A fast dynamic programming multi-objective knapsack problem

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Abstract

This work addresses... The Multidi Objective knapsack programming. The dynamic programming method... The data structure...

1 Introduction

2 The Multiobjective Knapsack Problem

A general multiobjective optimization problem can be described as a vector function f that maps a decision variable (solution) to a tuple of m objectives. Formally:

$$\begin{aligned} &\min/\max\ \boldsymbol{y} = f(\boldsymbol{x}) = \big(f_1(\boldsymbol{x}), f_2(\boldsymbol{x}), \ldots, f_m(\boldsymbol{x})\big) \\ &\text{subject to } \boldsymbol{x} \in X \end{aligned}$$

where $\boldsymbol{x}$ is the *decision variable*, X denotes the set of feasible solutions, and $\boldsymbol{y}$ is the *objective vector* where each objective has to be minimized (or maximized).

Considering two decision vectors $\boldsymbol{a}, \boldsymbol{b} \in X$, $\boldsymbol{a}$ is said to *dominate* $\boldsymbol{b}$ if, and only if $\boldsymbol{a}$ is at least as good as $\boldsymbol{b}$ in all objectives and better than $\boldsymbol{b}$ in at least one objective. For shortening we will say that $\boldsymbol{a}$ dominates $\boldsymbol{b}$ by saying $dom(\boldsymbol{a}, \boldsymbol{b})$. Formally:

$$dom(\boldsymbol{a}, \boldsymbol{b}) = \begin{cases} \forall i \in \{1, 2, \ldots, m\} : f_i(\boldsymbol{a}) \geq f_i(\boldsymbol{b}) \text{ and} \\ \exists j \in \{1, 2, \ldots, m\} : f_j(\boldsymbol{a}) > f_j(\boldsymbol{b}) \end{cases}$$

A feasible solution $\boldsymbol{a} \in X$ is called *efficient* if it is not dominated by any other feasible solution. The set of all efficient solutions of a multiobjective optimization problem is known as *Pareto optimal*. Solving a multiobjective problem consists in giving its Pareto optimal set.

An instance of a multiobjective knapsack problem (MOKP) with m objectives consists of an integer capacity $W > 0$ and n items. Each item i has a

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positive weight w_i and nonnegative integer profits $p_i^1, p_i^2, \ldots, p_i^m$. Each profit p_i^k represents the contribution of the i-th item for k-th objective. A solution is represented by a set $\boldsymbol{x} \subseteq \{1, \ldots, n\}$ containing the indexes of the items included in the solution. A solution is feasible if the total weight included in the knapsack does not exceed its capacity. Formally the definition of the problem is:

$$\begin{aligned} \max\ & f(\boldsymbol{x}) = \big(f_1(\boldsymbol{x}), f_2(\boldsymbol{x}), \ldots, f_m(\boldsymbol{x})\big) \\ \text{subject to}\ & w(\boldsymbol{x}) < W \\ & \boldsymbol{x} \subseteq \{1, \ldots, n\} \end{aligned}$$

where

$$f_j(\boldsymbol{x}) = \sum_{i \in \boldsymbol{x}} p_i^j$$

$$w(\boldsymbol{x}) = \sum_{i \in \boldsymbol{x}} w_i$$

The MOKP is considered a $\mathcal{NP}$-Hard problem since it is a generalization of the well-known $0-1$ knapsack problem, in which $m = 1$. It is quite difficult to determine the Pareto optimal set for the MOKP, especially for high dimension instances, in which the Pareto optimal set tends to grow exponentially. Even for the bi-objective case, small problems may prove intractable. For this reason we are interested in developing efficient methods for handling large solution sets, which may bring tractability to previously intractable instances.

