Algorithms and Data Structures

Fundamental Data Structure: Lists

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Content of this Lecture

- ADT List
- Using Arrays
- Linked Lists
- Double-linked Lists
- Iterators

Lists

- Very often, we want to manage a list of „things“
  - A list of customer names that have an account on a web site
  - A list of windows that are visible on the current screen
  - A list of IDs of students enrolled in a course
- Lists are fundamental: We have things and lists of things
- Difference between a list and a set?
  - Lists: Have order, may contain duplicates
- Note: A list has “an” order; this doesn’t imply that we require that the elements of a list have a particular order
  - Like lexicographic / numerical / …
- But when dealing with list, one usually wants to maintain the given order

Representing Lists

- We already discussed an ADT for a list

```
type list( T)
operators
  isEmpty:  list → bool;
  add:      list x T → list;
  delete:   list x T → list;
  contains: list x T → bool;
  length:   list → integer;
```

- Today, we require slightly different operations
  - Positions are counted starting from 1
  - **`insert(L,t,p)`:** Add element t at pos p of L; if p=|L|+1, add t to L
  - **`delete(L,p)`:** Delete element at position p of list L
  - **`search(L,t)`:** Return first pos of t in L if $t \in L$; return 0 otherwise
- We also require that the order of the existing elements in the list is not changed by any of these operations

Representing Lists

- How can we implement this ADT?

```
type list( T)
import integer, bool;
operators
  isEmpty: list → bool;
  insert:  list x integer x T → list;
  delete:  list x int → list;
  search:  list x T → integer;
  length:  list → integer;
```

- We shall discuss three options
  - Arrays
  - Linked-Lists
  - Double-Linked lists
- We assume to store values of constant size (e.g. real)

Just a Snapshot

- Of course, there are many more issues
  - If the list gets too large to fit into main memory
  - If the list contains complex objects and should be searchable by different attributes (first name, last name, age, ...)
  - If the list is stored on different computers, but should be accessible through a single interface
  - If different users can access and modify the list at the same time
  - If the list contains lists as elements (nested lists)
  - ...

Just a Snapshot

- Of course, there are many more issues
  - If the list gets too large to fit into main memory
    - See databases, caching, operating systems
  - If the list contains complex objects and should be searchable by different attributes (first name, last name, age, ...)
    - See databases; multidimensional index structures
  - If the list is stored on different computers, but should be accessible through a single interface
    - See distributed algorithms, cloud-computing, peer-2-peer
  - If different users can access and modify the list at the same time
    - See databases; transactions; parallel/multi-threaded programming
  - If the list contains lists as elements (nested lists)
    - See trees and graphs
  - ...

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Lists through Arrays

- Probably the simplest method
  - Fix a maximal number of elements **max_length**
  - Access elements by their offset within the array

```
class list {
  size: integer;
  a: array[1..max_length]

  func void init() {
    size := 0;
  }
  func bool isEmpty() {
    if (size=0)
      return true;
    else
      return false;
    end if;
  }
}
```

Insert, Delete, Search

Problem!

```
func void insert (t real, p integer) {
  if size = max_length then
    return ERROR;
  end if;
  if p!=size+1 then
    if (size<p) or (p<1) then
      return ERROR;
    end if;
    for i := size-1 downto p do
      A[i+1] := A[i];
    end for;
  end if;
  A[p] := t;
  size := size + 1;
}
```

```
func void delete(p integer) {
  if (size<p) or (p<1) then
    return ERROR;
  end if;
  for i := p .. size-1 do
    A[i] := A[i+1];
  end for;
  size := size - 1;
}
```

```
func int search(t real) {
  for i := 1 .. size do
    if A[i]=t then
      return i;
    end if;
  end for;
  return 0;
}
```

- Complexity (worst-case)?
  - Insert: O(n)
  - Delete: O(n)
  - Search: O(n)

