

# HW help: proving polyhedrality

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How to prove a set  $S \subseteq \mathbb{R}^n$  is polyhedral?

One, often difficult way, is to write  $S = \{x: Ax \leq b, Fx = d\}$ . Usually we have tricks:

## Tricks

① Finite intersection of polyhedral sets is still polyhedral

② "lift-and-project"  $S \subseteq \mathbb{R}^n$ , then "lift" to a larger space  $\mathbb{R}^{n+k}$  by introducing new variables, and  $S = P(T)$ ,  $T$  a polyhedral set in  $\mathbb{R}^{n+k} \Rightarrow S$  polyhedral  
 $\uparrow$  project

(Why does this work?  $T$  polyhedral  $\Rightarrow \exists \{z_1, \dots, z_m\} \in \mathbb{R}^{n+k}$  s.t.  $T = \text{conv}(\{z_1, \dots, z_m\})$   
 $\Rightarrow S = \text{conv}(\{Pz_1, \dots, Pz_m\}) \Rightarrow S$  polyhedral)  
 $\downarrow$  via ③

③  $\text{conv}(V)$ ,  $|V| < \infty$ , is polyhedral. So find vertices  $V$   
 finite set.

Why?  $\text{conv}(\{v_1, \dots, v_m\}) = \left\{ x = \sum_{i=1}^m t_i \vec{v}_i : \sum t_i = 1, t_i \geq 0 \right\}$   
 encode as

$$\begin{bmatrix} I & -A \end{bmatrix} \begin{bmatrix} x \\ t \end{bmatrix} = \begin{bmatrix} 0 \end{bmatrix} \quad \text{where } t = \begin{bmatrix} t_1 \\ \vdots \\ t_m \end{bmatrix}$$

$A = [v_1, \dots, v_m]$   $n \times m$  matrix

encode as

$$\begin{bmatrix} 0 & \mathbf{1}^T \end{bmatrix} \begin{bmatrix} x \\ t \end{bmatrix} = 1$$

encode as

$$\begin{bmatrix} 0 & I \end{bmatrix} \begin{bmatrix} x \\ t \end{bmatrix} \geq \begin{bmatrix} 0 \end{bmatrix}$$

so  $\begin{bmatrix} x \\ t \end{bmatrix}$  is polyhedral, so projection to just  $x$  is polyhedral  
 (trick ②)

In my quick justifications, I used ③ to prove ② (in fact I used an "iff" version of ③), and used ② to prove ③ ... so not an actual proof)

From "Actually Doing it: Polyhedral Computations and its Applications"  
by Jesus A. De Loera, UC Davis, Math dept

**Definition 2.1.2** The set of solutions of a system of linear inequalities is called a polyhedron. In its general form a polyhedron is then a set of the type

$$P = \{x \in \mathbb{R}^d : \langle c_i, x \rangle \leq \beta_i\} \quad \{x : Ax \leq b, Fx = d\}$$

for some non-zero vectors  $c_i$  in  $\mathbb{R}^d$  and some real numbers  $\beta_i$ .

**Definition 2.1.3** A polytope is the convex hull of a finite set of points in  $\mathbb{R}^d$ .  
conv(v) vertices ✓

**Theorem 2.3.1 (Weyl-Minkowski theorem)** Every polytope is a polyhedron.  
Every bounded polyhedron is a polytope.

**Lemma 2.2.1** Given a bounded polyhedron  $P = \{x \in \mathbb{R}^n \mid Ax \leq b\}$  then

- There is a translation of  $P$  that can be represented with in the form  $\{x \in \mathbb{R}^n : A'x \leq b' \text{ and } x \geq 0\}$ .
- There is a polyhedron of the form  $Q = \{x \in \mathbb{R}^q : Bx = c, x \geq 0\}$  such that
  - the coordinate-erasing linear projection

"lift-and-project"

$$\pi : \mathbb{R}^q \longrightarrow \mathbb{R}^n : x = (x_1, \dots, x_n, \dots, x_q) \mapsto \pi(x) = (x_1, \dots, x_n)$$

provides a bijection between  $Q$  and  $P$ .

- The bijection implies that  $\pi(Q \cap \mathbb{Z}^q) = P \cap \mathbb{Z}^n$ .
- $y \in Q$  is an extreme point if and only if  $\pi(x)$  is an extreme point of  $P$
- $\dim(Q) = \dim(P)$ .