Considering two solutions $\boldsymbol{x}, \boldsymbol{y} \subseteq \{1, \ldots, n\}$, $\boldsymbol{y}$ is called *extension* of $\boldsymbol{x}$, denoted as $ext(\boldsymbol{y}, \boldsymbol{x})$, iff $\boldsymbol{x} \subseteq \boldsymbol{y}$. Any set $\boldsymbol{e} \subseteq \{1, \ldots, n\}$ such that $\boldsymbol{x} \cap \boldsymbol{e} = \emptyset$ is called an *extender* of $\boldsymbol{x}$. If $w(\boldsymbol{x}) + w(\boldsymbol{e}) \leq W$ than $\boldsymbol{e}$ is called a *feasible extender* of $\boldsymbol{x}$. A solution $\boldsymbol{x}$ is called *deficient* if it has available space to fit one or more item, i.e., $w(\boldsymbol{x}) + min\{w_i : i \notin \boldsymbol{x}\} \leq W$. We say $\boldsymbol{x}$ *knapsack-dominates* $\boldsymbol{y}$, denoted as $dom_k(\boldsymbol{x}, \boldsymbol{y})$, if $\boldsymbol{x}$ dominates $\boldsymbol{y}$ and does not weight more than $\boldsymbol{y}$. Formally:

$$dom_k(\boldsymbol{x}, \boldsymbol{y}) = \begin{cases} dom(\boldsymbol{x}, \boldsymbol{y}) & \text{and} \\ w(\boldsymbol{x}) \leq w(\boldsymbol{y}) \end{cases}$$

Figure 1 illustrates the concept of kanpsack-dominance for a problem with $m = 1$. Any solution in the cross-hatched area knapsack-dominates the marked solution.

Theorem 1. *Consider two solutions $\boldsymbol{x}, \boldsymbol{y} \subseteq \{1, \ldots, n\}$ and $\boldsymbol{e}$ an extender for $\boldsymbol{x}$ and $\boldsymbol{y}$. If $dom_k(\boldsymbol{x}, \boldsymbol{y})$ than $dom_k(\boldsymbol{x} \cup \boldsymbol{e}, \boldsymbol{y} \cup \boldsymbol{e})$.*

Proof. □

3 The Dynamic Programming Algorithm

Paragrafo de introducao da secao, justificando toda a explicacao que segue..

The dynamic programming algorithm is based on the classical Nemhauser-Ullmann (NU) algorithm proposed in [6]. concept of domination for knapsack problems, proposed by Weingartner and Ness [10].

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(6,1)

Some property of the MOKP will be explored... three filter as optimizations for the trying to reduce the number of solutions handled on mid stages of the algorithm... We will introduce the dynamic programming (DP) algorithm for the MOKP proposed in [1].

3.1 Item ordering

An important issue in the MOKP is the ordering of items. It is well-known that good solutions of knapsack problems are generally composed of items with good profit-weight relation. Therefore, the prioritization of those items tends to lead to better solutions.

For the multi-objective case however, there is no such natural measure. The method proposed by [1] will be considered in this work, in which is based on the ranking of items from their cost-benefit measures in the various objectives. Those measures will give us several items orders which will be useful during the development of the algorithm.

We denote $\mathcal{O}^j$ the set of items ordered by ascending order of cost-benefit function regarding objective j. Formally, considering the function $cb^j(a) = p_a^j/w_a$ the cost-benefit function of item a regarding objective j, order $\mathcal{O}^j$ can be defined by:

$$\mathcal{O}^j = (o_1^j, \ldots, o_n^j), \qquad cb^j(o_1^j) \leq cb^j(o_2^j) \leq \cdots \leq cb^j(o_n^j)$$

Let r_i^j be the rank of item i in order $\mathcal{O}^j$. $\mathcal{O}^{sum}$, $\mathcal{O}^{min}$ and $\mathcal{O}^{max}$ denotes an order according to inscreasing values of the sum, minimum and maximum ranks. Formally:

$$\begin{aligned} r_i^j &= \max\{k \mid o_k^j = i\} \\ r_i^{sum} &= \textstyle\sum_{j=1}^{m} r_i^j \\ \mathcal{O}^{sum} &= (o_1, \ldots, o_n), \qquad r_{o_1}^{sum} \leq \cdots \leq r_{o_n}^{sum} \end{aligned}$$

The orders $\mathcal{O}^{min}$ and $\mathcal{O}^{max}$ are conceived in the same manner as $\mathcal{O}^{sum}$, except for the fact that $\frac{r_i^{sum}}{m}$ is added up in their ranks as tie breaking criteria. The notation $\mathcal{O}(\boldsymbol{s})$ will be used to denote the respective ordering of a $\boldsymbol{s}$ restrict set of items.

