Performance Prediction for GPU-based HPC codes

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Abstract—The purpose of this report is to performance prediction for GPU-based HPC code. The prediction of execution time is divided into two part. One is Kernel execution time. Another is memory copy of a randomly filled floating-point array from CPU host to GPU device (H2D) and from GPU device to CPU host (D2H). We will use R programming language to explore the correlation between different values and find the prediction model.

Keywords—C, Canny Edge Detector, CUDA, GPU, Prediction.

I. INTRODUCTION

The purpose of this report is to analyze the accelerating by designing a serial code and then using CUDA method to speedup. In Canny Edge Detector, there are several kernels, such as convolution, Suppressed, and Hysteresis. In those kernels, the program needs to get and calculate the value of each pixels. They take up execute time the most, thus by executing on GPUs device, we can effectively to observe the reducing of execution time. In addition, in order to prediction, we divided our execution time into three parts: kernel execution time, memory copy from host to device and memory copy from device to host. So, first we implemented the features detection. This kernel will detect different numbers of features along with different block sizes and image sizes. We tested program with several image sizes which are 256, 512, 1024, 2048, 4096, 8192, 10240. This time we only use one sigma 0.6. And, the most important part is block size. We tested with block sizes: 8,12,16,20,24,28,32. We have run 50 times to get a solid result. Moreover, for each kernel time, we considered that float point computation times and global memory access times to predict our kernel time.

II. IMPLEMENTATION

A. Canny edge detector with CUDA

We used this function to detect the edge of image. In order to effectively reduce our execution time. We used GPUs to implement our accelerating. On GPUs device, we can use global memory and shared memory to speed up our program. In convolution part, especially, we used shared memory. Other kernels used only global memory. We only used shared memory for convolution because during the calculation, some data were out of the bound that the threads need to access global memory to get the data what it needs. We can restore our data to shared memory to enhance their speed. For other part of functions, like Suppressed and Hysteresis, we only use global memory for implementation. This is due to there is not calculation which needs the data from different block.

B. Corner Detection on GPU

In the Corner features detection kernel, the first step is to evaluate cornerness values for each pixel by computing the Z matrix over the neighborhood window (7x7) which is a similar step to convolution. After computing the cornerness values, we need to implement an algorithm which helps us

to find the location of the maximum cornerness in a window of BLOCKSIZExBLOCKSIZE.

The following is how to implement:

- 1. Set the thread block dimensions as dimBlock(BLOCKSIZE,BLOCKSIZE)
- 2. Using shared memory, have threads in a thread block find the location of the maximum cornerness value. Have thread-0 of the block write the location of the maximum corner to the list of corners.

The following figure is an example of corners output with image-size is 256x256 and Block-Size is 8.

corneress

Index:6512	31 25 4	J:32 J:80 J:112 J:159
Index:6512	25 4	J:112
	4	
Indovidado II		J:159
index: 1103	3	
Index:931 :	-	J:163
Index:978	3	J:210
Index:1021	3	J:253
Index:11267	44	J:3
Index:13119	51	J:63
Index:13890 :	54	J:66
Index:15212	59	J:108
Index:10132	39	J:148
Index:16312	63	J:184
Index:15055 :	58	J:207
Index:10237	39	J:253
Index:16650 :	65	J:10
Index:21055	82	J:63
Index:24389	95	J:69
Index:22905	89	J:121
Index:22928 :	89	J:144
Index:16568	64	J:184
Index:20427	79	J:203
Index:19936	77	J:224
Index:32522 :	127	J:10
Index:31807 :	124	J:63

It appears that the indexes and the threads location I and J.

C. Testing

After We implemented the code. We tested the program on GPU in several sizes of images and Block-Sizes. Image sizes {256, 512, 1024, 2048, 4096, 8192, and 10240} x {BLOCKSIZES 8, 12, 16, 20, 24, 28, and 32} x {sigma 0.6}. In addition, in order to get a solid execution time, we executed the program in 30 times for each parameter.

