Numerische
Mathematik

Tight linear envelopes for splines

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Summary. A sharp bound on the distance between a spline and its B-spline control polygon is derived. The bound yields a piecewise linear envelope enclosing spline and polygon. This envelope is particularly simple for uniform splines and splines in Bernstein-Bézier form and shrinks by a factor of 4 for each uniform subdivision step. The envelope can be easily and efficiently implemented due to its explicit and constructive nature.

Mathematics Subject Classification (1991): 65D07

1 Introduction & main results

B-splines are routinely used as approximating functions and to represent geometry for numerical calculations. A central feature that allows reasoning about these nonlinear piecewise polynomials is the fact that the spline is closely outlined by its *control polygon*, line segments connecting the control points. The efficiency of many applications, for example rendering or intersection testing, hinges on a tight quantitative estimate of the maximal distance of the spline from its control polygon. However, apart from the convex hull and the min–max bound, no quantitive estimates exist to date.

This paper shows how to bound the maximal distance between a spline and its B–spline control polygon in terms of second differences of the control points and linear interpolants of special piecewise convex and nonnegative splines that depend only on the knot sequence. The bounds yield piecewise linear envelopes consisting of a positive and a negative offset of the control polygon (cf. Figs. and 2).
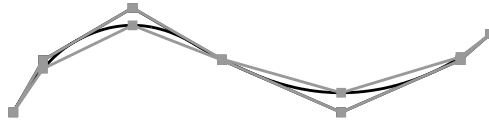

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Fig. 1. A quadratic spline over the knot sequence $(0, 0, 0, 1, 3, 4, 7, 8, 8, 8)$ with control points $(0, 2, 4, 2, 0, 2, 3)$. The spline is shown in black. The envelope according to Corollary 1 is shown in grey. The control polygon is covered by the envelope

Fig. 2. A quintic spline and its envelope according to Corollary 1 . The knots are at $(0, 1, 3, 7)$ and the control points are $(0, 1, 2, 0, 0, 1, 1, 0)$. The first and last knot have multiplicity 6

In most practical applications, the control points are varied far more often than the knot sequence. For fixed knot sequences, including the important cases of Bernstein polynomials and uniform splines, the bounds require only forming the scalar product of a precomputed vector with the vector of second differences of the control points.

The resulting envelopes are as narrow as possible at each control point whose abscissa maps into a convex or concave polynomial piece: one envelope polyline matches that control point while the other touches the spline. This ensures quadratic convergence of the envelope to the spline under subdivision since the number of inflection points of a spline stays fixed while more and more pieces without inflection points are produced.

2 Notation

A piecewise polynomial p of degree d is in B–Spline form if

$$p(t) = \sum_{k \in \mathbb{Z}} b_k N_k^d(t)$$

where the *control points* b_k are real numbers and the B–Spline basis functions N_k^d are defined recursively based on a nondecreasing sequence of real numbers, the *knots* t_k [dB93]. We may assume that p is at least continuous since otherwise we can treat p as two seperate splines. This implies that any knot can appear with multiplicity at most d, except for the first and last knot which can have multiplicity $d + 1$. Therefore the *Greville abscissae* t_k^*,

$$t_k^* := \frac{1}{d} \sum_{i=k+1}^{k+d} t_i,$$

are distinct. We denote the line segment from (t_k^*, a_1) to (t_{k+1}^*, a_2) by

$$\mathcal{L}_k(t \mid a_1, a_2) = a_1 \frac{t_{k+1}^* - t}{t_{k+1}^* - t_k^*} + a_2 \frac{t - t_k^*}{t_{k+1}^* - t_k^*}.$$

and the linear interpolant of the function f at t_k^* and t_{k+1}^* as $\mathcal{L}_k(f) = \mathcal{L}_k(\,\cdot\mid f(t_k^*), f(t_{k+1}^*))$. The *control polygon* $\ell(p)$ of p is the piecewise linear interpolant to control points (t_k^*, b_k). That is, over the interval $[t_k^*, t_{k+1}^*]$ the k–th piece $\ell_k(p)$ of the control polygon $\ell(p)$ is given by

$$\ell_k(p)(t) = \mathcal{L}_k(t \mid b_k, b_{k+1}).$$

We abbreviate $\ell(p)$ to ℓ if the spline p is understood from the context.

