

DATASTRUCTURES

BIT

```
const int N=1e5+10;
int n, BIT[N];

void UpdateIndex(int idx,int value) // O(logn) - Updates value of a certain index
{
    for(;idx<=n;idx+=idx&-idx)
        BIT[idx]+=value;
}

int GetAns(int idx) // O(logn) - Gets answer in range [1:idx]
{
    int sum = 0;
    for(;idx;idx--=(idx&-idx))
        sum+=BIT[idx];
    return sum;
}
```

SEGMENT TREE

```
// Size = 4*N as there are children to the leaves
// The array a[N] is 1-based so as to correspond to the Build function

const int N=1e5+10;
int n,a[N]; // Don't redefine n
int seg[4*N],lazy[4*N];

void propagate(int ind, int l, int r) // Propagate contains l,r for summation values
{
    seg[ind*2+1]|=lazy[ind];
    seg[ind*2+2]|=lazy[ind];
    lazy[ind*2+1]|=lazy[ind];
    lazy[ind*2+2]|=lazy[ind];
    lazy[ind]=0;
}

void Build(int ind=0, int l=1, int r=n) // O(nlogn)
{
    if(l==r)
    {
        seg[ind]=a[l];
        return;
    }
    int mid=l+(r-l)/2;
    Build(ind*2+1,l,mid);
    Build(ind*2+2,mid+1,r);
    seg[ind]=seg[ind*2+1]|seg[ind*2+2];
}
```

```

void UpdateIndex(int qind,int qv, int ind=0, int l=1, int r=n) // O(logn)
{
    if(qind<l || qind>r)    return;
    if(l==r)
    {
        seg[ind]|=qv;
        return;
    }
    int mid=l+(r-l)/2;
    UpdateIndex(qind,qv,ind*2+1,l,mid);
    UpdateIndex(qind,qv,ind*2+2,mid+1,r);
    seg[ind]=seg[ind*2+1]|seg[ind*2+2];
}

void UpdateRange(int qx, int qy, int qv, int ind=0, int l=1, int r=n) // O(logn) with
lazy propagation
{
    if(r<qx || l>qy)    return;
    if(l>=qx && r<=qy)
    {
        seg[ind]|=qv;
        lazy[ind]|=qv;
        return;
    }
    int mid=l+(r-l)/2;
    propagate(ind, l, r);
    UpdateRange(qx,qy,qv,ind*2+1,l,mid);
    UpdateRange(qx,qy,qv,ind*2+2,mid+1,r);

    seg[ind]=seg[ind*2+1]|seg[ind*2+2];
}

int GetAns(int qx, int qy, int ind=0, int l=1, int r=n, bool isRangeProblem=0) // O(logn)
{
    if(r<qx || l>qy)    return 0;
    if(l>=qx && r<=qy)    return seg[ind];
    int mid=l+(r-l)/2;
    if(isRangeProblem)    propagate(ind, l, r);
    int r1=GetAns(qx,qy,ind*2+1,l,mid,isRangeProblem);
    int r2=GetAns(qx,qy,ind*2+2,mid+1,r,isRangeProblem);
    return r1|r2;
}

```

SPARSE TABLE

```

const int N = 1e5 + 5;
const int logN = 20;
int n, a[N], ST[logN][N], LOG[N];

void Build() // O(nlogn) - Builds the sparse table (min sparse table)
{
    LOG[0] = -1;
    for(int i=0;i<n;i++) // Computes floor(log2(i)) for all values
        ST[0][i] = i, LOG[i+1] = LOG[i] + !(i & (i+1));
}

```

```

    for(int j=1; (1<<j) <= n; j++)
        for(int i=0; (i +(1<<j)) <= n; i++)
        {
            int x = ST[j-1][i];
            int y = ST[j-1][i + (1<<(j-1))];

            ST[j][i] = (a[x] <= a[y] ? x : y); // Changes according to the operation
        }
}

int GetAns(int l, int r) // O(1) - Queries for the values from l to r
{
    int g = LOG[r-l+1];
    int x = ST[g][l];
    int y = ST[g][r - (1<<g) + 1];
    return (a[x] <= a[y] ? x : y); // Changes according to the operation
}

