

LINEAR ALGEBRA Cliff Notes

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- We'll work in \mathbb{R}^n usually (or an isomorphic space like $\mathbb{R}^{m \times n}$) and use its vector space properties.

- Measure "size" w/ a norm $\|\cdot\|$. Let V be a vector space

Def a norm $\|\cdot\| : V \rightarrow \mathbb{R}_+$ (\mathbb{R}_+ = non-negative \mathbb{R} numbers)

satisfies ① $\|x\| = 0$ iff $x = 0$

② $\forall x \in F, \|\alpha x\| = |\alpha| \cdot \|x\|$, F is a field, for us \mathbb{R} or \mathbb{C}

③ $\|x+y\| \leq \|x\| + \|y\|$ "triangle ineq." or " Δ -ineq."

- Measure "distance" w/ metric induced by norm

$$d(x, y) = \|x - y\|.$$

Fact: reverse Δ -ineq: $\|x\| = \|x + y - y\| \leq \|x + y\| + \|y\|$
 $\|x\| = \|x - y + y\| \leq \|x - y\| + \|y\|$

so ① $\|x + y\| \geq \|x\| - \|y\|$

and ② $\|x + y\| \geq \|y\| - \|x\|$ by interchanging x, y

③ $\|x - y\| \geq \|x\| - \|y\|$ ④ $\|x - y\| \geq \|y\| - \|x\|$

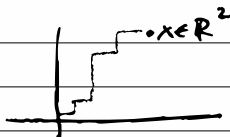
Ex: Euclidean norm/metric, $\|x\|_2 = \sqrt{\sum_{i=1}^n |x_i|^2}$ for $x \in \mathbb{R}^n$

"distance as the crow flies"

Ex: Mahalanobis $\|x\|_V = \sqrt{x^T V x}$ for a matrix $V \succ 0$ pos. def.

Ex: Taxi-cab / Manhattan distance (i.e. ℓ_1 norm)

$$\|x\|_1 = \sum |x_i|$$



Inner product: For a vector space (V, F) , $\langle \cdot, \cdot \rangle : V \times V \rightarrow F$

is bi-linear (well, almost, it have \mathbb{C} -numbers)

"angles"

Usually use Euclidean inner product $\langle x, y \rangle = x^T y = \sum x_i y_i$

which induces $\|\cdot\|_2$ ($\|x\|_2^2 = \langle x, x \rangle$)

Fact: $\|\cdot\|_p$ ($p \neq 2$) norms not induced by any inner product

if $V = \mathbb{C}^n$, two options: ① $F = \mathbb{R}$, then output of $\langle \cdot, \cdot \rangle$ must be \mathbb{R}

x^* means \overline{x}^T ← transpose
 ↗ adjoint

So $\langle x, y \rangle = \text{Re}(x^* y)$

② $F = \mathbb{C}$, then $\langle x, y \rangle = x^* y$

Positive definite matrix

$A > 0$ means A is p.d., meaning $\forall x \neq 0, \langle x, Ax \rangle > 0$
 Usually we implicitly impose that $A = A^*$
 in which case $A > 0$ iff all eigenvalues $\lambda_i > 0$
 ↗ "squeggly >"

and pos. semidefinite (psd) iff all $\lambda_i \geq 0$.

Recall a matrix is diagonalizable if $A = VDV^{-1}$

$A = \begin{bmatrix} 0 & 1 \\ 0 & 0 \end{bmatrix}$ not diagonalizable.

V invertible
 D diag.

If A is normal ($AA^* = A^*A$) \Rightarrow it is unitarily diagonalizable
 $A = VDV^{-1} = VDV^*, \quad V^{-1} = V^*$

If A is self-adjoint / Hermitian ($A = A^*$) \Rightarrow it is normal
 and $D_{ii} = \lambda_i \in \mathbb{R}$

Matrix norms

• Frobenius / Hilbert-Schmidt $\|A\|_F = \sqrt{\sum_{i,j} |A_{ij}|^2} = \langle A, A \rangle$

where $\langle A, A \rangle = \text{tr}(A^*A) = \text{vec}(A)^* \text{vec}(A)$

$\text{vec} : \mathbb{R}^{m \times n} \rightarrow \mathbb{R}^{mn \times 1}$
 $\text{mat} : \mathbb{R}^{mn \times 1} \rightarrow \mathbb{R}^{m \times n}$ } inverses and adjoints of each other

