

# CptS 451- Introduction to Database Systems

## SQL as a Query Language - part1 (DMS ch-5)

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*World Class. Face to Face.*





# SQL = Structured Query Language

Standard language for querying and manipulating relational data

- Query capabilities of SQL are similar to those in *relational algebra*
- Many standards: SQL92, SQL2, SQL3, SQL99
- SQL language has several aspects:
  - ✓ Data Definition Language (DDL)
    - CREATE TABLE, ALTER TABLE, DROP TABLE
  - Query Language
    - SELECT
  - Data Manipulation Language(DML)
    - INSERT, DELETE, UPDATE
  - Triggers and Advanced Integrity Constraints



# What is special about SQL?

- You describe ***what*** you want
- The job of the DBMS is to figure out ***how*** to compute what you want efficiently.

# Topics

- SQL as a Query Language
  - Select queries
  - Set operations: UNION, ...
  - Aggregation, Group by

## The basic form of a SQL query is *select-from-where*

**SELECT** desired attributes

Project out  
everything not in  
the final answer

**FROM** one or more tables

Every table you  
want to join,  
together

**WHERE** condition on the rows of  
the tables

All the join and  
selection conditions

# SQL as a Query Language

```
SELECT A1, A2, ..., An  
FROM   R1, R2, ..., Rm  
WHERE  conditions;
```

- **Example:** Emp(ssn, ename, dno, sal),  
Dept(dno, dname, mgr),  
Proj(proj\_id, ptitle, startdate, enddate, numEmp),  
ProjEmp(proj\_id, ssn, begindate)

Query 1: “Find employees’ names who work in department 132.”

```
SELECT  ename  
FROM    Emp  
WHERE   dno=132;
```

Query 2: “Find the manager of the Marketing department.”

```
SELECT  mgr  
FROM    Dept  
WHERE   dname = 'Marketing';
```

# SQL vs Relational Algebra

```
SELECT A1, A2, ..., An  
FROM   R1, R2, ..., Rm  
WHERE conditions;
```

- Equivalent relational algebra expression:

$$\Pi_{A1, \dots, An} (\sigma_{\text{cond}} (R1 \times R2 \times \dots Rm))$$

- Difference:
  - Relational algebra uses set semantics
  - Most SQL operators uses bag semantics
    - However, SQL set operators use set semantics
    - Set operators are applied on query results

# “Select” Clause

- Specify attributes to project onto (different from the “selection” operator in the relational algebra)
- Use star \* to denote all attributes:  

```
SELECT *
FROM Emp
WHERE ename = 'Jack' AND sal > 50K;
```

Emp(ssn,ename,dno,salary)

ssn	Ename	Dno	sal
111-111-1111	Jack	111	81K
<del>222-111-2222</del>	<del>Alice</del>	<del>111</del>	<del>70K</del>
<del>333-111-3333</del>	<del>Lisa</del>	<del>222</del>	<del>32K</del>
<del>444-111-4444</del>	<del>Tom</del>	<del>333</del>	<del>56K</del>
<del>555-111-5555</del>	<del>Mary</del>	<del>333</del>	<del>65K</del>
<del>666-111-6666</del>	<del>Jack</del>	<del>444</del>	<del>45K</del>



# Here is a way to think about how the query might be implemented

1. Imagine a *tuple variable* ranging over each tuple of the relation mentioned in FROM.
2. Check if the “current” tuple satisfies the WHERE clause.
3. If so, output the attributes/expressions of the SELECT clause using the components of this tuple.



A	B	C
A1	B1	C1
A2	B2	C2
A3	B3	C3
A4	B4	C4
A5	B5	C5
A6	B6	C6
A7	B7	C7

```
SELECT A, B
FROM R
WHERE A = 'A3';
```

A	B
A3	B3

# “Select” Clause

- Single Relation vs. Multi Relation Queries
- Single relation:  

```
SELECT *
FROM Emp
WHERE ename = 'Jack' AND sal > 50K;
```
- Multiple relations:
  - Can use relation prefix (especially when we need to disambiguate attribute names)  

```
SELECT *
FROM Emp, Dept
WHERE Emp.dno = Dept.dno;
```

Emp(ssn,ename,dno,salary)

ssn	Ename	dno	sal
111-111-1111	Jack	111	81K
222-111-2222	Alice	111	70K
333-111-3333	Lisa	222	32K
444-111-4444	Tom	333	56K
555-111-5555	Mary	333	65K

