### Data Structure and Algorithms

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### List

- A collection of items of the same type
- A flexible structure that can grow and shrink on demand

- Elements can be:
  - Inserted
  - Accessed
  - Deleted

at any position!

#### List

- A list is a sequence of zero or more elements of a given type
- Represented by a comma-separated sequence of elements

- Where n ≥ 0 and each ai is of type element\_type
- Number of elements n determines the length of the list
  - If  $n \ge 1$ 
    - > a1 is the first element
    - > an is the last element
  - If n = 0
    - List has no elements (empty list)

### List

- Elements of a list can be linearly ordered according to their position
  - $a_i$  precedes  $a_{i+1}$  for i = 1,2,3...n-1
  - $a_i$  follows  $a_{i-1}$  for i = 2,3,4...n
- Element a<sub>i</sub> is at position i

### Properties of Lists

- Can have a single element
- Can have no elements

- Can be list of lists
- Can be concatenated together
- Can be split into sub-lists

### Basic Operations

- Create the list
  - The list is initialized to an empty state
- Determine whether the list is empty
  - Determine whether the list is full
- Find the size of the list.
- Destroy, or clear, the list
- Insert an item in the list at the specified location
- Delete an item from the list at the specified location
- Replace an item at the specified location with another item
- Retrieve an item from the list at the specified location
- Search the list for a given item
- Traverse (iterate through) the elements of the list

### Basic Operations

- INSERT(x, p, L)
  - Insert x at position p in list L
  - If list L has no position p, the result is undefined
- RETRIEVE(p, L)
  - Return the element at position p on list L
- LOCATE(x, L)
  - Return the position of x on list L
- DELETE(p, L)
  - Delete the element at position p on list L

### **Basic Operations**

- MAKENULL(L)
  - Causes L to become an empty list and returns position END(L)
- NEXT(p, L)
  - Return the position following p on list L
- PREVIOUS(p,L)
  - Return the position preceding position p on list L
- FIRST(L)
  - Returns the first position on the list L
- PRINTLIST(L)
  - Print the elements of L in order of occurrence
- And more ...

#### List as a Data Structure

- We know the ADT of the list, how to implement it?
- Create a List class, containing at least the following function members
  - Constructor
  - isEmpty()
  - insert()
  - delete()
  - print()
- What are other function members?
  - isFull(), listSize(), retrieve(), replace(), search(), clearList(),

#### List as a Data Structure

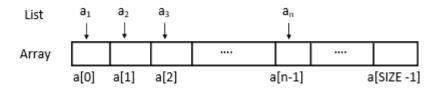
- Implementation involves
  - Defining data members
  - Defining function members from design phase

- In terms of implementation, there are two possible approaches
  - Array-based Implementation of lists
  - Linked list using pointers-based implementation of lists

# Array-Based Implementation of Lists

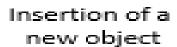
- An array is a viable choice for storing list elements
  - Elements are sequential
  - Array is a commonly available data type
  - Algorithm development is easy

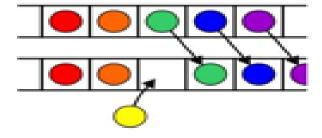
Normally sequential orderings of list elements match with array indices



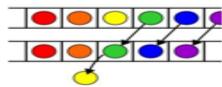
## Implementing Operations

- Constructor
  - Static array allocated at compile time
- isEmpty
  - Check if size == 0
- traverse/ print
  - Use a loop from 0<sup>th</sup> element to size − 1
- insert
  - Shift elements to right of insertion point
- delete
  - Shift elements back









### Operations Code in C++

```
#include <iostream>
using namespace std;
class List {
private:
   int items[10];  // Static array to store list elements
int currentSize;  // Current number of items in the list
const int maxSize;  // Maximum capacity of the list (fixed size)
public:
List() : maxSize(10), currentSize(0) {}
void clearList() {
       currentSize = 0;
```

## Empty and Full Operation

```
bool isEmpty() const {
    return currentSize == 0;
}

// Checks if the list is full
bool isFull() const {
    return currentSize == maxSize;
}
```

# Inserts an item at the end of the list

```
void insert(int item) {
    if (isFull()) {
    cout << "List is full. Cannot insert item." << endl;
        return;
    }
    items[currentSize++] = item;
}</pre>
```

## Deletes the first occurrence of an item from the list

```
void deleteItem(int item) {
     int index = -1;
     for (int i = 0; i < currentSize; ++i) {
       if (items[i] == item) {
          index = i;
          break;
     if (index == -1) {
       cout << "Item "" << item << "' not found in list." << endl;
       return:
// Shift elements to the left to fill the gap
     for (int i = index; i < currentSize - 1; ++i) {
       items[i] = items[i + 1];
     --currentSize; } }
```

### Prints all items in the list

```
void print() const {
    for (int i = 0; i < currentSize; ++i) {
        cout << items[i] << " ";
    }
    cout << endl;
}

// Returns the current size of the list
int listSize() const {
    return currentSize;
}</pre>
```

# Retrieves the item at a specific index

```
int retrieve(int index) const {
    if (index >= 0 && index < currentSize) {
        return items[index];
    } else {
        cerr << "Index out of range." << endl;
        return -1; // Could also throw an exception
    }
}</pre>
```

// cerr is an output stream in C++ that's used to output error messages.

# Searches for an item and returns its index

```
int search(int item) const {
    for (int i = 0; i < currentSize; ++i) {
       if (items[i] == item) {
          return i;
     return -1;
```

# Replaces the item at a specific index with a new item

```
void replace(int index, int newItem) {
    if (index >= 0 && index < currentSize) {
      items[index] = newItem;
    } else {
      cerr << "Index out of range." << endl;
```

# List Class with Static Arrays - Problems

- Stuck with "one size fits all"
  - Could be wasting space
  - Could run out of space

Better to have instantiation of a list by specifying the capacity (i.e., size)

Consider creating a List class with dynamically-allocated array

# Dynamic Allocation of List Class

- Changes required in data members
  - Eliminate constant declaration for CAPACITY/SIZE
  - Data member to store capacity specified by client program

- Little or no changes required for many function members
  - isEmpty()
  - display()
  - delete()
  - insert()

# Dynamic Allocation of List Class

```
Now possible to specify different sized lists
        cin >> maxListSize;
        List aList1 (maxListSize);
        List alist2 (500):
aList1
            mySize
        myCapacity 1024
                                                                     1023
        myArrayPtr
aList2
            mySize
        myCapacity 500
        myArrayPtr
```

# Dynamic Allocation of List Class

```
class List {
private:
  int* items; // Pointer to dynamically allocated array
  int currentSize; // Current number of items in the list
  int maxSize; // Maximum capacity of the list
public:
  // Constructor with a default size
  List(int size = 10) : maxSize(size), currentSize(0) {
    items = new int[maxSize]; // Dynamically allocate array of given size
  // Destructor to free dynamically allocated memory
  ~List() {
    delete[] items;
```

### Any Question So Far?

