## **Algorithms**

**Lecture 4: Greedy Algorithm** 

Anxiao (Andrew) Jiang

CH 16. Greedy Algorithms

16.3 Huffman Code

Symbol	a	b	$\mathcal{C}$	d	e	f
Probability	0.45	0.13	0.12	0.16	0.09	0.05

How to represent the symbols using bits, to minimize the average number of bits needed?

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How to represent the symbols using bits, to minimize the average number of bits needed?

Assume: we use "Fixed Length Code (FLC)"

1-bit codewords: 0, 1 (not enough)

2-bit codewords: 00, 01, 10, 11 (not enough)

3-bit codewords: 000, 001, 010, 011, 100, 101, 110, 111 (enough)

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Variable Length Code (VLC): the codewords can have different lengths.

Prefix Code: no codeword is the prefix of another codeword.

Let's study Variable Length Prefix Code (VLPC).

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Text:  $a \ a \ b \ f \ e \ d \ c \ a \ \cdots$ 

Compression (encoding): turn a text to bits

Decompression (decoding): turn a bit sequence back to text

a a

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Text:  $a \ a \ b \ f \ e \ d \ c \ a \cdots$ 

Compression (encoding): turn a text to bits

00101110011111000 ...

Decompression (decoding): turn a bit sequence back to text

 $a \quad a \quad b \quad f \quad e \quad d \quad c \quad a \quad \dots$ 

Symbol	a	b	С	d	е	f	
Probability	0.45	0.13	0.12	0.16	0.09	0.05	
FLC	000	001	010	011	100	101	3 bits per symbol
VLPC	0	101	100	111	1101	1100	2.24 bits per symbol

## Average codeword length for VLPC:

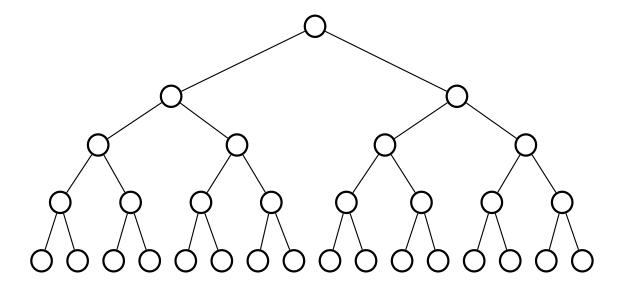
 $1 \times 0.45 + 3 \times 0.13 + 3 \times 0.12 + 3 \times 0.16 + 4 \times 0.09 + 4 \times 0.05 = 2.24$  bits/symbol

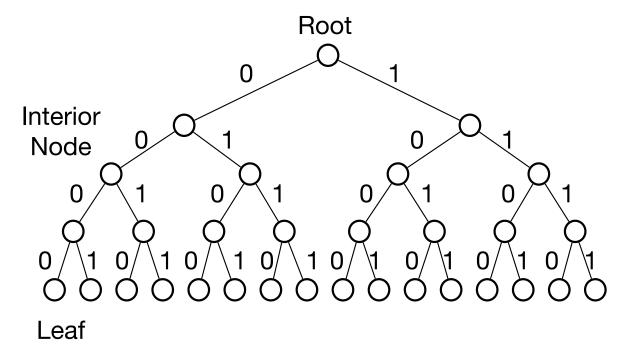
Input: n symbols  $s_1, s_2, \cdots, s_n$ . For  $i=1,2,\cdots,n$ , the symbol  $s_i$  has probability  $f_i$ .

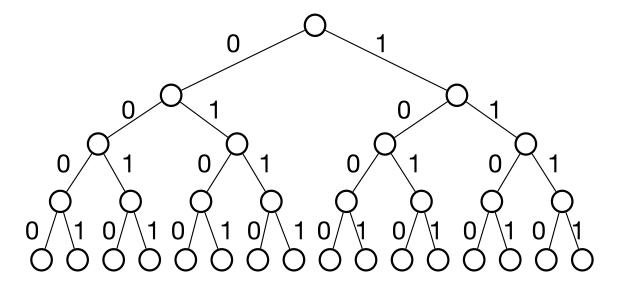
Output: Design a prefix code for the n symbols such that the average codeword length is minimized.

