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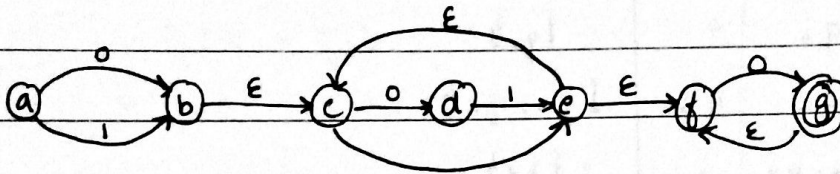
ROLL NO: 8669

CLASS: SEIT

AT TUTORIAL NO: 2.

Q.1] Convert $(0+1)(01)^*(0+E)$ into NFA with E-moves and obtain DFA.

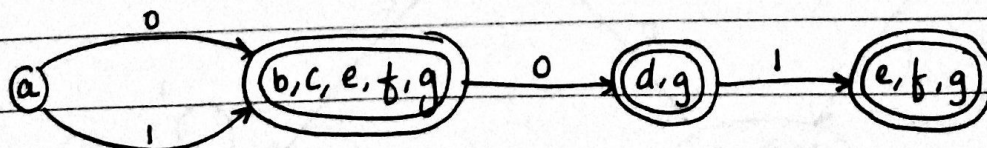
Ans:



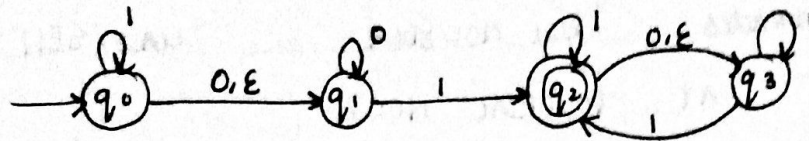
states	0	1	E
→ a	b	b	-
b	-	-	c
c	d	-	c
d	-	e	-
e	-	-	c, f
f	g	-	g
* g			

States for DFA:

states	0	1
a	b, c, e, f, g	b, c, e, f, g
b, c, e, f, g	d, g	-
d, g	-	e, f, g



Q.2] Construct DFA for the following NFA with ϵ -moves.

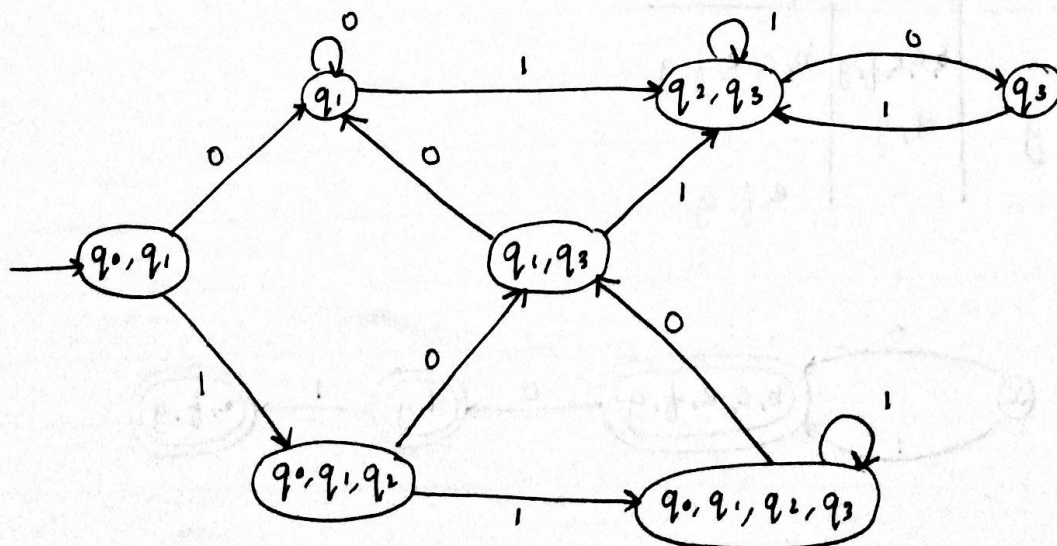


Ans:

