### CS3310 – Operating Systems Lab

## Assignment Report - Lab 4

# **SURE MANOJ KUMAR - CS12B028**

#### Exercise 1:

```
    (i) – Allocating an array of Env's of NENV Size
    envs=(struct Env*)boot_alloc(NENV*sizeof(struct Env));
    (ii) – Mapping to user Read-only at UENVS
    boot_map_region(kern_pgdir, UENVS, NENV * sizeof(struct Env), PADDR(envs), PTE_U |
```

After writing the above code, when I ran the kernel, the function check\_kern\_pgdir() gave succeeded as result.

PTE P); //Read-only permissions "PTE U | PTE P"

#### Exercise 2:

```
(i) - env_init()
```

adding the environments to the env\_free\_list linked list in the same order as of *envs* array and marking them as free.

Allocating a page directory and initialising the kernel portion for the environment

```
(iii) – Region alloc()
Allocating and mapping physical memory for the environment
      //making va and (va+len) page – aligned.
      uint32 t a= ROUNDDOWN((uint32 t)va,PGSIZE); //rounding va down
      uint32 t b= ROUNDUP((uint32 t)(va+len), PGSIZE); // rounding (va+len) up
      struct PageInfo* p;
      uint32 ti;
      for(i=a;i< b;i++)
            p = page alloc(0); // allocating a page
            if (p = = NULL) // panicking if the page can not be allocated
                   panic("region alloc failed");
            page insert(e->env pgdir,p,(void *)i,PTE U|PTE W);
            //allocating this page for under e->env pgdir with read-only permissions.
      }
(iv) - load_icode()
parsing elf binary and loading it into user environment space of new environment
      lcr3(PADDR(e->env pgdir));
      struct Elf* elfheader = (struct Elf*)binary;
      struct Proghdr* ph1,*ph2;
      ph1=(struct Proghdr*)(elfheader->e phoff+binary);
      ph2=ph1+elfheader->e phnum;
      e->env tf.tf eip=elfheader->e entry;
      char* start, *content;
      for(;ph1 < ph2;ph1 + +)
          region alloc(e,(void*)ph1->p_va,ph1->p_memsz);
          memmove((void*)ph1->p va,(char*)binary+ph1->p offset,ph1->p filesz);
          memset((void*)(ph1->p va+ph1->p filesz),0,ph1->p memsz - ph1->p filesz);
      lcr3(PADDR(kern pgdir));
      region alloc(e,(void*)(USTACKTOP-PGSIZE),PGSIZE);
(v) - env_create()
allocating an environment an loading the elf binary into it.
      struct Env* env;
      int a =env alloc(&env, 0); //env alloc returns 0 on successfull allocation
      if (a != 0)
```

```
panic("env create failed."); //panic if env alloc returns non-zero(failed)
      env->env parent id=0; //setting parent's id to zero.
      env->env type=type;
      load icode(env,binary,size); //loading the elf binary into the environment.
(vi) – env_run()
running given environment in user mode.
      if(curenv!=e)
            //changing the 'curenv' status to runnable from running if it was already
            //in running state.
             if((curenv!=NULL) && curenv->env status==ENV RUNNING)
                   curenv->env status=ENV RUNNABLE;
             curenv=e;
             curenv->env status=ENV RUNNING; //setting env status to env running
             curenv->env runs++; //updating the env runs
             lcr3(PADDR(curenv->env pgdir)); //switching to the env pgdir address space
      }
      env pop tf(&curenv->env tf); //popping the current environment's trap frame.
Exercise 3:
             Reading Part.
Exercise 4:
alltraps:
pushl %ds
pushl %es
pushal
movl $GD KD, %eax //loading GD KD into es and ds
movw %ax, %es
movw %ax, %ds
pushl %esp //pushl esp and calling trap
call trap
trap init() - initialising the idt to point to each of the entries in trapentry.S
      int i;
      //initialising the idt
      for (i = 0; i \le 48; i++)
             SETGATE(idt[i], 0 , GD KT, handlers[i], 0);
      SETGATE(idt[T BRKPT], 0, GD KT, handlers[T BRKPT], 3); //for break point test to
                                                               //work
      SETGATE(idt[T SYSCALL], 0, GD_KT, handlers[T_SYSCALL], 3);//added for
exercise 7
```

Trap numbers from 8 to 15 generate an error code and others do not. So, for them we push the corresponding number and for others we push zero in place of error code.

