

Linear Programming

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Lecture 6 Interior Point Method

Outline:

1. Motivation

2. Basic concepts

3. Primal affine scaling algorithm

4. Dual affine scaling algorithm

● Motivation

- **Simplex method** works well in general, but **suffers from exponential-time** computational **complexity**.
- Klee-Minty example shows simplex method may have to **visit every vertex** to reach the optimal one.
- **Total complexity** of an iterative algorithm
= **# of iterations** \times **# of operations in each iteration**

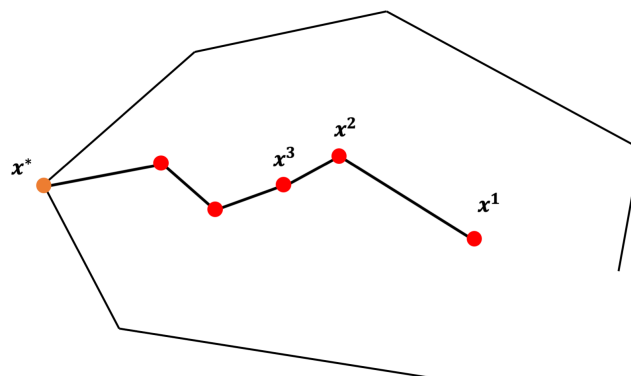
● Simplex method

- Simple operations: only check adjacent extreme points.
- May take many iterations: Klee-Minty example.

Question: **any fix?**

● Karmarkar's (interior point) approach

- Basic idea: **approach optimal solutions from the interior** of the feasible domain.



- Take **more complicated operations** in each iteration to find a better moving direction.
- Require much **fewer iterations**.

● General scheme of an interior point method

-An iterative method that moves in the interior of the feasible domain.

Step 1: Start with an interior solution.

Step 2: If current solution is good **enough**, STOP! Otherwise,

Step 3: Check all directions for improvement and move to a better interior solution.
Go to Step 2.

● Interior movement (iteration)

-Given a current interior feasible solution \mathbf{x}^k , we have

$$A\mathbf{x}^k = \mathbf{b}, \mathbf{x}^k > \mathbf{0}.$$

-An interior movement has a **general** format

$$\mathbf{x}^{k+1} := \mathbf{x}^k + \alpha \mathbf{d}_x^k$$

$$\begin{cases} \alpha \geq 0 : & \text{step-length} \\ \mathbf{d}_x^k \in R^n : & \text{moving direction} \end{cases}$$

● Key knowledge

1. Who is in the interior?

-Initial solution.

2. How do we know a current solution is optimal?

-**Optimality condition.**

3. How to move to a new solution?

-Which direction to move?(**good feasible direction**)

-How far to go?(**step-length**)

● Q1 - Who is in the interior?

Standard form LP

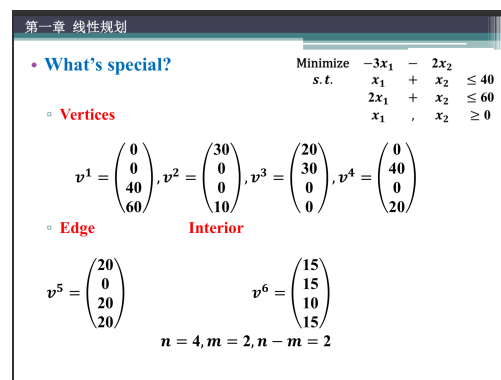
$$\begin{aligned} \min \quad & \mathbf{c}^T \mathbf{x} \\ \text{(LP)} \quad & \text{s.t. } A\mathbf{x} = \mathbf{b} \\ & \mathbf{x} \geq \mathbf{0} \end{aligned}$$

Who is at the vertex?

Who is on the edge?

Who is on the boundary?

Who is in the interior?