The first and second *divided* differences of b are

$$b_i' = \frac{b_i - b_{i-1}}{t_i^* - t_{i-1}^*} \qquad b_i'' = (d - 1)\frac{b_i' - b_{i-1}'}{t_{i+d-1} - t_i}.$$

They are the B–spline control points of the first and second derivatives of p. The *centered* second differences of b are defined as

$$\Delta_2 b_i = b_{i+1}' - b_i'.$$

Only finitely many B–spline basis functions have support on some part of $[t_k^*, t_{k+1}^*]$. Let $\underline{k}$ and $\overline{k}$ be the index of the first and the last B–spline basis function that are nonzero on some part of $[t_k^*, t_{k+1}^*]$. There will be more than $d + 1$ such basis functions if $[t_k^*, t_{k+1}^*]$ contains one or more knots in its interior. Since $t_k^* \leq t_{k+d} \leq t_{k+d+1} \leq t_{k+1}^*$, $\underline{k}$ and $\overline{k}$ are related to k by $\underline{k} \leq k < \overline{k}$.

3 Tight envelopes for arbitrary knot sequences

The key to deriving the bounds in this paper is to factor the difference between a spline and its B–spline control polygon into two parts: the second differences of the control polygon and nonnegative splines β_{ki} that only depend on the knot sequence. The β_{ki} are piecewise convex and therefore well suited for piecewise linear approximations. This observation is made precise in the following theorem.

Theorem 1 *Let* $t \in [t_k^*, t_{k+1}^*]$ *and functions* β_{ki} *be defined as*

(3.1) $$\beta_{ki} = \begin{cases} \sum_{j=i}^{\overline{k}}(t_j^* - t_i^*)N_j^d & i > k \\ \sum_{j=\underline{k}}^{i}(t_i^* - t_j^*)N_j^d & i \leq k. \end{cases}$$

The difference between the spline p and its B–spline control polygon ℓ is

(3.2)
$$p - \ell = \sum_{i=\underline{k}}^{\overline{k}} \Delta_2 b_i \, \beta_{ki}.$$

The functions β_{ki} are nonnegative and convex.

Proof. We write $p - \ell_k$ over $[t_k^*, t_{k+1}^*]$ as

$$\sum_i b_i \, \alpha_{ki} = \sum_i b_i \left(N_i^d(t) - \mathcal{L}_k(t \mid \delta_{ik}, \delta_{i,k+1}) \right),$$

where $\delta_{ik} = 1$ if $i = k$ and 0 otherwise. Only the α_{ki} with $\underline{k} \leq i \leq \overline{k}$ can be nonzero on $[t_k^*, t_{k+1}^*]$.

We show that $\alpha_{ki} = \Delta_2 \beta_{ki}$: the partition of unity $\sum_i N_i^d = 1$ implies that $\sum_i \alpha_{ki} = 0$ and the linear precision of B–Splines, $\sum_i t_i^* N_i^d = t$ implies on the interval $[t_k^*, t_{k+1}^*]$ that $\sum_i t_i^* \alpha_{ki} = 0$. Hence, for any i, $\sum_{j=\underline{k}}^{\overline{k}} (t_j^* - t_i^*) \alpha_{kj} = 0$.

For $i > k$,

$$\beta_{ki} = \sum_{j=i+1}^{\overline{k}} (t_j^* - t_i^*) N_j^d - \sum_{j=\underline{k}}^{\overline{k}} (t_j^* - t_i^*) \alpha_{kj} = \sum_{j=\underline{k}}^{i} (t_i^* - t_j^*) \alpha_{kj}.$$

Hence, for any i, $\beta_{ki} = \sum_{j=\underline{k}}^{i} (t_i^* - t_j^*) \alpha_{kj}$ which allows us verify that $\Delta_2 \beta_{ki} = \alpha_{ki}$. Summation by parts then completes the proof of (3.2),

$$p - \ell_k = \sum_{i=\underline{k}}^{\overline{k}} b_i \, \alpha_{ki} = \sum_{i=\underline{k}}^{\overline{k}} b_i \, \Delta_2 \beta_{ki} = \sum_{i=\underline{k}}^{\overline{k}} \Delta_2 b_i \, \beta_{ki}.$$

The functions β_{ki} are nonnegative since their B–spline coefficients are nonnegative by (3.1).