• Operator norm

$\|A\|_{a \rightarrow b} := \sup_{x \neq 0} \frac{\|Ax\|_b}{\|x\|_a} = \sup_{\|x\|_a = 1} \|Ax\|_b$

Submultiplicative: $\|Ax\|_b \leq \|A\|_{a \rightarrow b} \cdot \|x\|_a$

the most common operator norm is the "spectral norm"
using $\|\cdot\|_a = \|\cdot\|_b = \|\cdot\|_2$ (Euc. norm)

$\|A\|_{2 \rightarrow 2}$, usually written $\|A\|_2$ or just $\|A\|$

⚠ in Matlab, $\text{norm}(A)$ is spectral norm by default
in Python, $\text{norm}(A)$ is Frobenius norm by default

$\|A\|_2$ calculated via SVD: $A = U \Sigma V^*$
 $\Sigma = \text{diag}(\sigma_1, \dots, \sigma_r)$ then $\|A\|_2 = \|\sigma\|_\infty$

$$\|A\|_F = \|\sigma\|_2$$

If A is $n \times n$, $\|A\|_2$ requires $O(n^3)$ flops
 $\|A\|_F$ requires $O(n^2)$ flops

In finite dimensions (like \mathbb{R}^n), all norms are equivalent,
which means for norms $\|\cdot\|_a, \|\cdot\|_b$, $\exists c, C > 0$
s.t.
 $\forall x \in \mathbb{R}^n, \quad c \|x\|_a \leq \|x\|_b \leq C \cdot \|x\|_a$

Equivalent \neq the same. Only means they induce the
same topology

Open sets: $U \subseteq \mathbb{R}^n$ is open if $\forall x \in U, \exists \varepsilon > 0$ s.t.

the open ball $B_\varepsilon(x) \subseteq U$

Closed set: $F \subseteq \mathbb{R}^n$ closed if F^c is open

or, equiv., if it is sequentially closed, $(x_n) \in F$

$x_n \rightarrow x \Rightarrow x \in F$. "Contains its limit points"

Compact: since in \mathbb{R}^n , means closed and bounded

or, equiv., if $(x_n) \in K$, K compact, then \exists a
convergent subseq. of (x_n) w/ limit in K

More misc. notes:

A symmetric matrix X is positive semi-definite ($X \succeq 0$)
iff \exists matrix G s.t. $X = GG^T$ (or GG^* if \mathbb{C})

(if $X \succ 0$, then G is square & invertible and called the Cholesky factor of X)

" \succeq " induces a partial order on S^n , called the "Loewner order"
not a total order:
not all elements comparable

(but if they are, then acts nicely, e.g.
 $A \succeq B, B \succeq C \Rightarrow A \succeq C$ etc.)

to say $A \succeq B$ means $A - B \succeq 0$

⚠ If A has eigenvalues λ_i
and B has eigenvalues σ_i
then $A - B$ doesn't necessarily have eigenvalues $\lambda_i - \sigma_i$

trace

$$\text{tr}(A) = \sum_{i=1}^n A_{ii} \quad \text{is sum of diagonal entries} \\ = \sum \lambda_i \quad \text{sum of eigenvalues}$$

$$\text{tr}(ABC) = \text{tr}(BCA) = \text{tr}(CAB) \quad \text{"cyclic property"}$$

tricks:

- Compute $\text{tr}(A^T B)$, $A, B \in \mathbb{R}^{n \times n}$

$$\text{BAD: } \begin{cases} (1) \text{ compute } C = A^T B & O(n^3) \text{ flops} \\ (2) \text{ find } \text{tr}(C) & O(n) \text{ flops} \end{cases}$$

$$\text{BETTER: } \text{tr}(A^T B) = \langle A, B \rangle = \langle \text{vec}(A), \text{vec}(B) \rangle \\ = \text{vec}(A)^T \text{vec}(B) \\ O(n^2) \text{ flops}$$

