

Analysis, Anti-Analysis, Anti-Anti-Analysis: An Overview of the Evasive Malware Scenario

Marcus Botacin¹, Vitor Falcão da Rocha⁽¹⁾, André Grégio^{1,2},
Paulo Lício de Geus¹

¹Institute of Computing - UNICAMP
{marcus,vitorf,paulo}@lasca.ic.unicamp.br

²Federal University of Paraná (UFPR)
gregio@inf.ufpr.br

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Topics

- 1 Part I
 - Analysis and Anti-analysis
- 2 Part II
 - Tricks and detection methods
- 3 Part III
 - Tests and Results
- 4 Part IV
 - Concluding Remarks

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Arms-Race

Researchers Reverse Engineer Latest CryptoBit Ransomware to Decrypt Files

By GoldSparrow in [Computer Security](#)

User Rating: ★★★★★ (1 votes, average: 5.00 out of 5)



Figura: Enigma: <https://tinyurl.com/kydgvve>

PetrWrap Crypto Ransomware Blocks Security Researchers From Reverse Engineering Code Samples

JP Buntinx March 16, 2017 [News, Security](#)

Figura: Themerkle: <https://tinyurl.com/kasuxcr>

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APIs and direct calls.

Attackers I

Listagem 1: Attackers.

```
1 do_malicious();
```

Analysts I

Listagem 2: Analysts.

```
1 if do_malicious is present;  
2   detect()
```

APIs and direct calls.

Attackers II

Listagem 3: Attackers.

```
1  if !is_debug():  
2      do_malicious();
```

Analysts II

Listagem 4: Analysts.

```
1  if is_debug() is present;  
2      detect();
```

APIs and direct calls.

Attackers III

Listagem 5: Attackers.

```
1 typedef struct _PEB {  
2 BYTE Reserved1 [2];  
3 BYTE BeingDebugged;  
4 BYTE Reserved2 [1];
```

Attackers III

Listagem 6: Analysts.

```
1 mov eax, [fs:0x30]  
2 mov     eax, [eax+0x0c]
```


APIs and direct calls.

Analysts III

Listagem 7: Analysts.

```
1 def check():
2     if instruction in ['mov', 'movsx', 'movzx']:
3         if 'fs:0x30' in op2:
4             self.found_op1 = op1
5             self.found_keyword = True
6         if self.found_keyword:
7             if instruction in ['cmp', 'cmpxchg', '
8                 mov', ...]:
9                 if '[' + self.found_op1 + '+0xc]' or
10                    ...:
11                     print "Detected!"
```

Anti-Disassembly.

Technique	Description	Detection
PUSH POP MATH	PUSH and POP a value on/from the stack instead of using a direct MOV	Detect a sequence of PUSH and POP on/from a register.
PUSH RET	PUSH a value on the stack and RET to it instead of the ordinary return.	Detect a sequence of PUSH and RET
LDR address resolving	Get loaded library directly from the PEB instead of using a function call	Check memory access referring the PEB offset.
Stealth API import	Manually resolving library imports instead of directly importing them.	Check for a sequence of access/compares of PEBs offsets.
NOP sequence	Breaks pattern matching by implanting NO-Operations	Detect a sequence of NOPs within a given window
Fake Conditional	Create an always-taken branch	Check for branch-succeeded instructions which set branch flags
Control Flow	Changing control flow within an instruction block	Check for the PUSH-RET instruction sequence
Garbage Bytes	Hide data as instruction code	Check for branch-preceded data

Anti-Debug.

Technique	Description	Detection
Known Debug API	Call a debug-check API	Check for API imports
Debugger Fingerprint	Check the presence of known debugger strings	Check known strings inside the binary
NtGlobalFlag	Check for flags inside the PEB structure	Check for access on the PEB offset
IsDebuggerPresent	Check the debugger flag on the PEB structure	Check access to PEB on the debugger flag offset
Hook Detection	Verify whether a function entry point is a JMP instruction	Check for a CMP instruction having JMP opcode as an argument
Heap Flags	Check for heap flags on the PEB	check for heap checks involving PEB offsets
Hardware Breakpoint	Check whether hardware breakpoint registers are not empty	Check for access involving the debugger context
SS Register	Insert a check when interruptions are disabled	Check for SS register's POPs
Software Breakpoint	Check for the INT3 instruction	Check for CMP with INT3
SizeOfImage	Change code image field	Check for PEB changes.