The convexity of the β_{ki} over $[t_k^*, t_{k+1}^*]$ follows from the convexity of their B–spline control polygons: for $i > k$, the part of the control polygon of β_{ki} that influences β_{ki} over $[t_k^*, t_{k+1}^*]$ lies on the function $\max\{\cdot - t_i^*, 0\}$ while for $i \leq k$ it lies on $\max\{t_i^* - \cdot, 0\}$. In both cases, the control polygon of β_{ki}, and hence β_{ki}, is nonnegative and convex.

The characterization of $p - \ell$ from Theorem 1 gives us a piecewise linear envelope for p that improves in practice on the convex hull, for example for the splines in Figs. 1 and 2.

Corollary 1 *Over the interval $[t_k^*, t_{k+1}^*]$, the spline p is enveloped by linear offsets of the control polygon ℓ,*

$$(3.3) \qquad \ell + \mathcal{L}_k \left(\sum_{i=\underline{k}}^{\overline{k}} \Delta_2^- b_i \, \beta_{ki} \right) \leq p \leq \ell + \mathcal{L}_k \left(\sum_{i=\underline{k}}^{\overline{k}} \Delta_2^+ b_i \, \beta_{ki} \right),$$

where $\Delta_2^+ b_i = \max\{\Delta_2 b_i, 0\}$ and $\Delta_2^- b_i = \min\{\Delta_2 b_i, 0\}$ and $\mathcal{L}_k(f)$ the linear interpolant to f at t_k^* and t_{k+1}^*.

Proof. We have from Theorem 1

$$p - \ell = \sum_{i=\underline{k}}^{\overline{k}} \Delta_2 b_i \, \beta_{ki} = \sum_{i=\underline{k}}^{\overline{k}} \Delta_2^+ b_i \, \beta_{ki} + \sum_{i=\underline{k}}^{\overline{k}} \Delta_2^- b_i \, \beta_{ki}.$$

The positivity of the β_{ki} implies that the first sum on the right–hand side is positive and the second is negative and that therefore

$$\sum_{i=\underline{k}}^{\overline{k}} \Delta_2^- b_i \, \beta_{ki} \leq p - \ell \leq \sum_{i=\underline{k}}^{\overline{k}} \Delta_2^+ b_i \, \beta_{ki}.$$

Since the β_{ki} are convex over $[t_k^*, t_{k+1}^*]$, they can be bounded linearly to yield (3.3).

3.1 Implementation

The envelopes of Corollary 1 depend on the control points b_i, their second differences $\Delta_2 b_i$, and the values $\beta_{ki}(t_k^*)$. In evaluating either of the sums in equation (3.3) for some $t \in \{t_k^*, t_{k+1}^*\}$, it suffices to sum over the subset $\mathcal{I}(t)$,

$$(3.4) \qquad\qquad\qquad \mathcal{I}(t) = \{\, i : \beta_{ki}(t) > 0 \,\},$$

of the integers between $\underline{k}$ and $\overline{k}$. $\mathcal{I}(t)$ is determined by equation (3.1): let $t \in [t_k^*, t_{k+1}^*]$ and let $t_l < t \leq t_{l+1}$ be the knots immediately to the left and right of t (see Fig. 3). Since $N_j^d(t)$ can only be non-zero if $l - d \leq j \leq l$ and since $\beta_{k,l-d}(t) = \beta_{kl}(t) = 0$,

$$\mathcal{I}(t) = \{\, i : l - d < i < l \text{ for } t_l < t \leq t_{l+1} \,\}.$$