Properties

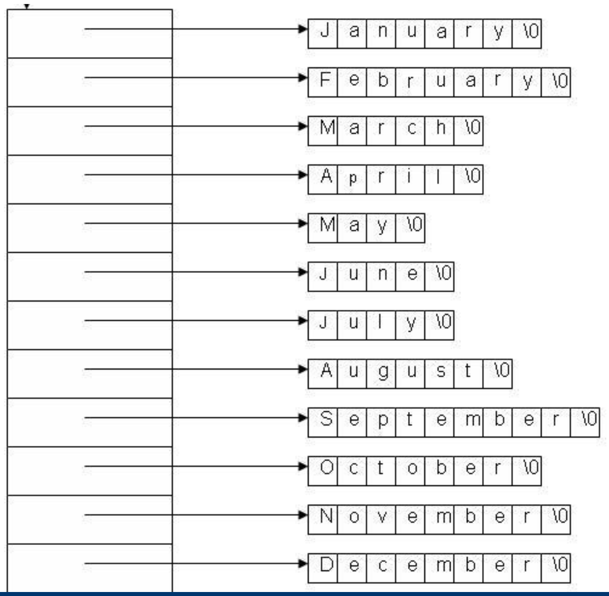
- Source of complexity: We can jump to position p, but we need to move O(n) elements to insert/delete an item
  - If all positions appear with the same probability, we expect n/2 operations on average for each insert/delete (still O(n))
- Unbalanced: Adding to the end of an array costs O(1), adding to the start costs O(n) operations
- Disadvantages
  - If `max_length` too small, we run into errors
  - If `max_length` too large, we waste space
- Help: Dynamic arrays (with other disadvantages)
  - Think of Java vectors ...

Arrays of Strings

- We assumed that every element of the list requires constant space
  - Thus, elements are one-after-the-other in main memory
  - Element at position p can be access directly by computing the address of the memory cell
- What happens for other data types, e.g. strings?

Arrays of Strings

- We assumed that every element of the list requires constant space
  - Thus, elements are one-after-the-other in main memory
  - Element at position p can be access directly by computing the address of the memory cell
- What happens for other data types, e.g. strings?
  - Each string actually is a list itself
  - Thus, we are building a list of lists
  - Array A holds pointer to strings
  - Strings can be implemented in various ways (arrays, linked lists, ...)
  - A pointer requires constant space



Summary

	Array	Linked list	Double-linked l.
Insert	O(n) (or error)		
Delete	O(n)		
Search	O(n)		
Add to list	O(1) (or error)		
Space	Static, upfront		

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Linked Lists

- The static space allocation is a severe problem of arrays
- Alternative: Linked lists
  - An element of a list becomes a tuple (**value**, **next**)
  - **value** is the value of the element
  - **next** is a pointer to the next element in the list
- Disadvantage: O(n) additional space or all the pointers

```
class element {
  value: real;
  next: element;
}
```

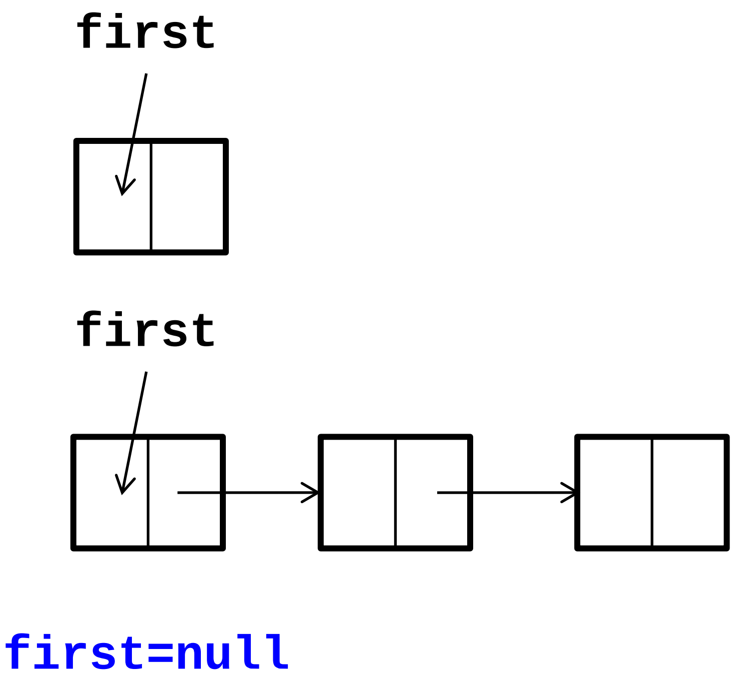
```
class list {
  first: element;

func void init() {
    first := null;
  }
func bool isEmpty() {
    if (first=null)
      return true;
    else
      return false;
    end if;
  }
}
```