Assume the codeword for symbol  $s_i$  has  $L_i$  bits.

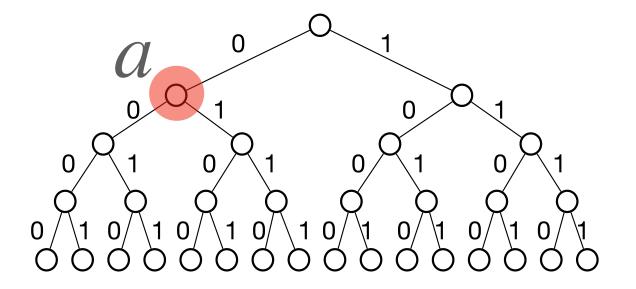
Average Codeword Length = 
$$\sum_{i=1}^{n} f_i L_i$$



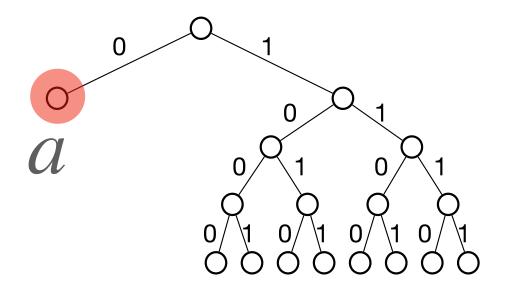




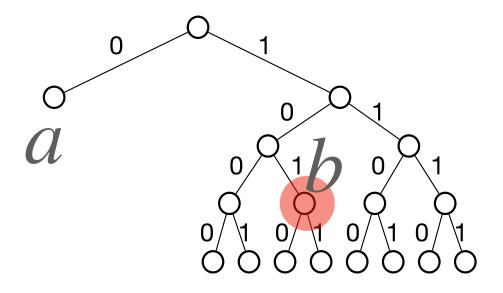
Symbol	a	b	$\mathcal{C}$	d	e	f
Probability	0.45	0.13	0.12	0.16	0.09	0.05
VLPC	0	101	100	111	1101	1100



