

Algorithms

Lecture 14: Linear Programming (Part 4)

Anxiao (Andrew) Jiang

CH 29.5 Initial Basic Feasible Solution

How to solve an LP:

- 1) Decide if LP is feasible
- 2) If LP is feasible, turn it into a slack form whose basic solution is feasible.
- 3) Run SIMPLEX algorithm.

CH 29.5 Initial Basic Feasible Solution

$$\begin{array}{ll}\text{maximize} & 2x_1 - x_2 \\ \text{s.t.} & \\ & 2x_1 - x_2 \leq 2 \\ & x_1 - 5x_2 \leq -4 \\ & x_1, x_2 \geq 0\end{array}$$

Is this LP feasible?

That is, does this LP have a solution to x_1 and x_2 that satisfies all the constraints?

Note that feasibility has nothing to do with the objective function.

How to solve an LP:

- 1) Decide if LP is feasible
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CH 29.5 Initial Basic Feasible Solution

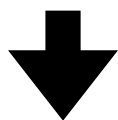
maximize $2x_1 - x_2$

s.t.

$$2x_1 - x_2 \leq 2$$

$$x_1 - 5x_2 \leq -4$$

$$x_1, x_2 \geq 0$$



Slack-form LP:

$$z = 2x_1 - x_2$$

$$x_3 = 2 - 2x_1 + x_2$$

$$x_4 = -4 - x_1 + 5x_2$$

Basic solution: $x_1 = 0, x_2 = 0, x_3 = 2, x_4 = -4$

Is this LP feasible?

That is, does this LP have a solution to x_1 and x_2 that satisfies all the constraints?

Note that feasibility has nothing to do with the objective function.

The slack-form LP does not have a feasible basic solution.

Does it mean the LP is infeasible? Not at all.

How to solve an LP:

- 1) Decide if LP is feasible
- 2) If LP is feasible, turn it into a slack form whose basic solution is feasible.
- 3) Run SIMPLEX algorithm.

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Is this LP feasible?

Whether this LP is feasible or not,
we can always add a “helper” to make it feasible.

How to solve an LP:

- 1) Decide if LP is feasible
- 2) If LP is feasible, turn it into a slack form whose basic solution is feasible.
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CH 29.5 Initial Basic Feasible Solution

Original LP:

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Auxiliary LP:

$$\begin{array}{ll} \text{maximize} & \text{[redacted]} \\ \text{s.t.} & \\ & 2x_1 - x_2 - x_0 \leq 2 \\ & x_1 - 5x_2 - x_0 \leq -4 \\ & x_0, x_1, x_2 \geq 0 \end{array}$$

Is this LP feasible?

Whether this LP is feasible or not,
we can always add a “helper” to make it feasible.

x_0 is our “helper”.

This auxiliary LP is ALWAYS feasible. Why?

How to solve an LP:

- 1) Decide if LP is feasible
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CH 29.5 Initial Basic Feasible Solution

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Auxiliary LP:

$$\begin{array}{ll}\text{maximize} & \text{[redacted]} \\ \text{s.t.} & \\ & 2x_1 - x_2 - x_0 \leq 2 \\ & x_1 - 5x_2 - x_0 \leq -4 \\ & x_0, x_1, x_2 \geq 0\end{array}$$

Is this LP feasible?

Whether this LP is feasible or not,
we can always add a “helper” to make it feasible.

x_0 is our “helper”.

This auxiliary LP is ALWAYS feasible. Why?

Because we can always make x_0 sufficiently large to satisfy all constraints.

For example, here we can make $x_0 = 4$, and a feasible solution is

$$x_0 = 4, x_1 = 0, x_2 = 0$$

How to solve an LP:

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CH 29.5 Initial Basic Feasible Solution

Original LP:

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Auxiliary LP:

$$\begin{array}{ll}\text{maximize} & x_0 \\ \text{s.t.} & \\ & 2x_1 - x_2 - x_0 \leq 2 \\ & x_1 - 5x_2 - x_0 \leq -4 \\ & x_0, x_1, x_2 \geq 0\end{array}$$

Is this LP feasible?