Two criteria for a point \mathbf{x} to be an interior feasible solution:

1. $A\mathbf{x} = \mathbf{b}$ (every linear constraint is satisfied)

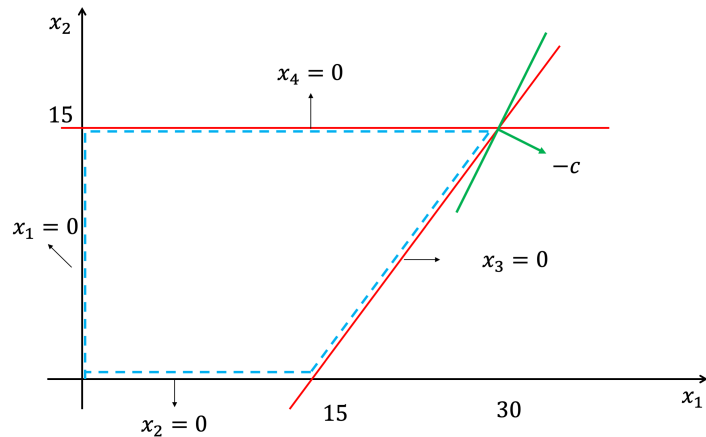
2. $\mathbf{x} > \mathbf{0}$ (every component is positive)

Comments:

1. On a hyperplane $H = \{x \in R^n | a^T x = \beta\}$. (**every point is interior relative to H since H does not have the concept of “boundary”!**)
2. For the first orthant $K = \{x \in R^n | x \geq 0\}$. (**only those $x > 0$ are interior relative to K .**)

Example:

$$\begin{array}{ll} \min & -2x_1 + x_2 \\ \text{s.t.} & x_1 - x_2 \leq 15 \\ & x_2 \leq 15 \\ & x_1, x_2 \geq 0 \end{array}$$



How to find an initial interior solution?

Like the simplex method, we have

-**Big M** method

-**Two-phase** method

(to be discussed later in “Part Two”)

● Q2 - How do we know a current solution is optimal?

Basic concept of optimality:

A current feasible solution is optimal if and only if “no feasible direction at this point is a good direction.”

In other words, **“every feasible direction is not a good direction to move!”**

Feasible direction

In an interior-point method, a **feasible direction** at a current solution is a direction that allows it to take a **small movement** while **staying to be interior feasible**.

Observations:

$$x^{k+1} = x^k + \alpha d_x^k, \quad Ax^k = b, \quad x^k > 0.$$

There is no problem to stay interior if the step-length is small enough.

To maintain feasibility, we need

$$Ax^{k+1} = b \iff Ax^k + \alpha Ad_x^k = b \Rightarrow Ad_x^k = 0,$$

i.e.

$d_x^k \in \mathcal{N}(A)$: null space of A .

Good direction:

In an interior-point method, a **good direction** at a current solution is a direction that leads to a new solution with a **lower objective value**.

Observations:

$$\mathbf{c}^T \mathbf{x}^{k+1} \leq \mathbf{c}^T \mathbf{x}^k \iff \mathbf{c}^T \mathbf{x}^k + \alpha \mathbf{c}^T \mathbf{d}_x^k \leq \mathbf{c}^T \mathbf{x}^k \Rightarrow \mathbf{c}^T \mathbf{d}_x^k \leq 0.$$

Optimality check

Principle: “no feasible direction at this point is a good direction.”

At a current solution, we check that:

$$\text{No } \mathbf{d}_x^k \in \mathbb{R}^n \text{ with } A\mathbf{d}_x^k = \mathbf{0} \text{ can make } \mathbf{c}^T \mathbf{d}_x^k < 0.$$

● Q3 - How to move to a new solution?

1. Which direction to move?

- a good, feasible direction.

“Good” requires $\mathbf{c}^T \mathbf{d}_x^k \leq 0$.

“Feasible” requires $A\mathbf{d}_x^k = \mathbf{0}$, $\mathbf{d}_x^k \in \mathcal{N}(A)$: null space of A .

Question: any suggestion?

