

# Optimizing Prosumer Policies in Periodic Double Auctions Inspired by Equilibrium Analysis

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TCS Research

# Motivation

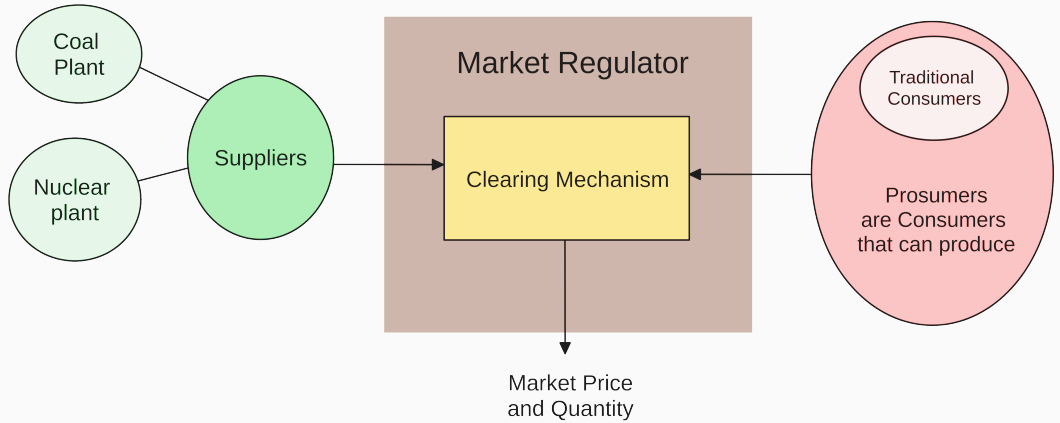
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- The Periodic Double Auction<sup>1</sup> is made up of **Suppliers** and **Consumers**.
- **Objective** is to find **Policies** for the consumers to **Minimize the Cost**.
- Application of PDA is that it models a **Day Ahead Electricity Market**.
- Hence, the developed policies can be used in the Market.

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<sup>1</sup>It is denoted as PDA.

- Day Ahead Electricity Market has two components: Suppliers and Cosumers.
- When Consumers are allowed to sell small quantities, they are called as **Prosumers**.

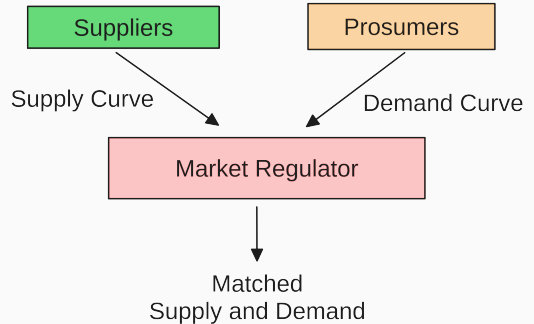


# Introduction

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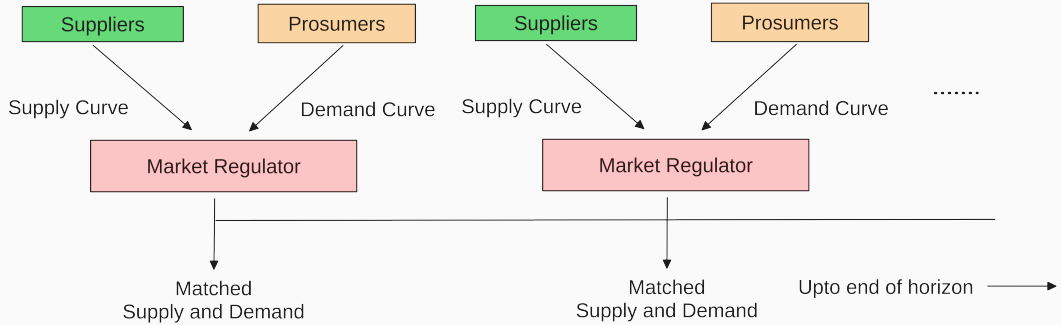
# Double Auction

- **Double Auction** is where the Suppliers and Prosumers are strategic.
- A Market Regulator matches the asks by supplier with the bids by the Prosumers.
- Both Asks and Bids are the pair of price and quantity.



The PDA is the sequence of Double Auctions and it has a finite horizon.