In the figure 2, it shows the average execution time in each parameter.

execution time(MicroSecond)								
BlockSize	Size Image-Hight Image_Width Sigma i_o_time kernel_time communication_time endToEndNolandO endToEndwith							endToEndwithlandO
8	256	256	0.6	1664	317	795	394535	391065
8	512	512	0.6	3683	942	2739	371915	362081
8	1024	1024	0.6	11553	2531	5873	387782	383232
8	2048	2048	0.6	31420	7165	17651	371036	301018
8	4096	4096	0.6	83210	54308	57252	603029	453749
8	8192	8192	0.6	277440	244661	214003	1466363	1022242
8	10240	10240	0.6	483865	295274	359343	2166675	1381534
12	256	256	0.6	716	272	470	222455	221153
12	512	512	0.6	1665	755	3095	277933	159817
12	1024	1024	0.6	8737	2662	4903	235379	244986
12	2048	2048	0.6	29571	10194	15384	285331	240212
12	4096	4096	0.6	32416	49322	23141	392203	352184
12	8192	8192	0.6	264705	267556	209200	1644761	1219253
12	10240	10240	0.6	424019	383380	294074	2053407	1372857
16	256	256	0.6	1527	435	662	436495	433750
16	512	512	0.6	3419	1192	2356	486852	469087
16	1024	1024	0.6	9275	4246	5791	511293	490480
16	2048	2048	0.6	39828	26777	20528	554129	458458
16	4096	4096	0.6	84981	43184	59393	618851	469431
16	8192	8192	0.6	266443	174774	252485	1642102	1203434
16	10240	10240	0.6	385434	304631	311565	2055560	1389587
20	256	256	0.6	974	15566	3538	294464	251342
20	512	512	0.6	3119	15721	3022	381569	275737
20	1024	1024	0.6	22475	17297	7533	375599	333727
20	2048	2048	0.6	45750	11165	13242	435272	362727
20	4096	4096	0.6	89172	44742	62721	663469	480882
20	8192	8192	0.6	258094	195301	214106	1529829	1120852
20	10240	10240	0.6	398766	330364	339009	2139809	1483318
24	256	256	0.6	852	7718	3382	314090	321495
24	512	512	0.6	2977	1241	2625	420554	411548
24	1024	1024	0.6	7096	4154	5629	427279	408947
24	2048	2048	0.6	24824	17371	15143	392919	392020
24	4096	4096	0.6	78888	45960	53370	661098	507868

Figure 2.

The above executions will help us to develop performance prediction models for execution time prediction.

D. Performance Modeling

256 0.6

512 0.6 3215

2048 0.6 39916

512 0.6 3007

1024 0.6 9113

4096 0.6 58321

0.6 277558

0.6

0.6 85308

0.6

0.6 429577

0.6 283911

Our application's runtime is governed by the time taken by GPU kernels, H2D communications, D2H communications, and any CPU computations. Especially, entire code is on GPUs, Gaussian kernel evaluation which is on CPU is way too small for any meaningful modeling. So, we will not consider CPU computations time. For performance modeling, we have several modeling.

1. Modeling CPU-GPU Communications

In this part, we have to implement a simple benchmark that performs memory copy of a randomly filled floating-point array from CPU host to GPU device (H2D) and from GPU device to CPU host (D2H). For a given array size, we had performed 30 statistical executions to compute the average H2D time and D2H time for these array size: 1 KB, 2 KB, 4 KB, 8 KB, 16 KB, 32 KB, 64 KB, 128 KB, 256 KB, 512 KB, 1 MB, 2 MB, 4 MB, 8 MB, 16 MB, 32 MB, 64 MB, 128 MB, 256 MB, 512 MB.

