Conjugate Gradient method

Sunday, March 21, 2021 8:49 AM

Fast algo. for smooth unconstrained optim. grad descent, Nesterov

- Newton (See Bayd + Vondubegh, self-concordant analysis)

- Nonlinear conjugate gradient

· (linear) Conjugate Gradient (these notes follow Nocedal + Wright text and Anne Greenbaum's monograph) 143tong: Hestenes, Stiefel '50s (nonlinea: Fletcher + Reeves '60s)

(G solves Ax=6 if A>0, Ais nxn (naively: O(n3) flops) approximately

Motivation / link to optimization:

Min $\frac{1}{2} \parallel \tilde{A} \times -\tilde{b} \parallel^2$ $\varphi(x) = \frac{1}{2} \times \tilde{A} + \frac{1}{2} \tilde{b}$ $= \frac{1}{2} \times \tilde{A} \times -\tilde{b} \times \tilde{A} \times -\tilde{b} \times \tilde{A} \times + \frac{1}{2} \tilde{b} \times \tilde$

Py(x) = Ax-b, so find x st Ax=6

1 If you do want to solve least-squares, don't form A=ATA. Instead use LSQR

One idea to mon p(x) is coordinate descent / alternating minimization Converge in n steps

Conjugate Directions Ep. 3 are emjugate directions if they are A-orthogonal

Notation:
$$\langle P_i, A P_j \rangle = 0$$
 if $i \neq j$ (Fact: $\langle X_i y \rangle_A = x^T A y$ is an inner product)

Notation: $\langle x | A | y \rangle = \langle x_i, A y \rangle = \langle A x_i, y \rangle$
 $= x^T A y = \langle A x_i y \rangle$

So, $\langle P_i | A | P_j \rangle = 0$ if $i \neq j$

If we have $\{P_i, J_{i=0}^{n-1} A - e r t h u g}$, it's a basis.

Ex: P_i are exponentative: if A is symm, its exponentiative are arthing.

 $\langle P_i | A | P_j \rangle = \lambda_i \langle P_i | P_j \rangle = 0$ if $i \neq j$
 $A_j P_j$

Goal: $A_j P_j = \lambda_j \langle P_i | P_j \rangle = 0$ if $i \neq j$
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Ax* = 6

20 < pr | 4 | x - x > = 0

claim
$$\Gamma_{k} = \langle \underbrace{P_{k}|A|x^{*}-X_{k}}\rangle = \langle \underbrace{P_{k}|b-AX_{k}}\rangle = \alpha_{k}$$

$$\Rightarrow \chi_{n} = x^{*}$$

Facts via any CD method,

Conjugate Grandient = C.D. method that cheaply compute PK

.
$$X_0$$
 anything, $r_0 = Ax_0 - b$, $p_0 = -r_0$

· Iterate

$$P_{k} = -r_{k} + \beta_{k} \cdot \beta_{k-1}, \quad \beta_{k} = \frac{\langle r_{k} | A | p_{k-1} \rangle}{\langle p_{k-1} | A | p_{k-1} \rangle}$$

$$X_{k+1} = X_{k} + A_{k} P_{k} \quad \text{chosen so } \langle p_{k} | A | p_{k-1} \rangle = 0$$

$$A_{k} = -\frac{\langle r_{k} | p_{k} \rangle}{\langle p_{k} | A | p_{k} \rangle}$$

Cost: one matrix-vector multiply per step

Convergence Regult (NW book)

$$\|x_{k} - x_{k}\|^{2} \leq 2 \cdot \left(\frac{\sqrt{k} - 1}{\sqrt{k} + 1}\right)^{k} \|x_{0} - x_{k}\|^{2}$$

K=K(A) is the condition # = max sing. value min. sity. value

Proeped issues exist (eg., stability (roundoff), not discussed here.

· Nonlinear Conjugate Gradient

Linear (G: min $\phi(x) := \frac{1}{2} < x |A| \times 7 - < b| \times 7$, $\nabla \phi(x) = Ax - b = r$ I'e., solve $Ax = b \Leftrightarrow hence name "linear"$ Nonlinear (G: min $\phi(x)$, $\phi(x)$ isn't quadrative, $\nabla \phi$ isn't linear $X_{K+1} = x_K + d_K p_K$, $d_K = argmin \phi(x_K + ap_K)$ Known in closed form

Number: XK+1 = XK + dK PK, dK ~ argmin 4(XK+ xPK) Not in closed form.

 $P_{K+1} = -\nabla \varphi(x_{K+1}) + \beta_{K+1} P_{K}$ $\Rightarrow \sum_{k=1}^{K} P_{K+1} P_{K}$ $\Rightarrow \sum_{k=1}^{K} P_{K}$ $\Rightarrow \sum_$

BEH has many choices, all of which reduce to < TK+1 | A | PK > IF & quadration Ex: BER = 11 79 KH 112 others too (Polyak-Ribiere ...)

Comments on non-lihear CG

- can be fast, but limited to unconstrained optimization Doesn't play well we constraints, even simple x > 0 on = s Die, if q is graduate, CG is essentially optimal among all 1st order methods. So y "almost quadratic", non-linear (G might do very well
- Finicky w, d, B terms. Hager + Thang have many advanced versions, some require many (eg. 20) parameters.
- Global convergence theory isn't great
- Nemirousky + Yudin proved there exist reasonable cux finitions for which non-linear CG is slower than grandient descent

- IMO, quasi-Newton methods about as fast, and more stable (simpler)

(ie., if you tuned all non-linear CG param. j'ust right, it might be faster, but parameters are problem dependent!)