A good direction

-Reduce the objective value

$$\mathbf{c}^T \mathbf{d}_x^k \leq 0 \quad \text{Candidate: } \mathbf{d}_x^k = -\mathbf{c} \text{ (negative gradient) (steepest descent)}$$

-Maintain feasibility

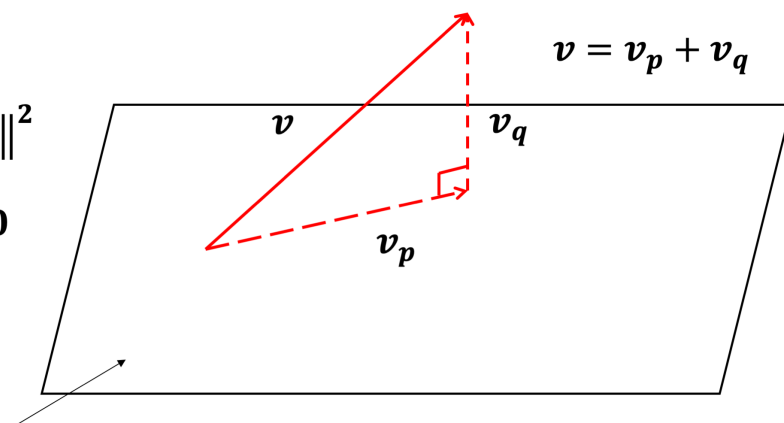
$$A\mathbf{d}_x^k = \mathbf{0} \quad \text{Candidate: projected negative gradient}$$

$$\mathbf{d}_x^k = [\mathbf{I} - A^T(AA^T)^{-1}A](-\mathbf{c})$$

Convex QP

$$\min_{\mathbf{v}_p} \|\mathbf{v}_p - \mathbf{v}\|^2$$

$$\text{s. t. } A\mathbf{v}_p = \mathbf{0}$$



$\mathcal{N}(M)$ = Null space of matrix M

$$\mathcal{N}(M) = \{\mathbf{x} \mid M\mathbf{x} = \mathbf{0}\}$$

$$\mathbf{v}_p = [\mathbf{I} - M^T(MM^T)^{-1}M]\mathbf{v}$$

$$\mathbf{v}_q = M^T(MM^T)^{-1}M\mathbf{v}$$

$$\mathbf{d}_x^k = [\mathbf{I} - \mathbf{A}^T(\mathbf{A}\mathbf{A}^T)^{-1}\mathbf{A}](\mathbf{-c}) = \mathbf{P}_x^k(\mathbf{-c}),$$

$$\mathbf{P}_x^k \triangleq \mathbf{I} - \mathbf{A}^T(\mathbf{A}\mathbf{A}^T)^{-1}\mathbf{A} : \text{null space projection matrix}$$

2. How far to go?

-To **satisfy every linear constraint**.

Since $\mathbf{A}\mathbf{d}_x^k = \mathbf{0}$, $\mathbf{d}_x^k \in \mathcal{N}(\mathbf{A})$: null space of \mathbf{A} , $\mathbf{A}\mathbf{x}^{k+1} = \mathbf{A}\mathbf{x}^k + \alpha\mathbf{A}\mathbf{d}_x^{k+1} = \mathbf{b}$, the step-length can be real number.

-To **stay to be an interior solution**, we need $\mathbf{x}^{k+1} > \mathbf{0}$.

How to choose step-length?

One easy approach: in order to keep $\mathbf{x}^{k+1} = \mathbf{x}^k + \alpha\mathbf{d}_x^k > \mathbf{0}$, we may use the “**minimum ratio test**” to determine the step-length.

Observation:

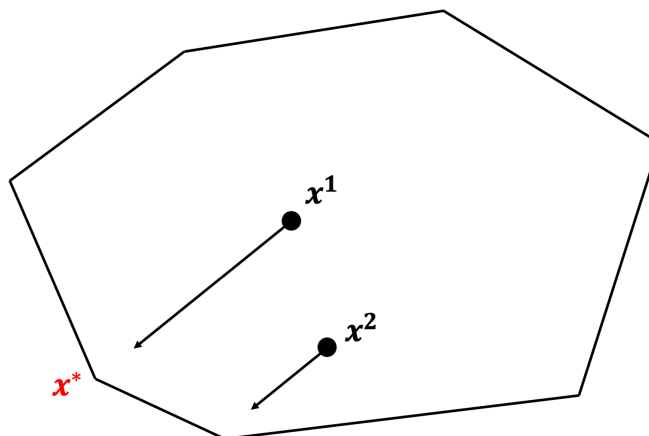
When \mathbf{x}^k is close to the boundary, the step-length may be very small.