```

TRIE

```

const int N = 1e5 + 5;
const int logN = 20;
int n, a[N], ST[logN][N], LOG[N];

void Build() // O(nlogn) - Builds the sparse table (min sparse table)
{
    LOG[0] = -1;
    for(int i=0;i<n;i++) // Computes floor(log2(i)) for all values
        ST[0][i] = i, LOG[i+1] = LOG[i] + !(i & (i+1));

    for(int j=1; (1<<j) <= n; j++)
        for(int i=0; (i +(1<<j)) <= n; i++)
        {
            int x = ST[j-1][i];
            int y = ST[j-1][i + (1<<(j-1))];

            ST[j][i] = (a[x] <= a[y] ? x : y); // Changes according to the operation
        }
}

int GetAns(int l, int r) // O(1) - Queries for the values from l to r
{
    int g = LOG[r-l+1];
    int x = ST[g][l];
    int y = ST[g][r - (1<<g) + 1];
    return (a[x] <= a[y] ? x : y); // Changes according to the operation
}

```

DYNAMIC PROGRAMMING

LCS

```
vector<int> vec1,vec2;
int dp[5005][5005];

int LCS(int i, int j)
{
    if(i == vec1.size() || j == vec2.size()) return 0;
    if(dp[i][j] != -1) return dp[i][j];
    if(vec1[i] == vec2[j])
        return dp[i][j] = LCS(i+1,j+1) + 1;
    return dp[i][j] = max(LCS(i,j+1) , LCS(i+1,j));
}
```

LIS $O(N^2)$

```
const int N = 1e5+10;
int dp[N];
vector<int> a,path;

int LIS()
{
    for(int i=0;i<N;i++) dp[i] = 1;
    for(int i=1;i<a.size();i++)
        for(int j=0;j<i;j++)
            if(a[i] > a[j] && dp[i] < dp[j]+1)
                dp[i] = dp[j] + 1;

    int ans = 0, mx = 0, idx = 0;

    for(int i=0;i<N;i++)
        if(dp[i] > mx)
            mx = dp[i],idx=i;

    ans = mx;

    for(int i=idx;i>=0;i--)
        if(dp[i] == mx)
            path.pb(a[i]),mx--;

    reverse(path.begin(),path.end());

    return ans;
}
```

LIS $O(N\log N)$

```
const int N = 1e5+5;

int n,a[N];
```

```

int getLIS()
{
    if(n<1) return 0;

    int len = 0;
    vector<int> LIS(n,0);

    for(int i=0;i<n;i++)
    {
        int idx = lower_bound(LIS.begin(),LIS.end(),a[i]) - LIS.begin();
        LIS[idx] = a[i];
        len = max(len,idx);
    }

    return len+1;
}

```

GRAPHS

DSU

```

const int N = 1e5+5;

int par[N];
int num[N];

for(int i=0;i<N;i++) // Initialization in main
{
    par[i] = i;
    num[i] = 1;
}

int FindSet(int n) // Use FindSet instead of par
{
    if(par[n] == n)
        return n;
    return par[n] = FindSet(par[n]);
}

void Union(int x,int y)
{
    x = FindSet(x);
    y = FindSet(y);
    if(x==y)
        return;
    par[x]=y;
    num[y] += num[x]; // num holds number of elements in a group
}

```

FLOYD

```

// Shortest path between all pairs of nodes, negative edges allowed, no negative cycles

int n; // Number of nodes
int adj[n][n]; // adj[u][v] = cost of edge from u to v

```



```

bool checkNegativeCycle()
{
    bool ret = false;
    for(int i=0;i<n;i++)
        ret = ret || (adj[i][i] < 0);
    return ret;
}

```

KRUSKALL MST

```

const int N = 1e5+5;

vector< pair< int, pair<int,int> > edges; // {cost, {u,v}}

int par[N];
int num[N];

for(int i=0;i<N;i++) // Initialization in main
{
    par[i] = i;
    num[i] = 1;
}

int FindSet(int n) // Use FindSet instead of par
{
    if(par[n] == n)
        return n;
    return par[n] = find_set(par[n]);
}

void Union(int x,int y)
{
    x = find_set(x);
    y = find_set(y);
    if(x==y)
        return;
    par[x]=y;
    num[y] += num[x]; // num holds number of elements in a group
}

int getMST() // Returns summation of edges of Minimum Spanning Tree and builds the tree
in DSU
{
    int sum = 0;
    sort(edges.begin(),edges.end());
    //reverse(edges.begin(), edges().end()); // Uncomment for Max. spanning tree
    for(auto e : edges)
    {
        int u = e.s.f; int v = e.s.s;
        if(FindSet(u) == FindSet(v))
            continue;
        Union(u,v);
        sum+=e.f;
    }
    return sum;
}

