



Algorithms: Design
and Analysis, Part II

Approximation Algorithms for NP-Complete Problems

A Greedy Knapsack Heuristic

Strategies for NP-Complete Problems

① identify computationally tractable special cases

example: knapsack instances with small capacity

[i.e., knapsack capacity $W = \text{polynomial in number of items } n$]

② heuristics — today

for knapsack

— pretty good greedy heuristic

— excellent dynamic programming heuristic

③ exponential time but better than brute-force search

example: $O(nW)$ -time dynamic programming vs. $O(2^n)$ brute-force search

Ideally: should provide a performance guarantee (i.e., "almost correct") for all (or at least many) instances.

Knapsack Revisited

Input: n items. Each has a positive value v_i and a size w_i . Also, knapsack capacity is W

Output: a subset $S \subseteq \{1, 2, 3, \dots, n\}$ that

maximizes $\sum_{i \in S} v_i$

Subject to $\sum_{i \in S} w_i \leq W$.

A Greedy Heuristic

Motivation: ideal items have big value, small size.

Step 1: Sort and reindex item so that

$$\frac{v_1}{w_1} \geq \frac{v_2}{w_2} \geq \frac{v_3}{w_3} \geq \dots \geq \frac{v_n}{w_n}$$

[i.e., nondecreasing
"bang-per-buck"]

Step 2: pack items in this order until one doesn't fit, then halt.

Example: $v_1=2$ $w_1=1$ \Rightarrow greedy gives $\{1, 2\}$
 $w=5$ $v_2=4$ $w_2=3$ [also optimal]
 $v_3=3$ $w_3=3$

Quiz

Consider a knapsack instance with $V_1 = 2$ $w_1 = 1$
 $V_2 = 1000$ $w_2 = 1000$
 $W = 1000$.

Question: What is the value of the greedy solution and the optimal solution, respectively?

(A) 2 and 1000

(C) 1000 and 1002

(B) 2 and 1002

(D) 1002 and 1002

A Refined Greedy Heuristic

Upshot: greedy solution can be arbitrarily bad relative to an optimal solution.

Fix: add

Step 3: return either the Step-2 solution, or the maximum value item, whichever is better.

Theorem: value of the 3-step greedy solution is always $\geq 50\%$ of value of an optimal solution.

[i.e., a " $\frac{1}{2}$ -approximation algorithm"]

[also runs in $O(n \log n)$ time]