For an efficient implementation, we precompute the $\mathcal{I}(t_k^*)$ and the values $\beta_{ki}(t_k^*)$, $i \in \mathcal{I}(t_k^*)$. Computing the lower and upper envelope from Corollary 1 for any p over the fixed knot sequence (t_k) then only requires computing one second difference and one scalar product with $d - 1$ terms for each Greville abscissa t_k^*.
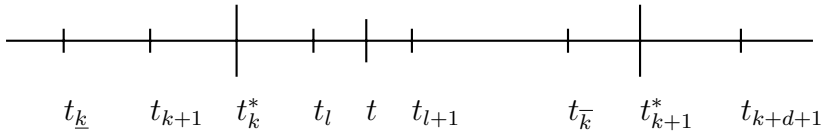
$$t_{\underline{k}} \qquad t_{k+1} \quad t_k^* \quad t_l \quad t \quad t_{l+1} \qquad t_{\overline{k}} \quad t_{k+1}^* \quad t_{k+d+1}$$

Fig. 3. For a given $t \in [t_k^*, t_{k+1}^*]$, the set $\mathcal{I}(t) = \{\, i : \beta_{ki}(t) > 0 \,\}$ is determined by the positions of the displayed knots relative to one another
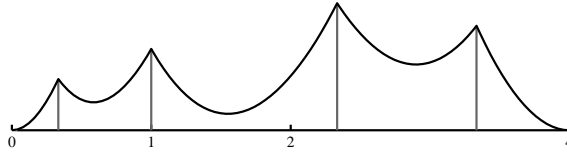

Fig. 4. The function z for cubic splines over the knot sequence $(0, 0, 0, 0, 1, 2, 4, 4, 4, 4)$. The grey lines indicate the position of the Greville abscissae. Knots are shown as ticks on the x–axis. The values of z range from 0 to 0.26

3.2 Sharpness and convergence

The envelope from Corollary 1 is sharp in the sense that the upper (lower) part of the envelope *equals* p at t_k^* if $\Delta_2 b_i \geq 0$ ($\Delta_2 b_i \leq 0$) for all $i \in \mathcal{I}(t_k^*)$. This is the case if the whole polynomial piece of p on which $p(t_k^*)$ lies is convex (concave). In particular, the envelope is sharp for splines of degree 2 since $\mathcal{I}(t_k^*)$ contains only one element.

Sharpness implies that, under subdivison, the width of the envelope shrinks quadratically in the distance of the knots: since the envelope consists of linear pieces, the maximum width is attained at the break points. Since a spline has only finitely many inflection points almost all pieces generated by subdivision are either concave or convex for which the width of the envelope equals the distance of the spline to the control polygon which is known to decrease quadratically [CS85, Dah86].

4 Coarse envelopes

The envelope in Corollary 1 requires the evaluation of $d - 1$ splines β_{ki} at each Greville abscissa t_k^*. We can reduce the computational cost to just one evaluation per Greville abscissa by subsuming the values of the β_{ki} into one spline ζ and estimating its values. The resulting envelopes, though very simple to compute, are in general considerably larger than the ones from Corollary 1. They do, however, give valuable insight into the asymptotic behavior of the envelope.

Theorem 2 *The difference between the spline p and its B–spline control polygon $\ell(p)$ over the interval $[t_k^*, t_{k+1}^*]$ is bounded by*

(4.5) $$|p - \ell_k(p)| \le \mathcal{M} \, |\zeta - \ell_k(\zeta)| = \mathcal{M} \, z.$$

where $\mathcal{M}(t) = \max \{ \, |\Delta_2 b_i| : i \in \mathcal{I}(t) \, \}$. The splines ζ,

(4.6) $$\zeta = \sum_j \zeta_j N_j^d \qquad where \qquad \zeta_j = jt_j^* + \sum_{i=j}^{\overline{k}} t_i^*,$$

and $z = \zeta - \ell(\zeta)$ are nonnegative and convex.

Proof. Using Theorem 1, we have

$$|p - \ell_k(p)| = \left| \sum_{i=\underline{k}}^{\overline{k}} \Delta_2 b_i \, \beta_{ki} \right| \le \mathcal{M} \left| \sum_{i=\underline{k}}^{\overline{k}} \beta_{ki} \right| = \mathcal{M} \sum_{i=\underline{k}}^{\overline{k}} \beta_{ki}.$$

Equation (4.5) follows if we can show that $\sum_i \beta_{ki} = z$.