```

LCA O(1)

```
int const N = 1e5+10;
int const nN = 2*N;
int id, ID[N], rID[N], First[N], ST[nN][logN], lg[nN];
vector<int>vec;

void dfs(int cur, int par)
{
    ID[cur] = id++;          // Creating ur own ID by euler tour.
    rID[ID[cur]] = cur;     // reverseID to get back to the original node name.
    vec.push_back(ID[cur]); // Push ur tour.
    for(auto x:adj[cur])
        if(x!=par)
            dfs(x, cur), vec.push_back(ID[cur]);
}

inline void Build()
{
    int cnt = 0, two = 1;
    while(two<nN) lg[two] = cnt++, two*=2;
    for(int i=3; i<nN; ++i) if(!lg[i]) lg[i] = lg[i-1]; //Create an array to get the log of
any number [nearest log].
    //the previous part should be done only one time! Put it before the test cases.

    int nN = sz(vec);
    memset(First, -1, sizeof First);

    for(int i=0; i<nN; ++i) // Fill the "First" array. remeber this array
contains the idx of the first appearence of an id.
        ST[i][0] = vec[i], (First[vec[i]]==-1? First[vec[i]]=i:0);

    // Bulding the sparse table.
    for(int i=1; (1<<i) <= nN; ++i)
        for(int j=0; j+(1<<i)-1<nN; ++j)
            ST[j][i] = min(ST[j][i-1], ST[j+(1<<(i-1))][i-1]);
}

inline int LCA(int a, int b)
{
    int L, R, idx, d, lca;
    a = ID[a]; b = ID[b];          // Go to ur domain.
    L = First[a], R = First[b];
    if(L>R) swap(L, R);
    d = R-L+1;
    d = lg[d];
    idx = R - (1<<d) + 1;
    lca = min(ST[L][d], ST[idx][d]);
    lca = rID[lca];
    a = rID[a]; b = rID[b];        //Back to main domain.
    return lca;
}
```

LCA O(logN)

```
const int N = 1e5+10;
const int logN = 20;
int ST[logN][N], depth[N];
```



```

void dfs(int cur, int par)
{
    ST[0][cur] = par;

    for(int i=1; i<logN; ++i)
        ST[i][cur] = ST[i-1][ST[i-1][cur]];

    for (auto x:adj[cur])
        if(x != par)
            depth[x]=depth[cur]+1, dfs(x, cur);
}

int KthAncestor(int u, int k)
{
    if(!k) return u;
    int d = depth[u] - k; // Get the depth of the wanted node.
    for (int j = logN-1; j >= 0; --j)
    {
        int nu = ST[j][u];
        if (!nu) continue;
        if (depth[nu] == d) return nu;
        else if (depth[nu] > d && u) u = nu;
    }

    return 0;
}

int LCA(int a, int b)
{
    if (depth[a] > depth[b]) a = KthAncestor(a, depth[a] - depth[b]); // make them
at the same depth.
    else if (depth[b] > depth[a]) b = KthAncestor(b, depth[b] - depth[a]);

    if (a == b) return a;

    for (int j = logN-1; j >= 0; --j)
        if (ST[j][b] != ST[j][a])
            a = ST[j][a], b = ST[j][b];

    return ST[0][a];
}

```

LONGEST WEIGHTED PATH

```

const int N = 1e5+5;
vector<pair<int,int> > adj[N];

ll dfs(int u=0, int par=-1)
{
    ll ret = 0;
    for(int i=0; i<adj[u].size(); i++)
        if(adj[u][i].f != par)
            ret = max(ret, dfs(adj[u][i].F, u) + adj[u][i].s);
    return ret;
}