Anti-VM.

Technique	Description	Detection
VM Fingerprint	Check for known strings, such as serial numbers	Check for known strings inside the binary
CPUID Check	Check CPU vendor	Check for known CPU vendor strings
Invalid Opcodes	Launch hypervisor-specific instructions	Check for specific instructions on the binary
System Table Checks	Compare IDT values	Look for checks involving IDT
HyperCall Detection	Platform specific feature	Look for specific instructions

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Sections

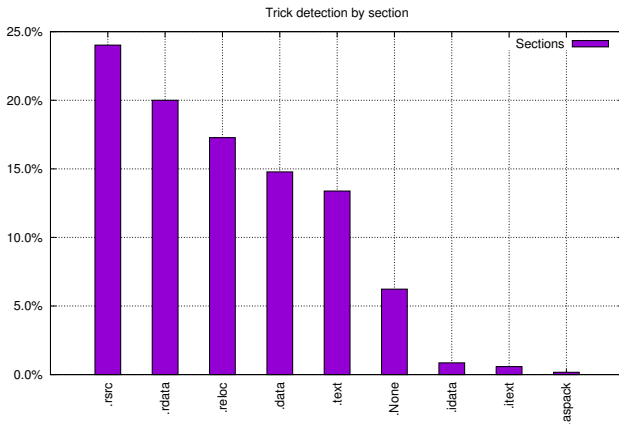


Figura: Tricks by section.

Sections

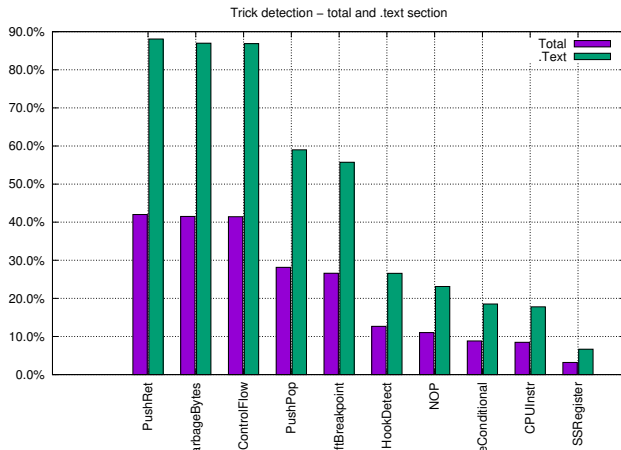


Figura: Tricks - total and .text section.

Packers

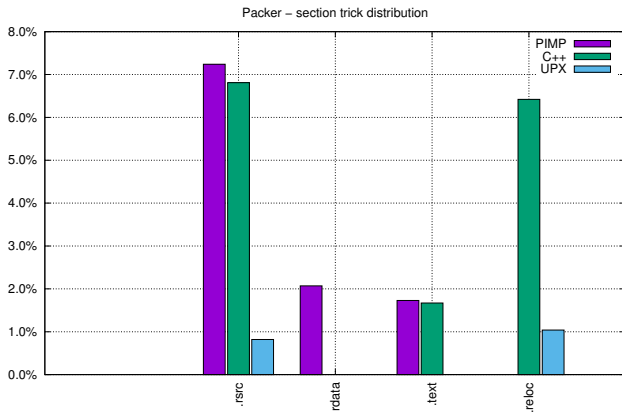


Figura: Packer distribution across binary sections.