Search

- Return the first element with value=t, or null if no such element exists
  - Note: Here we return the element, not the position of the element
  - Makes sense in linked lists: Access to properties in O(1)

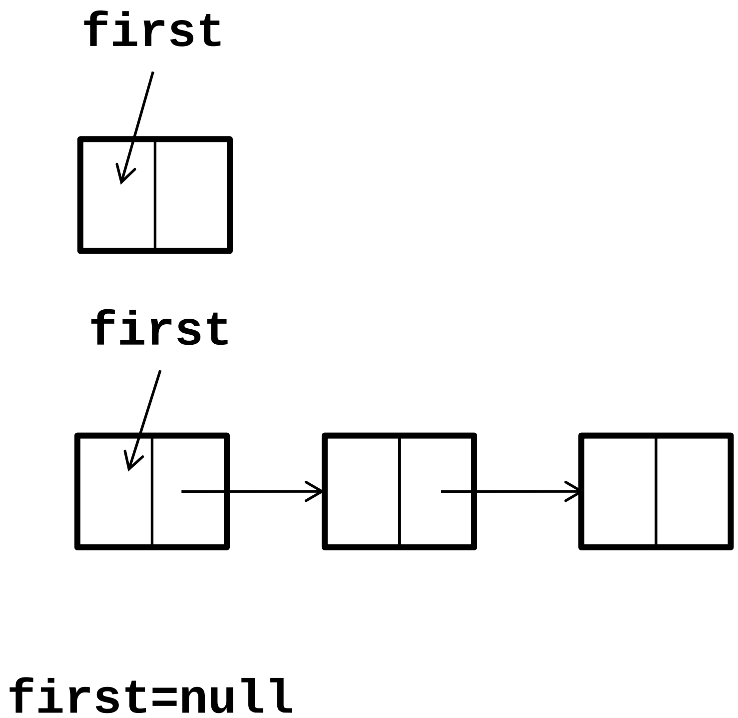
```
func element search(t real) {
  v := first;
  if v.value = t then
    return v;
  end if;
  while (v.next != null) do
    v := v.next;
    if (v.value  = t) then
      return v;
    end if;
  end while;
  return null;
}
```