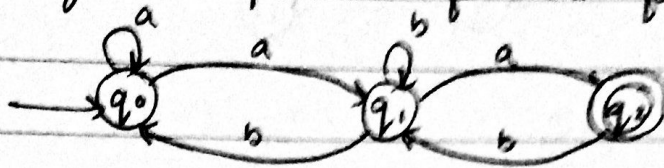
States	0	1	ϵ	ϵ -closure
q_0	q_1	q_0	q_1	$\{q_0, q_1\}$
q_1	q_1	q_2	-	$\{q_1\}$
q_2	q_3	q_2	q_3	$\{q_2, q_3\}$
q_3	-	q_2, q_3	-	$\{q_3\}$

Corresponding DFA:

States	0	1
$\rightarrow q_0, q_1$	q_1	q_0, q_1, q_2
q_1	q_1	q_2, q_3
q_0, q_1, q_2	q_1, q_3	q_0, q_1, q_2, q_3
q_2, q_3	q_3	q_2, q_3
q_1, q_3	q_1	q_2, q_3
q_0, q_1, q_2, q_3	q_1, q_3	q_0, q_1, q_2, q_3
q_3	-	q_2, q_3



Q.3] Construct regular expression for the following FA



Ans: $q_2 = q_1 a$ ——— (1)

$$q_1 = q_{0a} + q_{1b} + q_{2b} \text{ --- (2)}$$

$$q_0 = q_{0a} + q_{1b} + \epsilon \quad \text{--- (3)}$$

$$q_1 = q_{0a} + q_{1b} + q_{1ab} \quad [\text{from } ① \text{ \& } ②]$$

$$q_1 = q_0 a + q_1 (b + ab)$$

$$\therefore q_1 = q_0 a (b + ab)^*$$

[By arden's theorem]

$$q_0 = \varepsilon + q_0 a + (q_0 a (b + ab)^*) b$$

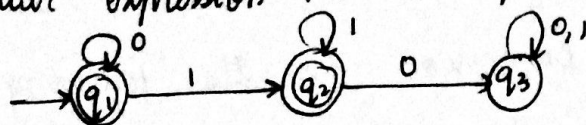
$$\therefore q_0 \in (a + a(b+ab)^*)b)^* \quad [\text{By arden's theorem}]$$

∴ Final state:

$$q_2 = q_1 a = (q_0 a [b + ab]^* a) a$$

$$= (a + a(b + ab)^* b)^* (a(b + ab)^* a)$$

Q.4] Find a regular expression RE corresponding to the following F.



Ans.

$$\begin{aligned}
 q_3 &= q_3 0 + q_3 1 + q_2 0 \\
 \therefore q_3 &= q_3 (0+1) + q_2 0 \quad \text{--- (1)} \\
 q_2 &= q_1 1 + q_2 1 \quad \text{--- (2)} \\
 q_1 &= q_1 0 + \epsilon \quad \text{--- (3)}
 \end{aligned}$$

$$q_3 = q_2 0 (0+1)^*$$

[By arden's theorem]

$$q_1 = \epsilon (0)^*$$

[By arden's theorem]

$$\begin{aligned}
 q_2 &= \epsilon (0)^* + q_2 1 \\
 &= 0^* 1^*
 \end{aligned}$$

[By arden's theorem]

$$\text{Final state : } q_1 + q_2 = 0^* + 0^* 1^*$$

Q.5] Using pumping lemma prove that the following language is not regular :-

$$(i) L = \{ww / w \in (0,1)^*\}$$

Ans: Assume L is regular, then

Pumping length = P .

$$\text{Let } S = 0^P 1 0^P 1$$

$$\text{Let } P = 4$$

$$S = \underbrace{00}_{x} \underbrace{00}_{y} \underbrace{100001}_z$$

$$xy^2z = 00 \ 0000 \ 100001 \notin L$$

$$\text{Here } |y| > 0 \text{ and } |xy| \leq P$$

$\therefore L$ is not regular.

$$(ii) L = \{a^n b^n c^n / n \geq 1\}$$

Ans: Assume L is regular, then

Pumping length = P

$$\text{Let } S = a^P b^P c^P, \quad P = 4$$

$$S = \underbrace{aa}_{x} \underbrace{aa}_{y} \underbrace{bbbb \ cccc}_z$$

$$xy^2z = aa \ aaaa \ bbbb \ cccc \notin L$$

$$\text{Here } |y| > 0 \text{ and } |xy| \leq P.$$

$\therefore L$ is not regular.