```

TARJAN SCC

```
vector<vector<int> > SCCs /* The components itself*/;
#define comps SCCs
vector<int> compIndex /* for each node, what is the index of the
component this node inside*/
,ind, lowLink;
stack<int> st;
vector<bool> inst;
vector<vector<int> > adj /*The intial graph*/;
int idx = 0; //must be intialized by zero;

void tarjanSCC(int i) {
    lowLink[i] = ind[i] = idx++;
    st.push(i);
    inst[i] = true;
    for (int j = 0; j < adj[i].size(); j++) {
        int k = adj[i][j];
        if (ind[k] == -1) {
            tarjanSCC(k);
            lowLink[i] = min(lowLink[i], lowLink[k]);
        } else if (inst[k]) {
            lowLink[i] = min(lowLink[i], lowLink[k]);
        }
    }
    if (lowLink[i] == ind[i]) {
        vector<int> comp;
        int n = -1;
        while (n != i) {
            n = st.top();
            st.pop();
            comp.push_back(n);
            inst[n] = 0;
            compIndex[n] = comps.size();
        }
        comps.push_back(comp);
    }
}

void SCC() {
    comps.clear();
    compIndex.resize(adj.size());
    ind.clear();
    ind.resize(adj.size(), -1);
    lowLink.resize(adj.size());
    inst.resize(adj.size());
    idx = 0; //must be intialized by zero;
    for (int i = 0; i < adj.size(); i++)
        if (ind[i] == -1)
            tarjanSCC(i);
}

int cntSrc /*the number of source components*/,
cntSnk /*the number of sink components*/;
vector<vector<int> > cmpAdj /*The new graph between components*/;
vector<int> inDeg, outDeg /*the in degree and out degree for each
component*/;
```

```

void computeNewGraph() {
    outDeg.clear();
    outDeg.resize(comps.size());
    inDeg.clear();
    inDeg.resize(comps.size());
    cntSrc = cntSnk = comps.size();
    cmpAdj.clear();
    cmpAdj.resize(comps.size());
    for (int i = 0; i < adj.size(); i++) {
        for (int j = 0; j < adj[i].size(); j++) {
            int k = adj[i][j];
            if (compIndex[k] != compIndex[i]) {
                cmpAdj[compIndex[i]].push_back(compIndex[k]);
                if (!(inDeg[compIndex[k]]++))
                    cntSrc--;
                if (!(outDeg[compIndex[i]]++))
                    cntSnk--;
            }
        }
    }
}

```

TREE DIAMETER

```

const int N = 1e5+5;
int sp[N];

int bfs(int u)
{
    queue<int> q;
    q.push(u);
    memset(sp, -1, sizeof(sp));
    sp[u] = 0;

    while(!q.empty())
    {
        u = q.front();
        q.pop();

        for(auto v : adj[u])
            if(sp[v] == -1)
                sp[v] = sp[u] + 1, q.push(v);
    }

    return u;
}

int calcTreeDiameter(int root)
{
    int u = bfs(root);
    int v = bfs(u);
    return sp[v];
}

```

BELLMAN FORD

```

// Shortest path from source to all nodes, negative edges allowed, no negative cycle
const int N = 1e5+5;

vector<pair<int,int> > adj[N];

```

```

int n,par[N]; // n is number of nodes

bool bellmanFord(int source)
{
    memset(par,-1,sizeof par);
    memset(dis, 0x3F, sizeof(dis));
    dis[source] = 0;

    bool updated = 1;

    for(int k=0;k<n && updated;k++)
    {
        updated = 0;
        for(int u = 1; u<=n;u++)
        {
            for(auto& e : edges[u])
            {
                int v = e.first;
                int c = e.second;

                if(dis[v] > dis[u] + c)
                {
                    dis[v] = dis[u] + c;
                    par[v] = u;
                    updated = 1;
                }
            }
        }
    }

    return updated; // Whether a negative cycle exist or not
}

```

BFS

```

const int N = 1e5+5;
vector<int> adj[N];
int sp[N];
void bfs()
{
    memset(sp,-1,sizeof sp);
    queue<int> q;

    q.push(s); //pushes start
    sp[s] = 0;

    while(!q.empty()){
        int u = q.front();
        q.pop();

        for(int i=0;i<adj[u].size();i++){
            int v = adj[u][i];
            if(sp[v] == -1){
                q.push(v);
                sp[v] = sp[u] + 1;
            }
        }
    }
}

```

DFS

```
const int N = 1e5+5;
vector<int> adj[N];

int visited[N];

void dfs(int u){
    visited[u] = 1;
    for(int i=0;i<adj[u].size();i++){
        int v = adj[u][i];
        if(!visited[v])
            dfs(v);
    }
}
```

DIJKSTRA

```
// Dijkstra SP for weighted graph + path build, No negative edges
const int N = 1e5+5;
vector< pair<int,int> > adj[N]; // adj[u] = {v,cost}
vector<int> path;
int sp[N]; int par[N];

void build(int u) {
    while (u != -1) {
        path.pb(u);
        u = par[u];
    }

    reverse(path.begin(),path.end()); //The path is reversed, fix it
}

int Dijkstra(int source, int destination)
{
    int u,v,c;
    for (int i = 0; i < n; ++i) // initialization
        sp[i] = oo, par[i] = -1;

    priority_queue< pair<int,int> , vector<pair<int,int> >, greater<pair<int,int> > > pq;

    pq.push( { 0, source } );
    sp[source] = 0;

    while (!pq.empty()) {
        u = pq.top().second;
        c = pq.top().first;
        pq.pop();

        if (sp[u] < c)
            continue;

        for (auto edge : adj[u]) {
            int to = edge.first;
            int cost = edge.second;

            if (c + cost < sp[to]) {
```

```

        sp[to] = c + cost;
        pq.push( { sp[to], to } );

        //assign (u) as parent of (to) to rebuild the path
        par[to] = u;
    }
}

build(destination); // Path is stored in path
return sp[destination];
}

