

# Sorting #3 - The Merge of the King

CptS 223 - Fall 2017 - Aaron Crandall



# Today's Agenda

- Announcements
- Thing of the day
- Sorting is fun! - Time for some merge sort

# Announcements



- Looks like our midterm #2 will be...
  - Any major objections?
  - Topics include: Hashing, Heaps, Sorting



# To answer the “Is heap sort stable” question:

## No.

- More specifically, consider sorting:
  - Input list: [21 20a 20b 12 11 8 7]  
It will come out: [7 8 11 12 20b 20a 21]
- There's also opportunity for much larger heaps to cause unstable behavior if equal elements are in different subtrees with well crafted parents.

# Playing with Oracle of Bacon



- All movie performers are really close to Kevin Bacon
  - <https://oracleofbacon.org/>
- This is a graph traversal problem finding shortest paths
  - Dijkstra's Algorithm
- We'll look at it more after Midterm #2

# TotD: Robots doing to-home grocery delivery?



**“These six-wheeled robots are about to start delivering food in the US”**

**Door to door delivery of groceries and small goods. Mostly autonomous, but always learning. Currently being tested in D.C., as it shown to my right. →**

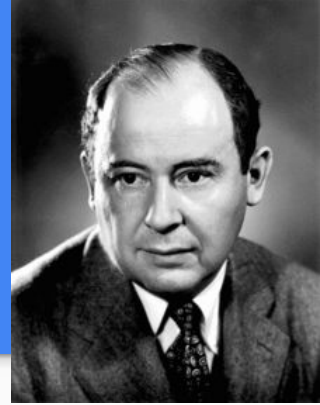


# Mergesort! - It's a way to come together

- Runs in  $O(N \log N)$  time
- Number of comparisons nearly optimal
- Fundamentally, it's merging two lists together
  - The lists are already sorted (by definition)
  - Results go in a 3rd list - beware space complexity!
    - Can be mitigated by using in-place swaps at cost of time
- FYI: Invented by John von Neumann in 1945
  - Was ranked as a top secret algorithm by the US military at the time

