

## Bases de dades NoSQL.

### Distributed architectures - consistent hashing

index: associa clau amb direcció física

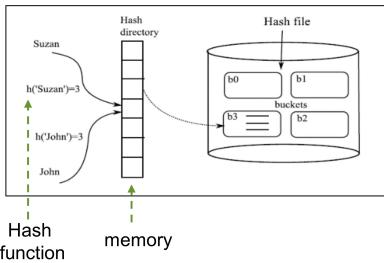
↳ Tree → complexitat de cerca logarítmica

↳ Hash tables → complexitat de cerca constant

✓ consultes ponctuals

✗ queries de rang

✗ raycasting collections (canviem ràpidament)



Què passa si les dades creixen molt?

→ Refresher de hash.

Distributed hashing challenges

- Dynamicity: grow and shrink rapidly
  - Distribution: Assign buckets to participating nodes
    - \*all the nodes should share the hash function for it to work

E.g.,  $h(x) = x \% \# \text{servers}$



- Adding a new server implies modifying  $h$ ...
  - Re-hashing all the objects
  - Communicating the new  $h$  to all servers

- Location of the hash directory: any access must go through the hash directory
  - Potential bottleneck

### Consistent hashing

Results of the most common queries are in **caches** (i.e., in-memory) of several servers

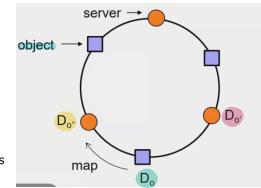
A dedicated **proxy** machine records which server stores which query results
 

- Queries are assigned to servers according to a hash function over the query

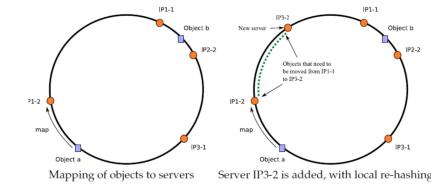
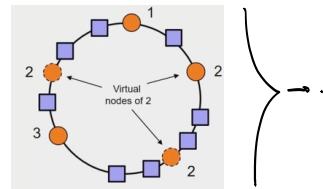
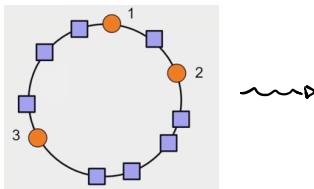
#### Copying with dynamicity:

- The hash function **never** changes

- Choose a very large domain  $D$  (address space) and map server IP addresses and object keys to such domain
- Organize  $D$  as a ring in clockwise order so each node has a successor
- Objects are assigned as follows:
  - For an object  $O$ ,  $f(O) = D_o$ 
    - Let  $D_{o'}$  and  $D_o$  be the two nodes in the ring such that
      - $D_{o'} < D_o \leq D_o'$
      - $O$  is assigned to  $D_{o'}$
  - Adding a new server is straightforward
    - It is placed in the ring and part of its successors objects transferred



### Examples



#### Further refinements:

- Assign to the same server several hash values (virtual servers) to balance load

- Lazy update of the hash directory

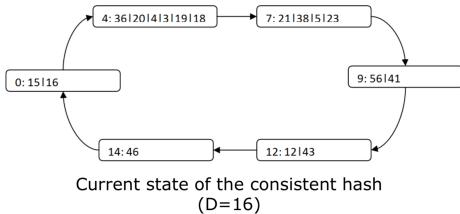
### Activity - consistent hashing

Objective: Understand how the consistent hash works

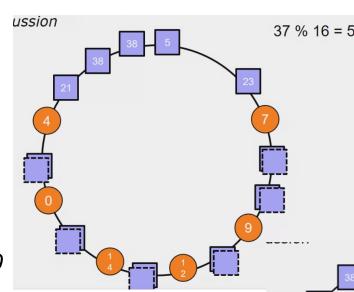
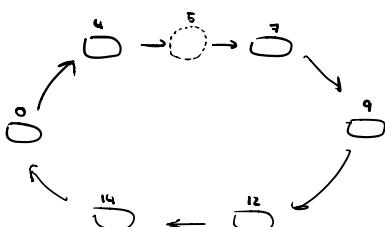
Tasks:

- (5') By pairs, solve the following exercise
  - What happens in the structure when we register a new server with IP address '37'? Draw the result.
- (5') Discussion

