DATASTRUCTURES

BIT

```
const int N=1e5+10;
int n, BIT[N];

void UpdateIndex(int idx,int value) // O(logn) - Updates value of a certain index
{
    for(;idx<=n;idx+=idx&-idx)
        BIT[idx]+=value;
}

int GetAns(int idx) // O(logn) - Gets answer in range [1:idx]
{
    int sum = 0;
    for(;idx;idx-=(idx&-idx))
        sum+=BIT[idx];
    return sum;
}</pre>
```

SEGMENT TREE

```
// Size = 4*N as there are children to the leaves
// The array a[N] is 1-based so as to correspond to the Build function
const int N=1e5+10;
int n,a[N]; // Don't redefine n
int seg[4*N],lazy[4*N];
void propagate(int ind, int 1, int r) // Propagate contains 1,r for summation values
{
        seg[ind*2+1]|=lazy[ind];
        seg[ind*2+2]|=lazy[ind];
       lazy[ind*2+1]|=lazy[ind];
        lazy[ind*2+2]|=lazy[ind];
       lazy[ind]=0;
}
void Build(int ind=0, int l=1, int r=n) // O(nlogn)
        if(l==r)
               seq[ind]=a[1];
               return;
        int mid=l+(r-l)/2;
       Build(ind*2+1,1,mid);
       Build (ind*2+2, mid+1, r);
        seg[ind] = seg[ind*2+1] | seg[ind*2+2];
}
```

```
void UpdateIndex(int qind,int qv, int ind=0, int l=1, int r=n) // O(logn)
        if(gind<l || gind>r)
                               return;
       if(l==r)
               seg[ind] |=qv;
               return;
       int mid=l+(r-l)/2;
       UpdateIndex(qind,qv,ind*2+1,l,mid);
       UpdateIndex(qind, qv, ind*2+2, mid+1, r);
        seg[ind] = seg[ind*2+1] | seg[ind*2+2];
void UpdateRange(int qx, int qy, int qv, int ind=0, int l=1, int r=n) // O(logn) with
lazy propagation
        if(r<qx || 1>qy)
                            return:
       if(1>=qx && r<=qy)
               seg[ind]|=qv;
               lazy[ind] |=qv;
               return;
       int mid=l+(r-l)/2;
       propagate(ind, l, r);
       UpdateRange (qx, qy, qv, ind*2+1, l, mid);
       UpdateRange(qx,qy,qv,ind*2+2,mid+1,r);
       seg[ind] = seg[ind*2+1] | seg[ind*2+2];
int GetAns(int qx, int qy, int ind=0, int l=1, int r=n, bool isRangeProblem=0) // O(logn)
       if(r<qx || 1>qy)
                               return 0;
       if(1>=qx && r<=qy)
                               return seg[ind];
       int mid=l+(r-l)/2;
       if(isRangeProblem)
                               propagate(ind, l, r);
       int r1=GetAns(qx,qy,ind*2+1,l,mid,isRangeProblem);
       int r2=GetAns(qx,qy,ind*2+2,mid+1,r,isRangeProblem);
       return r1|r2;
```

SPARSE TABLE

```
const int N = 1e5 + 5;
const int logN = 20;
int n, a[N], ST[logN][N], LOG[N];

void Build() // O(nlogn) - Builds the sparse table (min sparse table)
{
    LOG[0] = -1;
    for(int i=0;i<n;i++) // Computes floor(log2(i)) for all values
        ST[0][i] = i, LOG[i+1] = LOG[i] + !(i & (i+1));</pre>
```