Packers

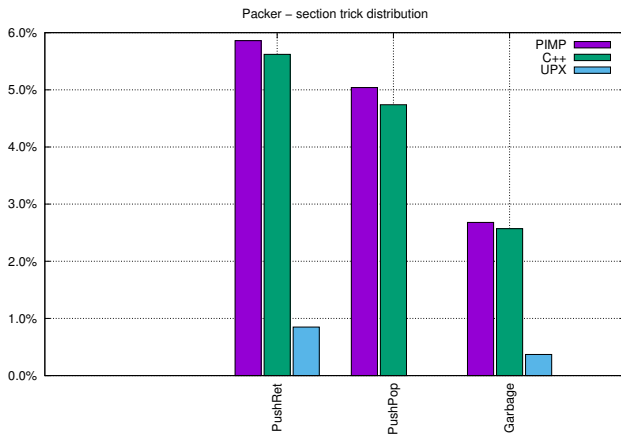


Figura: Tricks detected on distinct packers.

Packers

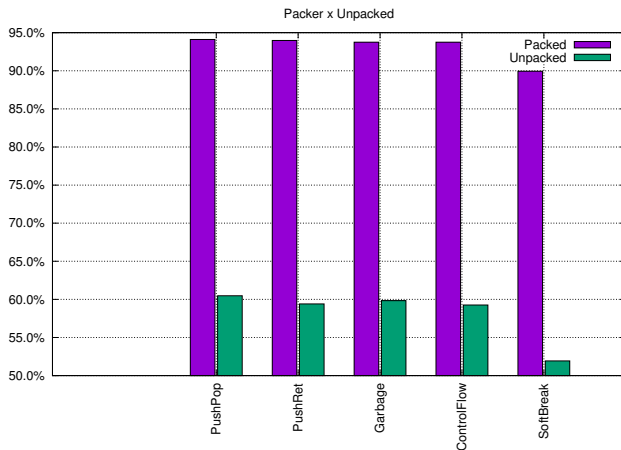


Figura: Packer influence on trick detection.

Malware and Goodware

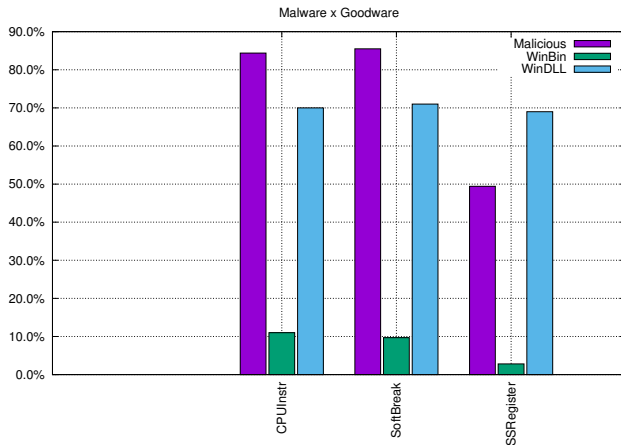


Figura: Tricks detection on malware and goodware.

Comparing Scenarios

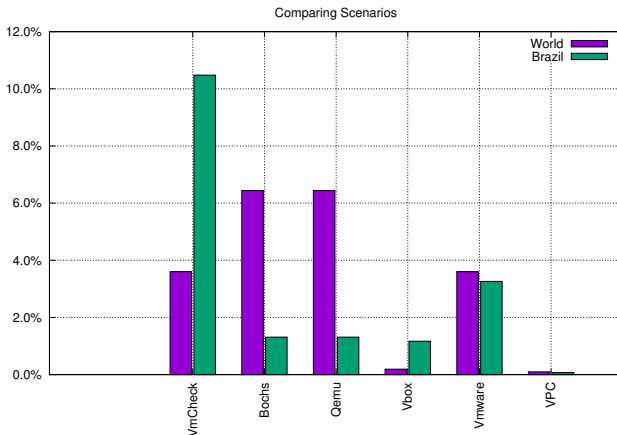


Figura: Comparing scenarios: PEframe detection.