# John von Neumann

<https://goo.gl/ORD7D>

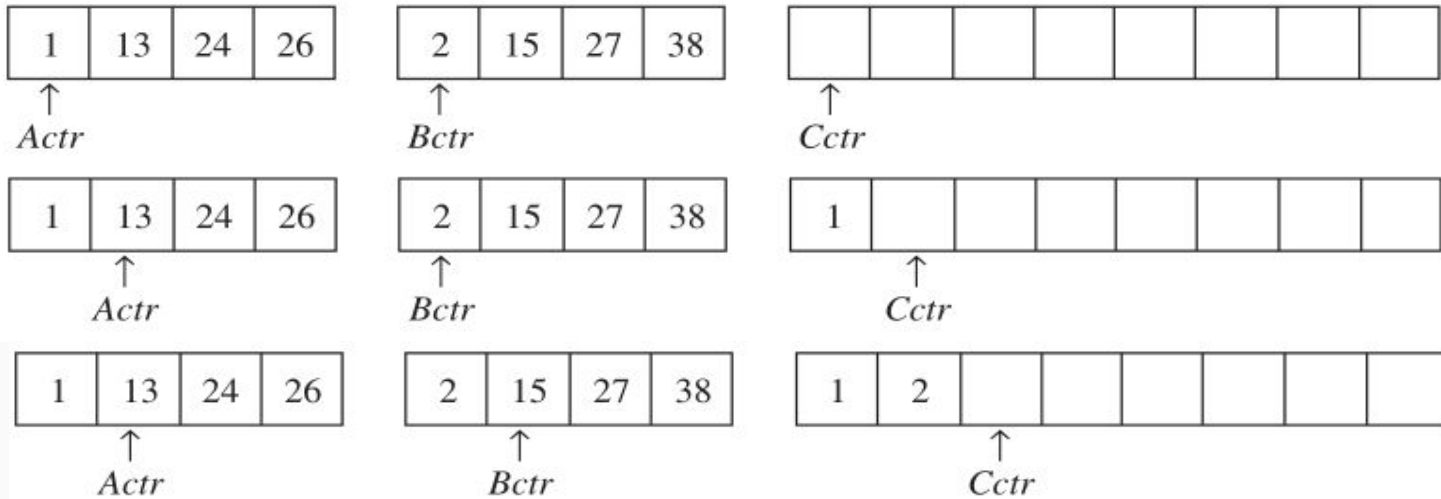


- "the last representative of the great mathematicians" like Euler, Gauss, Poincaré or Hilbert"
- Pioneer of operator theory to quantum mechanics
- Key dev in game theory
- Cellular automata, the universal constructor and the digital computer
- 150 papers: 60 math, 20 physics, 60 applied math - last one from his hospital bed became a book "The Computer and The Brain"
- Analysis of self replication work preceded DNA
- This is just the tip of the iceberg - we owe him so very much



# Top level view of merging lists in general

- Lists: A, B, C
- Counters: Actr, Bctr, Cctr



# And so on

1	13	24	26
---	----	----	----

↑  
*Actr*

2	15	27	38
---	----	----	----

↑  
*Bctr*

1	2	13					
---	---	----	--	--	--	--	--

↑  
*Cctr*

1	13	24	26
---	----	----	----

↑  
*Actr*

2	15	27	38
---	----	----	----

↑  
*Bctr*

1	2	13	15				
---	---	----	----	--	--	--	--

↑  
*Cctr*

1	13	24	26
---	----	----	----

↑  
*Actr*

2	15	27	38
---	----	----	----

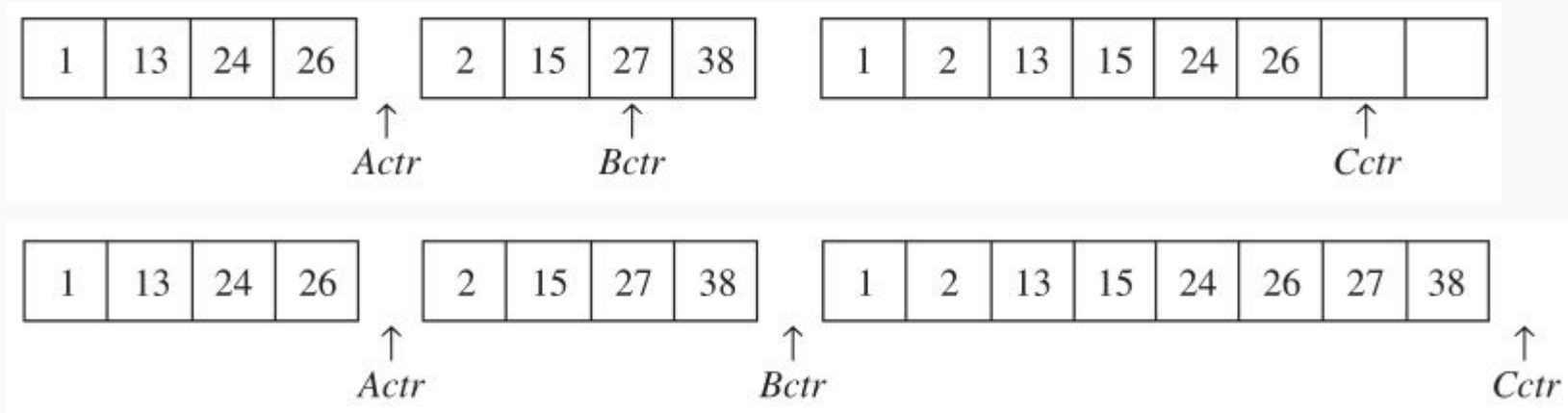
↑  
*Bctr*

1	2	13	15	24			
---	---	----	----	----	--	--	--

↑  
*Cctr*

# Short circuit if one of the lists ends early

- Remainder of B gets copied without comparisons into C when A runs out.



# Time to merge two lists

- Clearly linear: at most  $N-1$  comparisons are made
- The question is: how to get two sorted lists?