Whether this LP is feasible or not,
we can always add a “helper” to make it feasible.

x_0 is our “helper”. This auxiliary LP is ALWAYS feasible.

But how large does x_0 have to be to make this auxiliary LP feasible?

If x_0 can be 0, then no “help” is needed, which means the original LP is feasible.

If x_0 has to be greater than 0, then the original LP is not feasible.

Which case is it? Let's find out.

How to solve an LP:

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Is this LP feasible?



Auxiliary LP:

$$\begin{array}{ll}\text{maximize} & -x_0 \\ \text{s.t.} & \\ & 2x_1 - x_2 - x_0 \leq 2 \\ & x_1 - 5x_2 - x_0 \leq -4 \\ & x_0, x_1, x_2 \geq 0\end{array}$$

How small can x_0 be?

In particular, can x_0 be as small as 0?

How to solve an LP:

- 1) Decide if LP is feasible
- 2) If LP is feasible, turn it into a slack form whose basic solution is feasible.
- 3) Run SIMPLEX algorithm.

CH 29.5 Initial Basic Feasible Solution

Theorem: Let L be an LP

maximize $c_1x_1 + c_2x_2 + \cdots + c_nx_n$

s.t.

$$a_{1,1}x_1 + a_{1,2}x_2 + \cdots + a_{1,n}x_n \leq b_1$$

$$a_{2,1}x_1 + a_{2,2}x_2 + \cdots + a_{2,n}x_n \leq b_2$$

$$\vdots$$

$$a_{m,1}x_1 + a_{m,2}x_2 + \cdots + a_{m,n}x_n \leq b_m$$

$$x_1, x_2, \cdots, x_n \geq 0$$

Let L_{aux} be its auxiliary LP

maximize $-x_0$

s.t.

$$a_{1,1}x_1 + a_{1,2}x_2 + \cdots + a_{1,n}x_n - x_0 \leq b_1$$

$$a_{2,1}x_1 + a_{2,2}x_2 + \cdots + a_{2,n}x_n - x_0 \leq b_2$$

$$\vdots$$

$$a_{m,1}x_1 + a_{m,2}x_2 + \cdots + a_{m,n}x_n - x_0 \leq b_m$$

$$x_0, x_1, x_2, \cdots, x_n \geq 0$$

Then L is feasible if and only if the optimal objective value for L_{aux} is 0.

CH 29.5 Initial Basic Feasible Solution

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maximize $c_1x_1 + c_2x_2 + \cdots + c_nx_n$

s.t.

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$$a_{2,1}x_1 + a_{2,2}x_2 + \cdots + a_{2,n}x_n \leq b_2$$

$$\vdots$$

$$a_{m,1}x_1 + a_{m,2}x_2 + \cdots + a_{m,n}x_n \leq b_m$$

$$x_1, x_2, \cdots, x_n \geq 0$$

Let L_{aux} be its auxiliary LP

maximize $-x_0$

s.t.

$$a_{1,1}x_1 + a_{1,2}x_2 + \cdots + a_{1,n}x_n - x_0 \leq b_1$$


$$a_{2,1}x_1 + a_{2,2}x_2 + \cdots + a_{2,n}x_n - x_0 \leq b_2$$

$$\vdots$$

$$a_{m,1}x_1 + a_{m,2}x_2 + \cdots + a_{m,n}x_n - x_0 \leq b_m$$

$$x_0, x_1, x_2, \cdots, x_n \geq 0$$

Then L is feasible if and only if the optimal objective value for L_{aux} is 0.

Proof:  If L has a feasible solution (x_1, x_2, \cdots, x_n) ,
then L_{aux} has an optimal solution $(x_0 = 0, x_1, x_2, \cdots, x_n)$.