Question: **then what?**

Observations

If a current solution is **near the center** of the feasible domain (polyhedral set), it makes sense to move along the **steepest descent direction**.

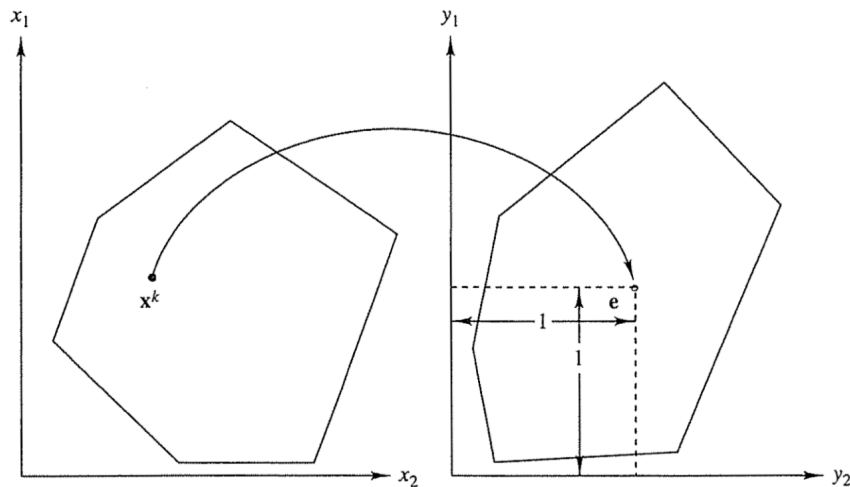
If a current solution is not **near the center**, we need to **re-scale** its coordinates to transform it to become “near the center”.



Question: **but how?**

● **Where is the center?**

We need to know where is the “**center**” of the **non-negative/first orthant**: $\{x \in \mathbb{R}^n \mid x \geq 0\}$.



If $x^k = e = (1, 1, \dots, 1)^T$, then

- (1) x^k is one-unit away from the boundary.
- (2) As long as $\alpha < 1$, $x^{k+1} > 0$. (if $|(d_x^k)_i| \leq 1, \forall i$)

Question: **If not, what to do?**

● **Concept of scaling**

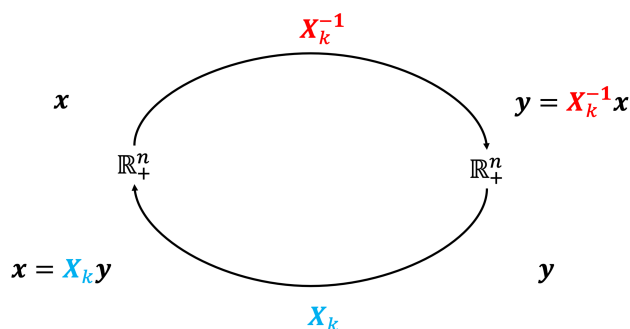
Scale x^k to be e .

Define a diagonal scaling matrix:

$$X_k = \text{diag}(x_i^k) = \begin{pmatrix} x_1^k & 0 & \dots & 0 \\ 0 & x_2^k & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & x_n^k \end{pmatrix}$$

then $X_k^{-1}x^k = e$.

