

# COMP1819

## Algorithms and Data Structures

Lecture 06: Trees

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# LEARNING DATA STRUCTURE & ALGORITHM IS IMPORTANT



## Content

- Review Lab 05
- General Trees
- Binary Trees
- Implementing Trees
- Tree Traversal Algorithms
- Reinforcement

# Lab 05

## 1. Improve Bubble Sort

The given code (lecture/github) always runs  $O(n^2)$  time even if the array is sorted. It can be optimized by stopping the algorithm if inner loop didn't cause any swap.

### Examples:

Input	Output
1 3 10 18	Bubble sort function stops after 1 pass.

### Hints

- Code for these searches given in the lecture slides

## 2. Compare Bubble sort, Selection sort and Insertion sort

Create a large array, compare the time complexity of three sort algorithms with these cases:

- Sorted list (ascending order)
- Sorted list (descending order)
- Random list

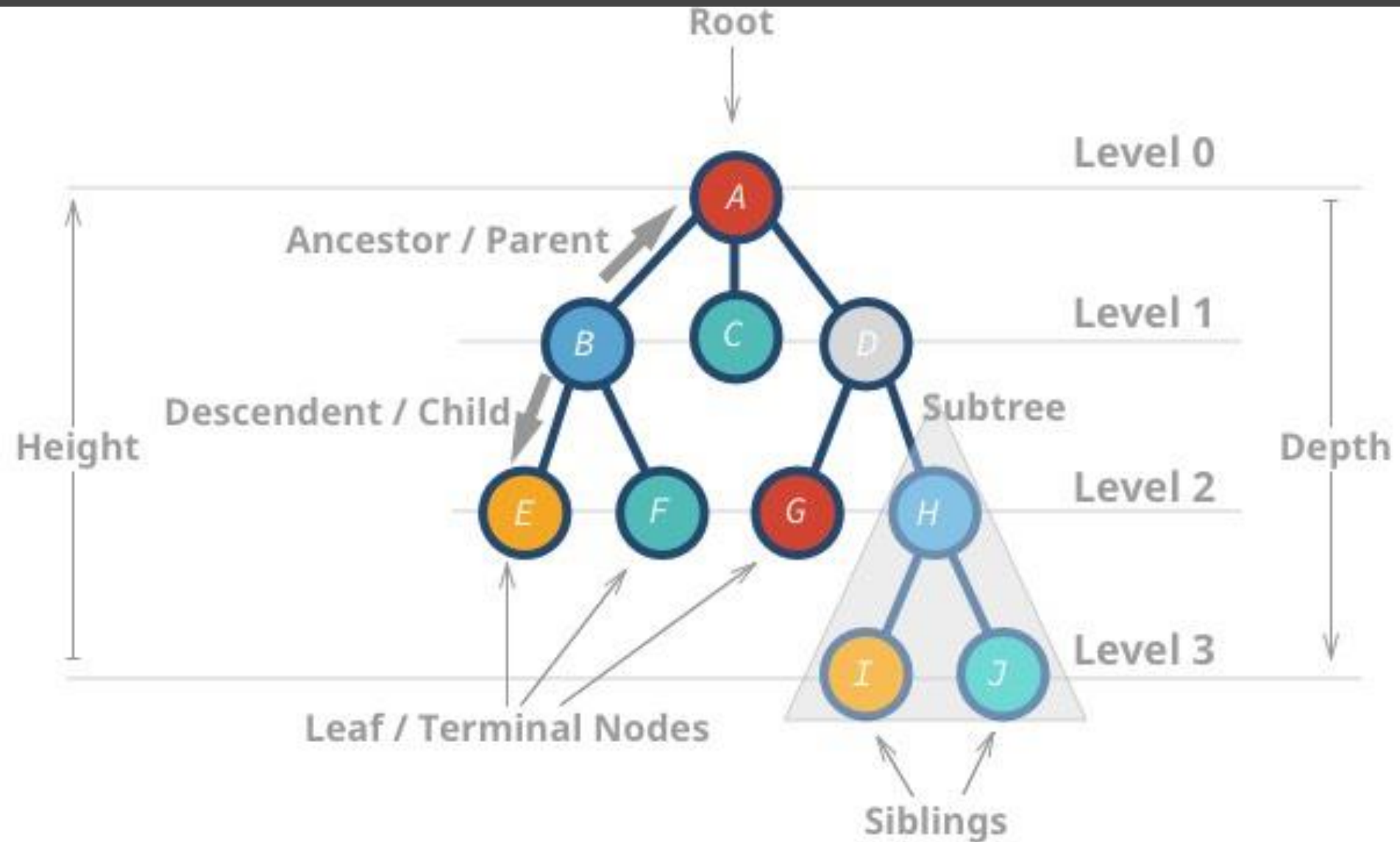
Which one is faster?