TRIE

```
const int N = 1e5 + 5;
const int logN = 20;
int n, a[N], ST[logN][N], LOG[N];
void Build() // O(nlogn) - Builds the sparse table (min sparse table)
{
    LOG[0] = -1;
    for(int i=0;i<n;i++) // Computes floor(log2(i)) for all values</pre>
        ST[0][i] = i, LOG[i+1] = LOG[i] + !(i & (i+1));
    for (int j=1; (1<<j) <= n; j++)</pre>
        for (int i=0; (i + (1<< j)) <= n; i++)</pre>
            {
                 int x = ST[j-1][i];
                 int y = ST[j-1][i + (1 << (j-1))];
                 ST[j][i] = (a[x] \le a[y] ? x : y); // Changes according to the operation
            }
}
int GetAns(int 1, int r) // O(1) - Queries for the values from 1 to r
    int g = LOG[r-l+1];
    int x = ST[q][1];
    int y = ST[g][r - (1 << g) + 1];
    return (a[x] \leftarrow a[y] ? x : y); // Changes according to the operation
```

}

DYNAMIC PROGRAMMING

LCS

```
vector<int> vec1, vec2;
int dp[5005][5005];
int LCS(int i, int j)
    if(i == vec1.size() || j == vec2.size()) return 0;
    if(dp[i][j] != -1) return dp[i][j];
    if(vec1[i] == vec2[j])
        return dp[i][j] = LCS(i+1,j+1) +1;
    return dp[i][j] = max(LCS(i,j+1), LCS(i+1,j));
}
                                       LIS O(N^2)
const int N = 1e5+10;
int dp[N];
vector<int> a,path;
int LIS()
    for (int i=0; i<N; i++) dp[i] = 1;</pre>
    for (int i=1; i < a.size(); i++)</pre>
        for (int j=0; j<i; j++)</pre>
             if(a[i] > a[j] \&\& dp[i] < dp[j]+1)
                 dp[i] = dp[j] + 1;
    int ans = 0, mx = 0, idx = 0;
    for (int i=0; i<N; i++)</pre>
        if(dp[i] > mx)
            mx = dp[i], idx=i;
    ans = mx;
    for (int i=idx; i>=0; i--)
        if(dp[i] == mx)
            path.pb(a[i]), mx--;
    reverse(path.begin(),path.end());
    return ans;
                                     LIS O(NlogN)
const int N = 1e5+5;
```

```
int n,a[N];
```

```
int getLIS()
{
    if(n<1) return 0;

    int len = 0;
    vector<int> LIS(n,oo);

    for(int i=0;i<n;i++)
    {
        int idx = lower_bound(LIS.begin(),LIS.end(),a[i]) - LIS.begin();
        LIS[idx] = a[i];
        len = max(len,idx);
    }

    return len+1;
}</pre>
```

GRAPHS DSU

```
const int N = 1e5+5;
int par[N];
int num[N];
for(int i=0;i<N;i++) // Initialization in main</pre>
       par[i] = i;
        num[i] = 1;
}
int FindSet(int n) // Use FindSet instead of par
        if(par[n] == n)
                return n;
        return par[n] = FindSet(par[n]);
void Union(int x,int y)
        x = FindSet(x);
        y = FindSet(y);
        if(x==y)
                return;
        par[x]=y;
        num[y] += num[x]; // num holds number of elements in a group
}
```