And, Calculate the bandwidth as (data size/time), plot H2D bandwidth vs. array-size (KB) and observe the trend. (Dependent variable: bandwidth, independent variable: data size). The figures 3 to 5 are H2D time and the figures 6 to 8 are D2H time. The unit of time is microsecond.

arraySize_KB	time_ms	bandwidth
1	11.70000	0.08547009
2	12.60000	0.15873016
4	13.46667	0.29702972
8	19.66667	0.40677967
16	56.50000	0.28318584
32	67.70000	0.47267358
64	86.46667	0.74016962
128	162.16667	0.78931138
256	304.89999	0.83961956
512	645.43335	0.79326549
1024	1013.29999	1.01055957
2048	2083.86670	0.98278839
4096	4047.03345	1.01209937
8196	8887.46680	0.92219754
16384	18273.09961	0.89661855
32768	29746.73242	1.10156637
65536	73312.42969	0.89392754
131072	137917.32810	0.95036644
262144	261403.17190	1.00283404
624200	543245 25000	0.96510370

Figure 3.

H2D bandwidth vs. arraySize

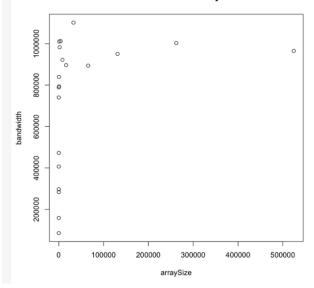


Figure 4.

Call: Im(formula = bandwidth ~ arraySize)

Residuals:

Min 1Q Median 3Q Max -602222 -231512 103842 203694 387298

Coefficients:

Estimate Std. Error t value Pr(>|t|) (Intercept) 687690.7993 75502.0145 9.108 0.0000000368 *** arraySize 0.8111 0.5577 1.454 0.163

Signif. codes: 0 '***' 0.001 '**' 0.01 '*' 0.05 '.' 0.1 ' ' 1

Residual standard error: 311300 on 18 degrees of freedom Multiple R-squared: 0.1051, Adjusted R-squared: 0.05542 F-statistic: 2.115 on 1 and 18 DF, p-value: 0.1631

Figure 5.

The formula for H2D bandwidth is: H2D_bandwidth = 687690.8+0.8111*array Size

arraySize_KB	time_ms	bandwidth
1	15.86667	0.06302521
2	15.06667	0.13274336
4	19.96667	0.20033389
8	25.30000	0.31620555
16	48.43333	0.33035099
32	89.13333	0.35901272
64	179.70000	0.35614914
128	262.43332	0.48774295
256	434.60001	0.58904739
512	906.09998	0.56505906
1024	1220.00000	0.83934426
2048	2955.00000	0.69306261
4096	4789.39990	0.85522197
8196	13376.66699	0.61270868
16384	18093.43359	0.90552188
32768	33889.73438	0.96690047
65536	64856.69922	1.01047387
131072	138962.00000	0.94322189
262144	309439.34380	0.84715795
524288	540365.93750	0.97024620

Figure 6.

D2H bandwidth vs. arraySize

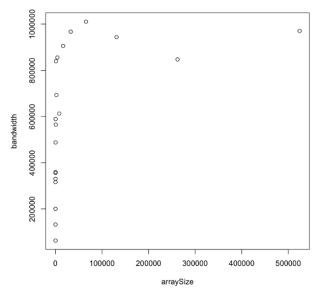


Figure 7.

Call:

Im(formula = bandwidth ~ arraySize)

Residuals:

Min 1Q Median 3Q Max -480744 -194126 16035 263684 393695

Coefficients:

Estimate Std. Error t value Pr(>|t|) (Intercept) 543768.0945 67834.5476 8.016 0.000000238 *** arraySize 1.1140 0.5011 2.223 0.0392 *

Signif. codes: 0 '***' 0.001 '**' 0.01 '*' 0.05 '.' 0.1 ' ' 1

Residual standard error: 279700 on 18 degrees of freedom (1 observation deleted due to missingness)

Multiple R-squared: 0.2154, Adjusted R-squared: 0.1718

F-statistic: 4.943 on 1 and 18 DF, p-value: 0.03925

Figure 8.