# Today



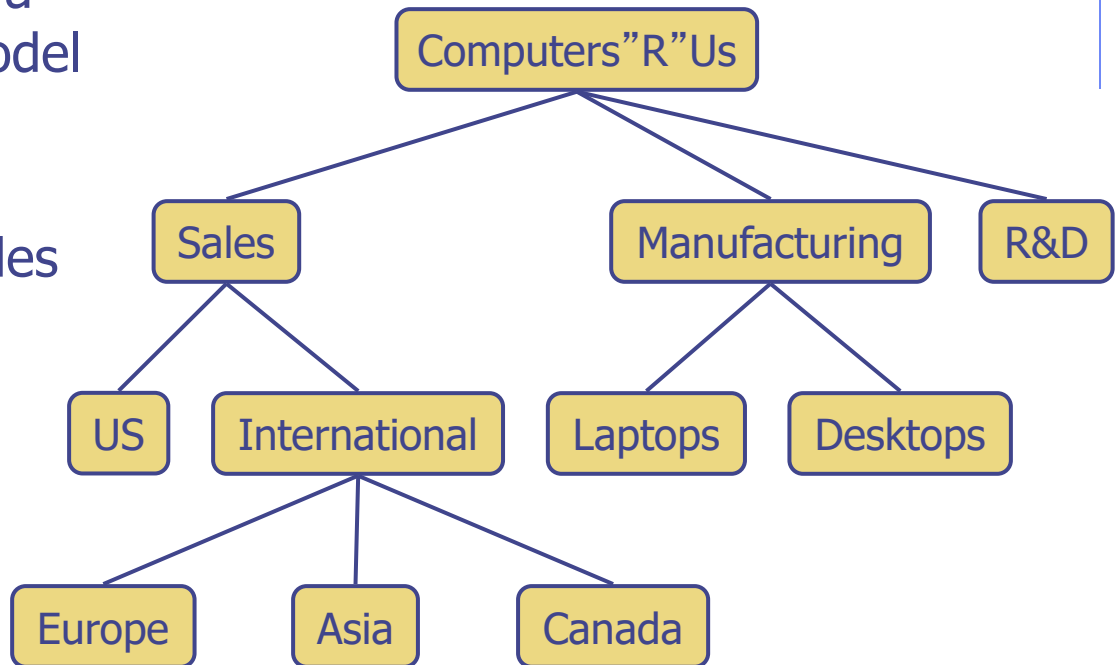
[Tree link](#)

# This one



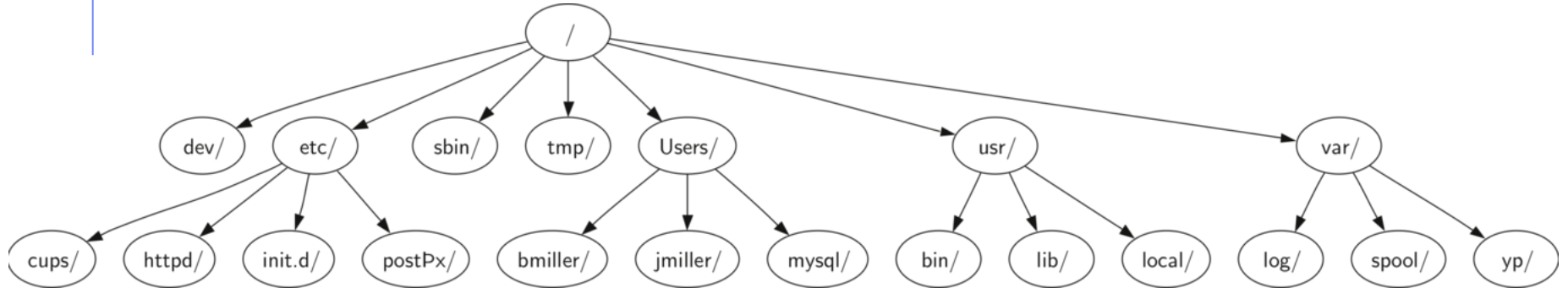
# What is a Tree

- ❑ In computer science, a tree is an abstract model of a hierarchical structure
- ❑ A tree consists of nodes with a parent-child relation
- ❑ Applications:
  - Organization charts
  - File systems
  - Programming environments