CH 29.5 Initial Basic Feasible Solution

Theorem: Let L be an LP

maximize $c_1x_1 + c_2x_2 + \cdots + c_nx_n$

s.t.

$$a_{1,1}x_1 + a_{1,2}x_2 + \cdots + a_{1,n}x_n \leq b_1$$

$$a_{2,1}x_1 + a_{2,2}x_2 + \cdots + a_{2,n}x_n \leq b_2$$

\vdots

$$a_{m,1}x_1 + a_{m,2}x_2 + \cdots + a_{m,n}x_n \leq b_m$$

$$x_1, x_2, \cdots, x_n \geq 0$$

Let L_{aux} be its auxiliary LP

maximize $-x_0$

s.t.

$$a_{1,1}x_1 + a_{1,2}x_2 + \cdots + a_{1,n}x_n - x_0 \leq b_1$$

$$a_{2,1}x_1 + a_{2,2}x_2 + \cdots + a_{2,n}x_n - x_0 \leq b_2$$

\vdots

$$a_{m,1}x_1 + a_{m,2}x_2 + \cdots + a_{m,n}x_n - x_0 \leq b_m$$

$$x_0, x_1, x_2, \cdots, x_n \geq 0$$

Then L is feasible if and only if the optimal objective value for L_{aux} is 0.

Proof:



If L_{aux} has an optimal solution $(x_0 = 0, x_1, x_2, \cdots, x_n)$,

then L has a feasible solution (x_1, x_2, \cdots, x_n) .

CH 29.5 Initial Basic Feasible Solution

Original LP:

$$\begin{array}{ll}\text{maximize} & 2x_1 - x_2 \\ \text{s.t.} & \\ & 2x_1 - x_2 \leq 2 \\ & x_1 - 5x_2 \leq -4 \\ & x_1, x_2 \geq 0\end{array}$$

Is this LP feasible?

Let's solve the Auxiliary LP to find out.

Auxiliary LP:

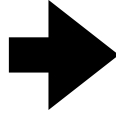
$$\begin{array}{ll}\text{maximize} & -x_0 \\ \text{s.t.} & \\ & 2x_1 - x_2 - x_0 \leq 2 \\ & x_1 - 5x_2 - x_0 \leq -4 \\ & x_0, x_1, x_2 \geq 0\end{array}$$

How small can x_0 be?

In particular, can x_0 be as small as 0?

Auxiliary LP:

$$\begin{array}{ll}\text{maximize} & -x_0 \\ \text{s.t.} & \\ & 2x_1 - x_2 - x_0 \leq 2 \\ & x_1 - 5x_2 - x_0 \leq -4 \\ & x_0, x_1, x_2 \geq 0\end{array}$$



$$z = -x_0$$

$$x_3 = 2 + x_0 - 2x_1 + x_2$$

$$x_4 = -4 + x_0 - x_1 + 5x_2$$

Auxiliary LP:

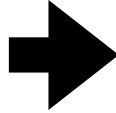
maximize $-x_0$

s.t.

$$2x_1 - x_2 - x_0 \leq 2$$

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$$z = -x_0$$

$$x_3 = 2 + x_0 - 2x_1 + x_2$$

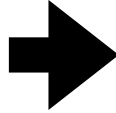
$$x_4 = -4 + x_0 - x_1 + 5x_2$$

Do a pivot:

- 1) Move x_0 to left.
- 2) Pick a basic variable whose constant term is the most negative, and move it to right.

Auxiliary LP:

$$\begin{array}{ll}\text{maximize} & -x_0 \\ \text{s.t.} & \\ & 2x_1 - x_2 - x_0 \leq 2 \\ & x_1 - 5x_2 - x_0 \leq -4 \\ & x_0, x_1, x_2 \geq 0\end{array}$$



$$\begin{array}{l}z = -x_0 \\ x_3 = 2 + x_0 - 2x_1 + x_2 \\ x_4 = -4 + x_0 - x_1 + 5x_2\end{array}$$

Most negative

Do a pivot:

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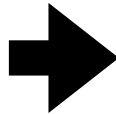
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s.t.