Search

- Return the first element with value=t, or null if no such element exists

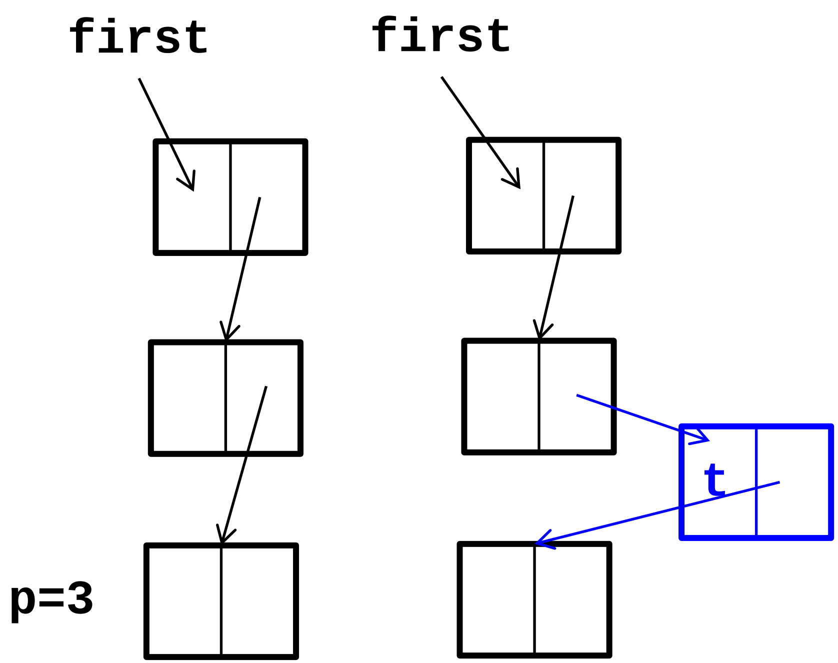
```
func element search(t real) {
  if first=null then
    return null;
  end if;
  v := first;
  if v.value = t then
    return v;
  end if;
  while (v.next != null) do
    v := v.next;
    if (v.value  = t) then
      return v;
    end if;
  end while;
  return null;
}
```



Insert

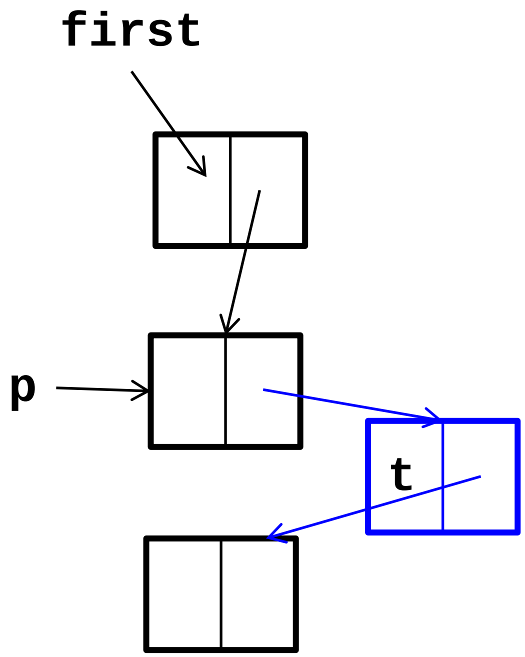
- Insert value t after the p'th element of the list

```
func void insert (t real, p integer) {
  v := first;
  if v=null then
    if p≠1 then
      return ERROR;
    else
      first := new element(t, null);
      return;
    end if;
  end if;
  new := new element (t, null);
  for i := 1 .. p-1 do
    if (v.next=null) then
      return ERROR;
    else
      v := v.next;
    end if;
  end for;
  new.next := v.next;
  v.next := new;
}
```



InsertAfter

- In linked lists, a slightly different operation also makes sense: We insert after element p, not at position p
  - E.g., we search an element p and want to insert a new element right after p
- No difference in arrays, but large difference for linked lists

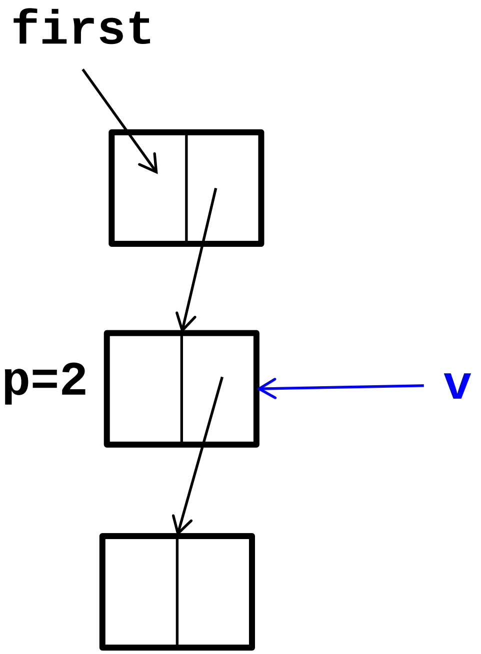



```
func void insertAfter (t real, p element) {
  new := new element (t, null);
  new.next := p.next;
  p.next := new;
}
```