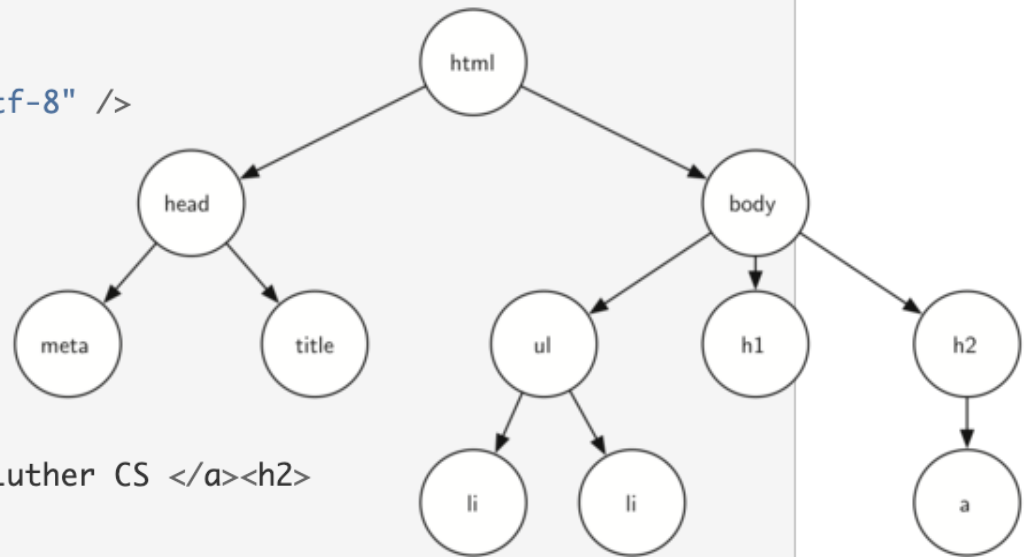


# Example: Unix File System Hierarchy



# Example: HTML webpage

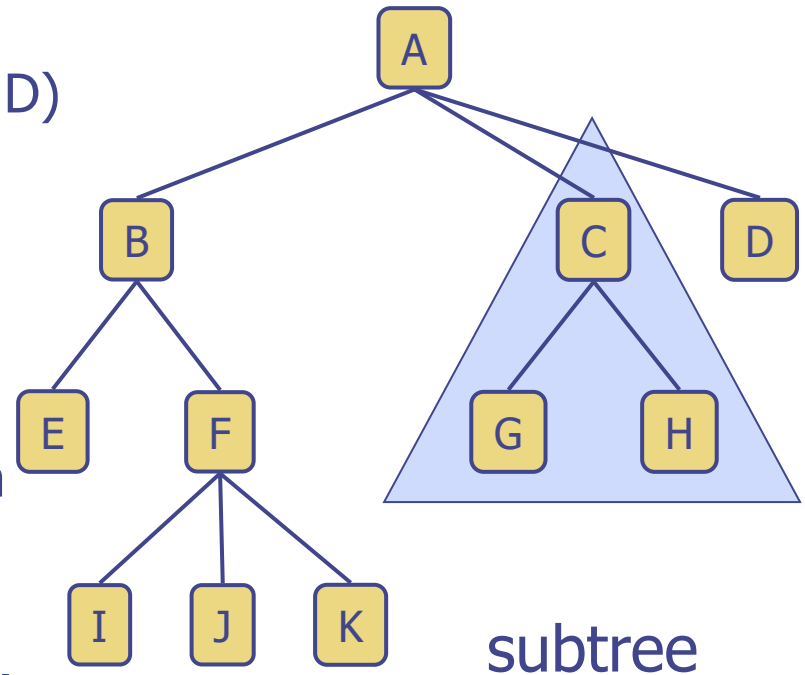
```
<html xmlns="http://www.w3.org/1999/xhtml"
      xml:lang="en" lang="en">
<head>
  <meta http-equiv="Content-Type"
        content="text/html; charset=utf-8" />
  <title>simple</title>
</head>
<body>
<h1>A simple web page</h1>
<ul>
  <li>List item one</li>
  <li>List item two</li>
</ul>
<h2><a href="http://www.cs.luther.edu">Luther CS </a></h2>
</body>
</html>
```