$$\sum_{i=\underline{k}}^{\overline{k}} \beta_{ki} = \sum_{i=\underline{k}}^{\overline{k}} \sum_{j=\underline{k}}^{i} (t_i^* - t_j^*) \alpha_{kj} = \sum_{j=\underline{k}}^{\overline{k}} \sum_{i=j}^{\overline{k}} (t_i^* - t_j^*) \alpha_{kj}$$

$$= \sum_{j=\underline{k}}^{\overline{k}} \left[\sum_{i=j}^{\overline{k}} t_i^* - (\overline{k} - j + 1) t_j^* \right] \alpha_{kj} = \sum_{j=\underline{k}}^{\overline{k}} \left[jt_j^* + \sum_{i=j}^{\overline{k}} t_i^* \right] \alpha_{kj}$$

(4.7) $$= \sum_{j=\underline{k}}^{\overline{k}} \zeta_j \, \alpha_{kj} = \sum_{j=\underline{k}}^{\overline{k}} \zeta_j N_j^d - \mathcal{L}_k(\,\cdot\, | \, \zeta_k, \zeta_{k+1}) = z.$$

It is easily checked that $\zeta_j'' > 0$. The functions ζ and $z = \zeta - \ell(\zeta)$ are therefore strictly convex on $[t_k^*, t_{k+1}^*]$. The coefficients of ζ and z are nonnegative.

Corollary 2 *The difference between p and its control polygon ℓ over the interval $[t_k^*, t_{k+1}^*]$ is bounded by*

$$|p - \ell| \le \mathcal{L}_k \left(\mathcal{M} z \right).$$

Proof. By Theorem 1 and the convexity of the β_{ki}, we have

$$|p - \ell| \le \sum_{i=\underline{k}}^{\overline{k}} |\Delta_2 b_i| \, \beta_{ki} \le \mathcal{L}_k \left(\sum_{i=\underline{k}}^{\overline{k}} |\Delta_2 b_i| \, \beta_{ki} \right).$$

The definition of $\mathcal{M}$ and equation (4.7) in the proof of Theorem 2 show that

$$\mathcal{L}_k \left(\sum_{i=\underline{k}}^{\overline{k}} |\Delta_2 b_i| \, \beta_{ki} \right) \le \mathcal{L}_k \left(\mathcal{M} \sum_{i=\underline{k}}^{\overline{k}} \beta_{ki} \right) = \mathcal{L}_k \left(\mathcal{M} z \right).$$
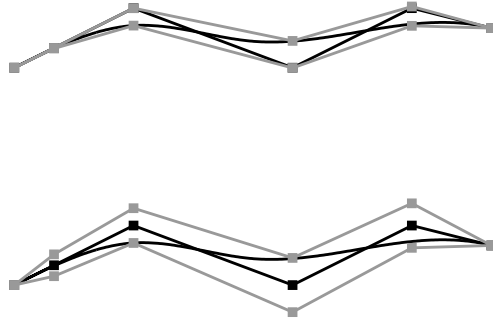
Fig. 5. A cubic spline and its envelope according to Corollary 1 *(top)* and Corollary 2 *(bottom)* over the knot sequence $(0, 0, 0, 0, 1, 2, 4, 4, 4, 4)$ with control points $(0, 1, 3, 0, 3, 2)$

4.1 Bounding with a quadratic spline

The same considerations as in Theorem 2 but with divided instead of centered second differences yield a quadratic bounding spline.

Corollary 3 *The difference between the spline p and its control polygon $\ell(p)$ over the interval $[t_k^*, t_{k+1}^*]$ is bounded by*

$$|p - \ell(p)| \leq \widehat{\mathcal{M}}(t) \, |\eta - \ell(\eta)|$$

where $\widehat{\mathcal{M}}(t) = \max \{ |b_i''| : i \in \mathcal{I}(t) \}$. *The function* $\eta = \sum \eta_j N_j^d$ *is a quadratic spline with* $\eta'' = 1$.