3.2 Nemhauser-Ullmann algorithm – the knapsack dominance

The algorithm considered in this work is based on the Nemhauser-Ullmann algorithm, a dynamic programming method for knapsack problems, which is presented by Algorithm 1

At k-th stage the algorithm receives the set S^{k-1} of solutions and generates the set S^k of solutions that correspond to subsets containing exclusively the first k items, i.e., $\forall \boldsymbol{x} \in S^k, \boldsymbol{x} \subseteq \{1, \ldots, k\}$.

Algorithm 1 Basic dynamic programming algorithm for MOKP

1: **function** DP($\boldsymbol{p}, \boldsymbol{w}, W$)
2: $S^0 = \{\emptyset\}$
3: **for** $k \leftarrow 1, n$ **do**
4: $S^k_* = S^{k-1} \cup \{\boldsymbol{x} \cup k \mid \boldsymbol{x} \in S^{k-1}\}$ ▷ solutions extension
5: $S^k = \{\boldsymbol{x} \mid \nexists \boldsymbol{a} \in S^k_* : dom_k(\boldsymbol{a}, \boldsymbol{x})\}$ ▷ partial dominance filter
6: **end for**
7: $P = \{\boldsymbol{x} \mid \nexists \boldsymbol{a} \in S^n : dom(\boldsymbol{a}, \boldsymbol{x})\}$ ▷ dominance filter
8: **return** P
9: **end function**

This is done by expanding S^{k-1} by adding a copy of each solution with the inclusion of k-th item (line 4). We will refer as *partial solutions* all the solutions handled by stages prior to n-th stage.

The clever part of it is that it uses the concept of knapsack dominance to filter solutions that will not lead to efficient solutions (line 5). Considering two partial solutions $\boldsymbol{x}, \boldsymbol{y} \in S^k$, if $\boldsymbol{x}$ is knapsack-dominated by $\boldsymbol{y}$ then we may discard $\boldsymbol{x}$ since all solutions generated from $\boldsymbol{x}$ will be dominated by those generated from $\boldsymbol{y}$.

3.3 Avoiding deficient solutions

The first optimization that can be made on Algorithm 1 is avoiding the generation of deficient solutions. At k-th stage all previous solution is copied to the new solution set without adding k-th item (line 4). However preserving solutions with a lot of space left, concerning the remaining itens, may lead to deficient solutions.

Theorem 2. *Considering the k-th stage of the algorithm and $\boldsymbol{x} \in S^{k-1}$. If $w(\boldsymbol{x}) + \sum_{i \in \{k,\dots,n\}} w_i \leq W$ than*

Considering the k-th stage, if a partial solution $\boldsymbol{x} \in S^{k-1}$ has enough space to fit all remaining items, i.e., $w(\boldsymbol{x}) + \sum_{i=k}^{n} w_i \leq W$, $\boldsymbol{x}$ may be discarded and only $\boldsymbol{x} \cup \{k\}$ keeped, once keeping $\boldsymbol{x}$ will certainly lead to deficient solutions.

3.4 Removing unpromissing solutions

Another optimization that can be applied on later stages is filtering unpromissing solutions by computing upper bounds for its objectives functions and comparing it with the set of available lower bounds. Considering a given k-th iteration, an upper(lower) bound of a partial solution is an upper(lower) limit each objective value can achieve, given its remaining capacity and the remaining items $(k+1, \dots, n)$. A solution can be discharged if its upper bound is dominated by an existing lower bound, since it will generate no efficient solution.

A lower bound of a solution can be computed by greedily filling the knapsack with respect to $\mathcal{O}^{max}$ which is a good quality order of items. Formally:

$$lb(\boldsymbol{x}, \boldsymbol{s}) = \boldsymbol{x} \cup \left\{ o_i \;\middle|\; w(\boldsymbol{x}) + \textstyle\sum_{j=1}^{i} w_{o_j} \leq W \right\}$$

where

$$(o_1, \ldots, o_k) = \mathcal{O}^{max}(\boldsymbol{s})$$

The upper-bound of a partial solution $\boldsymbol{x}$ is computed considering its available capacity and the remaining items on the current algorithm stage. To ensure the upper-bound its computation is done separately for. For the each j-th objective the reversed order $\mathcal{O}^j(\boldsymbol{s})$ on the set $\boldsymbol{s}$ of remaining items is considered to iteratively fill the remaining capacity of $\boldsymbol{x}$.