Dept(dno,dname, mgr)

dno	dname	mgr
111	HR	Alice
222	R&D	Lisa
333	Production	Mary

# “Select” Clause

```
SELECT *
FROM Emp, Dept
WHERE Emp.dno = Dept.dno;
```

Emp(ssn,ename,dno,salary)

ssn	Ename	dno	sal
111-111-1111	Jack	111	81K
222-111-2222	Alice	111	70K
333-111-3333	Lisa	222	32K
444-111-4444	Tom	333	56K
555-111-5555	Mary	333	65K

Dept(dno,dname, mgr)

dno	dname	mgr
111	HR	Alice
222	R&D	Lisa
333	Production	Mary

Emp.ename	Emp.dno	Emp.sal	Dept.dno	Dept.dname	Dept.mgr
Jack	111	81K	111	HR	Alice
Alice	111	70K	111	HR	Alice
Lisa	222	32K	222	R&D	Lisa
Tom	333	56K	333	Production	Mary
Mary	333	65K	333	Production	Mary

# Eliminate Duplicates

- “SELECT” does **not** automatically eliminate duplicates.  

```
SELECT dno
FROM Emp;
```

  - If there are more than 1 employee in the department 333, then 333 will appear more than once in the result.
- Use keyword **distinct** to explicitly remove duplicates

Emp(ssn,ename,dno,salary)

ssn	Ename	dno	sal
111-111-1111	Jack	111	81K
222-111-2222	Alice	111	70K
333-111-3333	Lisa	222	32K
444-111-4444	Tom	333	56K
555-111-5555	Mary	333	65K

```
SELECT distinct dno
FROM Emp;
```



dno
111
222
333



## “Select” Clause (cont.)

- You can rename the attributes in the result, using “as <new name>”

```
SELECT  ename, mgr as manager
FROM    Emp, Dept
WHERE   Emp.dno = Dept.dno AND
        manager='Alice';
```

Emp(ssn,ename,dno,salary)

ssn	Ename	dno	sal
111-111-1111	Jack	111	81K
222-111-2222	Alice	111	70K
333-111-3333	Lisa	222	32K
444-111-4444	Tom	333	56K
555-111-5555	Mary	333	65K

- You can create a new column and give it a constant value, in the SELECT clause

```
SELECT  ename, dno, 'temporary' as status
FROM    Emp
WHERE   dno = 111;
```

ename	dno	status
Jack	111	temporary
Alice	111	temporary

## “Select” Clause (cont.)

- You can use math in the SELECT clause

```
SELECT  eNaMe, sal*1.05 as newSalary
FROM    Emp
WHERE   ename='O''Fallon';
```

Case-insensitive, except inside quoted strings

Two single quotes inside a string  
= one apostrophe

ename	newSalary
O'Fallon	85.05K

# “FROM” clause

- Specify relations
- Renaming relations:
  - Use “**as**” to define “**variables**,” to disambiguate multiple references to the same relation
  - Example: “who has higher salary than their manager”

```

SELECT  E1.ename
FROM    Emp as E1, Dept D, Emp as E2
WHERE   E1.dno = D.dno AND
        D.mgr = E2.ename AND
        E1.sal > E2.sal;

```

E1: Emp(ssn,ename,dno,salary)

ssn	Ename	dno	sal
111-111-1111	Jack	111	81K
222-111-2222	Alice	111	70K
333-111-3333	Lisa	222	32K
444-111-4444	Tom	333	56K
555-111-5555	Mary	333	65K

Dept(dno,dname, mgr)

dno	dname	mgr
111	HR	Alice
222	R&D	Lisa
333	Production	Mary

E1: Emp(ssn,ename,dno,salary)

ssn	Ename	dno	sal
111-111-1111	Jack	111	81K
222-111-2222	Alice	111	70K
333-111-3333	Lisa	222	32K
444-111-4444	Tom	333	56K
555-111-5555	Mary	333	65K

# “WHERE” clause

- Specify conditions
- Optional
- Complex conditions:
  - AND, OR, NOT, ...
  - “Employees who work for Lisa and have a salary < 70K”

```
SELECT  ename
FROM    Emp, Dept
WHERE   Emp.dno=Dept.dno AND
        mgr = 'Lisa' AND
        sal < 70K;
```



# “WHERE” clause (cont.)

- String patterns:
  - **LIKE** keyword uses a regular expression to contain the pattern that the values are matched against
  - “**s LIKE p**”: string **s** matches pattern **p**
  - Pattern may include:
    - **% (percent)**: zero, one, or multiple occurrences of any character
      - `dname LIKE 'TOM %'`
        - » ‘TOM KERRY’, ‘TOM JOHNSON’, ‘TOM ‘ ...
    - **\_ (underbar)**: one-character wildcard
      - `dname LIKE 'a_c'`
        - » ‘abc’ ‘adc’ ‘azc’ ‘a9c’ ...