● **Transformation - affine scaling**



The transformation is

1. one-to-one
2. onto
3. invertible
4. boundary to boundary
5. interior to interior

$$\begin{array}{llll}
(7.1 \downarrow) & & & \\
\min \quad \mathbf{c}^T \mathbf{x} & \min \quad \mathbf{c}^T \mathbf{X}_k \mathbf{y} & \min \quad (\mathbf{X}_k^T \mathbf{c})^T \mathbf{y} & \min \quad (\mathbf{X}_k \mathbf{c})^T \mathbf{y} \\
\text{s.t.} \quad \mathbf{A} \mathbf{x} = \mathbf{b} & \iff \text{s.t.} \quad \mathbf{A} \mathbf{X}_k \mathbf{y} = \mathbf{b} & \iff \text{s.t.} \quad (\mathbf{A} \mathbf{X}_k) \mathbf{y} = \mathbf{b} & \iff \text{s.t.} \quad (\mathbf{A} \mathbf{X}_k) \mathbf{y} = \mathbf{b} \\
\mathbf{x} \geq \mathbf{0} & \mathbf{y} \geq \mathbf{0} & \mathbf{y} \geq \mathbf{0} & \mathbf{y} \geq \mathbf{0}
\end{array}$$

$$\mathbf{x}^k > \mathbf{0}$$

$$\mathbf{y}^k = \mathbf{e}$$

$$\mathbf{d}_y^k = [\mathbf{I} - \mathbf{X}_k \mathbf{A}^T (\mathbf{A} \mathbf{X}_k^2 \mathbf{A}^T)^{-1} \mathbf{A} \mathbf{X}_k] (-\mathbf{X}_k \mathbf{c})$$

$$\mathbf{y}^{k+1} = \mathbf{y}^k + \alpha_k \frac{\mathbf{d}_y^k}{\|\mathbf{d}_y^k\|}$$

$$\alpha_k = 0.99 \text{ (say)} \quad 0 < \alpha_k < 1$$

$$\mathbf{x}^{k+1} = \mathbf{X}_k \mathbf{y}^{k+1} = \mathbf{X}_k (\mathbf{y}^k + \alpha_k \frac{\mathbf{d}_y^k}{\|\mathbf{d}_y^k\|})$$

$$= \mathbf{X}_k \mathbf{y}^k + \alpha_k \mathbf{X}_k \frac{\mathbf{d}_y^k}{\|\mathbf{d}_y^k\|}$$

$$= \mathbf{x}^k + \frac{\alpha_k}{\|\mathbf{d}_y^k\|} \mathbf{d}_y^k$$

$$\mathbf{X}_k \mathbf{d}_y^k = \mathbf{d}_x^k = -\mathbf{X}_k [\mathbf{I} - \mathbf{X}_k \mathbf{A}^T (\mathbf{A} \mathbf{X}_k^2 \mathbf{A}^T)^{-1} \mathbf{A} \mathbf{X}_k] (\mathbf{X}_k \mathbf{c})$$

● Step-length in the transformed space

Minimum ratio test in the y-space

In order to make sure that $\mathbf{y}^{k+1} > \mathbf{0}$, we need

$$\mathbf{y}^k + \alpha_k \mathbf{d}_y^k > \mathbf{0}$$

Case 1: $\mathbf{d}_y^k \geq \mathbf{0}$, then $\alpha_k \in (0, +\infty)$. Typically, $\mathbf{d}_y^k = \mathbf{0}$, then the objective value is a constant in its feasible domain.

Case 2: $(\mathbf{d}_y^k)_i < 0$ for some i .

Then

$$\begin{aligned}
(\mathbf{y}^k)_i + \alpha_k (\mathbf{d}_y^k)_i > 0 &\iff (\mathbf{y}^k)_i > -\alpha_k (\mathbf{d}_y^k)_i \iff \alpha_k < \frac{(\mathbf{y}^k)_i}{-(\mathbf{d}_y^k)_i} = \frac{1}{-(\mathbf{d}_y^k)_i} \\
\Rightarrow \alpha_k &= \min_i \left\{ \frac{\alpha}{-(\mathbf{d}_y^k)_i} \mid (\mathbf{d}_y^k)_i < 0 \right\} \text{ for some } \alpha \in (0, 1).
\end{aligned}$$