FLOYD

```
// Shortest path between all pairs of nodes, negative edges allowed, no negative cycles
int n; // Number of nodes
int adj[n][n]; // adj[u][v] = cost of edge from u to v
```

```
void init()
    for (int i=0; i<n; i++)</pre>
         for (int j=0; j<n; j++)</pre>
              adj[i][j] = (i==j? 0 : oo); // Sets adj[u][u] = 0, adj[u][u] = oo if no edge
exists
void floyd()
    for (int i=0; i<n; i++)</pre>
         for (int j=0; j<n; j++)</pre>
              adj[i][j] = oo;
    for (int k=0; k<n; k++)</pre>
         for (int i=0; i<n; i++)</pre>
              for (int j=0; j<n; j++)</pre>
                  if(adj[i][k] + adj[k][j] < adj[i][j])
                       adj[i][j] = adj[i][k] + adj[k][j];
bool checkNegativeCycle()
    bool ret = false;
    for(int i=0;i<n;i++)</pre>
         ret = ret || (adj[i][i] < 0);
    return ret;
}
                                        IS BIPARTITE
// Shortest path between all pairs of nodes, negative edges allowed, no negative cycles
int n; // Number of nodes
int adj[n][n]; // adj[u][v] = cost of edge from u to v
void init()
    for (int i=0; i<n; i++)</pre>
         for (int j=0; j<n; j++)</pre>
              adj[i][j] = (i==j? 0 : oo); // Sets adj[u][u] = 0, adj[u][u] = oo if no edge
exists
void floyd()
    for (int i=0; i<n; i++)</pre>
         for (int j=0; j<n; j++)</pre>
              adj[i][j] = oo;
    for (int k=0; k<n; k++)</pre>
         for (int i=0; i<n; i++)</pre>
              for (int j=0; j<n; j++)</pre>
                  if(adj[i][k] + adj[k][j] < adj[i][j])
                       adj[i][j] = adj[i][k] + adj[k][j];
```

```
bool checkNegativeCycle()
    bool ret = false;
    for (int i=0; i<n; i++)</pre>
        ret = ret || (adj[i][i] < 0);
    return ret;
                                 KRUSKALL MST
const int N = 1e5+5;
vector< pair< int, pair<int, int> > edges; // {cost, {u,v}}
int par[N];
int num[N];
for(int i=0;i<N;i++) // Initialization in main</pre>
{
       par[i] = i;
       num[i] = 1;
int FindSet(int n) // Use FindSet instead of par
        if(par[n] == n)
               return n;
       return par[n] = find set(par[n]);
void Union(int x,int y)
       x = find set(x);
       y = find set(y);
        if(x==y)
               return;
       par[x]=y;
       num[y] += num[x]; // num holds number of elements in a group
int getMST() // Returns summation of edges of Minimum Spanning Tree and builds the tree
in DSU
    int sum = 0;
    sort(edges.begin(),edges.end());
    //reverse(edges.begin(), edges().end()); // Uncomment for Max. spanning tree
    for (auto e : edges)
        int u = e.s.f; int v = e.s.s;
        if(FindSet(u) == FindSet(v))
            continue;
        Union(u,v);
        sum+=e.f;
    return sum;
}
```

LCA O(1)

```
int const N = 1e5+10;
int const nN = 2*N;
int id, ID[N], rID[N], First[N], ST[nN][logN], lg[nN];
vector<int>vec;
void dfs(int cur,int par)
       rID[ID[cur]] = cur;  // reverseID to get back to the original node name.
       vec.push back(ID[cur]);
                                             //Push ur tour.
       for(auto x:adj[cur])
               if(x!=par)
               dfs(x,cur),vec.push back(ID[cur]);
}
inline void Build()
       int cnt = 0, two = 1;
    while(two<nN) lg[two] = cnt++,two*=2;</pre>
    for (int i=3; i< nN; ++i) if (!lq[i]) lq[i] = lq[i-1]; //Create an array to get the log of
any number [nearest log].
       //the previous part should be done only one time! Put it before the test cases.
       int nN = sz(vec);
       memset(First, -1, sizeof First);
       for (int i=0; i < nN; ++i)</pre>
                                     // Fill the "First" array. remeber this array
contains the idx of the first appearence of an id.
       ST[i][0] = vec[i], (First[vec[i]] == -1? First[vec[i]] = i:0);
       // Bulding the sparse table.
       for (int i=1; (1<<i) <= nN; ++i)</pre>
       for (int j=0; j+(1<<i) -1<nN; ++j)</pre>
               ST[j][i] = min(ST[j][i-1], ST[j+(1<<(i-1))][i-1]);
inline int LCA(int a,int b)
       int L,R,idx,d,lca;
       a = ID[a]; b = ID[b];
                                     // Go to ur domain.
       L = First[a], R =First[b];
       if(L>R) swap(L,R);
       d = R-L+1;
       d = lq[d];
       idx = R - (1 << d) + 1;
       lca = min(ST[L][d],ST[idx][d]);
       lca = rID[lca];
       a = rID[a]; b = rID[b];
                                                      //Back to main domain.
       return lca;
}
```

