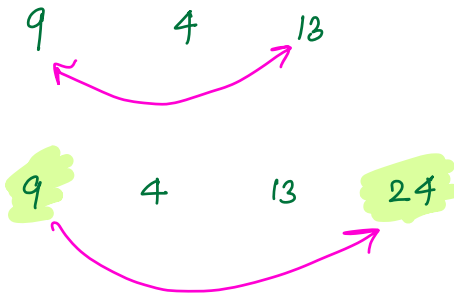
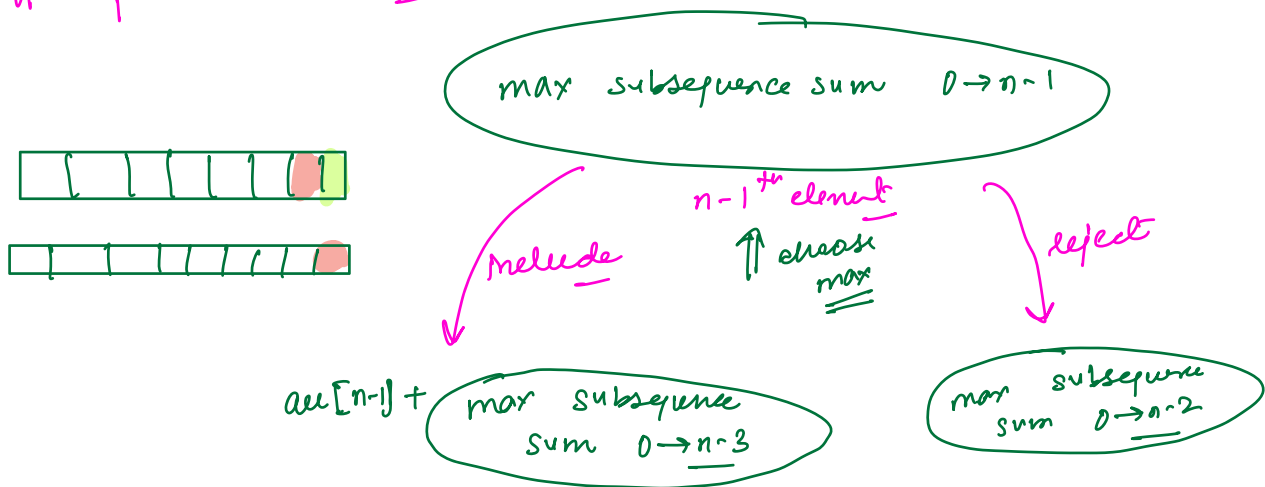


Q Given an array, Find maximum subsequence sum.  
can't pick adjacent elements

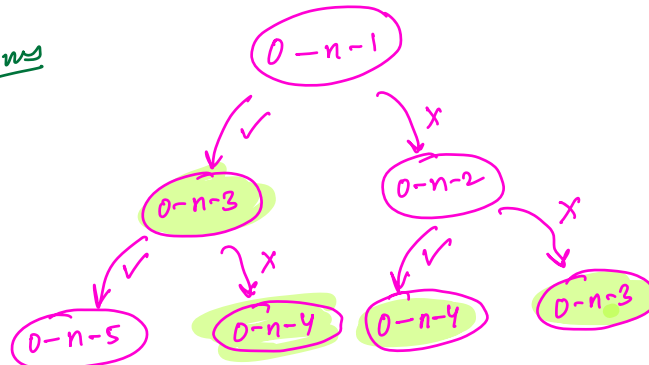


Bf/Basic idea: generate all valid subsequences

# optimal substructure



overlapping subproblems



$$\text{maxsum}(i) = \max(\text{arr}[i] + \text{maxsum}(i-2), \text{maxsum}(i-1))$$