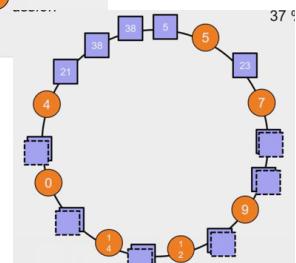
$$f(\text{key}) = \text{IP} \% 16$$



$$37 \% 16 = 5$$



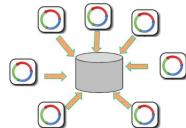
per mòdul



# KEY-VALUE ENHANCEMENTS DOCUMENT-ORIENTED DBS

## Applications Databases.

- SQL and relational databases played a key role as **integration mechanism between applications**
    - Multiple applications using a common integrated database
      - More complex
      - Changes by different apps need to be coordinated
      - Different apps have different performance needs, thus call for different index structures
      - Complex access policies



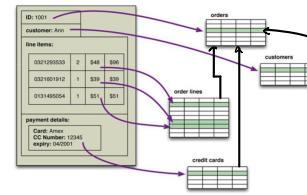
A different approach, treat your database as an **application database**

- ❑ An application database is only directly accessed by a single application, which makes it much **easier to maintain and evolve**
  - ❑ Interoperability concerns can now shift to the interfaces of the application:
    - During the 2000s we saw a shift to web services, where applications would communicate over HTTP
  - ❑ With a service you are able to use **richer data structures** (compared to SQL)
    - Usually represented as documents in XML or, more recently JSON

## Aggregate data models

- The relational model divides the information that we want to store into **tuples** (rows): this is a very simple structure for data
  - **Aggregate orientation** takes a different approach. It recognizes that often you want to operate on data in **units that have a more complex structure**
    - Think of it as a complex record that allows lists and other record structure to be nested inside
  - Key-value, document and column-family DBs can all be seen as aggregate-oriented databases
    - They differ in how they structure the aggregate and consequently how they allow for it to be accessed

- What is **good** about these models?
    - Dealing with aggregates makes it much easier for the databases to handle **operating on a cluster**, since the aggregate makes a natural unit for replication and sharding.
    - Also a natural unit to use for distribution
    - Also, it may help **solving the impedance mismatch problem**, i.e., the difference between the relational model and the in-memory data structures
    - The impedance mismatch is naturally solved in document-stores



An order, which looks like a single aggregate

## Structuring "value"

- ❑ Essentially, they are key-value stores
    - Same design and architectural features
  - ❑ The value is a document
    - XML (e.g., eXist)
    - JSON (e.g., MongoDB and CouchDB)
  - ❑ Tightly related to the Web
    - Normally, they provide RESTful HTTP APIs
  - ❑ So... what is the benefit of having documents?
    - ❑ New data model (collections and documents)
      - **New atom: from rows to documents**
    - ❑ Indexing

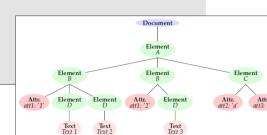
## Types of document-store

- ❑ JSON-like documents
    - MongoDB
    - CouchDB
  - ❑ JSON is a lightweight data interchange format
    - Brackets ([ ]) represent ordered lists
    - Curly braces ({} ) represent key-value dictionaries
      - Keys must be strings, delimited by quotes ("")
      - Values can be strings, numbers, booleans, lists, or key-value dictionaries
  - ❑ Natively compatible with JavaScript
    - Web browsers are natural clients for MongoDB / CouchDB  
<http://www.json.org/index.html>
  - ❑ XML-like documents
    - eXist, MarkLogic
      - Native XML extensions
    - Relational extensions for Oracle, PostgreSQL, etc.
      - Mapped to relational (impedance mismatch)
  - ❑ XML is a semistructured data model proposed as the standard for data exchange on the Web
    - Can be elegantly represented as a tree
      - Document: the root node of the XML document, denoted by "/"
      - Element: element nodes that correspond to the tagged nodes in the document
      - Attribute: attribute nodes attached to Element nodes
      - Text: text nodes, i.e., untagged leaves of the XML tree
  - ❑ Support Xpath, XQuery and XSLT
    - Xpath is a language for addressing portions of an XML document
      - Subtree selection
    - XQuery is a query language for extracting information from collections of XML documents
    - XSLT is a language for specifying transformations (from XML to XML)

## JSON Example

- **Definition:**  
*A document is an object represented with an unbounded nesting of array and object constructs.*

## XML Example



Labeled  
Unranked  
ambivalent

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