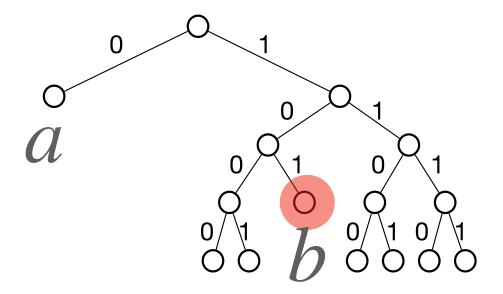
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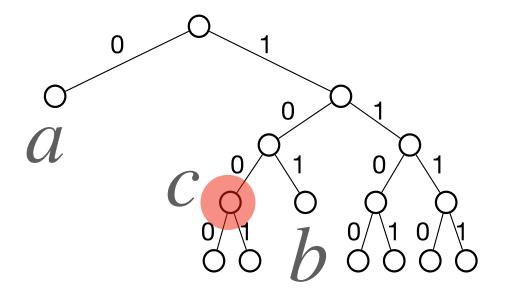
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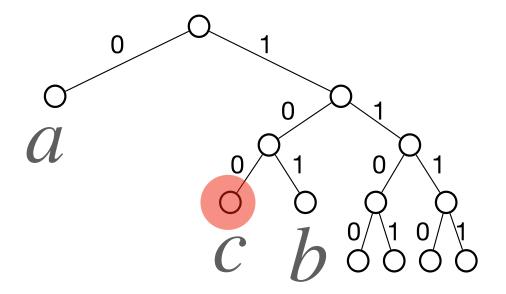
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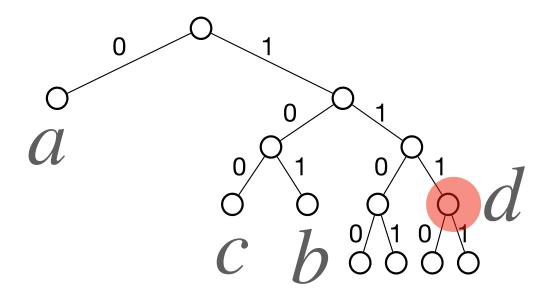
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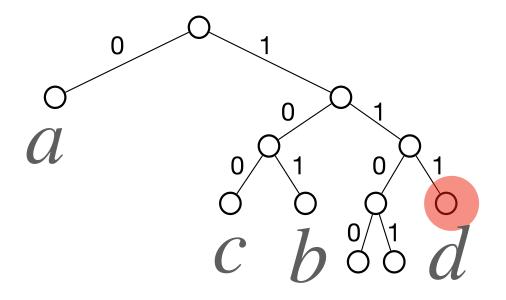
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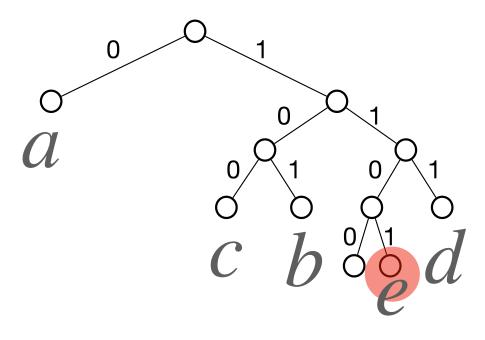
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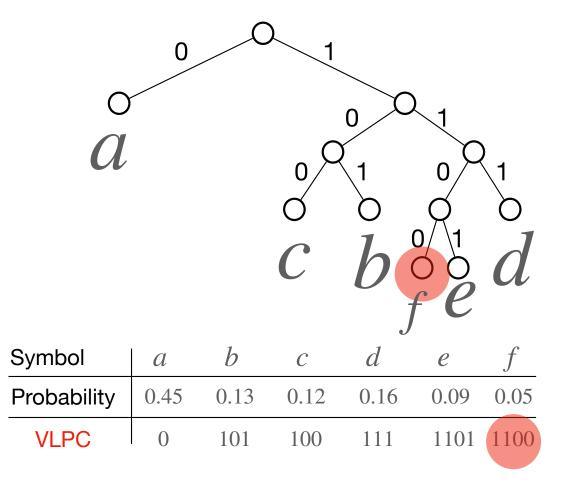
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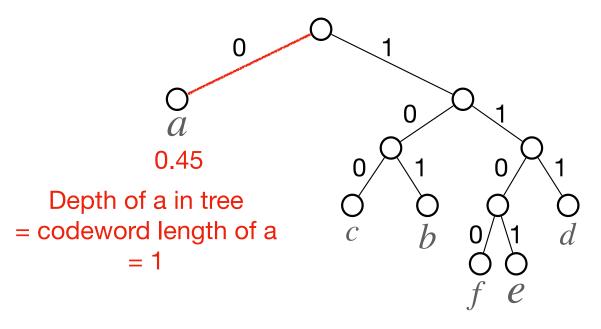