The formula for D2H bandwidth is: D2H bandwidth = 543768.1+1.1140*array Size For each of the Canny edge kernels and the cornerness kernel, we need to manually compute the number of floating-point computations and memory accesses per thread. When we multiply these with the total number of threads for that kernel, we can get total floating-point operations for the kernel (FLOPS_kernel) and total memory accesses (BYTES_kernel). Then, we explored the correlation between these two variables and the kernel runtime to decide a candidate mathematical model. (Independent variable: FLOPS_kernel, BYTES_kernel; Dependent variable: kernel time). For the approximation kernel, because we are testing for different values of BLOCKSIZE, the independent variables are FLOPS_kernel, BYTES_kernel, and BLOCKSIZE.

In the figure 9, it shows the FLOPS_kernel, BYTES_kernel, and kernel time in different image size and BLOCKSIZE. And the time unit is microsecond.

kernel_time(Microsecond)

FLOPS_kernel	BYTES_kernel	Kernel time	imagesize	blocksize
15104486	4539716	300	256	8
60963808	18575684	563	512	8
245433008	73971976	1776	1024	8
926885888	289379552	7468	2048	8
2261574400	661418656	44103	3072	8
4583325696	1800351360	56821	5120	8
4943773696	3209166848	158718	7680	8
15104486	2985540	297	256	16
60963808	12162372	820	512	16
245433008	49089028	2719	1024	16
926885888	189815488	10334	2048	16
2261574400	425587472	24727	3072	16
4583325696	529999408	67310	5120	16
4943773696	1984255872	176570	7680	16
15104486	2146436	467	256	32
60963808	8805956	1397	512	32
245433008	22295844	11021	1024	32
926885888	142522872	18643	2048	32
2261574400	302274928	39525	3072	32
4583325696	490202112	97976	5120	32
4943773696	1830021568	262450	7680	32

Figure 9.

Call:
Im(formula = Kernel.time ~ BYTES_kernel + FLOPS_kernel + blocksize)

Residuals:
Min 1Q Median 3Q Max
-66315 -18646 1315 12497 93630

Coefficients:
Estimate Std. Error t value Pr(>|t|)
(Intercept) -22996.576997398 12733.361375750 -1.806 0.0887 .

BYTES_kernel 0.000042797 0.000015612 2.741 0.0139 *
FLOPS_kernel 0.000014503 0.00006649 2.181 0.0435 *
blocksize 40.819234136 18.390685366 2.220 0.0403 *
--Signif. codes: 0 '*** 0.001 '** 0.01 '* 0.05 '.' 0.1 ' ' 1

Residual standard error: 33210 on 17 degrees of freedom
Multiple R-squared: 0.8148, Adjusted R-squared: 0.7822
F-statistic: 24.94 on 3 and 17 DF, p-value: 0.000001864

Figure 10.

Kernel time formula is

kernel_time = 22996.576997398+ BYTES kernel*0.000042797+

FLOPS_kernel*0.000014503+

BLOCKSIZE*40.81923

In figure 10, it shows the coefficients, p-values, and R^2 values and we can make the kernel time formula.

III. RESULTS AND ANALYSIS

By our prediction modeling, we predict the execution time as:

Execution Predict= Predicted H2D time + Predicted D2H time + Predicted kernel runtimes.

Predict the execution times for these configurations (all with Gaussian sigma=0.6):

- 1. Image Size 3072x 3072 with approximation algorithm BLOCKSIZE=8, 16, 32.
- 2.Image Size 5120x5120 with approximation algorithm BLOCKSIZE=8, 16, 32.
- 3.Image Size 7680x7680 with approximation algorithm BLOCKSIZE=8, 16, 32.

For the above 9 cases, we have a table called "Predicted H2D, D2H, and kernel times" (figure 11.) that gives the predicted H2D time, D2H time, kernel times, and CPU time. Also, another table called "Actual H2D, D2H, and kernel times for the actual times for H2D, D2H, kernels (figure 12.). These are obtained by executing code with the above configurations. The table (figure 13.) called "prediction errors for the tested configuration" that gives the prediction error for each of the above 9 configurations as:

Error = (predicted execution time - actual execution time)/actual execution time * 100%

Predicted H2D, D2H, kernel, and CPU times(second)