Proof. Using divided second differences b_i'', equation (3.2) becomes

$$p - \ell = \sum_{i=\underline{k}}^{\overline{k}} b_{i+1}'' \frac{t_{i+d} - t_{i+1}}{d - 1} \beta_{ki}.$$

Since, in analogy to the proof of Theorem 2,

$$\eta - \ell(\eta) = \sum_{i=\underline{k}}^{\overline{k}} \frac{t_{i+d} - t_{i+1}}{d - 1} \beta_{ki},$$

η is given by

$$\eta = \sum_{i=\underline{k}}^{\overline{k}} \sum_{j=\underline{k}}^{i} \frac{t_{i+d} - t_{i+1}}{d - 1} (t_i^* - t_j^*) N_j^d$$

$$= \sum_{j=\underline{k}}^{\overline{k}} \sum_{i=j}^{\overline{k}} \frac{t_{i+d} - t_{i+1}}{d - 1} (t_i^* - t_j^*) N_j^d.$$

It is now easy to verify that $\eta_j'' = 1$.

4.2 Implementation

The envelope from Corollary 2 requires $\mathcal{M}(t_k^*)$ and $z(t_k^*)$ for finite knot sequences $(t_0, \ldots, t_{m+d})$ and corresponding Greville abscissae $t_0^*, \ldots, t_m^*$.

At first sight, the definition of the coefficients ζ_j in equation (4.6) seems particular to the interval $[t_k^*, t_{k+1}^*]$ since it involves $\overline{k}$. However, we can extend the upper limit of summation in this definition to m, since only the values of $z = \zeta - \ell(\zeta)$ are needed and these are not affected by this change.

An efficient implementation of the bound of Corollary 2 stores the values t_k^*, $z(t_k^*)$, $\mathcal{I}(t_k^*)$ together with the knot sequence, so that only the maximal second difference $\mathcal{M}(t_k^*)$ needs to be recomputed whenever the control points of a spline change.

4.3 Sharpness

The proof of Theorem 2 contains only one inequality,

$$\left| \sum_{i \in \mathcal{I}(t)} \Delta_2 b_i \, \beta_{ki}(t) \right| \leq \mathcal{M}(t) \left| \sum_{i \in \mathcal{I}(t)} \beta_{ki}(t) \right|.$$

This is an equality exactly when all the second differences $\Delta_2 b_i, i \in \mathcal{I}(t)$, are equal to one another. The splines with this property are the linear polynomials and ζ.

Similarly, the bound from Corollary 3 is sharp for the linear polynomials and η, and therefore for all quadratic polynomials.

5 A bound for Bernstein polynomials

The B–splines N_j^d of degree d over the knot sequence $(t_0, \ldots, t_{2d})$ with $t_i = 1$ for $i > d$ and $t_i = 0$ else are called the Bernstein polynomials B_j^d of degree d. An explicit bound for the Bernstein polynomials was already given in [NPL99, Theorem 3.1]; we give an alternative proof for this theorem, showing how it can be derived from the more general exposition in this paper.

For Bernstein polynomials, the Greville abscissae are $t_i^* = i/d$ and the coefficients ζ_j are

$$\zeta_j = \frac{j(j+1) + d(d+1)}{2d}.$$

Thus,

$$z = \sum_j \zeta_j \alpha_{kj} = \frac{1}{d} \sum_j \binom{j}{2} \alpha_{kj} = \frac{1}{d} \left(\sum_j \binom{j}{2} B_j^d - \frac{k(2dt - k - 1)}{2} \right)$$

$$= \frac{d-1}{2} t^2 - \frac{k(2dt - k - 1)}{2d}$$

and

$$z(t_k^*) = \frac{k(d-k)}{2d^2}, \qquad z(t_{k+1}^*) = \frac{(k+1)(d-(k+1))}{2d^2}$$

at the endpoints of the interval $[t_k^*, t_{k+1}^*]$. The maximal value $z(t_k^*)$ over all k is taken on for $k = \lfloor d/2 \rfloor$. Since $\Delta_2 b_i = d(b_{i-1} - 2b_i + b_{i+1})$ the main theorem from [NPL99] follows:

Theorem 3 ([NPL99, Theorem 3.1]) *Let $p = \sum_{i=0}^d b_i B_i^d$ be a polynomial of degree d in Bernstein–Bézier form. The distance between p and its control polygon ℓ is uniformly bounded by*

$$|p - \ell| \leq \frac{\lfloor d/2 \rfloor \lceil d/2 \rceil}{2d} \max_{i=1}^{d-1} |\Delta_2 b_i|$$

where $\Delta_2 b_i = b_{i-1} - 2b_i + b_{i+1}$.