LCA O(logN)

```
const int N = 1e5+10;
const int logN = 20;
int ST[logN][N],depth[N];
```

```
void dfs(int cur, int par)
        ST[0][cur] = par;
        for (int i=1; i < logN; ++i)</pre>
               ST[i][cur] = ST[i-1][ST[i-1][cur]];
        for (auto x:adj[cur])
               if(x != par)
                       depth[x]=depth[cur]+1,dfs(x, cur);
}
int KthAncestor(int u, int k)
        if(!k) return u;
        int d = depth[u] - k; // Get the depth of the wanted node.
        for (int j = logN-1; j >= 0; --j)
               int nu = ST[j][u];
               if (!nu) continue;
               if (depth[nu] == d) return nu;
               else if (depth[nu] > d && u) u = nu;
       return 0;
int LCA(int a, int b)
        if (depth[a] > depth[b]) a = KthAncestor(a, depth[a] - depth[b]);  // make them
at the same depth.
       else if(depth[b] > depth[a]) b = KthAncestor(b, depth[b] - depth[a]);
        if (a == b) return a;
        for (int j = logN-1; j >= 0; --j)
               if (ST[j][b] != ST[j][a])
                       a = ST[j][a], b = ST[j][b];
       return ST[0][a];
}
                       LONGEST WEIGHTED PATH
const int N = 1e5+5;
vector<pair<int, int> > adj[N];
11 dfs(int u=0, int par=-1)
    11 \text{ ret} = 0;
    for (int i=0; i < adj [u] . size(); i++)</pre>
        if (adj[u][i].f != par)
            ret = max(ret, dfs(adj[u][i].F,u) + adj[u][i].s);
    return ret;
}
```

TARJAN SCC

```
vector<vector<int> > SCCs /* The components itself*/;
#define comps SCCs
vector<int> compIndex /* for each node, what is the index of the
component this node inside*/
, ind, lowLink;
stack<int> st;
vector<bool> inst;
vector<vector<int> > adj /*The intial graph*/;
int idx = 0; //must be intialized by zero;
void tarjanSCC(int i) {
    lowLink[i] = ind[i] = idx++;
    st.push(i);
    inst[i] = true;
    for (int j = 0; j < adj[i].size(); j++) {</pre>
        int k = adj[i][j];
        if (ind[k] == -1) {
            tarjanSCC(k);
            lowLink[i] = min(lowLink[i], lowLink[k]);
        } else if (inst[k]) {
            lowLink[i] = min(lowLink[i], lowLink[k]);
    if (lowLink[i] == ind[i]) {
        vector<int> comp;
        int n = -1;
        while (n != i) {
            n = st.top();
            st.pop();
            comp.push back(n);
            inst[n] = 0;
            compIndex[n] = comps.size();
        comps.push back(comp);
    }
}
void SCC() {
    comps.clear();
    compIndex.resize(adj.size());
    ind.clear();
    ind.resize(adj.size(), -1);
    lowLink.resize(adj.size());
    inst.resize(adj.size());
    idx = 0; //must be intialized by zero;
    for (int i = 0; i < adj.size(); i++)</pre>
        if (ind[i] == -1)
            tarjanSCC(i);
}
int cntSrc /*the number of source components*/,
cntSnk /*the number of sink components*/;
vector<vector<int> > cmpAdj /*The new graph between components*/;
vector<int> inDeg, outDeg /*the in degree and out degree for each
component*/;
```

```
void computeNewGraph() {
    outDeg.clear();
    outDeg.resize(comps.size());
    inDeq.clear();
   inDeg.resize(comps.size());
    cntSrc = cntSnk = comps.size();
    cmpAdj.clear();
    cmpAdj.resize(comps.size());
    for (int i = 0; i < adj.size(); i++) {</pre>
    for (int j = 0; j < adj[i].size(); j++) {</pre>
            int k = adj[i][j];
            if (compIndex[k] != compIndex[i]) {
                cmpAdj[compIndex[i]].push back(compIndex[k]);
                if (!(inDeg[compIndex[k]]++))
                    cntSrc--;
                if (!(outDeg[compIndex[i]]++))
                    cntSnk--;
            }
    }
    }
                                TREE DIAMETER
```

```
const int N = 1e5+5;
int sp[N];
int bfs(int u)
    queue<int> q;
    q.push(u);
    memset(sp,-1,sizeof(sp));
    sp[u] = 0;
    while(!q.empty())
        u = q.front();
        q.pop();
        for(auto v : adj[u])
            if(sp[v] == -1)
                sp[v] = sp[u] + 1,q.push(v);
    }
    return u;
}
int calcTreeDiameter(int root)
    int u = bfs(root);
    int v = bfs(u);
    return sp[v];
}
```

