

# Near-memory & In-Memory Detection of Fileless Malware

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# Agenda

- 1 Introduction
- 2 Proposed Solution
- 3 Evaluation
- 4 Conclusions

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# What Is Fileless Malware?

LILY HAY NEWMAN 02.09.17 07:22 PM

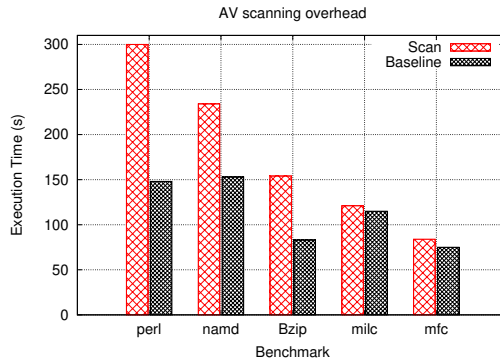
## Say Hello to the Super-Stealthy Malware That's Going Mainstream

Figure: Source: <https://www.wired.com/2017/02/say-hello-super-stealthy-malware-thats-going-mainstream/>



Figure: Source: <https://www.cyberscoop.com/kaspersky-fileless-malware-memory-attribution-detection/>

# The Cost of Scanning Memory



**Figure:** In-memory AV scans worst-case and best-case performance penalties.

# A Drawback for Current Security Solutions

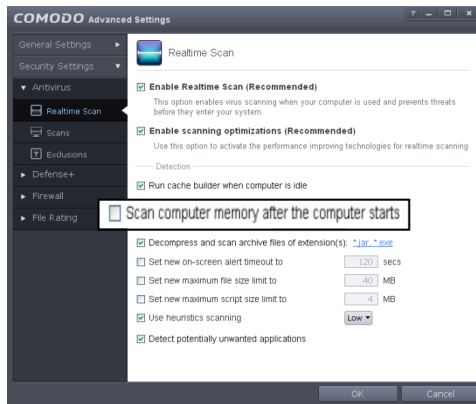
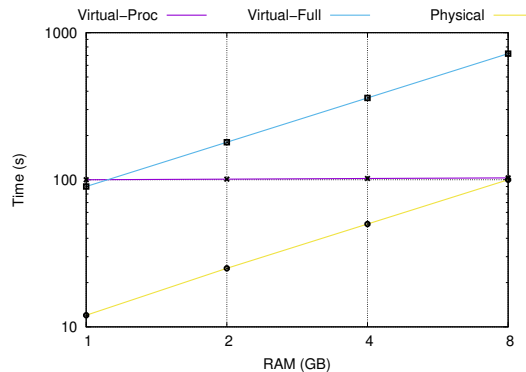


Figure: Default policy is not to scan memory.

# Why Not New AV implementations?



**Figure:** Memory dump time for distinct software-based techniques and memory sizes.

# Understanding Malware Detection Tasks

## Monitoring

- **You need to know:** When to inspect.

## Classifying

- **You need to know:** What to inspect.



## Can't We Rely on Page Faults?

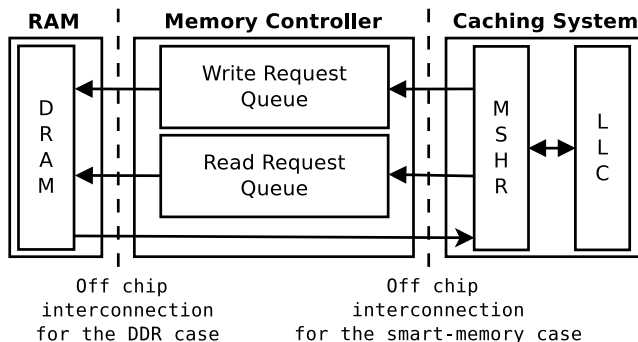
**Table: Blocking on Page Faults.** The performance impact is greater as more complex is the applied detection routine.

Benchmark	Cycles	PF	5K	10K	20K	30K
perf	187G	1,8M	4,74%	9,48%	18,96%	28,44%
mcf	69G	375K	2,72%	5,45%	10,89%	16,34%
milc	556G	1,2M	1,05%	2,10%	4,21%	6,31%
bzip	244G	170K	0,35%	0,69%	1,38%	2,08%
namd	491G	325K	0,33%	0,66%	1,32%	1,98%

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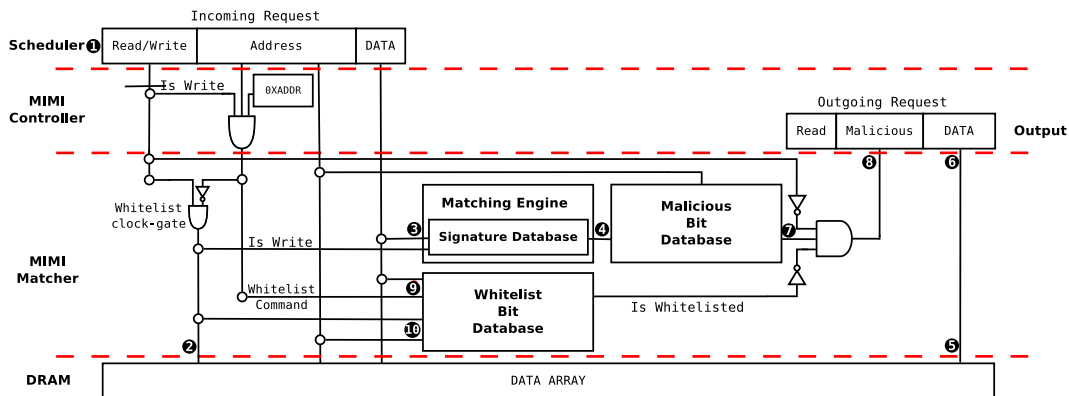
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## Observing Memory Accesses Patterns



**Figure: Write-to-Read window.** Read requests originated from the MSHR might overlap other memory-buffered read requests for any address, but must not overlap previous memory-buffered write requests for the same address.