● **Property 1**

Iteration in the x-space:

$$\begin{aligned}
\mathbf{x}^{k+1} &= \mathbf{X}_k \mathbf{y}^{k+1} = \mathbf{X}_k (\mathbf{e} + \alpha_k \mathbf{d}_y^k) \\
&= \mathbf{x}^k + \alpha_k \mathbf{X}_k \mathbf{d}_y^k \\
&= \mathbf{x}^k + \alpha_k \mathbf{X}_k \mathbf{P}_k (-\mathbf{X}_k \mathbf{c}) \\
&= \mathbf{x}^k + \alpha_k \mathbf{X}_k [\mathbf{I} - \mathbf{X}_k \mathbf{A}^T (\mathbf{A} \mathbf{X}_k^2 \mathbf{A}^T)^{-1} \mathbf{A} \mathbf{X}_k] (-\mathbf{X}_k \mathbf{c}) \\
&= \mathbf{x}^k + \alpha_k [\mathbf{X}_k - \mathbf{X}_k^2 \mathbf{A}^T (\mathbf{A} \mathbf{X}_k^2 \mathbf{A}^T)^{-1} \mathbf{A} \mathbf{X}_k] (-\mathbf{X}_k \mathbf{c}) \\
&= \mathbf{x}^k - \alpha_k [\mathbf{X}_k^2 \mathbf{c} - \mathbf{X}_k^2 \mathbf{A}^T (\mathbf{A} \mathbf{X}_k^2 \mathbf{A}^T)^{-1} \mathbf{A} \mathbf{X}_k^2 \mathbf{c}] \\
&= \mathbf{x}^k - \alpha_k \{ \mathbf{X}_k^2 [\mathbf{c} - \mathbf{A}^T (\mathbf{A} \mathbf{X}_k^2 \mathbf{A}^T)^{-1} \mathbf{A} \mathbf{X}_k^2 \mathbf{c}] \} \\
&= \mathbf{x}^k + \alpha_k \{ -\mathbf{X}_k^2 [\mathbf{c} - \mathbf{A}^T \underbrace{(\mathbf{A} \mathbf{X}_k^2 \mathbf{A}^T)^{-1} \mathbf{A} \mathbf{X}_k^2 \mathbf{c}}_{\mathbf{w}^k}] \} \\
&= \mathbf{x}^k + \alpha_k \{ -\mathbf{X}_k^2 [\mathbf{c} - \mathbf{A}^T \mathbf{w}^k] \} \\
&\quad \underbrace{\hspace{1.5cm}}_{\mathbf{d}_x^k} \\
&= \mathbf{x}^k + \alpha_k \mathbf{d}_x^k
\end{aligned}$$

● **Property 2**

Feasible direction in x-space:

$$\begin{aligned}
\mathbf{x}^{k+1} &= \mathbf{X}_k \mathbf{y}^{k+1} = \mathbf{X}_k (\mathbf{y}^k + \alpha_k \frac{\mathbf{d}_y^k}{\|\mathbf{d}_y^k\|}) \\
&= \mathbf{X}_k \mathbf{y}^k + \alpha_k \mathbf{X}_k \frac{\mathbf{d}_y^k}{\|\mathbf{d}_y^k\|} \\
&= \mathbf{x}^k + \frac{\alpha_k}{\|\mathbf{d}_y^k\|} \mathbf{X}_k \mathbf{d}_y^k \\
&= \mathbf{x}^k + \frac{\alpha_k}{\|\mathbf{d}_y^k\|} \mathbf{d}_x^k
\end{aligned}$$

Since $\mathbf{A} \mathbf{X}_k \mathbf{d}_y^k = \mathbf{0}$, $\mathbf{A} \mathbf{d}_x^k = \mathbf{0}$, i.e., $\mathbf{d}_x^k \in \mathcal{N}(\mathbf{A})$.