BELLMAN FORD

```
// Shortest path from source to all nodes, negative edges allowed, no negative cycle
const int N = 1e5+5;

vector<pair<int,int> > adj[N];
```

```
int n,par[N]; // n is number of nodes
bool bellmanFord(int source)
    memset(par,-1,sizeof par);
    memset(dis, 0x3F, sizeof(dis));
    dis[source] = 0;
    bool updated = 1;
    for(int k=0; k<n && updated; k++)</pre>
        updated = 0;
        for (int u = 1; u<=n;u++)</pre>
             for(auto& e : edges[u])
                 int v = e.first;
                 int c = e.second;
                 if(dis[v] > dis[u] + c)
                     dis[v] = dis[u] + c;
                     par[v] = u;
                     updated = 1;
            }
       }
    }
    return updated; // Whether a negative cycle exist or not
}
                                             BFS
const int N = 1e5+5;
vector<int> adj[N];
int sp[N];
void bfs()
    memset(sp,-1,sizeof sp);
    queue<int> q;
    q.push(s); //pushes start
    sp[s] = 0;
    while(!q.empty()){
        int u = q.front();
        q.pop();
        for (int i=0; i < adj [u] . size(); i++) {</pre>
             int v = adj[u][i];
             if(sp[v] == -1){
                 q.push(v);
                 sp[v] = sp[u] + 1;
             }
        }
    }
}
```

DFS

```
const int N = 1e5+5;
vector<int> adj[N];
int visited[N];
void dfs(int u) {
       visited[u] = 1;
       for(int i=0;i<adj[u].size();i++){</pre>
              int v = adj[u][i];
              if(!visited[v])
                     dfs(v);
       }
}
                                   DIJKSTRA
// Dijkstra SP for weighted graph + path build, No negative edges
const int N = 1e5+5;
vector< pair<int, int> > adj[N]; // adj[u] = {v,cost}
vector<int> path;
int sp[N]; int par[N];
void build(int u) {
       while (u != -1) {
              path.pb(u);
              u = par[u];
       }
       }
int Dijkstra(int source, int destination)
   int u, v, c;
   for (int i = 0; i < n; ++i) // initialization</pre>
       sp[i] = oo, par[i] = -1;
   priority queue< pair<int,int> , vector<pair<int,int> >, greater<pair<int,int> > > pq;
   pq.push( { 0, source });
   sp[source] = 0;
   while (!pq.empty()) {
       u = pq.top().second;
       c = pq.top().first;
       pq.pop();
       if (sp[u] < c)
           continue;
       for (auto edge : adj[u]) {
           int to = edge.first;
           int cost = edge.second;
           if (c + cost < sp[to]) {
```