$\text{int dp}[n]; \Rightarrow$  pick initial value acc to constraints  
 $\text{int maxsum}(\text{int arr}[], \text{int index})$  pick  
 $\{$   
 $\text{if}(\text{index} == 0) \text{ return } \max(\text{arr}[0], 0);$  something which  
 $\text{if}(\text{index} == 1) \text{ return } \max(\text{arr}[0], \text{arr}[1], 0);$  definitely  
 $\text{if}(\text{dp}[\text{index}] \neq -1)$  can't be  
 $\text{return dp}[\text{index}];$  your ans!  
 $\text{dp}[\text{index}] = \max(\text{arr}[i] + \text{maxsum}(\text{index}-2), \text{maxsum}(\text{index}-1));$   
 $\text{return dp}[\text{index}];$   
 $\}$

iterative

$\text{dp}[0] = \max(0, \text{arr}[0]);$   
 $\text{dp}[1] = \max(0, \text{arr}[0], \text{arr}[1]);$   
 $\text{for}(\text{int } i = 2; i < n; i++)$   
 $\{$   
 $\text{dp}[i] = \max(\text{arr}[i] + \text{dp}[i-2], \text{dp}[i-1]);$   
 $\}$   
 $\text{ans} = \text{dp}[n-1]$

Q2.  $2 \times N$  matrix - find max sum you can get from matrix.

can't pick adjacent elements  $\rightarrow$  horizontally  
 $\rightarrow$  diagonally

4	1	8	5	3	8
$\uparrow$	$\uparrow$	$\uparrow$	$\uparrow$	$\uparrow$	$\uparrow$
2	-3	8	5	1	6
4	1	5	2	3	8

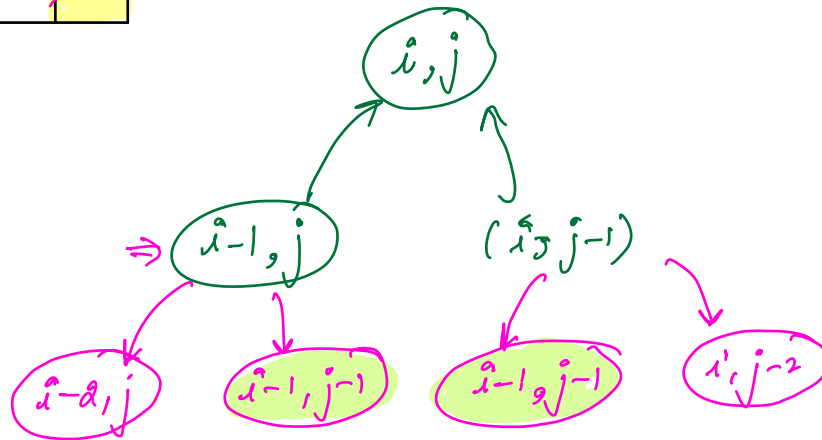
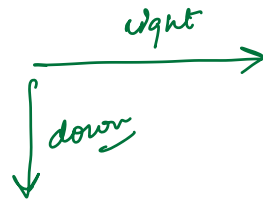
If you decide to pick a column  $\rightarrow$  can't pick adjacent column  
pick max element

3Q

	0	1	2	3
0	1	1	1	1
1	1			
2	1			
3	1			
4	1			

$n \times m$

No of ways to go from  $(0,0)$  to  $(n-1, m-1)$



$$\text{ways}(i, j) = \text{ways}(i-1, j) + \text{ways}(i, j-1)$$

$\text{dp}[n][m]$

$$\text{dp}[i][j] = \text{dp}[i-1][j] + \text{dp}[i][j-1]$$

$$\text{if } (i==0 \text{ || } j==0) \text{ dp}[i][j]=1;$$

T.C:  $O(n \times m)$

S.C:  $O(n \times m)$

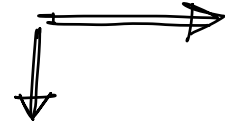
```
for (i=0; i<n; i++)
    for (j=0; j<m; j++)
        dp[i][j] = 0;
```