● **Property 3**

Good direction in x-space:

$$\begin{aligned}
 \mathbf{c}^T \mathbf{x}^{k+1} &= \mathbf{c}^T (\mathbf{x}^k + \alpha_k \mathbf{X}_k \mathbf{d}_y^k) \\
 &= \mathbf{c}^T \mathbf{x}^k + \alpha_k \mathbf{c}^T \mathbf{X}_k \mathbf{d}_y^k \\
 &= \mathbf{c}^T \mathbf{x}^k + \alpha_k \mathbf{c}^T \mathbf{X}_k [\mathbf{P}_k(-\mathbf{X}_k \mathbf{c})] \\
 &= \mathbf{c}^T \mathbf{x}^k - \alpha_k (-\mathbf{X}_k \mathbf{c})^T [\mathbf{P}_k(-\mathbf{X}_k \mathbf{c})] \\
 &= \mathbf{c}^T \mathbf{x}^k - \alpha_k \langle -\mathbf{X}_k \mathbf{c}, \mathbf{P}_k(-\mathbf{X}_k \mathbf{c}) \rangle \\
 &= \mathbf{c}^T \mathbf{x}^k - \alpha_k \|\mathbf{P}_k(-\mathbf{X}_k \mathbf{c})\|^2 \\
 &= \mathbf{c}^T \mathbf{x}^k - \alpha_k \|\mathbf{d}_y^k\|^2
 \end{aligned}$$

Hence, $\mathbf{c}^T \mathbf{x}^{k+1} \leq \mathbf{c}^T \mathbf{x}^k$, and $\mathbf{c}^T \mathbf{x}^{k+1} < \mathbf{c}^T \mathbf{x}^k$ if $\mathbf{d}_y^k \neq \mathbf{0}$.

Lemma 7.1 If $\exists \mathbf{x}^k \in P$, $\mathbf{x}^k > \mathbf{0}$ with $\mathbf{d}_y^k > \mathbf{0}$, then the standard LP is unbounded below.

● **Property 4**

Optimality check

Lemma 7.2 If there exists an $\mathbf{x}^k \in P^0$ with $\mathbf{d}_y^k = \mathbf{0}$, then every feasible solution of the linear programming problem (7.1) is optimal. (the objective value is a constant)

For $\mathbf{x}^k \in P^0 = \{\mathbf{x} \in R^n \mid \mathbf{A}\mathbf{x} = \mathbf{b}, \mathbf{x} \geq \mathbf{0}\}$, if $\mathbf{d}_y^k = -\mathbf{P}_k \mathbf{X}_k \mathbf{c} = \mathbf{0}$, then $\mathbf{X}_k \mathbf{c}$ falls in the orthogonal space of $\mathcal{N}(\mathbf{A}\mathbf{X}_k)$, i.e., $\mathbf{X}_k \mathbf{c} \in \text{row space of } (\mathbf{A}\mathbf{X}_k)$.

$$\Rightarrow \exists \mathbf{u}^k, \text{ s.t. } (\mathbf{A}\mathbf{X}_k)^T \mathbf{u}^k = \mathbf{X}_k \mathbf{c} \text{ or } (\mathbf{u}^k)^T \mathbf{A}\mathbf{X}_k = \mathbf{c}^T \mathbf{X}_k$$

$$\Rightarrow (\mathbf{u}^k)^T \mathbf{A} = \mathbf{c}^T$$

For any feasible solution \mathbf{x} ,

$$\mathbf{c}^T \mathbf{x} = (\mathbf{u}^k)^T \mathbf{A} \mathbf{x} = (\mathbf{u}^k)^T \mathbf{b}$$

Since $(\mathbf{u}^k)^T \mathbf{b}$ does not depend on \mathbf{x} , the value of $\mathbf{c}^T \mathbf{x}$ remains constant over P .

● **Property 5**

Lemma 7.3 If the linear programming problem (7.1) is bounded below and its objective function is not constant, then the sequence $\{\mathbf{c}^T \mathbf{x}^k \mid k = 1, 2, \dots\}$ is well-defined and strictly decreasing.