- The answer is:

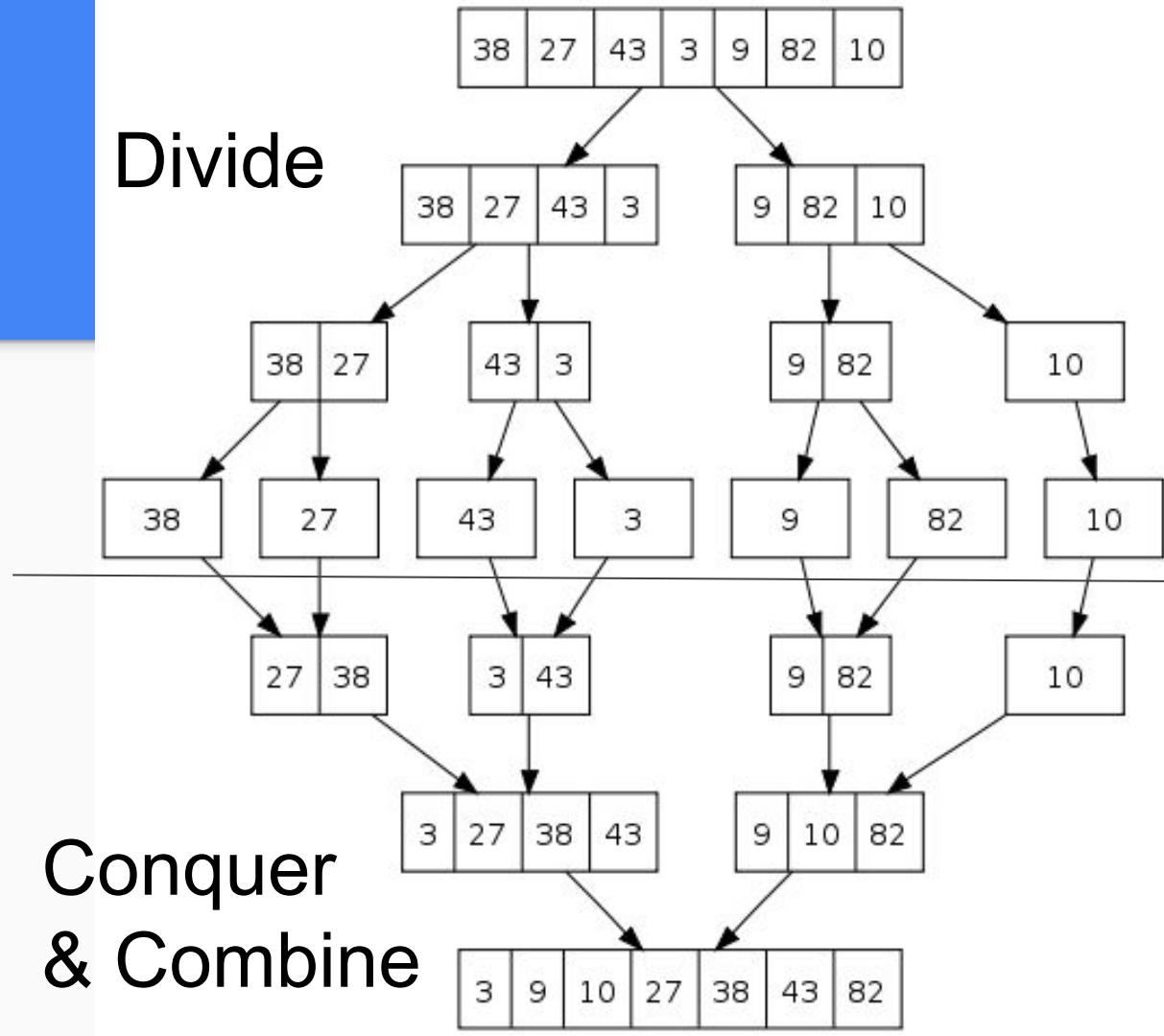
Recursively subdivide the original list and sort those subdivided lists

# For example

- To sort: [24,13,26,1,2,27,38,15]
- Recursively sort the first 4 and last 4 elements to get:  
[1,13,24,26,2,15,27,38]
- Merge the two halves to get: [1,2,13,15,24,26,27,38]
- This is a classic divide and conquer approach
  - Divide the problem into subproblems to be solved recursively
  - Conquer by patching together the solved pieces
- Is this algorithm a stable sort?