MATH

```
PRIME FACTORIZATION
const int MAX = 1e6+10;
int spf[MAX], power[MAX]
void Sieve() // O(nlogn)
        spf[1] = 1;
        for (int i = 2; i < MAX; i++) spf[i] = i;</pre>
        for (int i = 4; i < MAX; i += 2) spf[i] = 2;</pre>
        for (int i = 3; i < MAX; i++)</pre>
               if (spf[i] == i)
                       for (int j = i; j < MAX; j += i)</pre>
                               if (spf[j] == j)
                                       spf[j] = i;
}
void GetFactors(int num)
        int cnt = 1, last = spf[num], ans = 1;
       while (num != 1) {
               num = num / spf[num];
               if (last == spf[num]) {
                       cnt++;
                       continue;
               power[last] = cnt;
               cnt = 1;
               last = spf[num];
        }
}
                                 nCr WITH MOD
ll mod = 1e9+7;
const int N = 1e5+10;
11 fact[N];
void ComputeFactorial() // Don't forget to compute the factorial
    fact[0] = 1;
    for(ll i=1;i<N;i++)</pre>
        fact[i] = (fact[i-1] % mod * i% mod) % mod;
```

}

```
11 power(ll x, ll y)
    ll res = 1;
    x = x % mod;
    while (y > 0)
        if (y & 1)
           res = (res*x) % mod;
        y = y >> 1;
        x = (x*x) % mod;
    return res;
11 modInverse(ll n)
    return power(n, mod-2);
ll nCr(ll n, ll r)
   if (r==0)
      return 1;
    return (fact[n] * modInverse(fact[r]) % mod *
            modInverse(fact[n-r]) % mod) % mod;
}
                                          SIEVE
const int MAX = 1e6+10;
vector<bool> prime(MAX, true);
void Sieve() // O(nlog(logn))
{
        prime[0] = prime[1] = 0;
        for (int i = 2; i*i <= MAX; i++)</pre>
                if (prime[i])
                        for (int j = i * 2; j <= MAX; j += i)</pre>
                                prime[j] = 0;
}
                        SMALLEST PRIME FACTOR
const int MAX = 1e6+10;
int spf[MAX];
void Sieve() // O(nlogn)
        spf[1] = 1;
        for (int i = 2; i < MAX; i++) spf[i] = i;</pre>
        for (int i = 4; i < MAX; i += 2) spf[i] = 2;</pre>
        for (int i = 3; i < MAX; i++)</pre>
                if (spf[i] == i)
                        for (int j = i; j < MAX; j += i)</pre>
                                if (spf[j] == j)
                                        spf[j] = i;
}
```

EULER PHI

```
// Theorem: a^{h}(b) % b = 1, if a,b are coprimes (\gcd(a,b) == 1)
ll Phi(ll n) // O(sqrt(n)), returns Euler Phi of n
    ll result = n;
    for (int i=2; i*i<=n; i++)</pre>
        if(n%i == 0)
            result -= result/i;
        while (n\%i==0)
            n/=i;
    if(n>1)
        result -=result/n;
    return result;
ll phi[N];
void EulerTotient() // O(nlog(logn)), generates all Euler Phi's to N
    for (int i=0; i<N; i++) phi[i] = i;</pre>
    for (int i=2; i<N; i++)</pre>
        if(phi[i] == i)
            for(int j=i;j<N;j+=i) phi[j] -= phi[j]/i;</pre>
}
                               EXTENDED EUCLID
// O( log(max(a,b)) )
// Returns Bezout's coeffs of the smallest
// linear combination of s.a + t.b = gcd(a,b)
pair<int, int> ExtendedEuclid(int a, int b)
    if (b==0)
        return {1,0};
    pair<int, int> p = ExtendedEuclid(b,a%b);
    int s = p.first, t = p.second;
    return {t, s-t * (a/b)};
}
                         FASTPOWER WITH MOD
ll power(ll base, ll exp)
    ll ans = 1;
    base%=mod;
    while (exp>0)
        if(exp&1) ans = (ans*base)%mod;
        exp>>=1;
        base = (base*base) %mod;
    }
    return ans;
}
```