# Conditions in a “WHERE” clause

The following may appear in the WHERE condition

- constants of any supported type
- attribute names of the relation(s) used in the FROM.
- comparison operators: =, <>, <, >, <=, >=
- arithmetic operations: price\*2
- operations on strings (e.g., CONCAT for concatenation).
- lexicographic order on strings (lastname<'Norman').
- pattern matching: s LIKE p , s NOT LIKE p
- special operations for comparing dates and times.
- and combinations of the above using AND, OR, NOT, and parentheses
- Use relation prefix to disambiguate attribute names  

```
SELECT ename, dname, dept.dno
FROM   Emp, Dept
WHERE  Emp.dno = Dept.dno;
```

# Conditions in a “WHERE” clause

- What if an attribute value is unknown, or the attribute is inapplicable (i.e. is NULL)?

— Example:

```
SELECT ename, sal
FROM Emp
WHERE sal<=50K OR sal>50K;
```

Emp(ssn,ename,dno,salary)

ssn	Ename	dno	sal
111-111-1111	Jack	111	81K
222-111-2222	Alice	111	70K
333-111-3333	Lisa	222	NULL
444-111-4444	Tom	333	NULL
555-111-5555	Mary	333	65K

ename	sal
Jack	81K
Alice	70K
Mary	65K

Why???



# Conditions involving NULL evaluate to *unknown*, rather than *true* or *false*

Example condition	Evaluates to
'Tom' = 'Tom'	<b>true</b>
2 > 6	<b>false</b>
'Tom' = NULL	unknown
2 < NULL	unknown
<b>true</b> AND unknown	unknown
<b>true</b> OR unknown	<b>true</b>
<b>false</b> AND unknown	<b>false</b>
<b>false</b> OR unknown	unknown
unknown OR unknown	unknown

True-> 1  
False-> 0  
Unknown->1/2

A tuple only goes in the answer if its truth value for the WHERE clause is true.

# Conditions in a “WHERE” clause

- What if an attribute value is unknown, or the attribute is inapplicable?
  - Example:

```
SELECT ename, sal
FROM Emp
WHERE sal<=50K OR sal>50K;
```

*unknown* (circled around 'OR')

*unknown* (circled around 'sal<=50K')

*unknown* (circled around 'sal>50K')

ename	sal
Jack	81K
Alice	70K
Mary	65K

*unknown*

Emp(ssn,ename,dno,salary)

ssn	Ename	dno	sal
111-111-1111	Jack	111	81K
222-111-2222	Alice	111	70K
333-111-3333	Lisa	222	NULL
444-111-4444	Tom	333	NULL
555-111-5555	Mary	333	65K

# Dealing with NULL Values

Can test for NULL explicitly:

- IS NULL
- IS NOT NULL

```
SELECT ename, sal  
FROM Emp  
WHERE sal<=50000 OR sal>50000 OR sal is NULL;
```

The answer includes all employees!

# Ordering Output Tuples

```
SELECT *  
FROM   Emp  
WHERE  sal<=50000  
ORDER BY dno, sal desc, ename;
```

- First, order the tuples by dno (department).
- Within each department, order salaries from highest to lowest.
- For salary ties, use alphabetical order on the name.

# Ordering Output Tuples

```
SELECT *
FROM Emp
WHERE sal<=50000
ORDER BY dno, sal desc, ename;
```

By default, ORDER BY orders in ascending order. Use keyword “**desc**” for descending order.

- What if there are NULL values?

Emp(ssn,ename,dno,salary)

ssn	ename	dno	sal
111-111-1111	Jack	NULL	81K
222-111-2222	Alice	111	70K
333-111-3333	Lisa	222	32K
444-111-4444	Tom	333	56K
555-111-5555	Mary	NULL	65K

ssn	ename	dno	sal
111-111-1111	Jack	NULL	81K
555-111-5555	Mary	NULL	65K
222-111-2222	Alice	111	70K
333-111-3333	Lisa	222	32K
444-111-4444	Tom	333	56K

```
SELECT *
FROM Emp
WHERE sal<=50000 OR sal IS NULL
ORDER BY dno NULLS FIRST, sal desc, ename;
```

SQL:2003 standard



# Set Operations

- **Use the set semantics**
  - duplicates are eliminated in the result.
- **Example:**
- **Union:**  $\cup$  “Find employees who work either for the ‘Purchasing’ or the ‘HR’ department.”

```
(SELECT  ename FROM Emp, Dept
 WHERE Emp.dno=Dept.dno AND dname='Purchasing')
UNION
(SELECT  ename FROM Emp, Dept
 WHERE Emp.dno=Dept.dno AND dname='HR')
```