# Other example

Divide



Conquer  
& Combine

# Analyzing Mergesort

- Recursive analysis for this algorithm
- Assume  $N$  is a power of 2 (to make it easy)
- For  $N == 1$ , time to merge is  $O(1)$
- Time to mergesort  $N$  numbers is time to do two recursive mergesorts of size  $N/2$ , plus time to merge, which is linear (That  $N-1$  we saw earlier)
  - $T(1) = 1$
  - $T(N) = 2T(N/2) + N$

$$\frac{T(N)}{N} = \frac{T(N/2)}{N/2} + 1$$

$$\frac{T(N/2)}{N/2} = \frac{T(N/4)}{N/4} + 1$$

$$\frac{T(N/4)}{N/4} = \frac{T(N/8)}{N/8} + 1$$

$\vdots$

$$\frac{T(2)}{2} = \frac{T(1)}{1} + 1$$

This series sums up  $\log N$  times and mostly cancels out to:

$$\frac{T(N)}{N} = \frac{T(1)}{1} + \log N$$

Multiply by  $N$  on both sides for:

$$T(N) = N \log N + N = O(N \log N)$$



# Limits of Mergesort

- Although it's  $O(N \log N)$  time:
  - Requires linearly  $O(N)$  extra space for the merging
  - Has to copy to temp array and back with each layer of recursion
    - Can be avoided by swapping `a` and `tmpArray` as you alternate recursive calls
- Heavily dependent upon the wallclock costs of comparing elements and moving elements in the array - This is language implementation dependent
  - Java: comparisons are slow, moves are fast -> mergesort is default sort in the libraries
  - C++: copying can be expensive, comparing fast -> Quicksort default in libraries
  - C++11 has better copying semantics, so this might change, but it hasn't yet

# Our Book's Mergesort Implementation

```
template <typename Comparable>
void mergeSort( vector<Comparable> & a,
               vector<Comparable> & tmpArray, int left, int right )
{
    if( left < right )
    {
        int center = ( left + right ) / 2;
        mergeSort( a, tmpArray, left, center );
        mergeSort( a, tmpArray, center + 1, right );
        merge( a, tmpArray, left, center + 1, right );
    }
}
```

```
template <typename Comparable>
void merge( vector<Comparable> & a, vector<Comparable> & tmpArray,
           int leftPos, int rightPos, int rightEnd )
{
    int leftEnd = rightPos - 1;
    int tmpPos = leftPos;
    int numElements = rightEnd - leftPos + 1;

    // Main loop
    while( leftPos <= leftEnd && rightPos <= rightEnd )
        if( a[ leftPos ] <= a[ rightPos ] )
            tmpArray[ tmpPos++ ] = std::move( a[ leftPos++ ] );
        else
            tmpArray[ tmpPos++ ] = std::move( a[ rightPos++ ] );
}
```

```
while( leftPos <= leftEnd )    // Copy rest of first half
    tmpArray[ tmpPos++ ] = std::move( a[ leftPos++ ] );

while( rightPos <= rightEnd )  // Copy rest of right half
    tmpArray[ tmpPos++ ] = std::move( a[ rightPos++ ] );

// Copy tmpArray back
for( int i = 0; i < numElements; ++i, --rightEnd )
    a[ rightEnd ] = std::move( tmpArray[ rightEnd ] );
}
```

# Mergesort summary:

- Operates in  $O(N \log N)$  time - DOES NOT go faster for sorted data?
- Is stable
- Can take  $O(N)$  extra space unless you do extra work
- Invented by John Von Neumann in 1945
  - This name will come up again in your studies...
  - Was originally classified as Top Secret
  - Invented the Von Neumann computer architecture, which your computers are using now
- The sort used by Java because it's nearly optimal in comparisons