Delete

- Delete the p'th element of the list

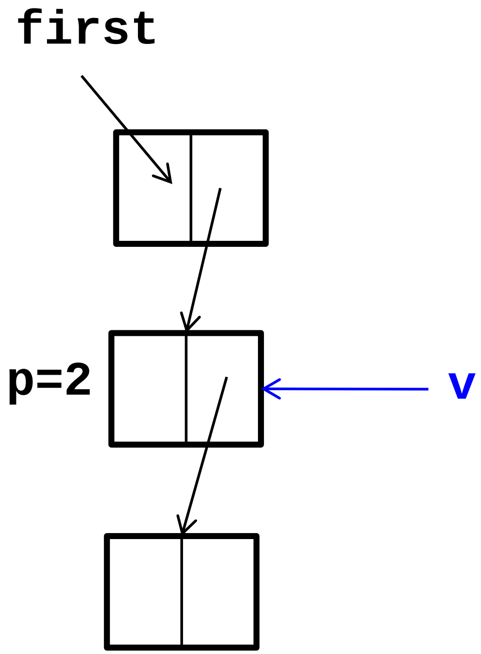
```
func void delete(t real, p integer) {
  v := first;
  if (v=null) or (p=0) then
    return ERROR;
  end if;
  for i := 1 .. p-1 do
    if (v.next=null) then
      return ERROR;
    else
      v := v.next;
    end if;
  end for;
  ? PROBLEM ?
}
```



Delete – Bug-free?

- Delete the p'th element of the list

```
func void delete(t real, p integer) {
  v := first;
  if (v=null) or (p=0) then
    return ERROR;
  end if;
  for i := 1 .. p-1 do
    last := v;
    if (v.next=null) then
      return ERROR;
    else
      v := v.next;
    end if;
  end for;
  last.next := v.next;
}
```

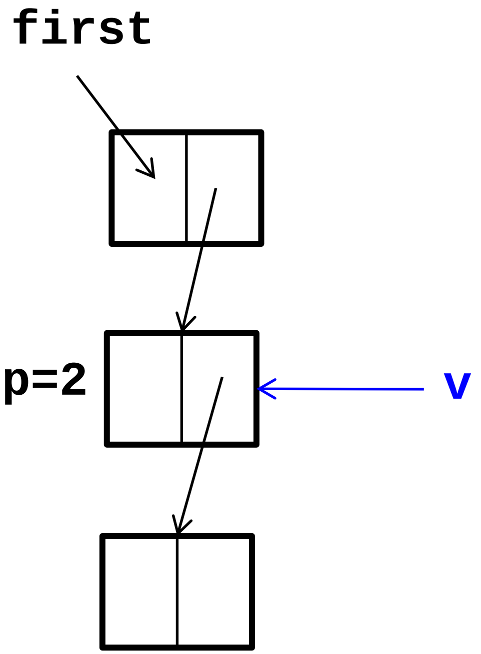


- What if p=1?