The schema of the  
SELECT results  
should be same

- **Intersect:**  $\cap$  “Find employees who work both for the ‘Purchasing’ or the ‘HR’ departments.”

```
(SELECT  ename FROM Emp, Dept
 WHERE Emp.dno=Dept.dno AND dname='Purchasing')
INTERSECT
(SELECT  ename FROM Emp, Dept
 WHERE Emp.dno=Dept.dno AND dname='HR')
```

# Set Operations

- **Except:** – “Find employees who work for the ‘Accounting’ department but not for the ‘Purchasing’ department.”

```
(SELECT  ename FROM Emp, Dept
 WHERE Emp.dno=Dept.dno AND dname='Accounting')
EXCEPT
(SELECT  ename FROM Emp, Dept
 WHERE Emp.dno=Dept.dno AND dname='Purchasing')
```

# Set Operations - Conserving Duplicates

- The UNION, INTERSECT, and EXCEPT operators use the set semantics, not bag semantics.
- To keep duplicates, use “**ALL**” after the operators:
  - UNION ALL, INTERSECT ALL, EXCEPT ALL
  - Example:

```
(SELECT ssn, name, “student” as standing FROM Student)
UNION ALL
(SELECT ssn, name, “TA” as standing FROM TeachingAssistant)
```

Student (ssn, name)

ssn	name
111	Tom
222	Jack
444	Mary

TA (ssno, name)

ssn	name
111	Tom
222	Jack
555	Alice

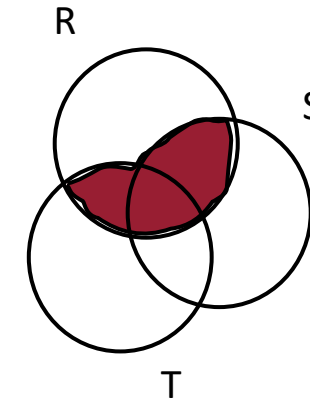
Result

ssn	name	standing
111	Tom	student
222	Jack	student
444	Mary	student
111	Tom	TA
222	Jack	TA
555	Alice	TA

# Set Operations - Example

- Relations:  $R(A)$ ,  $S(A)$ ,  $T(A)$
- Query: " $R \cap (S \cup T)$ "

```
SELECT R.A FROM R
intersect
( (SELECT S.A FROM S)
  union
  (SELECT T.A FROM T)
);
```



Solution-1

```
(SELECT R.A FROM R, S WHERE R.A=S.A)
union
(SELECT R.A FROM R, T WHERE R.A=T.A);
```

Solution-2

```
SELECT R.A
FROM   R, S, T
WHERE  R.A=S.A OR R.A=T.A;
```

Wrong!

- The SQL result becomes empty when T is empty

# Aggregations

- MIN, MAX, SUM, COUNT, AVG
  - input: collection of numbers/strings (depending on operation)
  - output: relation with a single attribute with a single row
- Example: “What is the minimum, maximum, average salary of employees in the ‘Marketing’ department”

```
SELECT MIN(sal), MAX(sal), AVG(sal)
FROM   Emp, Dept
WHERE  Emp.dno = Dept.dno and Dept.dname = Marketing ;
```

# Aggregations (cont.)

- Except “count,” all aggregations apply to a single attribute
- “Count” can be used on “\*”

```
SELECT Count(*) FROM Emp;
```

```
SELECT Count(ename) FROM Emp;
```

Emp(ssn,ename,dno,salary)

ssn	Ename	dno	sal
111-111-1111	Jack	111	81K
222-111-2222	Alice	111	70K
333-111-3333	Lisa	222	32K
444-111-4444	Tom	333	56K
555-111-5555	Mary	333	65K

Result : 5

# Duplication in Aggregations

- “What is the number of different dno’s in the Emp table”

```
SELECT count(dno)
FROM Emp;
```

Wrong! Since there can be duplicates

- Right Query:

```
SELECT count(DISTINCT dno)
FROM Emp;
```

Emp(ssn,ename,dno,salary)

ssn	Ename	dno	sal
111-111-1111	Jack	111	81K
222-111-2222	Alice	111	70K
333-111-3333	Lisa	222	32K
444-111-4444	Tom	333	56K
555-111-5555	Mary	333	65K

Result : 3

# GROUP BY Clause

- **GROUP BY** is used to apply aggregate function to a group of sets of tuples.
  - The aggregate function is applied to each group separately.
- **Example:** “For each department, list its total number of employees and total salary”

```
SELECT Dept.dno, SUM(sal), COUNT(ssn)
FROM Emp, Dept
WHERE Emp.dno = Dept.dno
GROUP BY Dept.dno;
```