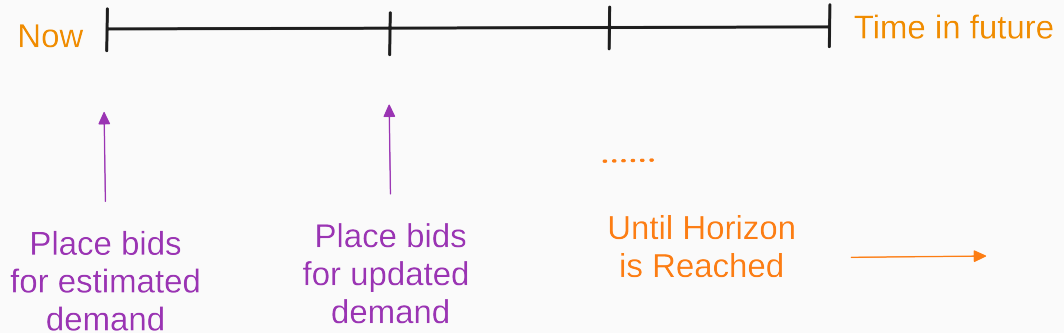


# Problem Statement

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The goal is to find **policies** for the **Prosumers** to minimize the **procurement cost**.

Update the demand estimation



- The PDA is modeled as a **Markov game**.
- A **Nash equilibrium** (NE) is formulated in an ideal setting<sup>2</sup>.
- Based on the NE analysis, an algorithm for the policy is developed.

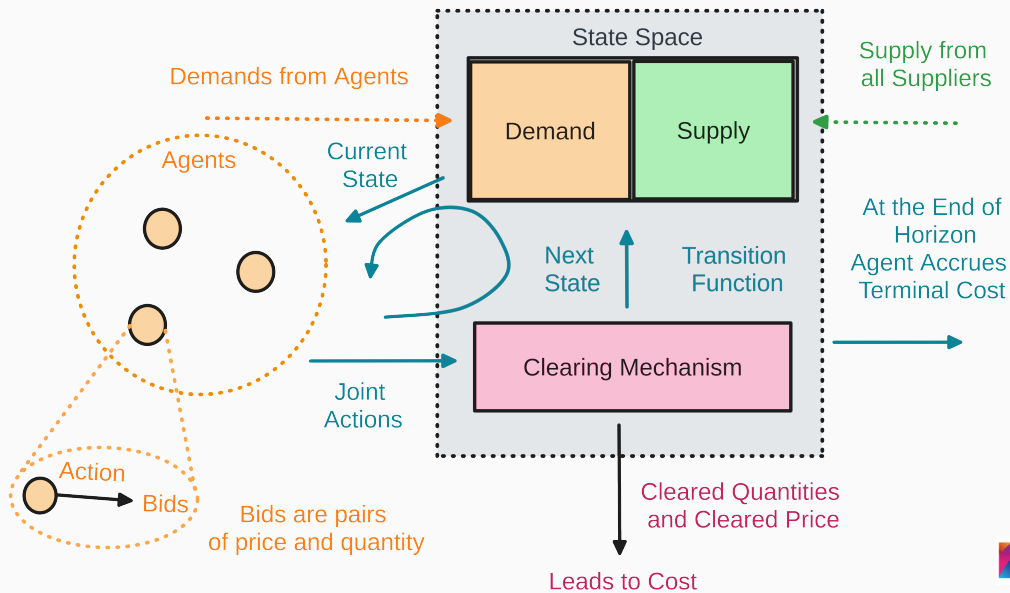
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<sup>2</sup>Some assumptions were made on the game to get to the NE.

## Modeling PDA as a Markov Game

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The PDA is modeled as a **Markov Game**  $\mathcal{M} = \langle N, A, S, P, C, H \rangle$





- The procurement cost is the **sum of costs** in each of the rounds.
- Starting from a state  $s \in S$ , the procurement cost is denoted as  $V(s)$ .
- $V(s)$  is also called as the Value function.

# Nash Equilibrium

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NE is a joint policy  $a \in A$ , in which no agent can **unilaterally deviate** from the allotted action.

- If there **are enough rounds** for players, then they would want to **utilize the rounds**.
- If rounds are not enough, then the players would bid at maximum possible price to **avoid paying very high price** outside the auction.
- The policy suggests the players **with higher requirement have more leverage**.