- Compute $\|A\|_F$, A is $n \times n$

$$\text{BAD: } \|A\|_F = \sqrt{\text{tr}(A^T A)} \quad \text{so compute } A^T A \dots O(n^3)$$

$$\text{BAD: Find SVD of } A, \|A\|_F^2 = \sum \sigma_i^2 \quad O(n^3)$$

$$\text{GOOD: } = \sqrt{\sum_i \sum_j A_{ij}^2} \quad O(n^2)$$

• Fancier trick. Let A be low-rank, $A = UV^T$

$$m \times \begin{bmatrix} \hat{A} \\ A \end{bmatrix} = m \times \begin{bmatrix} U \\ U \end{bmatrix} \times \begin{bmatrix} V \\ V \end{bmatrix}$$

For simplicity, let it be rank 1 and square: $k=1, m=n$

Suppose B is sparse, $n \times n$

Compute $\|A - B\|_F$ efficiently.

Naïve: $\sqrt{\sum_{i,j} (A_{i,j} - B_{i,j})^2}$ $O(n^2)$, and also $O(n^2)$ memory

Better: $\|A - B\|_F^2 = \textcircled{1} \|A\|_F^2 + \textcircled{2} \|B\|_F^2 - 2 \textcircled{3} \text{tr}(A^T B)$

$\textcircled{1} \|A\|_F^2 = \|UV^T\|^2 = \text{tr}((UV^T)^T UV^T)$
 $= \text{tr}(VU^T UV^T)$
 $= \text{tr}(U^T U V^T V)$ } cycle
 $\in \mathbb{R}$, costs $O(n)$ to compute ✓

$\textcircled{2} \|B\|_F^2$ costs $\text{nnz}(B)$ to compute (hopefully $\ll O(n^2)$)

$\textcircled{3} \text{tr}(A^T B) = \text{tr}(UV^T)^T B = \text{tr}(VU^T B)$
 $= \text{tr}(U^T B V)$

compute $w = BV$ (sparse matrix-vector multiply)

in $\text{nnz}(B)$ time, $\text{tr}(U^T w) = U^T w$
 $O(n)$ time

so $O(n + \text{nnz}(B))$ time and memory.

• Matrix inversion lemma (ie. Sherman-Morrison-Woodbury)

Easy (=quick) to invert a low-rank perturbation of a matrix for which you already know the inverse

$$(A + UV^T)^{-1} = A^{-1} - \frac{A^{-1}UV^T A^{-1}}{1 + V^T A^{-1}U}$$

and see wikipedia for generalizations

Basically exploiting the Schur Complement

• Hölder's inequality

$$|\langle x, y \rangle| \leq \|x\|_p \cdot \|y\|_q \quad \text{if } \underbrace{\frac{1}{p} + \frac{1}{q} = 1}_{\text{Hölder conjugates}}$$

eg., $p=q=2$ "Cauchy-Schwarz"

eg., $p=1, q=\infty$ or vice-versa

• Norms, inner products are continuous

$$(1) \lim \|x_n\| = \|\lim x_n\| \quad \text{if } \lim x_n \text{ exists}$$

$$(2) \lim \langle a, x_n \rangle = \langle a, \lim x_n \rangle \quad \text{if } \lim x_n \text{ exists}$$

• If $x_n \rightarrow x$ and $y_n \rightarrow y$, then

$$\lim_n \langle x_n, y_n \rangle = \langle x, y \rangle$$

proof: $y_n \rightarrow y \Rightarrow (y_n)$ bounded, i.e. $\|y_n\| \leq M \forall n$

$$\lim_n \langle x_n, y_n \rangle = \underbrace{\lim_n \langle x_n - x, y_n \rangle}_{\text{via (1)}} + \underbrace{\lim_n \langle x, y_n \rangle}_{\text{via (2)}} = \langle x, y \rangle$$

$$|\lim_n \langle x_n - x, y_n \rangle| = \lim_n |\langle x_n - x, y_n \rangle| \quad \text{via (1)}$$

$$\leq \lim_n \|x_n - x\| \cdot \|y_n\|$$

$$\leq M \cdot \lim_n \|x_n - x\| = 0$$