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Leaves of a subtree



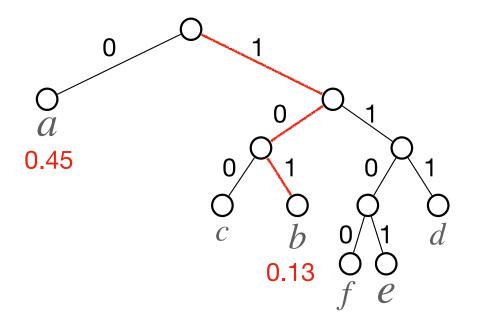


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Average codeword length



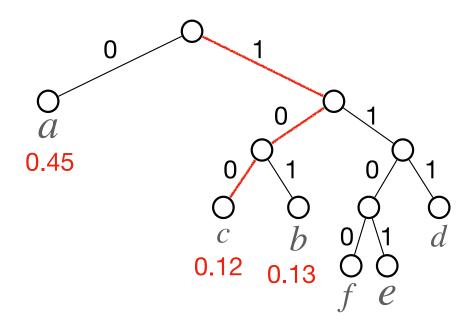
Average depth of leaves



Average codeword length



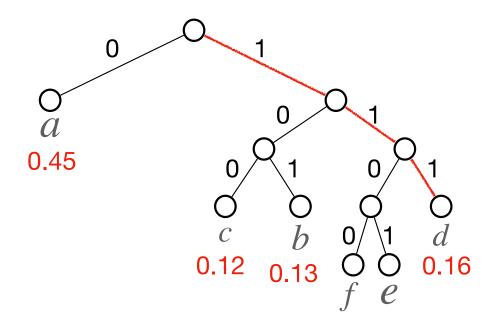
Symbol	a	b	С	d	e	f
Probability	0.45	0.13	0.12	0.16	0.09	0.05
VLPC	0	101	100	111	1101	1100



Average codeword length

Average depth of leaves

Symbol	a	b	C	d	e	f
Probability	0.45	0.13	0.12	0.16	0.09	0.05
VLPC	0	101	100	111	1101	1100

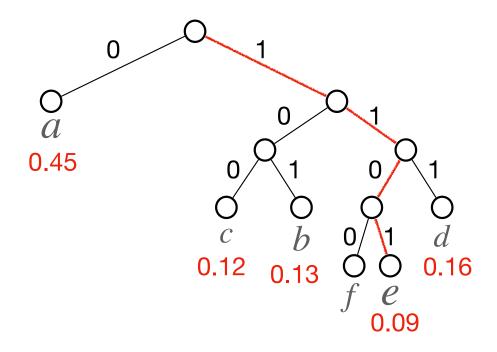


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Symbol	а	b	C	d	e	f
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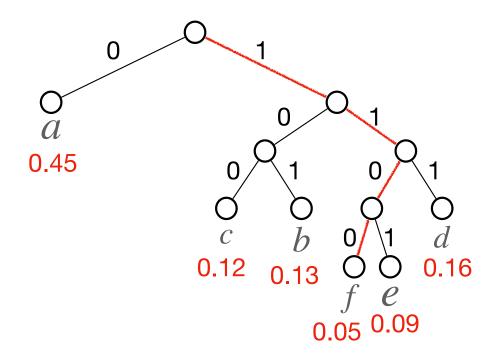


Average codeword length



Average depth of leaves

Symbol	a	b	С	d	e	f
Probability	0.45	0.13	0.12	0.16	0.09	0.05
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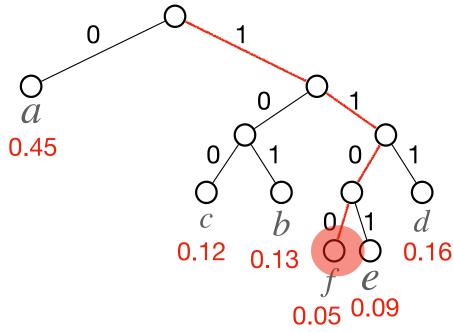
Symbol	а	b	C	d	e	f
Probability	0.45	0.13	0.12	0.16	0.09	0.05
VLPC	0	101	100	111	1101	1100

Average codeword length

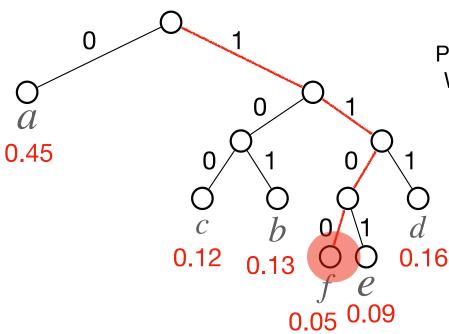


Average depth of leaves

Property 1 of optimal code: The (or one) symbol of lowest probability has the longest codeword.