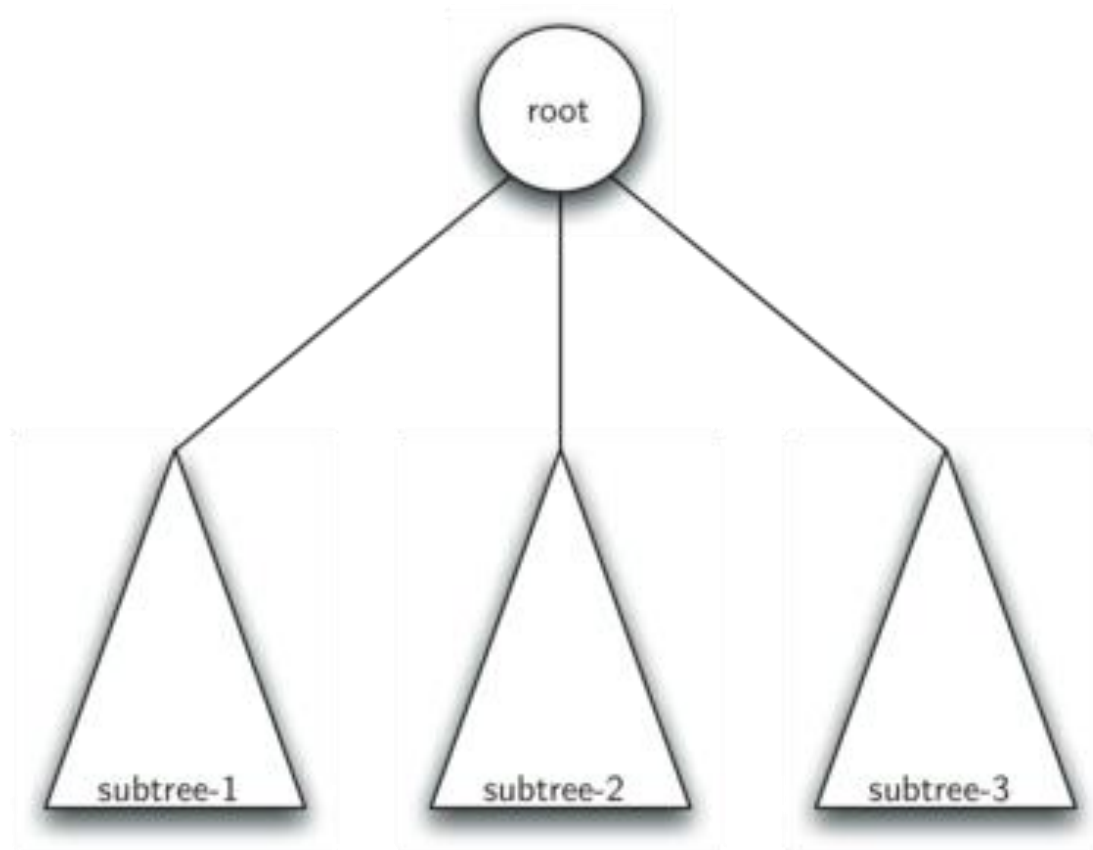


# Tree Terminology

- ❑ Root: node without parent (A)
- ❑ Internal node: node with at least one child (A, B, C, F)
- ❑ External node (a.k.a. leaf ): node without children (E, I, J, K, G, H, D)
- ❑ Ancestors of a node: parent, grandparent, grand-grandparent, etc.
- ❑ Depth of a node: number of ancestors
- ❑ Height of a tree: maximum depth of any node (3)
- ❑ Descendant of a node: child, grandchild, grand-grandchild, etc.
- ❑ Subtree: tree consisting of a node and its descendants



# Recursive definition



A tree is either empty or consists of a root and zero or more subtrees, each of which is also a tree. The root of each subtree is connected to the root of the parent tree by an edge.

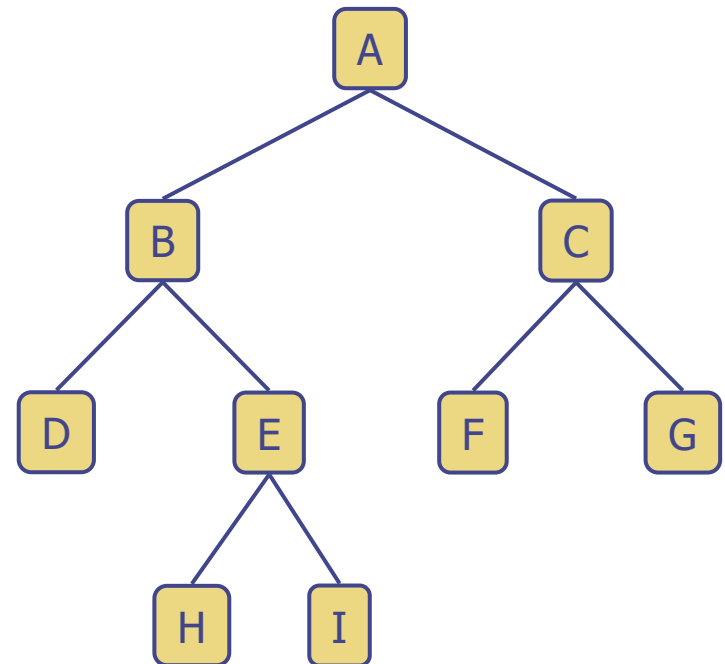
# Tree ADT

- We use positions to abstract nodes
- Generic methods:
  - Integer `len()`
  - Boolean `is_empty()`
  - Iterator `positions()`
  - Iterator `iter()`
- Accessor methods:
  - position `root()`
  - position `parent(p)`
  - Iterator `children(p)`
  - Integer `num_children(p)`
- ◆ Query methods:
  - Boolean `is_leaf(p)`
  - Boolean `is_root(p)`
- ◆ Update method:
  - element `replace(p, o)`
- ◆ Additional update methods may be defined by data structures implementing the Tree ADT

# Binary Trees

- A binary tree is a tree with the following properties:
  - Each internal node has at most two children (exactly two for **proper** binary trees)
  - The children of a node are an ordered pair
- We call the children of an internal node **left child** and **right child**
- Alternative recursive definition: a binary tree is either
  - a tree consisting of a single node, or
  - a tree whose root has an ordered pair of children, each of which is a binary tree