Algorithm 2 Upper-bound computation for a partial solution.

function $\mathrm{UB}_j(\boldsymbol{p}, \boldsymbol{w}, W, \boldsymbol{x}, \boldsymbol{s})$
  $w_{left} \leftarrow W - w(\boldsymbol{x})$
  $(o_1^j, \ldots, o_k^j) \leftarrow \mathcal{O}^j(\boldsymbol{s})$
  $u_j \leftarrow f_j(\boldsymbol{x})$
  $i \leftarrow 1$
  $l \leftarrow o_i^j$
  while $w_{left} \geq w_l$ **and** $i \leq k$ **do**
    $u_j \leftarrow u_j + p_l^j$
    $w_{left} \leftarrow w_{left} - w_l$
  end while
  if $i \leq k$ **then**
    $l \leftarrow o_i^j$
    $u_j \leftarrow u_j + \frac{w_{left}}{w_l}.p_l^j$
  end if
  return u_j
end function

Algorithm 3 Bazgan's DP algorithm for the MOKP

function BazDP$(\boldsymbol{p}, \boldsymbol{w}, W)$
  $S^0 = \{\emptyset\}$
  for $k \leftarrow 1, n$ **do**
    $S_*^k = \{\boldsymbol{x} \cup \{k\} \mid \boldsymbol{x} \in S^{k-1} \wedge (w(\boldsymbol{x}) + w_k \leq W)\}$
      $\cup \{\boldsymbol{x} \mid \boldsymbol{x} \in S^{k-1} \wedge (w(\boldsymbol{x}) + w_k + \ldots + w_n) > W\}$
    $S^k = \{\boldsymbol{x} \in S_*^k \mid (\nexists \boldsymbol{a} \in S_*^k)[dom_k(\boldsymbol{a}, \boldsymbol{x}) \vee dom(lb(\boldsymbol{a}), up(\boldsymbol{x}))]\}$
  end for
  return S^n
end function

4 The use of k-d tree

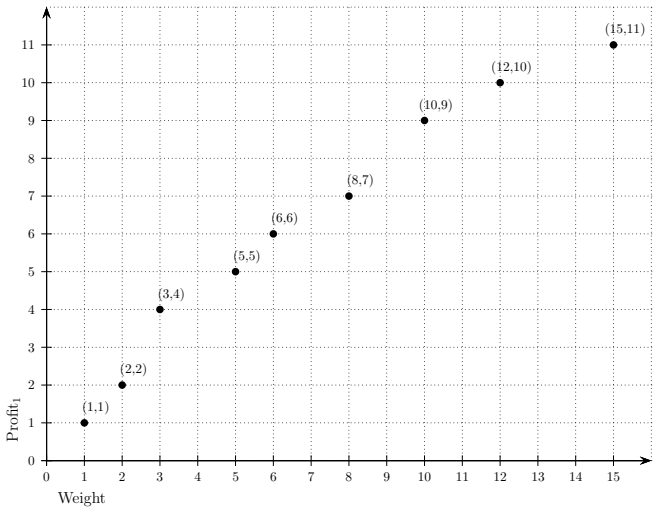
The k-d tree is a type of binary search tree for indexing multidimensional data with simple construction and low space usage. Despite its simplicity it efficiently supports operations like nearest neighbour search and range search [2]. For those reasons k-d tree is widely used on spacial geometry algorithms [8, 3], clustering [5, 4] and graphic rendering algorithms [7].

Like a standard binary search tree, the k-d tree subdivides data at each recursive level of the tree. Unlike a standard binary tree, that users only one key for all levels of the tree, the k-d tree uses k keys and cycles through these keys for successive levels of the tree.

Concerning it's efficiency, it is important to consider the number of dimensions k-d tree is indexing. As a general rule, a k-d tree is suitable for efficiently indexing of n elements if n is much greater than 2^k. Otherwise, when k-d tree are used with high-dimensional data, most of the elements in the tree will be evaluated and the efficiency is no better than exhaustive search [9].

Indexing the solutions and range operations.

Tends to increase the feasibility on problems with higher dimensions.



5 Computational experiments

- Base de dados utilizaca
- Parametros dos algoritmos
- Anlise dos resultados (comparao)

6 Conclusions and future remarks

- Concluses dos resultados
- Trabalhos futuros

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