This bound decreases by a factor of 4 under subdivision at $1/2$ [NPL99, Lemma 6.1].

6 A bound for uniform splines

A spline is *uniform* if the knots are all equidistant. Without loss of generality, we choose $t_k = k$. The uniform B–spline basis functions N_j^d are shifts of one another and we define $N^d = N_0^d$ so that $N_j^d(t) = N^d(t - j)$. The corresponding Greville abscissae are

$$t_k^* = \frac{1}{d} \sum_{i=k+1}^{k+d} i = k + \frac{d+1}{2}$$

and ζ has the B–spline coefficients

$$\zeta_j = j\left(j + \frac{d+1}{2}\right) + \sum_{i=j}^m \left(i + \frac{d+1}{2}\right)$$

$$= (m+1)\left(j + \frac{d+1}{2}\right) + \binom{j-m}{2}.$$

According to Equation (4.7), $z = \sum_j \zeta_j \alpha_{kj}$. For uniform splines this simplifies to

$$z = \sum_j \zeta_j \alpha_{kj} = \sum_{j=\underline{k}}^{\overline{k}} \left[(m+1)(j + \frac{d+1}{2}) + \binom{j-m}{2} \right] \alpha_{kj}$$

$$= \sum_j \binom{j-m}{2} \alpha_{kj} = \sum_j \binom{j-k}{2} N_j^d.$$

Regardless of the degree of p, z is a quadratic polynomial and has the monomial form

$$z(t) = \frac{1}{2} \left(t^2 - (t_k^* + t_{k+1}^*)t + \left(\frac{t_k^* + t_{k+1}^*}{2} \right)^2 + \frac{d-2}{12} \right).$$

Since z is a positive and convex function, it attains its maximum over $[t_k^*, t_{k+1}^*]$ at one of the endpoints of the interval. Its values there are

$$z(t_k^*) = z(t_{k+1}^*) = \frac{d+1}{24}.$$

This proves the following simplified version of Corollary 2:

Corollary 4 *Let $p = \sum_j b_j N_j^d$ be a uniform spline over the knot sequence $t_k = k$. Over the interval $[t_k^*, t_{k+1}^*]$, the difference between p and its control polygon ℓ is bounded by*

$$|p - \ell| \leq \frac{d+1}{24} \mathcal{L}_k (\mathcal{M}).$$

For uniform splines, the set $\mathcal{I}(t_k^*)$ from equation (3.4) is

$$\mathcal{I}(t_k^*) = \{ i : k - \lfloor d/2 \rfloor < i < k + \lfloor d/2 \rfloor \}.$$

Since $z(t_k^*) = (d+1)/24$ and $\mathcal{I}(t_k^*) = \{k\}$ for $d = 2$ or $d = 3$, Corollary 1 reduces for quadratic and cubic uniform splines to

(6.8)
$$\frac{d+1}{24} \mathcal{L}_k(\cdot \mid \Delta_2^- b_k, \Delta_2^- b_{k+1}) \leq p - \ell \leq \frac{d+1}{24} \mathcal{L}_k(\cdot \mid \Delta_2^+ b_k, \Delta_2^+ b_{k+1}).$$

7 Uniform refinement

For uniform splines, we consider the refinement of the knot sequence $t_k = k$ to the sequence $\hat{t}_k = k/2$. We now have two representations for p,

$$p(t) = \sum_k b_k N^d(t - k) = \sum_k \hat{b}_k N^d(2t - k),$$

where the new control points $\hat{b}_k$ are

(7.9)

$$\hat{b}_{2i} = 2^{-d} \sum_{j=0}^{\lceil d/2 \rceil} \binom{d+1}{2j} b_{i-j}, \qquad \hat{b}_{2i+1} = 2^{-d} \sum_{j=0}^{\lceil d/2 \rceil} \binom{d+1}{2j+1} b_{i-j}.$$