$$2x_1 - x_2 - x_0 \leq 2$$

$$x_1 - 5x_2 - x_0 \leq -4$$

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$$z = -x_0$$

$$x_3 = 2 + x_0 - 2x_1 + x_2$$

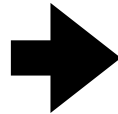
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Do a pivot:

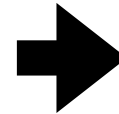
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Auxiliary LP:

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$$\begin{array}{l} z = -x_0 \\ x_3 = 2 + x_0 - 2x_1 + x_2 \\ x_4 = -4 + x_0 - x_1 + 5x_2 \end{array}$$



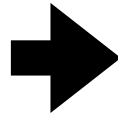
$$\begin{array}{l} z = -4 - x_1 + 5x_2 - x_4 \\ x_3 = 6 - x_1 - 4x_2 + x_4 \\ x_0 = 4 + x_1 - 5x_2 + x_4 \end{array}$$

Do a pivot:

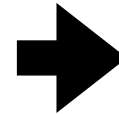
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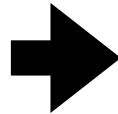
Positive

Do a pivot:

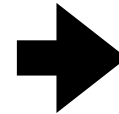
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Auxiliary LP:

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$$\begin{array}{l} z = -x_0 \\ x_3 = 2 + x_0 - 2x_1 + x_2 \\ x_4 = -4 + x_0 - x_1 + 5x_2 \end{array}$$



$$\begin{array}{l} z = -4 - x_1 + 5x_2 - x_4 \\ x_3 = 6 - x_1 - 4x_2 + x_4 \\ x_0 = 4 + x_1 - 5x_2 + x_4 \end{array}$$

Where does this 6 come from?

$$2 + 4 = 6$$

Do a pivot:

- 1) Move x_0 to left.
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Auxiliary LP:

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What if

$$\begin{array}{l} z = -x_0 \\ x_3 = -2 + x_0 - 2x_1 + x_2 \\ x_4 = -4 + x_0 - x_1 + 5x_2 \end{array}$$
$$\begin{array}{l} z = -4 - x_1 + 5x_2 - x_4 \\ x_3 = 6 - x_1 - 4x_2 + x_4 \\ x_0 = 4 + x_1 - 5x_2 + x_4 \end{array}$$

$-2 + 4 = 2$

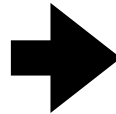
Still non-negative!

Do a pivot:

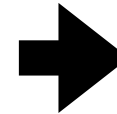
- 1) Move x_0 to left.
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Auxiliary LP:

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$$\begin{array}{l}z = -x_0 \\ x_3 = 2 + x_0 - 2x_1 + x_2 \\ x_4 = -4 + x_0 - x_1 + 5x_2\end{array}$$



$$\begin{array}{l}z = -4 - x_1 + 5x_2 - x_4 \\ x_3 = 6 - x_1 - 4x_2 + x_4 \\ x_0 = 4 + x_1 - 5x_2 + x_4\end{array}$$

All constant terms are non-negative.

Its basic solution is feasible,
So we can run the SIMPLEX Algorithm now!

Do a pivot:

- 1) Move x_0 to left.
- 2) Pick a basic variable whose constant term is the most negative, and move it to right.

Auxiliary LP:

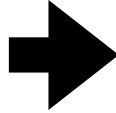
maximize $-x_0$

s.t.

$$2x_1 - x_2 - x_0 \leq 2$$

$$x_1 - 5x_2 - x_0 \leq -4$$

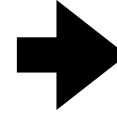
$$x_0, x_1, x_2 \geq 0$$



$$z = -x_0$$

$$x_3 = 2 + x_0 - 2x_1 + x_2$$

$$x_4 = -4 + x_0 - x_1 + 5x_2$$



$$z = -4 - x_1 + 5x_2 - x_4$$

$$x_3 = 6 - x_1 - 4x_2 + x_4$$

$$x_0 = 4 + x_1 - 5x_2 + x_4$$

Basic solution:

$$x_0 = 4, x_1 = 0, x_2 = 0, x_3 = 6, x_4 = 0$$

Auxiliary LP:

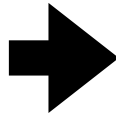
maximize $-x_0$

s.t.

$$2x_1 - x_2 - x_0 \leq 2$$

$$x_1 - 5x_2 - x_0 \leq -4$$

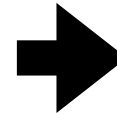
$$x_0, x_1, x_2 \geq 0$$



$$z = -x_0$$

$$x_3 = 2 + x_0 - 2x_1 + x_2$$

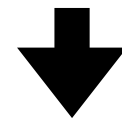
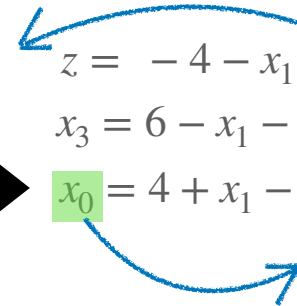
$$x_4 = -4 + x_0 - x_1 + 5x_2$$



$$z = -4 - x_1 + 5x_2 - x_4$$

$$x_3 = 6 - x_1 - 4x_2 + x_4$$

$$x_0 = 4 + x_1 - 5x_2 + x_4$$



Run
SIMPLEX
Algorithm

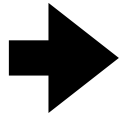
$$z = -x_0$$

$$x_2 = \frac{4}{5} - \frac{1}{5}x_0 + \frac{1}{5}x_1 + \frac{1}{5}x_4$$

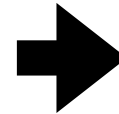
$$x_3 = \frac{14}{5} + \frac{4}{5}x_0 - \frac{9}{5}x_1 + \frac{1}{5}x_4$$

Auxiliary LP:

$$\begin{array}{ll}\text{maximize} & -x_0 \\ \text{s.t.} & \\ & 2x_1 - x_2 - x_0 \leq 2 \\ & x_1 - 5x_2 - x_0 \leq -4 \\ & x_0, x_1, x_2 \geq 0\end{array}$$



$$\begin{aligned}z &= -x_0 \\ x_3 &= 2 + x_0 - 2x_1 + x_2 \\ x_4 &= -4 + x_0 - x_1 + 5x_2\end{aligned}$$



$$\begin{aligned}z &= -4 - x_1 + 5x_2 - x_4 \\ x_3 &= 6 - x_1 - 4x_2 + x_4 \\ x_0 &= 4 + x_1 - 5x_2 + x_4\end{aligned}$$

Run
SIMPLEX
Algorithm

Original LP:

$$\begin{array}{ll}\text{maximize} & 2x_1 - x_2 \\ \text{s.t.} & \\ & 2x_1 - x_2 \leq 2 \\ & x_1 - 5x_2 \leq -4 \\ & x_1, x_2 \geq 0\end{array}$$

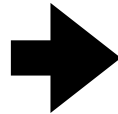
It is feasible!

$$\begin{aligned}z &= -x_0 \\ x_2 &= \frac{4}{5} - \frac{1}{5}x_0 + \frac{1}{5}x_1 + \frac{1}{5}x_4 \\ x_3 &= \frac{14}{5} + \frac{4}{5}x_0 - \frac{9}{5}x_1 + \frac{1}{5}x_4\end{aligned}$$

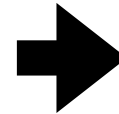
Optimal Objective Value = 0

Auxiliary LP:

$$\begin{array}{ll}\text{maximize} & -x_0 \\ \text{s.t.} & \\ & 2x_1 - x_2 - x_0 \leq 2 \\ & x_1 - 5x_2 - x_0 \leq -4 \\ & x_0, x_1, x_2 \geq 0\end{array}$$



$$\begin{array}{l}z = -x_0 \\ x_3 = 2 + x_0 - 2x_1 + x_2 \\ x_4 = -4 + x_0 - x_1 + 5x_2\end{array}$$



$$\begin{array}{l}z = -4 - x_1 + 5x_2 - x_4 \\ x_3 = 6 - x_1 - 4x_2 + x_4 \\ x_0 = 4 + x_1 - 5x_2 + x_4\end{array}$$

Run
SIMPLEX
Algorithm

Now we need to solve the Original LP. But how?