- Applications:
  - arithmetic expressions
  - decision processes
  - searching



# Properties of Proper Binary Trees

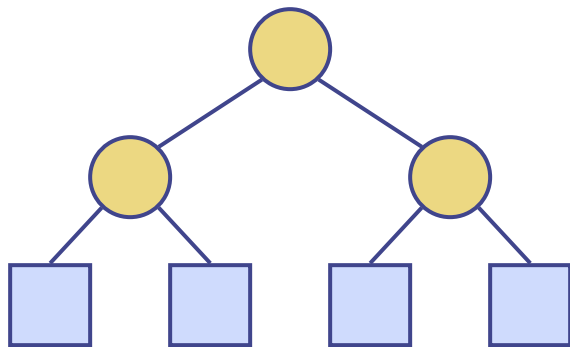
## □ Notation

$n$  number of nodes

$e$  number of  
external nodes

$i$  number of internal  
nodes

$h$  height



## ◆ Properties:

■  $e = i + 1$

■  $n = 2e - 1$

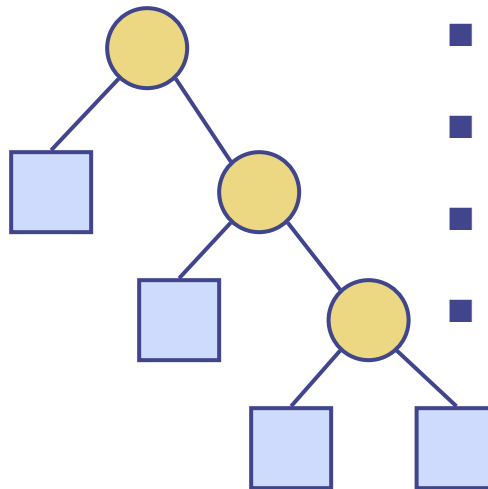
■  $h \leq i$

■  $h \leq (n - 1)/2$

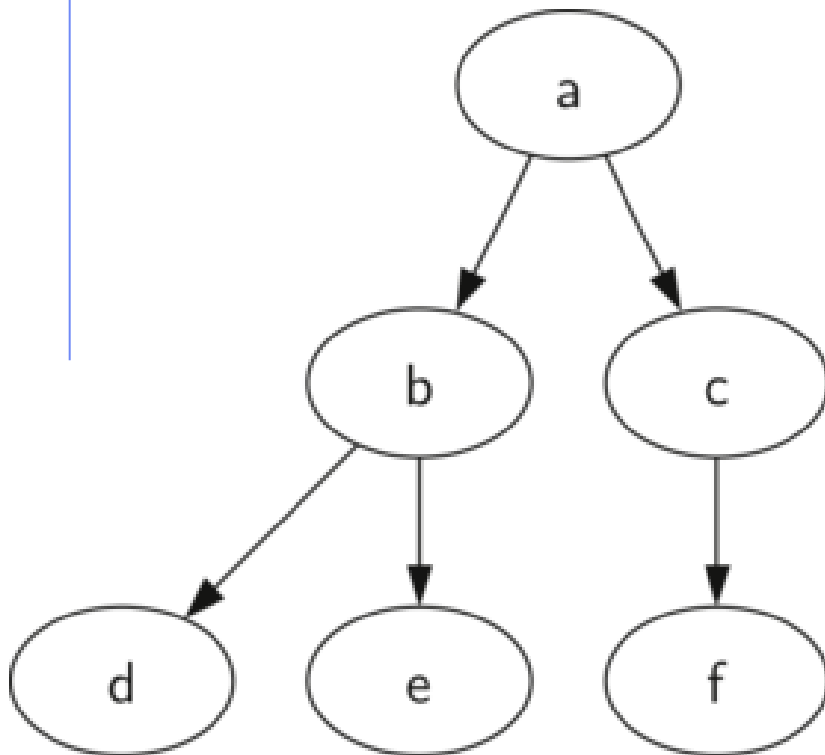
■  $e \leq 2^h$

■  $h \geq \log_2 e$

■  $h \geq \log_2 (n + 1) - 1$



# Tree by a list of lists



```
myTree = ['a',  #root
          ['b',  #left subtree
           ['d', [], []],
           ['e', [], []] ],
          ['c',  #right subtree
           ['f', [], []],
           [] ]
          ]
```

```

1  '''
2      This is an example of a binary tree data structure created with
3      python lists as the underlying data structure.
4  '''
5
6  def BinaryTree(r):
7      return [r, [], []]
8
9  def insertLeft(root,newBranch):
10     t = root.pop(1)
11     if len(t) > 1:
12         root.insert(1,[newBranch,t,[]])
13     else:
14         root.insert(1,[newBranch, [], []])
15     return root
16
17  def insertRight(root,newBranch):
18     t = root.pop(2)
19     if len(t) > 1:
20         root.insert(2,[newBranch,[],t])
21     else:
22         root.insert(2,[newBranch,[],[]])
23     return root

```

```

24
25  def getRootVal(root):
26     return root[0]
27
28  def setRootVal(root,newVal):
29     root[0] = newVal
30
31  def getLeftChild(root):
32     return root[1]
33
34  def getRightChild(root):
35     return root[2]
36
37
38  r = BinaryTree(3)
39  insertLeft(r,4)
40  insertLeft(r,5)
41  insertRight(r,6)
42  insertRight(r,7)
43  l = getLeftChild(r)
44  print(l)
45

```

**insertLeft:** first obtain the (possibly empty) list that corresponds to the current left child. Then add the new left child, installing the old left child as the left child of the new one.