Symbol	а	b	$\boldsymbol{\mathcal{C}}$	d	e	f
Probability	0.45	0.13	0.12	0.16	0.09	0.05
VLPC	0	101	100	111	1101	1100

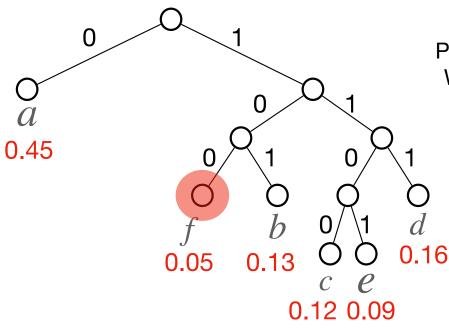


Symbol	а	b	$\boldsymbol{\mathcal{C}}$	d	e	$\int f$
Probability	0.45	0.13	0.12	0.16	0.09	0.05
VLPC	0	101	100	111	1101	1100

Property 1 of optimal code: The (or one) symbol of lowest probability has the longest codeword.

Proof: If not, then we can switch its codeword With another symbol, and get a better code.

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has the longest codeword.

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That will be a contradiction.

Example: Assume f does not have the longest codeword.

Average codeword length =  $1 \times 0.45 + 3 \times 0.13 + 4 \times 0.12 + 3 \times 0.16 + 4 \times 0.09 + 3 \times 0.05$ 

Property 1 of optimal code: The (or one) symbol of lowest probability has the longest codeword.

Proof: If not, then we can switch its codeword With another symbol, and get a better code.

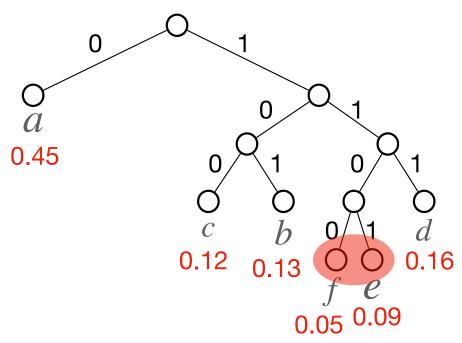
That will be a contradiction.

Example: Assume f does not have the longest codeword.

Now switch a with c.

Average codeword length = 
$$1 \times 0.45 + 3 \times 0.13 + 4 \times 0.12 + 3 \times 0.16 + 4 \times 0.09 + 3 \times 0.05$$
  
After switch, average codeword length =  $1 \times 0.45 + 3 \times 0.13 + 3 \times 0.12 + 3 \times 0.16 + 4 \times 0.09 + 4 \times 0.05$ 

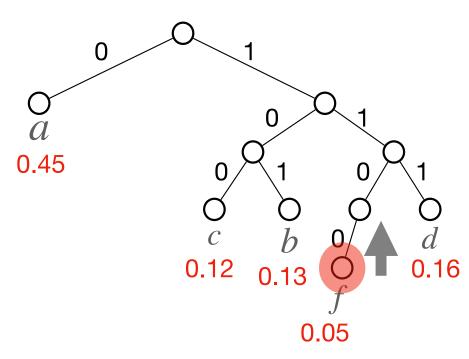
 $(4 \times 0.12 + 3 \times 0.05) - (3 \times 0.12 + 4 \times 0.05) = (4 - 3)(0.12 - 0.05) > 0$  Switch makes code better!



Symbol	a	b	C	d	e	f
Probability	0.45	0.13	0.12	0.16	0.09	0.05
VLPC	0	101	100	111	1101	1100

Property 2 of optimal code:

The symbol of lowest probability and longest codeword has a sibling leaf node.