Emp(ssn,ename,dno,salary)

ssn	Ename	dno	sal
111-111-1111	Jack	111	81K
222-111-2222	Alice	111	70K
333-111-3333	Lisa	222	32K
444-111-4444	Tom	333	56K
555-111-5555	Mary	333	65K

Dept(dno,dname, mgr)

dno	dname	mgr
111	HR	Alice
222	R&D	Lisa
333	Production	Mary

Result

dname	Sum(sal)	Count(ename)
HR	151K	2
R&D	32K	1
Production	121K	2



# GROUP BY Clause (cont.)

- Standard SQL: “SELECT” attributes must appear in Group-by attributes.
- The following queries cannot group the tuples.

```
SELECT dname, Emp.dno, SUM(sal), COUNT(ssn)
FROM Emp, Dept
WHERE Emp.dno = Dept.dno
GROUP BY Emp.dno;
```

```
SELECT dname, SUM(sal), COUNT(ssn)
FROM Emp, Dept
WHERE Emp.dno = Dept.dno
GROUP BY Emp.dno;
```

## GROUP BY Clause (cont.)

- Do the following queries return the same result?

```
SELECT dno  
FROM Emp  
GROUP BY dno;
```

```
SELECT distinct dno  
FROM Emp;
```

# HAVING Clause

- **HAVING** clause used along with GROUP BY clause to select some groups.
  - We can't define conditions on aggregate results in the WHERE clause
  - Syntax: **HAVING** aggregate\_function(column\_name) operator value
- Predicate in having clause applied after the formation of groups.
- **Example:** “List the department name, total salary, and number of employees for all departments with more than 2 employees.”

```
SELECT dname, SUM(sal), COUNT(ssn)
FROM Emp, Dept
WHERE Emp.dno = Dept.dno
GROUP BY dname
HAVING COUNT(ssn)>2;
```

Emp(ssn,ename,dno,salary)

ssn	Ename	dno	sal
111-111-1111	Jack	111	81K
222-111-2222	Alice	111	70K
333-111-3333	Lisa	222	32K
444-111-4444	Tom	333	56K
555-111-5555	Mary	333	65K

Dept(dno,dname, mgr)

dno	dname	mgr
111	HR	Alice
222	R&D	Lisa
333	Production	Mary

Result

dname	Sum(sal)	Count(ename)
HR	151K	2
Production	121K	2



# A General SQL Select Query

- “For each employee that works in two or more departments, print the total salary of his/her managers.”

```
SELECT ssn, ename, count(*)  
FROM Emp, Dept  
WHERE Emp.dno=Dept.dno  
GROUP BY ssn, ename  
HAVING count(*) > 1  
ORDER BY ssn,ename;
```

Find employees that works in  
two or more departments

```
SELECT E1.ssn,E1.ename, sum(E2.sal)  
FROM Emp as E1, Dept, Emp as E2  
WHERE E1.dno = Dept.dno AND E2.ename = Dept.mgr  
GROUP BY E1.ssn, E1.ename  
HAVING count(*) > 1  
ORDER BY E1.ssn, E1.ename;
```

For those  
employees, find  
their managers  
and calculate the  
sum of the  
managers' salaries.



# A General SQL Select Query

- For each employee that works in two or more departments, print the total salary of his/her managers.

```
SELECT E1.ssn, E1.ename, SUM(E2.sal)
FROM Emp E1, Dept, Emp E2
WHERE E1.dno = Dept.dno AND E2.ename = Dept.mgr
GROUP BY E1.ssn,E1.ename
HAVING count(distinct(Dept.dno)) > 1
ORDER BY E1.ssn,E1.ename;
```

5  
1  
2  
3  
4  
6



# A General SQL Query

- For each employee that works in two or more departments, print the total salary of his/her managers. Assume each dept has one manager.

```
SELECT E1.ssn, E1.ename, SUM(E2.sal)
FROM Emp E1, Dept, Emp E2
WHERE E1.dno = Dept.dno AND E2.ename = Dept.mgr
GROUP BY E1.ssn, E1.ename
HAVING count(distinct(Dept.dno)) > 1
ORDER BY E1.ssn, E1.ename;
```

5  
1  
2  
3  
4  
6

## Execution steps:

- Step 1: tuples are formed (Cartesian product)
- Step 2: tuples satisfying the conditions are chosen
- Step 3: groups are formed
- Step 4: groups are eliminated using “Having”
- Step 5: the aggregates are computed for the select line, flattening the groups
- Step 6: the output tuples are ordered and printed out.