# BinaryTree ADT

- ❑ The BinaryTree ADT extends the Tree ADT, i.e., it inherits all the methods of the Tree ADT
- ❑ Additional methods:
  - position **left**(p)
  - position **right**(p)
  - position **sibling**(p)
- ❑ Update methods may be defined by data structures implementing the BinaryTree ADT

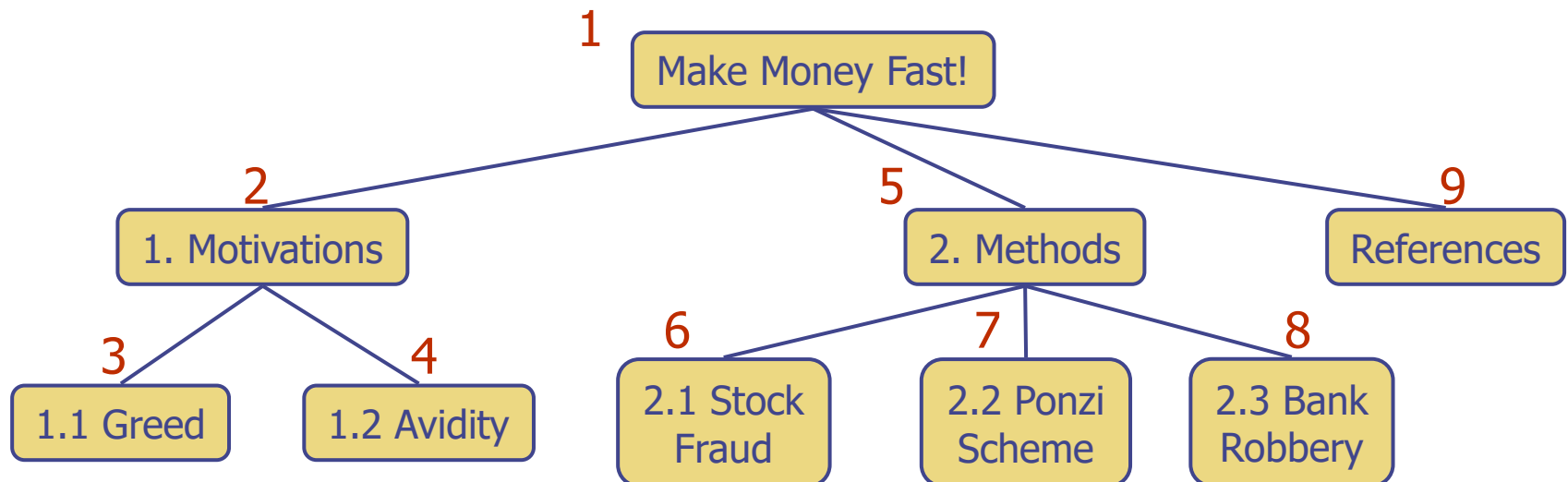
# ADT Binary Tree Class

```
1  class BinaryTree:
2      def __init__(self, rootObj):
3          self.key = rootObj
4          self.leftChild = None
5          self.rightChild = None
6
7      def insertLeft(self, newNode):
8          if self.leftChild == None:
9              self.leftChild = BinaryTree(newNode)
10         else:
11             t = BinaryTree(newNode)
12             t.leftChild = self.leftChild
13             self.leftChild = t
14
15     def insertRight(self, newNode):
16         if self.rightChild == None:
17             self.rightChild = BinaryTree(newNode)
18         else:
19             t = BinaryTree(newNode)
20             t.rightChild = self.rightChild
21             self.rightChild = t
22
23     def getRightChild(self):
24         return self.rightChild
25
26     def getLeftChild(self):
27         return self.leftChild
28
29     def setRootVal(self, obj):
30         self.key = obj
31
32     def getRootVal(self):
33         return self.key
34
```

# Preorder Traversal

- ❑ A traversal visits the nodes of a tree in a systematic manner
- ❑ In a preorder traversal, a node is visited before its descendants
- ❑ Application: print a structured document

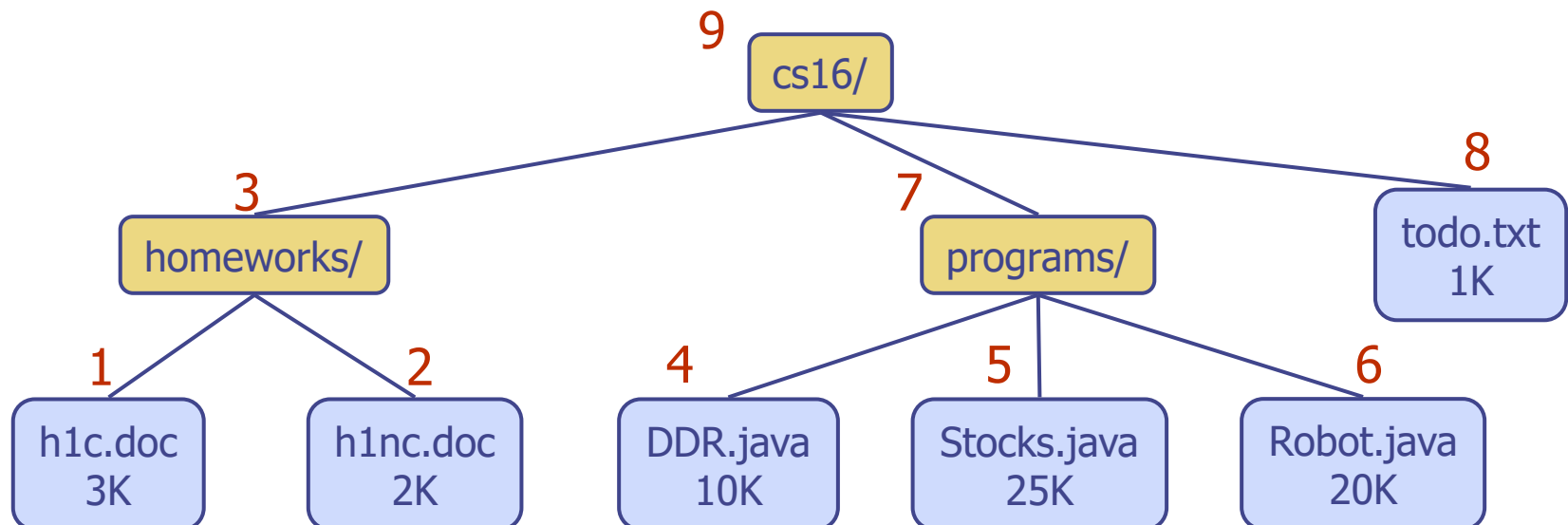
**Algorithm** *preOrder*( $v$ )  
*visit*( $v$ )  
**for each** child  $w$  of  $v$   
*preorder* ( $w$ )