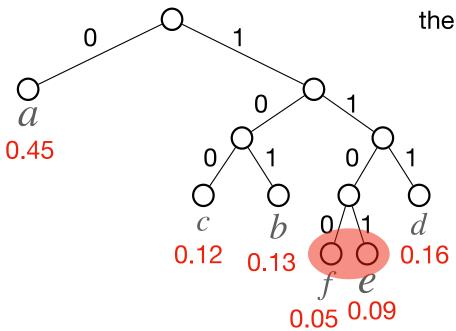
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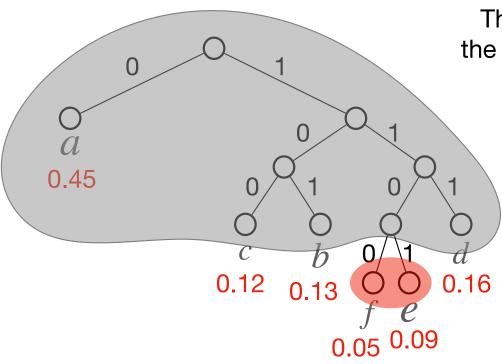
Proof: If not, we can move the codeword up to make it shorter.

Property 3 of optimal code:

There exists an optimal code where the two symbols of lowest probabilities are siblings.



Symbol	a	b	C	d	e	f
Probability	0.45	0.13	0.12	0.16	0.09	0.05
VLPC	0	101	100	111	1101	1100



Property 3 of optimal code:

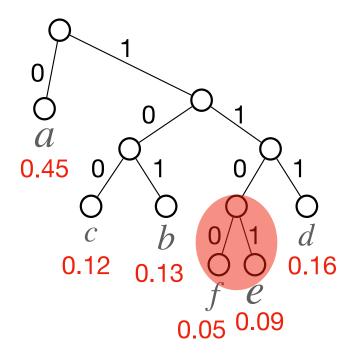
There exists an optimal code where the two symbols of lowest probabilities are siblings.

So we can build an optimal tree this way: first put the two symbols of lowest probabilities as siblings. Then figure out the rest of the tree.

Symbol	а	b	$\mathcal{C}$	d	e	f
Probability	0.45	0.13	0.12	0.16	0.09	0.05

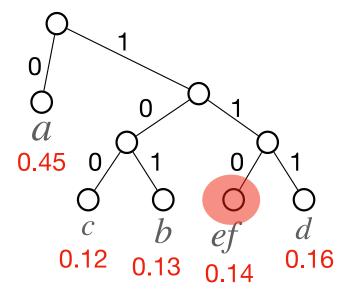
Idea: Once we put the two symbols are siblings, see them as one symbol (node) and combine their probabilities.

## Original Huffman Code



What is the relationship between the two Huffman codes?

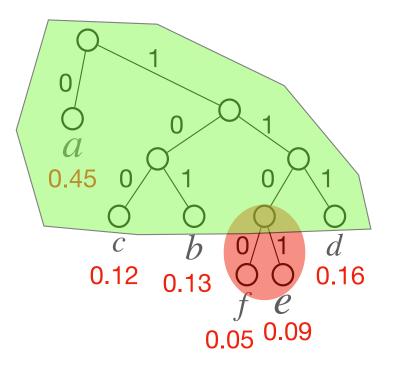
## New Huffman Code



Symbol	а	b	C	d	e f	
Probability	0.45	0.13	0.12	0.16	0.09 0.05	

Symbol	а	b	С	d	ef	
Probability	0.45	0.13	0.12	0.16	0.14	

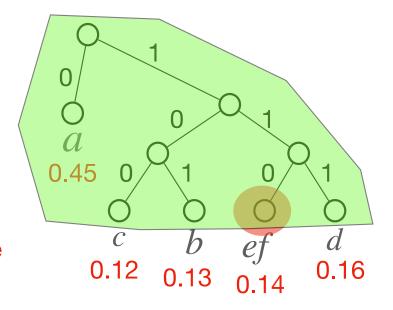
## Original Code



What is the relationship between the two Huffman codes?