Comparing Scenarios

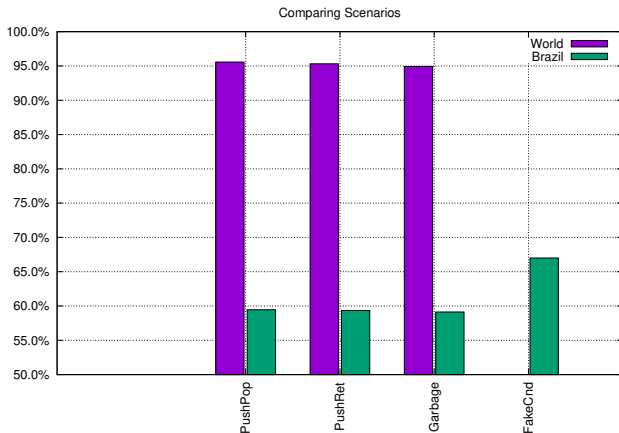


Figura: Comparing scenarios: Tricks detection.

Comparing Scenarios

More about the brazilian scenario

- Uma Visão Geral do Malware Ativo no Espaço Nacional da Internet entre 2012 e 2015 - SBSEG 2015
Marcus Botacin, André Grégio, Paulo Lício de Geus
(<http://siaiap34.univali.br/sbseg2015/anais/WFC/artigoWFC02.pdf>)

Trick Blocking

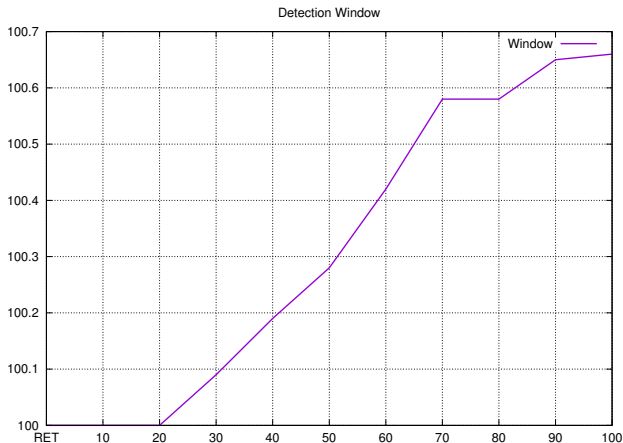


Figura: Evaluating block window effect on trick detection.

Trick Alignment

Tabela: Evaluating the occurrence of misaligned tricks.

Trick	Aligned	Unaligned
CPU	182	287
FakeJMP	63	203

Compiler-based evasion

Tabela: Compilation-based evasion.

ShellCode	Unarmored	ROPinjector
1 ¹	4/57	0/57
2 ²	15/58	0/57
3 ³	9/57	0/54
4 ⁴	7/58	0/54
5 ⁵	9/53	0/53

¹<http://shell-storm.org/shellcode/files/shellcode-898.php>

²<http://shell-storm.org/shellcode/files/shellcode-874.php>

³<http://shell-storm.org/shellcode/files/shellcode-627.php>

⁴<http://shell-storm.org/shellcode/files/shellcode-568.php>

⁵<http://shell-storm.org/shellcode/files/shellcode-714.php>

AV Detection

Shellcode	SC1		SC2		SC3	
Technique	W/o Trick	W/ Trick	W/o Trick	W/ Trick	W/o Trick	W/ Trick
Fakejmp	10/58	6/57	20/58	17/58	15/58	10/57
PushRet		7/57		17/58		10/58
NOP		6/57		17/57		10/58

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Summary

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- Arms race.
- Static detection can help detection.
- Tricks can be improved.
- Static analysis is theoretically limited.

Concluding Remarks

Handling anti-analysis dynamically

- Análise Transparente de Malware com Suporte por Hardware - SBSEG 2016

Marcus Botacin, Paulo Lício de Geus, André Grégio
(<http://sbseg2016.ic.uff.br/pt/files/anais/completos/ST8-3.pdf>)

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Contato:

marcus@lasca.ic.unicamp.br
vitorf@lasca.ic.unicamp.br
paulo@lasca.ic.unicamp.br
gregio@inf.ufpr.br