IS PRIME

bool IsPrime(int n)

```
if(n<=1) return false;</pre>
   if(n==2) return true;
   if(n%2==0) return false;
   for(int i=3 ; i*i<=n ; i +=2) {</pre>
       if(n%i==0)
               return false;
    return true;
}
                                     MATH MISC
ll gcd(ll a, ll b) { return (b == 0 ? a : gcd(b, a % b)); } // O(log(max(a,b)))
ll lcm(ll a, ll b) { return ((a*b) / gcd(a, b)); } // O(log(max(a,b)))
11 modInverse(ll a, ll m) // O(logm)
{
    return power(a, m-2, m);
}
ll nCr(ll n, ll r) // O(r) - Call the function as nCr(n, min(r, n-r)) for better
performance
{
    if(n<r) return 0;</pre>
    if(r==0) return 1;
    return n*nCr(n-1, r-1)/r;
}
// Factorization - O(sqrt(n))
vector<11>factors;
void GetFactors(ll n) {
        for (int i = 1; i*i <= n; i++) {</pre>
                if (n%i == 0) {
                        if (i != n / i)
                                factors.push_back(i), factors.push_back(n / i);
                                 factors.push back(i);
                }
       }
}
//Mod of very large numbers - O(n)
//\text{rule:} xy \pmod{a} \equiv ((x \pmod{a} * y) \pmod{a})
11 mod(string num)
{
    int res = 0;
    for (int i = 0; i < num.size(); i++)</pre>
         res = (res*10 + num[i] - '0') %mod;
    return res;
}
```

MATRIX POWER

```
const int mod = 1e9+7;
const int MatSize = 2;
struct Matrix
        long long mat[MatSize][MatSize];
};
Matrix MatMul (Matrix &a , Matrix &b)
        Matrix res;
    for(int i = 0 ; i<MatSize ; i++)</pre>
        for(int j = 0 ; j<MatSize ; j++)</pre>
        res.arr[i][j] = 0;
        for(int k = 0 ; k<MatSize ; k++)</pre>
             long long sum = a.arr[i][k] * b.arr[k][j];
             sum%=mod;
             res.arr[i][j] +=sum;
             res.arr[i][j]%=mod;
        }
   return res;
}
Matrix MatIdentity()
        Matrix res;
        for(int i = 0 ; i<MatSize ; i++)</pre>
        for(int j = 0 ; j<MatSize ; j++)</pre>
        res.arr[i][j] = (i==j);
        return res;
}
```

```
Matrix MatPower (Matrix a , long long x)
       Matrix res = MatIdentity();
       while(x)
       if (x&1)
              res = MatMul(res , a);
       a = MatMul(a,a);
       x>>=1;
       return res;
}
                                   TO BINARY
string toBinary(int num)
       string tmp = "";
       if(!num) tmp = "0";
       while (num) tmp+= ( (num&1) + '0'), num>>=1;
       reverse(tmp.begin(),tmp.end());
       return tmp;
}
                                    DIGIT SUM
int DigitSum(11 x)
       int ret = 0;
       while (x) ret+= (x%10), x/=10;
       return ret;
}
// DigitSum(a+b) = DigitSum(DigitSum(a) + DigitSum(b))
// DigitSum(a-b) = DigitSum(DigitSum(a)-DigitSum(b))
// DigitSum(a*b) = DigitSum(DigitSum(a)*DigitSum(b))
                               IS PALINDROME
bool Is_palindrome(string &p){
```

```
int s = 0 , e = p.size() - 1;
       while(s < e) {
       if(p[s] != p[e])return false;
       s++ , e--;
       return true;
}
```