● **Property 6**

We may define

$$\begin{aligned}\mathbf{w}^k &\equiv (\mathbf{A} \mathbf{X}_k^2 \mathbf{A}^T)^{-1} \mathbf{A} \mathbf{X}_k^2 \mathbf{c} && \text{(dual estimate)} \\ \mathbf{r}^k &\equiv \mathbf{c} - \mathbf{A}^T \mathbf{w}^k && \text{(reduced cost)}\end{aligned}$$

If $\mathbf{r}^k \geq \mathbf{0}$, then \mathbf{w}^k is dual feasible, and

$$(\mathbf{x}^k)^T \mathbf{r}^k = (\mathbf{X}_k \mathbf{y}^k)^T \mathbf{r}^k = (\mathbf{X}_k \mathbf{e})^T \mathbf{r}^k = \mathbf{e}^T \mathbf{X}_k^T \mathbf{r}^k = \mathbf{e}^T \mathbf{X}_k \mathbf{r}^k$$

becomes the duality gap, i.e., $\mathbf{c}^T \mathbf{x}^k - \mathbf{b}^T \mathbf{w}^k = \mathbf{e}^T \mathbf{X}_k \mathbf{r}^k$.

Therefore, if $\mathbf{r}^k \geq \mathbf{0}$ and $\mathbf{e}^T \mathbf{X}_k \mathbf{r}^k = 0$, then $\mathbf{x}^k \leftarrow \mathbf{x}^*$, $\mathbf{w}^k \leftarrow \mathbf{w}^*$.

● **Property 7 Moving direction and reduced cost**

$$\begin{aligned}\mathbf{d}_y^k &= [\mathbf{I} - \mathbf{X}_k \mathbf{A}^T (\mathbf{A} \mathbf{X}_k^2 \mathbf{A}^T)^{-1} \mathbf{A} \mathbf{X}_k] (-\mathbf{X}_k \mathbf{c}) \\ &= -[\mathbf{X}_k \mathbf{c} - \mathbf{X}_k \mathbf{A}^T (\mathbf{A} \mathbf{X}_k^2 \mathbf{A}^T)^{-1} \mathbf{A} \mathbf{X}_k^2 \mathbf{c}] \\ &= -\mathbf{X}_k [\mathbf{c} - \mathbf{A}^T (\mathbf{A} \mathbf{X}_k^2 \mathbf{A}^T)^{-1} \mathbf{A} \mathbf{X}_k^2 \mathbf{c}] \\ &= -\mathbf{X}_k [\mathbf{c} - \mathbf{A}^T \mathbf{w}^k] \\ &= -\mathbf{X}_k \mathbf{r}^k\end{aligned}$$

● **Key steps of primal affine scaling algorithm**

Step 1 Set $k \leftarrow 0$, $\epsilon > 0$, $0 < \alpha < 1$,
find $\mathbf{x}^0 > \mathbf{0}$ and $\mathbf{A} \mathbf{x}^0 = \mathbf{b}$.

Step 2 Compute

$$\begin{aligned}\mathbf{w}^k &= (\mathbf{A} \mathbf{X}_k^2 \mathbf{A}^T)^{-1} \mathbf{A} \mathbf{X}_k^2 \mathbf{c} \\ \mathbf{r}^k &= \mathbf{c} - \mathbf{A}^T \mathbf{w}^k \\ \text{If } \mathbf{r}^k &\geq \mathbf{0} \text{ and } \mathbf{e}^T \mathbf{X}_k \mathbf{r}^k \leq \epsilon, \\ \text{then STOP! } &\mathbf{x}^k \leftarrow \mathbf{x}^*, \mathbf{w}^k \leftarrow \mathbf{w}^*. \\ \text{Otherwise,}\end{aligned}$$

Step 3 Compute $\mathbf{d}_y^k = -\mathbf{X}_k \mathbf{r}^k$.

If $\mathbf{d}_y^k \geq \mathbf{0}$ (but $\neq \mathbf{0}$), then STOP! Unbounded below.

If $\mathbf{d}_y^k = \mathbf{0}$, then STOP! $\mathbf{x}^k \leftarrow \mathbf{x}^*$ (the objective value is a constant).

Otherwise,

Step 4 Find $\alpha_k = \min_i \left\{ \frac{\alpha}{-(\mathbf{d}_y^k)_i} \mid (\mathbf{d}_y^k)_i < 0 \right\}$.

$$\mathbf{x}^{k+1} = \mathbf{x}^k + \alpha_k \mathbf{X}_k \mathbf{d}_y^k$$

$k := k + 1$.

Go to Step 2.