If we optimize the new code, we also optimize the old code (and vide versa).

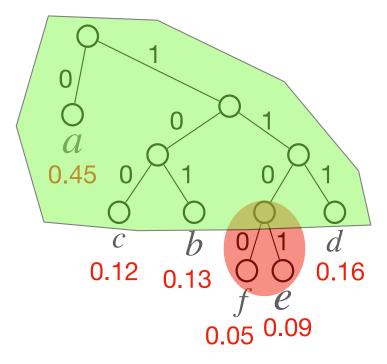
## New Code



Symbol	a	b	C	d	e f	
Probability	0.45	0.13	0.12	0.16	0.09 0.05	

Symbol	а	b	С	d	ef	
Probability	0.45	0.13	0.12	0.16	0.14	

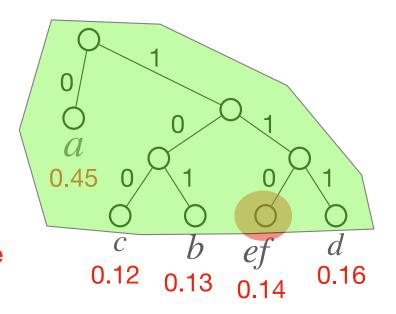
## Original Code



## **New Code**

What is the relationship between the two Huffman codes?

If we optimize the new code, we also optimize the old code (and vide versa).

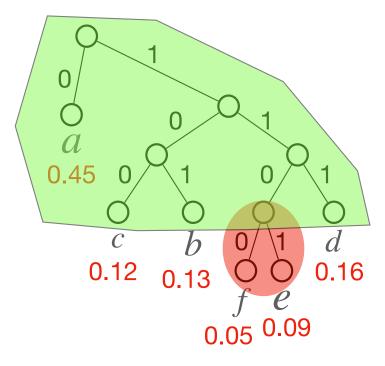


Average codeword length of original code =  $1 \times 0.45 + 3 \times 0.13 + 3 \times 0.12 + 3 \times 0.16 + 4 \times 0.09 + 4 \times 0.05$ Average codeword length of new code =  $1 \times 0.45 + 3 \times 0.13 + 3 \times 0.12 + 3 \times 0.16 + 3 \times (0.09 + 0.05)$ 

Symbol	а	b	С	d	e	f	
Probability	0.45	0.13	0.12	0.16	0.09	0.05	

Symbol	а	b	C	d	ef	
Probability	0.45	0.13	0.12	0.16	0.14	

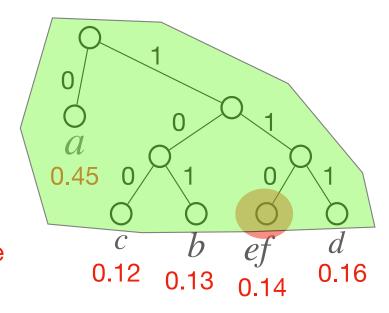
## **Original Code**



# New Code

What is the relationship between the two Huffman codes?

If we optimize the new code, we also optimize the old code (and vide versa).



They differ by  $f_e + f_f$ 

Average codeword length of original code = 
$$1 \times 0.45 + 3 \times 0.13 + 3 \times 0.12 + 3 \times 0.16 + 4 \times 0.09 + 4 \times 0.05$$
  
Average codeword length of new code =  $1 \times 0.45 + 3 \times 0.13 + 3 \times 0.12 + 3 \times 0.16 + 3 \times (0.09 + 0.05)$ 

Symbol	а	b	C	d	e	f	
Probability	0.45	0.13	0.12	0.16	0.09	0.05	

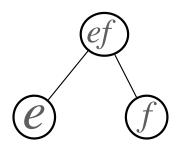
Symbol	а	b	C	d	ef	
Probability	0.45	0.13	0.12	0.16	0.14	

## Idea of Greedy Algorithm:

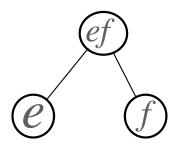
- 1) Make the two symbols of lowest probabilities siblings.
- 2) Combine them into one symbol, and repeat the above process.