```

MATH

PRIME FACTORIZATION

```

const int MAX = 1e6+10;
int spf[MAX], power[MAX]
void Sieve() // O(nlogn)
{
    spf[1] = 1;
    for (int i = 2; i < MAX; i++) spf[i] = i;
    for (int i = 4; i < MAX; i += 2) spf[i] = 2;
    for (int i = 3; i < MAX; i++)
        if (spf[i] == i)
            for (int j = i; j < MAX; j += i)
                if (spf[j] == j)
                    spf[j] = i;
}

void GetFactors(int num)
{
    int cnt = 1, last = spf[num], ans = 1;
    while (num != 1) {
        num = num / spf[num];
        if (last == spf[num]) {
            cnt++;
            continue;
        }
        power[last] = cnt;
        cnt = 1;
        last = spf[num];
    }
}

```

nCr WITH MOD

```

ll mod = 1e9+7;
const int N = 1e5+10;
ll fact[N];
void ComputeFactorial() // Don't forget to compute the factorial
{
    fact[0] = 1;
    for (ll i=1; i<N; i++)
        fact[i] = (fact[i-1]%mod * i%mod)%mod;
}

```

```

11 power(11 x, 11 y)
{
    11 res = 1;
    x = x % mod;
    while (y > 0)
    {
        if (y & 1)
            res = (res*x) % mod;
        y = y>>1;
        x = (x*x) % mod;
    }
    return res;
}

11 modInverse(11 n)
{
    return power(n, mod-2);
}

11 nCr(11 n, 11 r)
{
    if (r==0)
        return 1;

    return (fact[n]* modInverse(fact[r]) % mod *
            modInverse(fact[n-r]) % mod) % mod;
}

```

SIEVE

```

const int MAX = 1e6+10;
vector<bool> prime(MAX, true);

void Sieve() // O(nlog(logn))
{
    prime[0] = prime[1] = 0;
    for (int i = 2; i*i <= MAX; i++)
        if (prime[i])
            for (int j = i * 2; j <= MAX; j += i)
                prime[j] = 0;
}

```

SMALLEST PRIME FACTOR

```

const int MAX = 1e6+10;
int spf[MAX];

void Sieve() // O(nlogn)
{
    spf[1] = 1;
    for (int i = 2; i < MAX; i++) spf[i] = i;
    for (int i = 4; i < MAX; i += 2) spf[i] = 2;
    for (int i = 3; i < MAX; i++)
        if (spf[i] == i)
            for (int j = i; j < MAX; j += i)
                if (spf[j] == j)
                    spf[j] = i;
}

```

EULER PHI

```
// Theorem:  $a^{\phi(b)} \% b = 1$ , if a,b are coprimes ( $\gcd(a,b) == 1$ )
ll Phi(ll n) //  $O(\sqrt{n})$ , returns Euler Phi of n
{
    ll result = n;
    for(int i=2; i*i<=n; i++)
    {
        if(n%i == 0)
            result -= result/i;
        while(n%i==0)
            n/=i;
    }
    if(n>1)
        result -= result/n;
    return result;
}

ll phi[N];
void EulerTotient() //  $O(n \log(\log n))$ , generates all Euler Phi's to N
{
    for(int i=0; i<N; i++) phi[i] = i;
    for(int i=2; i<N; i++)
        if(phi[i] == i)
            for(int j=i; j<N; j+=i) phi[j] -= phi[j]/i;
}
```

EXTENDED EUCLID

```
//  $O(\log(\max(a,b)))$ 
// Returns Bezout's coeffs of the smallest
// linear combination of s.a + t.b = gcd(a,b)
pair<int,int> ExtendedEuclid(int a,int b)
{
    if(b==0)
        return {1,0};
    pair<int, int> p = ExtendedEuclid(b,a%b);
    int s = p.first, t = p.second;
    return {t, s-t * (a/b)};
}
```