Delete – Bug-free

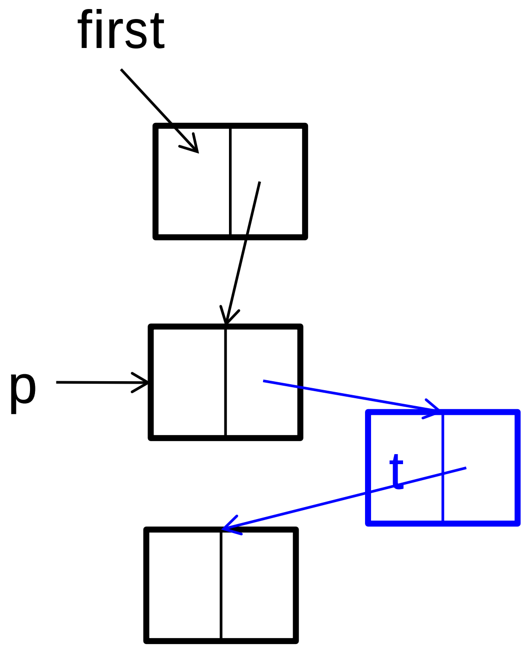
- Delete the p'th element of the list

```
func void delete(t real, p integer) {
  v := first;
  if (v=null) or (p=0) then
    return ERROR;
  end if;
  if p=1 then
    first := v.next;
    return;
  end if;
  for i := 1 .. p-1 do
    last := v;
    if (v.next=null) then
      return ERROR;
    else
      v := v.next;
    end if;
  end for;
  last.next := v.next;
}
```



DeleteThis

- In linked lists, a slightly different operation also makes sense: We delete element p, not that at position p
  - Again: We search an element p and then want to delete exactly p
- Big problem
  - If we have p, we cannot directly access the predecessor q of p (the q with **`q.next=p`**)
  - We need to go through the entire list to find p
  - **`deleteThis`** has the same complexity as **`delete`**, the additional information (p) does not help



Two More Issues

- Show me the list

```
func String print() {
  if (first=null) then
    return "";
  end if;
  tmp := "";
  while (v≠null) do
    tmp := tmp+v.value;
    v := v.next;
  end for;
  return tmp;
}
```

- What happens to deleted elements p?
  - In most languages, the space occupied by p remains blocked
  - These languages offer an explicit "dispose" which you should use
  - Java: "Dangling" space is freed automatically by garbage collector
    - After some (rather unpredictable) time

Summary

	Array	Linked list	Double-linked l.
Insert	O(n)	O(n)	
InsertAfter	O(n)	O(1)	
Delete	O(n)	O(n)	
DeleteThis	O(n)	O(n)	
Search	O(n)	O(n)	
Add to list	O(1)	O(1)	
Space	Static	n+1 add. pointers	

How?

Double-Linked List

- Simple idea: Every element holds a pointer to the next and to the previous element
- List holds pointer to first and last element
- Advantages
  - `deleteThis` can be implemented in O(1)
  - Concatenation of lists can be implemented in O(1)
  - Addition/removal of last element can be implemented in O(1)
- Disadvantages
  - Requires more space
    - Beware of the space necessary for a pointer on a 64bit machine
  - Slightly more complicated operations

Summary

	Array	Linked list	Double-linked l.
Insert	O(n)	O(n)	O(n)
InsertAfter	O(n)	O(1)	O(1)
Delete	O(n)	O(n)	O(n)
DeleteThis	O(n)	O(n)	O(1)
Search	O(n)	O(n)	O(n)
Add to start of list	O(n)	O(1)	O(1)
Add to end of list	O(1)	O(n)	O(1)
Space	Static	n+1 add. pointers	2(n+1) add. point.

- Can we do any better?
  - Yes – if we sort the list on the searchable value
  - Yes – if we know which elements are searched most often

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Example

- Assume we have a list of customers with their home addresses
- We want to know how many customers we have per city

Meier	Berlin
Müller	Hamburg
Meyer	Dresden
Michel	Hamburg
Schmid	Berlin
Schmitt	Hamburg
Schmidt	Wanne-Eikel
Schmied	Hamburg

⇨

Berlin	2
Hamburg	4
Dresden	1
Wanne-Eikel	1

Using a List

- Assume we have a list of cities (**`groups`**) with elements of class **`counts`**, and this list has an operation **`increment(city)`**

```
class counts {
  count: integer;
  city: String;
}

class customer{
  name: String
  city: String;
}
```

```
func void group_by( customers list;
                    groups: list) {
  if customers.isEmpty() then
    return;
  end if;
  c : customer;
  for i:= 1 .. customers.size do
    c := customers.search( i);
    groups.increment( c.city);
  end for;
  print groups;
}
```