Example: Symbol a b c d e f Probability 0.45 0.13 0.12 0.16 0.09 0.05

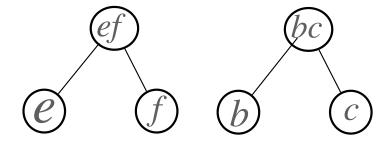
Example: Symbol a b c d e f Probability 0.45 0.13 0.12 0.16 0.09 0.05



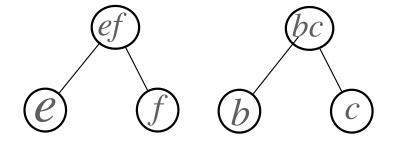
Example:	Symbol	a	b	$\mathcal{C}$	d	ef	
	Probability	0.45	0.13	0.12	0.16	0.14	



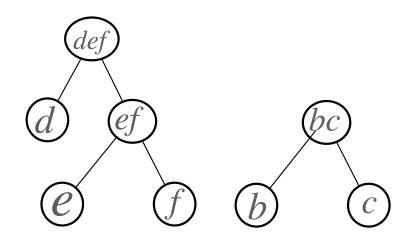
Example:	Symbol	a	b	C	d	ef
	Probability	0.45	0.13	0.12	0.16	0.14



Example:	Symbol	а	bc	d	ef
	Probability	0.45	0.25	0.16	0.14

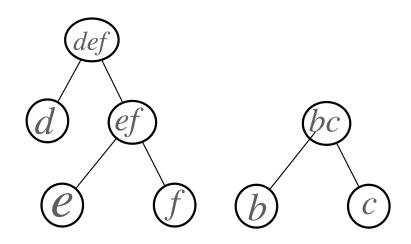


Example: Symbol a bc d efProbability 0.45 0.25 0.16 0.14



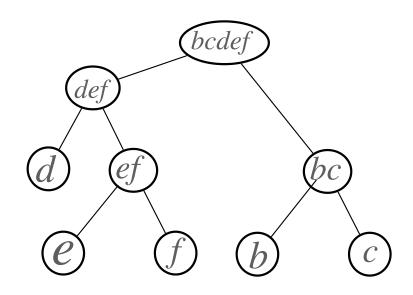
Example: Symbol a bc def

Probability 0.45 0.25 0.3

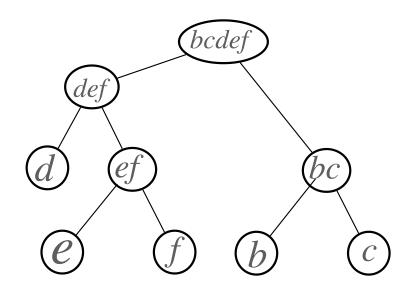


Example:

Symbol	а	bc	def
Probability	0.45	0.25	0.3

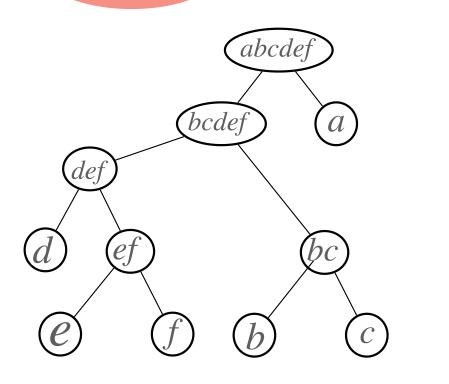


Example: Symbol *a bcdef*Probability 0.45 0.55



Example:

Symbol	a	bcdef
Probability	0.45	0.55



Example: Symbol *a bcdef*Probability 0.45 0.55