FASTPOWER WITH MOD

```
ll power(ll base, ll exp)
{
    ll ans = 1;
    base%=mod;
    while(exp>0)
    {
        if(exp&1) ans = (ans*base)%mod;
        exp>>=1;
        base = (base*base)%mod;
    }
    return ans;
}
```


IS PRIME

```
bool IsPrime(int n)
{
    if(n<=1) return false;
    if(n==2) return true;
    if(n%2==0) return false;

    for(int i=3 ; i*i<=n ; i +=2){
        if(n%i==0)
            return false;
    }
    return true;
}
```

MATH MISC

```
ll gcd(ll a, ll b) { return (b == 0 ? a : gcd(b, a % b)); } // O(log(max(a,b)))
ll lcm(ll a, ll b){ return ((a*b) / gcd(a, b)); } // O(log(max(a,b)))
```

```
ll modInverse(ll a, ll m) // O(logm)
{
    return power(a,m-2,m);
}
```

```
ll nCr(ll n, ll r) // O(r) - Call the function as nCr(n,min(r,n-r)) for better
performance
{
    if(n<r) return 0;
    if(r==0) return 1;
    return n*nCr(n-1,r-1)/r;
}
```

```
// Factorization - O(sqrt(n))
vector<ll>factors;
void GetFactors(ll n){
    for (int i = 1; i*i <= n; i++) {
        if (n%i == 0) {
            if (i != n / i)
                factors.push_back(i), factors.push_back(n / i);
            else
                factors.push_back(i);
        }
    }
}
```

```
//Mod of very large numbers - O(n)
//rule:  $xy \pmod a \equiv ((x \pmod a) * y) \pmod a$ 
ll mod(string num)
{
    int res = 0;
    for (int i = 0; i < num.size(); i++)
        res = (res*10 + num[i] - '0') %mod;
    return res;
}
```

```

// Derangements - O(n)
ll fact(int n)
{
    ll ret = 1;
    for(int i=n;i>1;i--) ret = ret* 1LL * i;
    return ret;
}

ll pww(ll p){return p&1 ? -1 : 1;}

ll Derangement(ll n)
{
    ll tmp = fact(n), sum=0;
    for(int i=0;i<=n;i++) sum+= (pww(i) * tmp) / fact(i);
    return sum;
}

// isPowerOfTwo - O(1)

bool isPowerOfTwo(ll x) {return x && (!(x&(x-1)))};

```

MATRIX POWER

```

const int mod = 1e9+7;

const int MatSize = 2;
struct Matrix
{
    long long mat[MatSize][MatSize];
};

Matrix MatMul(Matrix &a , Matrix &b)
{
    Matrix res;
    for(int i = 0 ; i<MatSize ; i++)
        for(int j = 0 ; j<MatSize ; j++)
        {
            res.arr[i][j] = 0;
            for(int k = 0 ; k<MatSize ; k++)
            {
                long long sum = a.arr[i][k] * b.arr[k][j];
                sum%=mod;
                res.arr[i][j] +=sum;
                res.arr[i][j]%=mod;
            }
        }
    return res;
}

Matrix MatIdentity()
{
    Matrix res;
    for(int i = 0 ; i<MatSize ; i++)
        for(int j = 0 ; j<MatSize ; j++)
            res.arr[i][j] = (i==j);
    return res;
}

```

```

Matrix MatPower(Matrix a , long long x)
{
    Matrix res = MatIdentity();
    while(x)
    {
        if(x&1)
        {
            res = MatMul(res , a);
        }
        a = MatMul(a,a);
        x>>=1;
    }
    return res;
}

```

TO BINARY

```

string toBinary(int num)
{
    string tmp = "";
    if(!num) tmp = "0";
    while(num) tmp+=( (num&1)+'0'),num>>=1;
    reverse(tmp.begin(),tmp.end());
    return tmp;
}

```

DIGIT SUM

```

int DigitSum(ll x)
{
    int ret = 0;
    while(x) ret+=(x%10),x/=10;
    return ret;
}

// DigitSum(a+b) = DigitSum(DigitSum(a)+DigitSum(b))
// DigitSum(a-b) = DigitSum(DigitSum(a)-DigitSum(b))
// DigitSum(a*b) = DigitSum(DigitSum(a)*DigitSum(b))

```

IS PALINDROME

```

bool Is_palindrome(string &p){
    int s = 0 , e = p.size() - 1;
    while(s < e){
        if(p[s] != p[e])return false;
        s++ , e--;
    }
    return true;
}

```