

Reasoning behind *ocamlts*

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June 6, 2016

1 Introduction

The purpose of *ocamlts* is to demonstrate abilities of writing basic OCaml programs and share my knowledge of optimization. The purpose of this document is to show some formal logic behind it.

Writing a complete formal proof of *ocamlts* is beyond the scope of that demo, however I would like to explain some reasoning behind it, that in future can serve to write a proper proof of optimality.

2 Optimal solution

2.1 Definitions

Let $path'$ be a path for problem instance P and p_{min} be a distance of $path'$. Let's assume $path'$ fulfills following properties

- 1 $path'$ is a Hamiltonian cycle
- 2 p_{min} is of globally minimal distance

Properties 1 and 2 define optimal solution to instance P of a *Traveling Salesman Problem*.

2.2 Program properties

For convenience main search function `optimize` is renamed to f .

$f(path, d, ub)$ returns (p, d')

Let's assume f be a function that accepts a partial path, it's distance d and best known *upper bound* ub . It returns a tuple of path p and distance d' associated to p .

Function f fulfills following properties

- 3 Final path will be a Hamiltonian cycle (source code line 4)
- 4 f performs a depth first search (recursive calls)
- 5 if branch is cut it implies that it doesn't contain solution better than best known solution (source code line 12)
- 6 optimal solution is guaranteed to be returned at the end
- 7 final solution is optimal solution to TSP

Some additional reasoning to justify points 5, 6 and 7.

- (5) explored branch of $f(path, d, ub)$ will result with distance $d' \geq ub$
- (6) suppose optimal solution p_{min} is not returned, this contradicts either property 4 not all solutions have been enumerated either property 5 optimal solution has been in one of excluded branches
- (7) because $path'$ is a Hamiltonian cycle, if not then contradicts property 3 and p_{min} is of minimal overall length, if not contradicts property 6, if solution that is a Hamiltonian cycle of minimal overall distance is not a solution to *Traveling Salesman Problem* then it contradicts properties 1 and 2

3 Source code

```
1 let rec optimize(dom, dist, path, cost, ub, visit) =
2   if dom = []
3   then
4     if hamiltonian(1, path, visit)
5     then
6       (cost, path)
7     else
8       (-1, [])
9   else
10    if length(hd(dom)) = 1
11    then
12      if cost + hd(hd(dist)) <= ub
13      then
14        optimize(tl(dom), tl(dist), insert(path, hd(hd(dom))),
15                hd(hd(dist)) + cost, ub, visit)
16      else
17        (-1, [])
18    else
19      let (pub, ppath) = optimize(tl(dom), tl(dist), insert(
20                                path, hd(hd(dom))), hd(hd(dist)) + cost, ub, visit)
21      in
22        if pub != -1
23        then
24          if pub <= ub
25          then
26            let (qub, qpath) = optimize(tl(hd(dom)) :: tl(dom)
27                                      , tl(hd(dist)) :: tl(dist), path, cost, pub,
28                                      visit)
29            in
30              if qub != -1
31              then
32                if pub < qub
33                then
34                  (pub, ppath)
35                else if qub < pub
36                then
37                  (qub, qpath)
38                else
39                  (qub, qpath)
40              else
41                (pub, ppath)
42            else
43              optimize(tl(hd(dom)) :: tl(dom), tl(hd(dist)) ::
44                        tl(dist), path, cost, ub, visit)
45          else
46            optimize(tl(hd(dom)) :: tl(dom), tl(hd(dist)) :: tl(
47                      dist), path, cost, ub, visit) ;;
```