

2) For a schedule to be recoverable, a commit must happen after all of the data is written. Reads can happen after a commit, but not writes.

T_1 can commit after it writes A

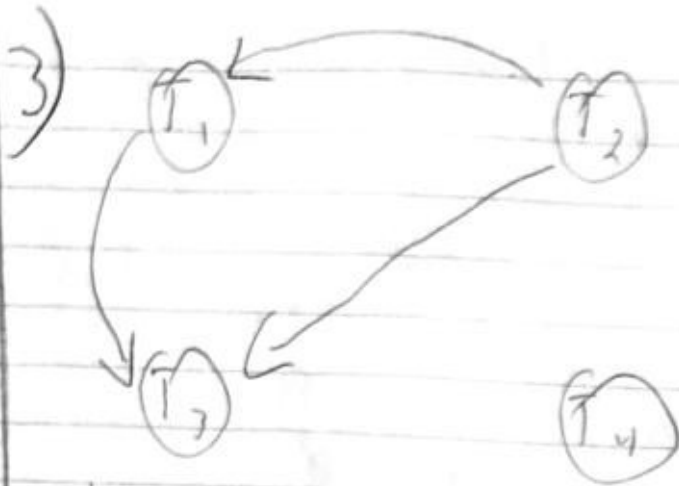
T_2 does not need to commit

T_3 can commit after it writes B

T_4 can commit after it writes C.

schedule table with commit:

T_1	T_2	T_3	T_4
read(A)	read(A)		
		read(B)	
read(D) write(A)	read(B)		
			read(C) write(C)
		read(A) write(B)	
		commit()	read(D)



This schedule is serializable because there are no cycles in the precedence graph.