#### Exercise 5 & 6:

trap dispatch(): call the appropriate function depending on the trapno in the trapframe.

```
int retval = 0;
      switch(tf->tf trapno)
             case T PGFLT: //pagefault
                    page fault handler(tf);
                    return;
             case T BRKPT: //break point - for exercise 6
                    monitor(tf); //monitor function is already implemented .
                    return;
             case T SYSCALL: //system call – for exercise 7
                    retval= syscall(tf->tf regs.reg eax,tf->tf regs.reg edx,tf-
       >tf regs.reg ecx,tf->tf regs.reg ebx,tf->tf regs.reg edi,tf->tf regs.reg esi);
                    tf->tf regs.reg eax = retval;
                    return;
      default: //we do not handle any other trap
             print trapframe(tf);
             if (tf->tf cs == GD KT)
                    panic("unhandled trap");
             else {
                    env destroy(curenv);
                    return;
             }
      }
Exercise 7:
sys_cputs():
      char *va;
      pte t *p;
      //checking whether the user has permissions or not
      for (va = (char *)s; va < s + len; va + = PGSIZE){
             p = pgdir walk(curenv->env pgdir, (void *)va, 0);
             if (p && (*p & PTE U) && (*p & PTE P))
                    continue;
             env destroy(curenv); //if user dont have necessary permissions we destroy
             return; //that environment and return, otherwise print the string(as given)
      }
```

```
syscall.c()
```

```
int retval = 0;
      //calls the appropriate function based on the syscallno
      switch(syscallno){
             case SYS cputs:
                    user mem assert(curenv, (void *)a1, a2, PTE U | PTE P);
                    //the above is for checking if the curenv has the required permissions
                    //or not. Asked in Exercise 9.
                    sys cputs((char *)a1, a2);
                    retval = a2;
                    break:
             case SYS cgetc:
                    retval = sys cgetc();
                    break;
             case SYS getenvid:
                    retval = sys getenvid();
                    break;
             case SYS env destroy:
                    retval = sys env destroy(a1);
                    break;
             default:
                    -E INVAL;
             return retval;
Exercise 8:
libmain.c
      thisenv = (struct Env *)envs + ENVX(sys getenvid());
      //thisenv point to environment's structure in envs array
Exercise 9 & 10:
page fault handler(): panicking if the page fault occurs in kernel mode.
if ((tf->tf cs \& 3)!= 3) //checking the last two bits of tf cs
      panic("kernel page fault");
user mem check():
      //checking if the environment has the permission to read [va,va+len] with
      //permissions perm | PTE P
      uint32 t ia;
      perm |= PTE P;
      for ( ia = (uint32 t)va; ia < (uint32 t)va + len; ia + +){
             if (ia > ULIM){
                    user mem check addr = ia;
                    return -E FAULT;
```

```
}
pte_t *p = pgdir_walk(env->env_pgdir, (void *)ia, 0);
if (p==NULL){
    user_mem_check_addr = ia;
    return -E_FAULT;
}
if ((*p & perm) != perm){
    user_mem_check_addr = ia;
    return -E_FAULT;
}
return -E_FAULT;
}
```

## Question 1:

Because, the hardware does not distinguish errors ,for which exception is called.some times we push the error code and some times we do not .if all the exceptions we delivered to the same handler, we lose this feature.

## Question 2:

No, we dont have to do anything to make the code work properly. If the chice is given to the user, then there is a security vulnerability as exceptions are handled with kernel privilages and he can run anything he wants, which should not happen. It gives int 13 because there is a violation in the privilages.

#### Question 3:

In trap\_init() we set the dpl to 3 for T\_BRKPT so that user can generate breakpoint trap and it will generate breakpoint trap ,if the dpl is set to 0,user can not generate it himself and it will give general protection fault.

## Question 4:

The mechanisms enforce permissions, they create gate so that user can make system calls through the exceptions depending on what exceptions we allow for the user to generate (depending on how you set the permissions in trap\_init using SETGATE). This is useful as a protection to not execute the malicious code written by user.

#### Exercise 9 – Question:

Because user process do not have permissions to access the code in lib/libmain.c.