# MPNE-BBS : An Algorithm Based on Nash Equilibrium Analysis

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## Algorithm 1 MPNE-BBS

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```
1: totalDmd[]  $\leftarrow$  netDmdPredict(time)
2: selfDmd[]  $\leftarrow$  indvDmdPredictor(time)
3: for hour in [1,...,23] do
4:   deliverySlot  $\leftarrow$  time + hour
5:   unclearedAsks[]  $\leftarrow$  Auction(time-1,deliverySlot)
6:   estBidPrice = getEstBidPrice(time,deliverySlot,unclearedAks,selfDmd,totalDmd)
7:   adjEstBidPrice = getAdjEstBidPrice(estBidPrice) To account for inaccurate SC
8:   bidList = getBids(adjEstBidPrice)
9:   Auction(time,deliverySlot)  $\leftarrow$  bidList
```

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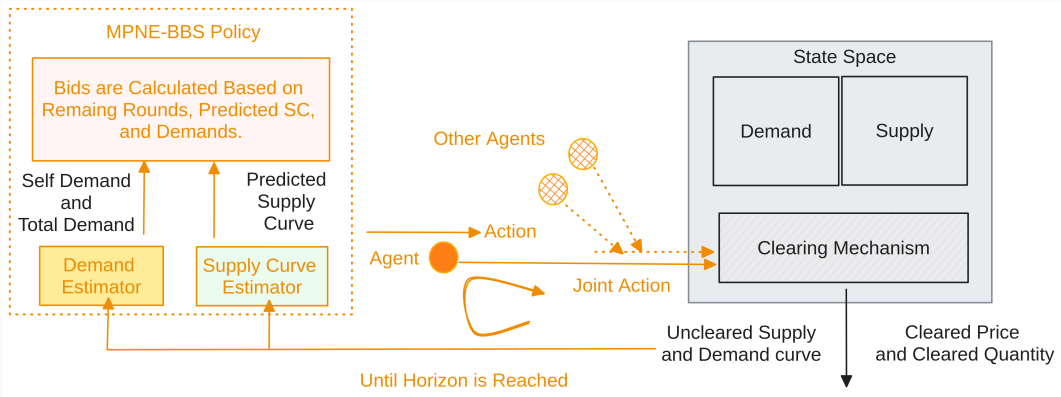
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**Algorithm 2** getEstBidPrice (time,deliverySlot,unclearedAks,selfDmd,totalDmd)

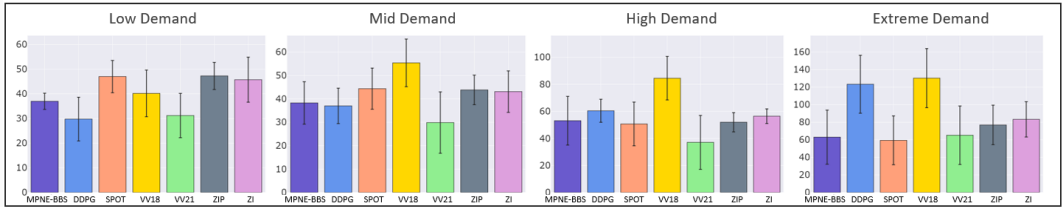
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```
1: if unclearedAsks is not empty then
2:   supplyQty  $\leftarrow$  sum(unclearedAsks.quant)
3:   if supplyQty > totalDmdForDSlot then
4:      $p_{u_h} \leftarrow$  Lowest price at which totalDmd[deliverySlot] can be cleared
5:     qtyReqForOthers  $\leftarrow$  totalDmd[deliverySlot] - selfDmd[deliverySlot]
6:      $p_{v_h^b} \leftarrow$  Lowest price at which qtyReqForOthers can be cleared
7:   else
8:      $p_{u_h}, p_{v_h^b} \leftarrow$  Maximum Ask Price
9:      $v_h^0 \leftarrow \max\{1, u_h - \text{hour} + 1\}$  To get remaining rounds
10:    return estBidPrice =  $\max\{p_{v_h^0}, p_{v_h^b}\}$  Adjust price to account for remaining rounds
11:  else
12:    clearedPrices[]  $\leftarrow$  Auction( $t, t + \text{hour}$ )  $\forall t < \text{currentTime}$ 
13:  return estBidPrice =  $\max\{\text{clearedPrices}\}$ 
```

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# Results



**Figure 1:** Wholesale Cost Comparison in 7-Player Games<sup>3</sup>

<sup>3</sup>Chandlekar et al. [2],Chandlekar et al. [1],Chowdhury et al. [3],Ghosh et al. [5],Cliff [4]



- A Nash Equilibrium for a PDA modeled as Markov game is devised.
- This Equilibrium helps in explaining the policy of the players in different scenarios.
- Devised an algorithm based on the Nash Equilibrium Analysis, results show that it works well in practice.
- In future work, aim is to extend the analysis to partial observable setting and see how it can help in devising algorithms.




For further discussion, please visit my poster presentation.

# Thank You

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