Original LP:

$$\begin{array}{ll}\text{maximize} & 2x_1 - x_2 \\ \text{s.t.} & \\ & 2x_1 - x_2 \leq 2 \\ & x_1 - 5x_2 \leq -4 \\ & x_1, x_2 \geq 0\end{array}$$

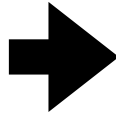
It is feasible!

$$\begin{array}{l}z = -x_0 \\ x_2 = \frac{4}{5} - \frac{1}{5}x_0 + \frac{1}{5}x_1 + \frac{1}{5}x_4 \\ x_3 = \frac{14}{5} + \frac{4}{5}x_0 - \frac{9}{5}x_1 + \frac{1}{5}x_4\end{array}$$

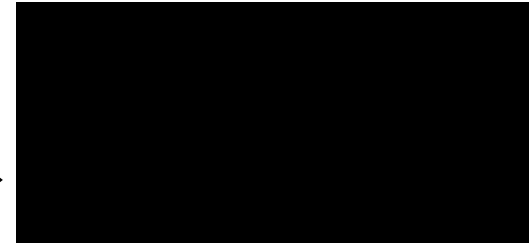
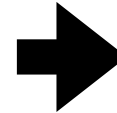
Optimal Objective Value = 0

Auxiliary LP:

$$\begin{array}{ll}\text{maximize} & -x_0 \\ \text{s.t.} & \\ & 2x_1 - x_2 - x_0 \leq 2 \\ & x_1 - 5x_2 - x_0 \leq -4 \\ & x_0, x_1, x_2 \geq 0\end{array}$$



$$\begin{array}{l}z = -x_0 \\ x_3 = 2 + x_0 - 2x_1 + x_2 \\ x_4 = -4 + x_0 - x_1 + 5x_2\end{array}$$



Now we need to solve the Original LP. But how?

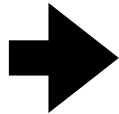
Original LP:

$$\begin{array}{ll}\text{maximize} & 2x_1 - x_2 \\ \text{s.t.} & \\ & 2x_1 - x_2 \leq 2 \\ & x_1 - 5x_2 \leq -4 \\ & x_1, x_2 \geq 0\end{array}$$

$$\begin{array}{l}z = -x_0 \\ x_2 = \frac{4}{5} - \frac{1}{5}x_0 + \frac{1}{5}x_1 + \frac{1}{5}x_4 \\ x_3 = \frac{14}{5} + \frac{4}{5}x_0 - \frac{9}{5}x_1 + \frac{1}{5}x_4\end{array}$$

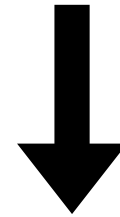
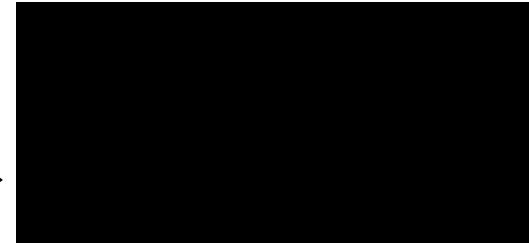
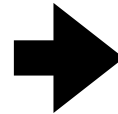
Auxiliary LP:

$$\begin{array}{ll}\text{maximize} & -x_0 \\ \text{s.t.} & \\ & 2x_1 - x_2 - x_0 \leq 2 \\ & x_1 - 5x_2 - x_0 \leq -4 \\ & x_0, x_1, x_2 \geq 0\end{array}$$



$$z = -x_0$$

$$\begin{array}{l}x_3 = 2 + x_0 - 2x_1 + x_2 \\ x_4 = -4 + x_0 - x_1 + 5x_2\end{array}$$



Now we need to solve the Original LP. But how?