# Postorder Traversal

- In a postorder traversal, a node is visited after its descendants
- Application: compute space used by files in a directory and its subdirectories

**Algorithm** *postOrder*(*v*)  
  **for each** child *w* of *v*  
    *postOrder* (*w*)  
  *visit*(*v*)

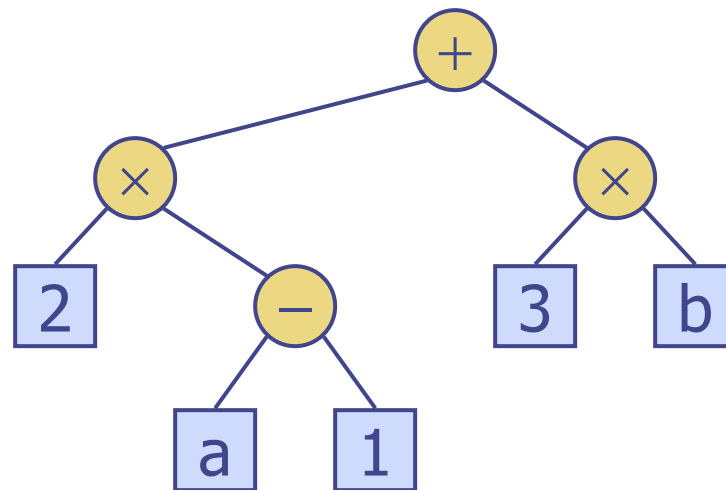


# Preoder & Postorder for Binary Tree

```
35     def preorder(self):
36         print(self.key)
37         if self.leftChild:
38             self.leftChild.preorder()
39         if self.rightChild:
40             self.rightChild.preorder()
41
42     def postorder(self):
43         if self.leftChild:
44             self.leftChild.postorder()
45         if self.rightChild:
46             self.rightChild.postorder()
47         print(self.key)
```

# Arithmetic Expression Tree

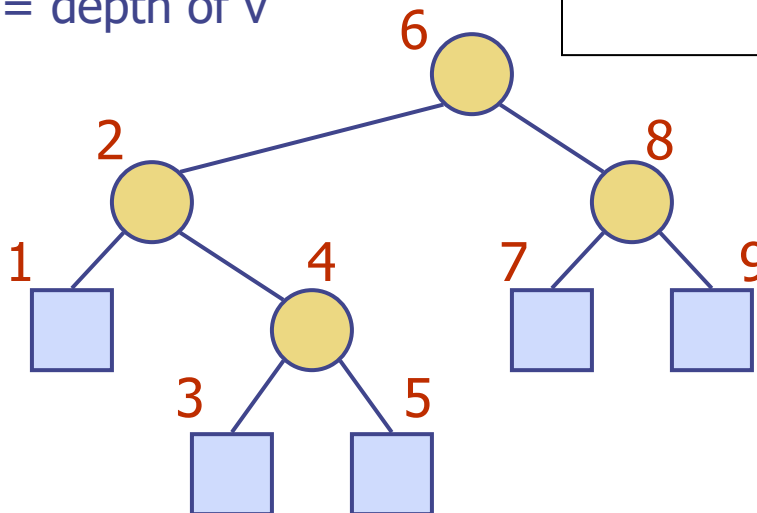
- Binary tree associated with an arithmetic expression
  - internal nodes: operators
  - external nodes: operands
- Example: arithmetic expression tree for the expression  $(2 \times (a - 1) + (3 \times b))$



# Inorder Traversal

- In an inorder traversal a node is visited after its left subtree and before its right subtree
- Application: draw a binary tree
  - $x(v)$  = inorder rank of  $v$
  - $y(v)$  = depth of  $v$

