 26$

If a word can be present multiple times!

Node d

char data;

~~bool isEnd;~~

Node child[26];

};

int freq;

how many words end at a particular node  
(freq > 0);

```
void insert ( Node root, string word)
```

```
{
```

```
    curr = root;
```

```
    L = word.length();
```

```
    for (i=0; i<L; i++)
```

```
    {
```

```
        int index = word[i] - 'a';
```

```
        if ( curr->children[index] == null)
```

```
        {
```

```
            curr->children[index] = new Node( word[i]);
```

```
        }
```

```
        curr = curr->children[index];
```

```
    }
```

```
    curr->freq++;
```

```
    // curr->isEnd = true;
```

```
}
```

```
bool search ( Node root, string word)
{
```

```
    curr = root;
```

```
    L = word.length();
```

```
    for (i=0; i<L; i++)
```

```
    {
```

```
        int index = word[i] - 'a';
```

```
        if ( curr->children[index] == null)
```

```
        {
```

```
            return false;
```

```
        }
```

```
        curr = curr->children[index];
```

```
    }
```

```
    if ( curr->freq > 0) return true;
    return false;
```

```
// End
```

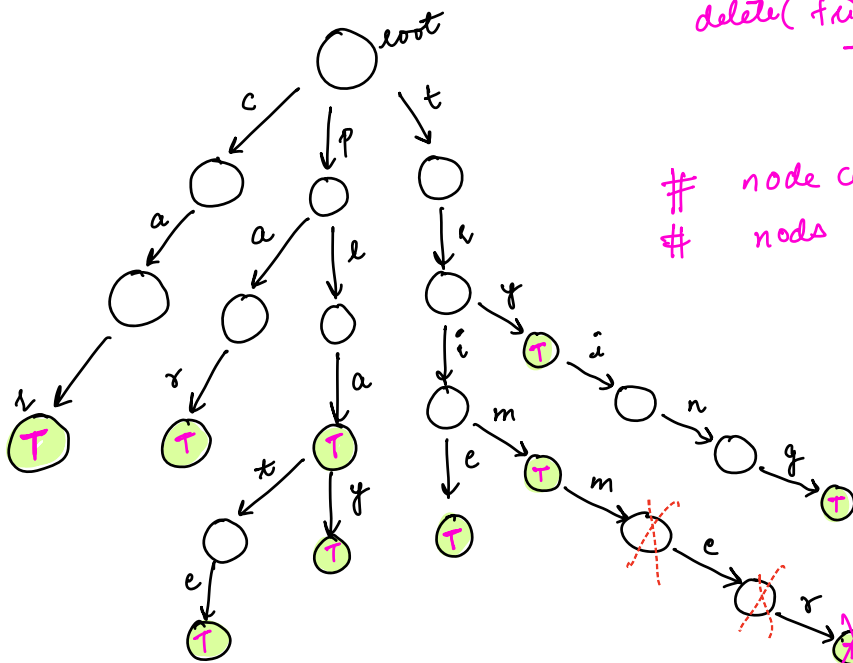
```
}
```

# deletion

last node = root ~~is not~~ m

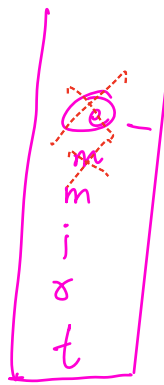
delete(trimmer)  
par try

# node completes word  
# node has children



~~False~~, fig--

$S.C: O(1)$



isEnd = false

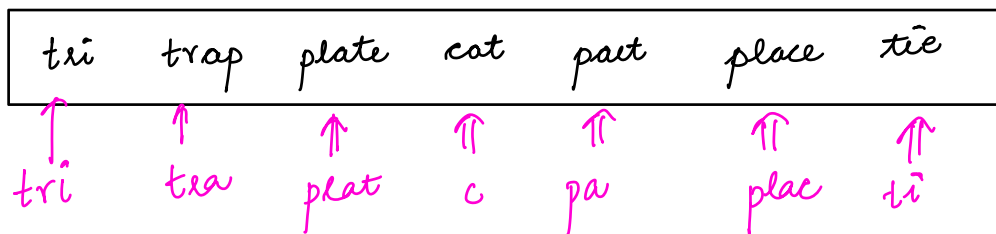
cur  
st.top().children[cur.data - 'a'] = null;

# get the last node which  
you can't delete at all

# # shortest unique prefix

Q Find shortest unique prefix to represent each word.

Note: Assume that no word is prefix of another  
In other words, the representation is always possible.



# "\_" → how many words is which it is prefix

tc: creation  $O(n \times k)$   
 for 1 query =  $O(k)$   
 for T.C.  $O(n \times k)$