#

	0	1	2	3
0	1	1	1	1
1	1	0	1	0
2	0	1	1	1
3	1	0	1	1
4	1	1	1	1

1 → open  
0 → closed

no of ways

 $(0,0) \rightarrow (n-1, m-1)$ 


	0	1	2	3
0	1	1	1	1
1	1	0	1	0
2	0	1	1	1
3	1	0	1	1
4	1	1	1	1

if (mat[i][j] == 0)  
dp[i][j] = 0;

dp[i][j] = dp[i-1][j] + dp[i][j-1];

if (mat[i][j] == 0) return 0;

#

	0	1	2	3	4
0	2	1	3	7	3
1	3	3	4	8	3
2	0	5	1	2	1
3	0	2	3	1	3
4	4	1	5	2	3

cost



min cost to go from

 $(0,0) \rightarrow (n-1, m-1)$ 

mincost(i, j) = cost[i][j] + min(mincost(i-1, j), mincost(i, j-1));

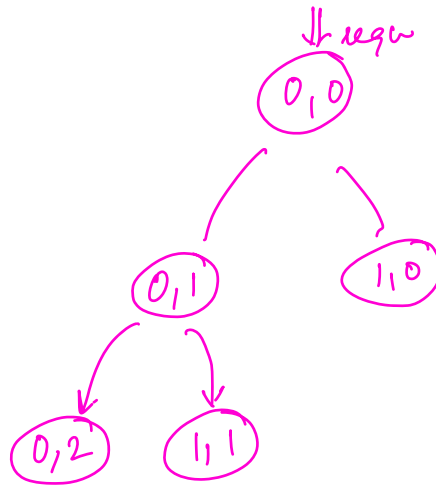
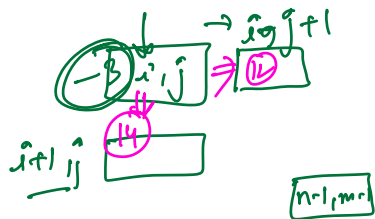
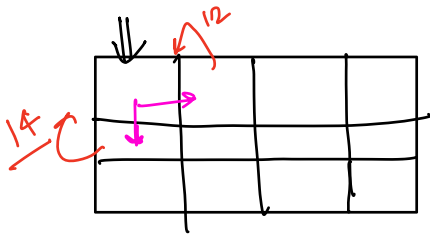
Q

# Dragons & princess

	0	1	2	3
0	-3	2	4	-5
1	-6	5	-4	6
2	-15	-7	5	-2
3	2	10	-3	-4

minimum  
initial energy  
should prince  
start with ?

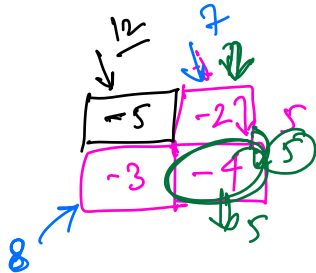
chamber → dragon  
energy diek  
die →  $\leq 0$  energy



$$\text{energy}(i,j) + \text{arr}[i][j] = \min(\text{energy}(i,j+1), \text{energy}(i+1,j))$$

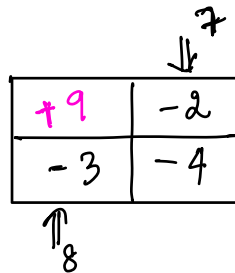
energy  
required  
to enter  
& reach  $(n-1, m-1)$

$$\text{energy}(i,j) = \min(\text{energy}(i,j+1), \text{energy}(i+1,j)) - \text{arr}[i][j]$$



$$5 - (-2) = 7$$

$$7 - (-5) = 12$$



$$\min(7, 8) = 7$$

$$= 7 - 9$$

$$= -2 \Rightarrow 1$$

$$dp[i][j] = \max(\min(dp[i+1][j], dp[i][j+1]) - arr[i][j], 1)$$

$$dp[n-1][m-1] = \max(1 - arr[n-1][m-1], 1)$$

$$\text{for } (i = n-1; i \geq 0; i--)$$

$$\quad \text{for } (j = m-1; j \geq 0; j--)$$

↓

↓

↓