# Malware Identification based on Near- and In-Memory Evaluation (MINI-ME)



**Figure: MINI-ME Architecture.** MINI-ME is implemented within the memory controller.

# Handling Notifications via Page Faults

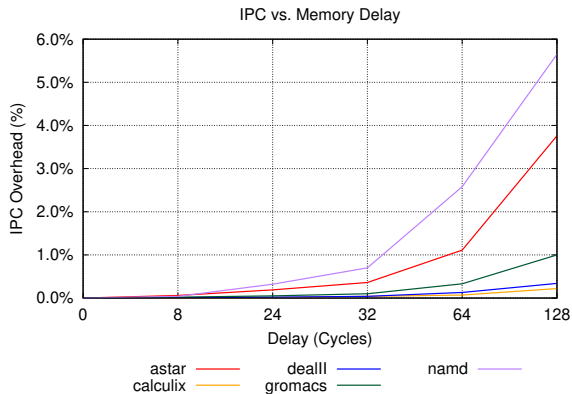
```
1 void __do_page_fault(...) {  
2     // Original Code  
3     if (X86_PF_WRITE) ...  
4     if (X86_PF_INSTR) ...  
5     // Added Code  
6     if (X86_MALICIOUS) ...
```

**Code 1: Modified PF handler.** Malicious bit is set when suspicious pages are mapped.

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# How Much Performance Overhead is Acceptable?



**Figure: MINI-ME database overhead.** Delays of up 32 cycles impose less than 1% of IPC overhead.

# Signatures as the Detection Mechanism

```
1  if(IsDebuggerPresent()){  
2      evade()
```

Code 2: C code

```
1  mov  eax, [fs:0x30]  
2  mov  eax, [eax+0x2]  
3  jne  0 <evade>
```

Code 3: ASM code

```
1  64 8b 04 25 30 00 00  
2  67 8b 40 02  
3  75 e1
```

Code 4: Instructions Bytes



## Signature Size Definition

**Table: Signature Generation.** Signatures (%) detected as false positives for each signature size and memory dump size.

		Memory Size			
		1 GB	2 GB	4 GB	8 GB
Signature Size	8 B	8.65%	9.92%	10.18%	11.45%
	16 B	3.06%	3.32%	3.32%	3.32%
	32 B	0.00%	0.00%	0.00%	0.00%
	64 B	0.00%	0.00%	0.00%	0.00%

# Matching Mechanism Definition

**Table: Matching Techniques.** FP rates for multiple signature sizes and techniques.

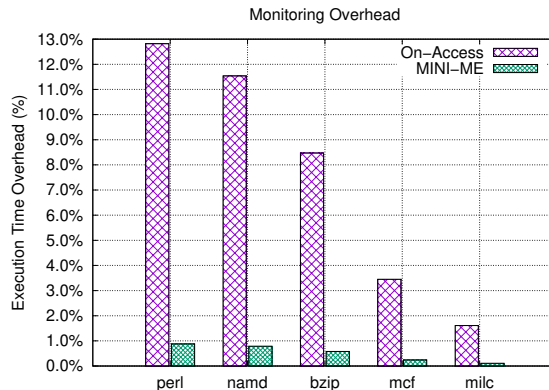
		Signature size			
		8 B	16 B	32 B	64 B
Match. Tech.	Dir. Mapped Table	8.33%	3.15%	0.00%	0.00%
	Signature Tree	8.33%	3.15%	0.00%	0.00%
	Bloom Filter	8.41%	3.47%	0.00%	0.00%

## Matching Policies Definition

**Table: Scan Policies.** FP rate for multiple signature sizes and policies.

		Signature size			
		8 B	16 B	32 B	64 B
Scan Policy	Whole Memory	8.33%	3.15%	0.00%	0.00%
	Mapped Pages	0.06%	0.01%	0.00%	0.00%
	Whitelist	0.00%	0.00%	0.00%	0.00%
	Code-Only	0.01%	0.00%	0.00%	0.00%

# MINI-ME in Practice



**Figure: Monitoring Overhead.** MINI-ME imposes a smaller overhead while still checking more pages than an on-access solution.

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# Conclusions

## Challenges & Lessons

- Fileless malware is a growing hard-to-detect class of threats.
- Traditional AntiViruses (AVs) impose significant performance overhead to perform memory scans.
- In-memory and Near-memory AVs helps reducing AV's performance overheads.
- The more complex the matching mechanism, the greater the performance overhead.
- MINI-ME as platform for future developments.

## Questions & Comments.

### Contact

- **mfbotacin@inf.ufpr.br**
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### Additional Material

- <https://github.com/marcusbotacin/In.Memory>