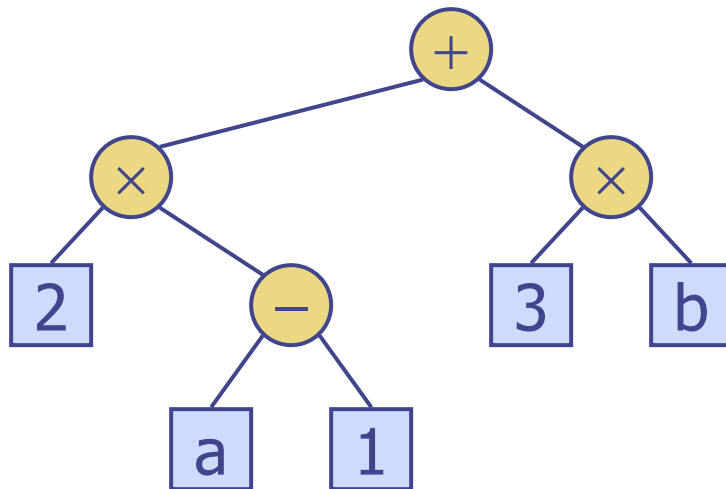
**Algorithm** *inOrder*( $v$ )  
**if**  $v$  **has a left child**  
    *inOrder* (*left* ( $v$ ))  
    *visit*( $v$ )  
**if**  $v$  **has a right child**  
    *inOrder* (*right* ( $v$ ))





# Print Arithmetic Expressions

- Specialization of an inorder traversal
  - print operand or operator when visiting node
  - print “(“ before traversing left subtree
  - print “)” after traversing right subtree



**Algorithm** *printExpression(v)*

**if** *v* **has a left child**

*print*“( ’ ’ )

*inOrder* (*left*(*v*))

*print*(*v.element* ())

**if** *v* **has a right child**

*inOrder* (*right*(*v*))

*print* (“)’ ’ )

$((2 \times (a - 1)) + (3 \times b))$

# Evaluate Arithmetic Expressions

- Specialization of a postorder traversal
  - recursive method returning the value of a subtree
  - when visiting an internal node, combine the values of the subtrees

**Algorithm** *evalExpr(v)*

**if** *is\_leaf*(*v*)

**return** *v.element* ()

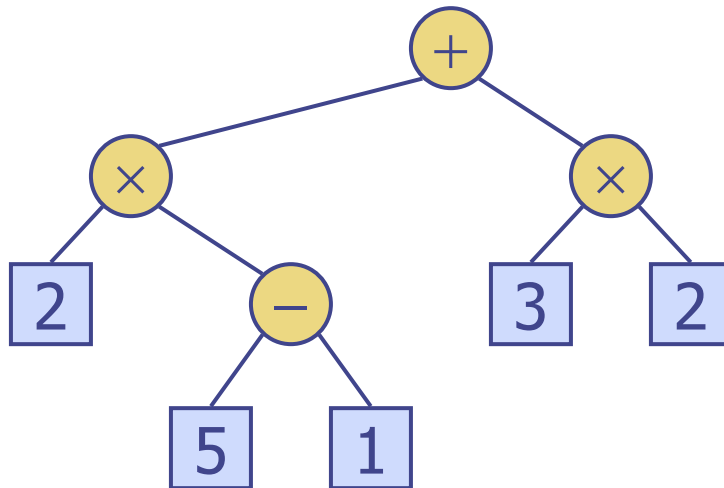
**else**

*x*  $\leftarrow$  *evalExpr*(*left* (*v*))

*y*  $\leftarrow$  *evalExpr*(*right* (*v*))

$\diamond \leftarrow$  operator stored at *v*

**return** *x*  $\diamond$  *y*



# Analysis of Algorithms

- What is the best/worst/average case?  
When?
- Big-O performance

# Reinforcement

# Discussion 2

Draw the tree structure resulting from the following set of tree function calls:

```
>>> r = BinaryTree(3)
>>> insertLeft(r,4)
[3, [4, [], []], []]
>>> insertLeft(r,5)
[3, [5, [4, [], []], []], []]
>>> insertRight(r,6)
[3, [5, [4, [], []], []], [6, [], []]]
>>> insertRight(r,7)
[3, [5, [4, [], []], []], [7, [], [6, [], []]]]
>>> setRootVal(r,9)
>>> insertLeft(r,11)
[9, [11, [5, [4, [], []], []], []], [7, [], [6, [], []]]]
```

# Self check:

```
x = BinaryTree('a')
insertLeft(x, 'b')
insertRight(x, 'c')
insertRight(getRightChild(x), 'd')
insertLeft(getRightChild(getRightChild(x)), 'e')
```

Which of the answers is the correct representation of the tree?

- ☐ A. ['a', ['b', [], []], ['c', [], ['d', [], []]]]
- ☐ B. ['a', ['c', [], ['d', ['e', [], []], []], ['b', [], []]]]
- ☐ C. ['a', ['b', [], []], ['c', [], ['d', ['e', [], []], []]]]
- ☐ D. ['a', ['b', [], ['d', ['e', [], []], []], ['c', [], []]]]

Using binary tree data structure created with python lists.

# Discussion

Consider the following list of integers: [1,2,3,4,5,6,7,8,9,10]. Show the binary search tree resulting from inserting the integers in the list.



# Discussion

Trace the algorithm for creating an expression tree for the expression  $(4*8)/6-3(4*8)/6-3$ .

# Quick overview

- Trees
- Binary Tree with Python list
- Binary Tree with ADT class
- Binary Tree for parsing and evaluating

# Extra reading

- Binary Search Tree (BTS)
- Balanced BTS



# Next week