Original “search” searches a value, not the i’th value

Our Lists Lack Functionality

```
class list {
  func void init() …
  func bool isEmpty() …
  func valueAt( i integer) { … }
  …
}
```

- Complexity
  - O(1) in arrays
  - O(n) in (double-)linked lists

Using a List 2

- Assume we have a list of cities (**`groups`**) with elements of class **`counts`**, and this list has an operation **`increment(city)`**

```
class counts {
  count: integer;
  city: String;
}

class customer{
  name: String
  city: String;
}
```

```
func void group_by( customers list) {
  if customers.isEmpty() then
    return;
  end if;
  groups: list( counts);
  c : customer;
  for i:= 1 .. customers.size do
    c := customers.valueAt( i);
    groups.increment( c.city);
  end for;
  print groups;
}
```

Complexity?

- We run once through costumers: O(n)
- For linked lists, this gives $O(n^2)$ in total
  - Only O(n) for arrays, but these had other problems
- Not satisfactory: We are doing unnecessary work
  - We only need to follow pointers – but driven by the client
  - Our data type “list” has no state, i.e., no internal “current” position
  - Without in-list state, the state (variable i) must be managed outside the list, and the list must be put to the right state again for every operation (`valueAt`)
  - Remedy: Stateful list ADT

Stateful Lists

```
type list( T)
import
  integer, bool;
operators
  isEmpty:  list → bool;
  insert:   list x integer → list;
  delete:   list x T → list;
  …
```

⇓

```
type stateful_list( T)
import
  integer, bool;
operators
  isEmpty:    list → bool;
  setState:   list x integer → list;
  insertHere: list → list;
  deleteHere: list → list;
  getNext:    list → T;
  search:     list x T → integer;
  size:       list → integer;
```

- Impl: List holds an internal pointer **p_current**
  - This is the state
- **p_current** can be set to position i using **setState()**
- **insertHere** inserts after **p_current**, **deleteHere** deletes **p_current**
- **getNext()** returns element at position **p_current** and increments **p_current** by 1

Using Stateful Lists

```
func void group_by( customers stateful_list) {
  if customers.isEmpty() then
    return;
  end if;
  groups: list( counts);
  c : customer;
  customers.setState(0);
  for i:= 1 .. customers.size do
    c := customers.getNext();
    groups.increment( c.city);
  end for;
  print groups;
}
```

- **`customers.setState(0)`** sets **`p_current`** before element 1
- **`getNext()`** can be implemented in O(1) using arrays or linked lists

Iterators

- **`stateful_lists`** allow to manage one state per list
- What if multiple threads want to read the list concurrently?
  - Every thread needs its own pointer
  - These pointers cannot be managed easily in the (one and only) list itself
- Iterators
  - An iterator is an object created by a list which holds list state
    - One **`p_counter`** per iterator
  - Multiple iterators can operate independently on the same list
  - Implementation of iterator depends on implementation of list, but can be kept secret from the client

Using an Iterator

```
func void group_by( customers stateful_list) {
  if customers.isEmpty() then
    return;
  end if;
  groups: list( counts);
  c : customer;
  it := customers.getIterator();
  while it.hasNext() do
    c := it.getNext();
    groups.increment( c.city);
  end while;
  print groups;
}
```

```
class iterator_for_linked_list (T) {
  p_current: T;

  func iterator init( l list) {
    p_current := l.getFirst();
  }

  func bool hasNext() {
    return (p_current ≠ null);
  }

  func T getNext() {
    if p_current = null then
      return ERROR;
    end if;
    tmp := p_current;
    p_current := p_current.next;
    return tmp;
  }
}
```

Take Home Message

- Finding robust ADTs that can remain stable for many applications is an art
  - See the complexity of standardization processes, e.g. Java community process
  - Growing trend to standardize ADTs / APIs
- Different implementations of an ADT yield different complexities of operations
- Therefore, one needs to look "behind" the ADT if efficient implementations for specific operations are required