	H2D	D2H	Kernel	Execution_predict
3072x3072+ BLOCKSIZE = 8	0.01325	0.0166343	0.055320	0.0852043
3072x3072+ BLOCKSIZE = 16	0.01325	0.0166343	0.067483	0.0973673
3072x3072+ BLOCKSIZE = 32	0.01325	0.0166343	0.192938	0.2228223
5120x5120+ BLOCKSIZE = 8	0.03613	0.044732	0.028388	0.10925
5120x5120+ BLOCKSIZE = 16	0.03613	0.044732	0.075838	0.1567
5120x5120+ BLOCKSIZE = 32	0.03613	0.044732	0.198382	0.279244
7680x7680+ BLOCKSIZE = 8	0.0784	0.09474	0.044932	0.218072
7680x7680+ BLOCKSIZE = 16	0.0784	0.09474	0.123012	0.296152
7680x7680+ BLOCKSIZE = 32	0.0784	0.09474	0.309238	0.482378

Figure 11.

Actual H2D, D2H, kernel, and CPU times(second)

	H2D	D2H	Kernel	Execution_predict
3072x3072+ BLOCKSIZE = 8	0.009657566	0.012346	0.044103	0.066106566
3072x3072+ BLOCKSIZE = 16	0.009657566	0.012346	0.056821	0.078824566
3072x3072+ BLOCKSIZE = 32	0.009657566	0.012346	0.158718	0.180721566
5120x5120+ BLOCKSIZE = 8	0.0300756	0.047185	0.024727	0.1019876
5120x5120+ BLOCKSIZE = 16	0.0300756	0.047185	0.067310	0.1445706
5120x5120+ BLOCKSIZE = 32	0.0300756	0.047185	0.176570	0.2538306
7680x7680+ BLOCKSIZE = 8	0.05595713	0.07510	0.039525	0.17058213
7680x7680+ BLOCKSIZE = 16	0.05595713	0.07510	0.097976	0.22903313
7680x7680+ BLOCKSIZE = 32	0.05595713	0.07510	0.262450	0.39350713

Figure 12.

prediction errors for the tested configuration

	H2D%	D2H%	Kernel%	Execution_predict
3072x3072+ BLOCKSIZE = 8	37	35	25	29
3072x3072+ BLOCKSIZE = 16	37	35	19	24
3072x3072+ BLOCKSIZE = 32	37	35	22	23
5120x5120+ BLOCKSIZE = 8	20	-5	15	7
5120x5120+ BLOCKSIZE = 16	20	-5	13	8
5120x5120+ BLOCKSIZE = 32	20	-5	12	10
7680x7680+ BLOCKSIZE = 8	40	26	14	28
7680x7680+ BLOCKSIZE = 16	40	26	26	29
7680x7680+ BLOCKSIZE = 32	40	26	18	23

Figure 13.

As we seen, H2D and D2H execution time will not change by different BLOCKSIZE because two execution time only execute the memory copy. And CPU time are too small to predict. So, kernel time will be the important part focus on. In the tables, for kernel execution time, we got 12% to 26% prediction error which is not too bad due to the formula of kernel time whose multiple R-squared value is up to 0.8148 and adjusted R-squared also up to 0.783. For H2D and D2H formula, multiple R-squared is only around 0.1 to 0.2. And the way to enhance our accuracy of H2D and D2H, we may need more sample data. In addition, in prediction error table, the range of percentage is from -5 to 40 because the accuracy of prediction is affected by many conditions. Sometimes, the execution time is slower than our expectations due to the computing resources are competed.

IV. CONCLUSION

One of the strengths of statistics machine learning method is that allows us to do large scale of prediction which is not able to practically test on any devices due to limitation of computing resource. But on the other hand, this is also one of weakness of statistics machine meaning method. In order to getting an accurate prediction, we need to collect as much as possible data or sample for our prediction. If we are not able to collect enough data, it may affect our accuracy of result. The most challenging thing is calculation of global access in the kernel which also uses shared memory. Because not each thread access global memory, we need to apply software method of counting instead of manual counting that helps us to understand more about how threads work .