Since $a_i'' = \Delta_2 a_{i-1}$, we can use the B–spline representation of p'' together with equation (7.9) to relate the centered second differences of the new control polygon to those of the old control polygon as

$$2^d \Delta_2 \hat{b}_{2i} = \sum_j \binom{d-1}{2j-1} \Delta_2 b_{i-j}, \qquad 2^d \Delta_2 \hat{b}_{2i+1} = \sum_j \binom{d-1}{2j} \Delta_2 b_{i-j}.$$

The equality $\sum_j \binom{d-1}{j} = 2^{d-1}$ and the symmetry of the binomial coefficients imply

$$\sum_j \binom{d-1}{2j} = \sum_j \binom{d-1}{2j-1} = 2^{d-2},$$

so that

$$\max_i \left| \Delta_2 \hat{b}_i \right| \leq \frac{1}{4} \max_i |\Delta_2 b_i|.$$

This ensures that the bound from Corollary 4 converges quadratically to zero under repeated uniform refinement. Figure 6 illustrates this for a cubic uniform spline.

7.1 Examples

For quadratic splines, uniform refinement is called Chaikin's algorithm and

$$\hat{b}_{2i} = 2^{-2}(3b_{i-1} + b_i), \qquad \hat{b}_{2i+1} = 2^{-2}(b_{i-1} + 3b_i).$$

This yields

$$\Delta_2 \hat{b}_{2i} = \Delta_2 \hat{b}_{2i-1} = \frac{1}{4} \Delta_2 b_{i-1}.$$
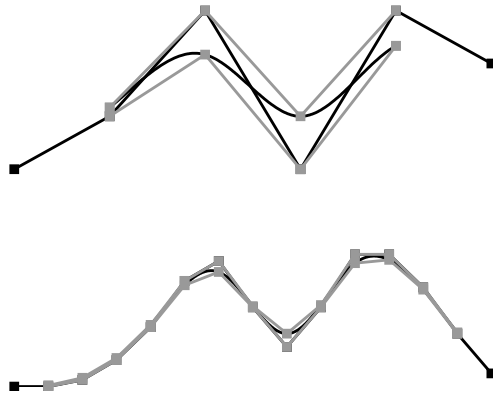
Fig. 6. A uniform cubic spline with control points $(0, 1, 3, 0, 3, 2)$. Shown are the envelope according to equation (6.8) *(top)* and the envelope after one step of uniform refinement *(bottom)*

Since every second difference decreases by a factor of four, subsequent envelopes are contained in one another.

Similarly, for cubic splines we have

$$\hat{b}_{2i} = 2^{-3}(b_{i-2} + 6b_{i-1} + b_i), \qquad \hat{b}_{2i+1} = 2^{-3}(4b_{i-1} + 4b_i).$$

and therefore

$$\Delta_2 \hat{b}_{2i} = \frac{1}{4} \Delta_2 b_{i-1}, \qquad \Delta_2 \hat{b}_{2i+1} = \frac{1}{8}(\Delta_2 b_{i-1} + \Delta_2 b_i).$$

8 Conclusion

We considered envelopes for spline *functions*. The extension to parametric curves can be computed as follows. First, at each control point, compute an axis–aligned box from the envelopes of each coordinate function. Secondly, for any two consecutive control points, form the convex hull of their boxes. The union of the convex hulls is guaranteed to contain the curve by the linearity of the envelopes.

The envelopes developed in this paper can be generalized to tensor-product splines and other bivariate bases commonly used in CAGD. The construction follows the principles outlined in this paper: rewrite $p - \ell$ in terms of second differences in a new basis and estimate the range of the basis functions at the Greville abscissae. Separating the vector of second differences of p into a negative and a positive part yields envelopes as piecewise offsets of the control polygon.

Added in print: In [Rei99], Reif gives an elegant and explicit formula for the bound in Corollary 3: for a spline p of degree d the difference between

p and its control polygon $\ell(p)$ is bounded by

$$|p - \ell(p)| \leq \frac{\sup_j \sigma_j^2}{2d} \max_i \left| b_i'' \right|$$

where σ_j^2 is the variance of d consecutive knots. Lutterkort's thesis [Lut00] generalizes the approach taken in this paper to a large class of refinable functions, e.g. bounding subdivision curves that do not have a closed representation. For applications of the spline bounds to pathfinding and one-sided approximation see [LP99].

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