
Rule 0.4 *(Refactoring Loops with Repeated Computations)*

<pre> [...] contract A { [...] function f(pds) { [...] for(init; cond; update) { stmts[expr] } stmts' } [...] } </pre>	=	<pre> [...] contract A' { [...] function f(pds) { [...] if (cond') { T local = expr; for(init; cond; update) { stmts[local] } } stmts' } [...] } </pre>
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where

$expr$ is an expression that is loop-invariant (yields the same result in every iteration);

T is the type of the computed expression $expr$;

$local$ is a local variable of type T used to cache the computation result;

$stmts[expr]$ represents statements inside the loop that use $expr$;

$stmts[local]$ represents the same statements with $expr$ replaced by $local$;

$init$, $cond$, and $update$ are the loop initialization, condition, and update expressions;

$cond'$ is the condition evaluated after $init$ to determine if the loop will execute;

$stmts'$ represents statements following the loop.

provided

$expr$ does not depend on the loop variable or any value modified within the loop;

$expr$ is side-effect free (does not modify state or call external functions);

The value of $expr$ remains constant throughout all loop iterations;

No variable in $expr$ is modified between the assignment to $local$ and the loop execution;

$cond'$ accurately represents the condition for loop entry (typically $cond$ evaluated after $init$);

The evaluation of $cond'$ does not have side effects.

Invariant:

Let s_i and s'_i be the initial state of A and A' , respectively.

Let s_f and s'_f be the state reached by A and A' , respectively, after $A.f()$ and $A'.f()$ are executed from s_i and s'_i , respectively.

Then, the coupling invariant is

$$\forall s_i, s'_i . (s_i = s'_i) \rightarrow (s_f = s'_f)$$