Original LP:

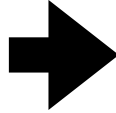
$$\begin{array}{ll}\text{maximize} & 2x_1 - x_2 \\ \text{s.t.} & \\ & 2x_1 - x_2 \leq 2 \\ & x_1 - 5x_2 \leq -4 \\ & x_1, x_2 \geq 0\end{array}$$

$$z = -x_0$$

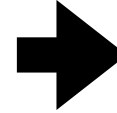
$$\begin{array}{l}x_2 = \frac{4}{5} - \frac{1}{5}x_0 + \frac{1}{5}x_1 + \frac{1}{5}x_4 \\ x_3 = \frac{14}{5} + \frac{4}{5}x_0 - \frac{9}{5}x_1 + \frac{1}{5}x_4\end{array}$$

Auxiliary LP:

$$\begin{array}{ll}\text{maximize} & -x_0 \\ \text{s.t.} & \\ & 2x_1 - x_2 - x_0 \leq 2 \\ & x_1 - 5x_2 - x_0 \leq -4 \\ & x_0, x_1, x_2 \geq 0\end{array}$$



$$\begin{array}{l}z = -x_0 \\ x_3 = 2 + x_0 - 2x_1 + x_2 \\ x_4 = -4 + x_0 - x_1 + 5x_2\end{array}$$



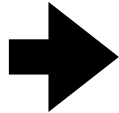
$$\begin{array}{l}z = -x_0 \\ x_2 = \frac{4}{5} - \frac{1}{5}x_0 + \frac{1}{5}x_1 + \frac{1}{5}x_4 \\ x_3 = \frac{14}{5} + \frac{4}{5}x_0 - \frac{9}{5}x_1 + \frac{1}{5}x_4\end{array}$$

Original LP:

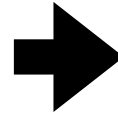
$$\begin{array}{ll}\text{maximize} & 2x_1 - x_2 \\ \text{s.t.} & \\ & 2x_1 - x_2 \leq 2 \\ & x_1 - 5x_2 \leq -4 \\ & x_1, x_2 \geq 0\end{array}$$

Auxiliary LP:

$$\begin{array}{ll}\text{maximize} & -x_0 \\ \text{s.t.} & \\ & 2x_1 - x_2 - x_0 \leq 2 \\ & x_1 - 5x_2 - x_0 \leq -4 \\ & x_0, x_1, x_2 \geq 0\end{array}$$



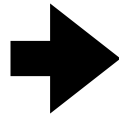
$$\begin{array}{l}z = -x_0 \\ x_3 = 2 + x_0 - 2x_1 + x_2 \\ x_4 = -4 + x_0 - x_1 + 5x_2\end{array}$$



$$\begin{array}{l}z = -x_0 \\ x_2 = \frac{4}{5} - \frac{1}{5}x_0 + \frac{1}{5}x_1 + \frac{1}{5}x_4 \\ x_3 = \frac{14}{5} + \frac{4}{5}x_0 - \frac{9}{5}x_1 + \frac{1}{5}x_4\end{array}$$

Original LP:

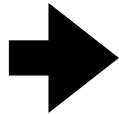
$$\begin{array}{ll}\text{maximize} & 2x_1 - x_2 \\ \text{s.t.} & \\ & 2x_1 - x_2 \leq 2 \\ & x_1 - 5x_2 \leq -4 \\ & x_1, x_2 \geq 0\end{array}$$



$$\begin{array}{l}z = 2x_1 - x_2 \\ x_3 = 2 - 2x_1 + x_2 \\ x_4 = -4 - x_1 + 5x_2\end{array}$$

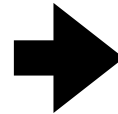
Auxiliary LP:

$$\begin{array}{ll}\text{maximize} & -x_0 \\ \text{s.t.} & \\ & 2x_1 - x_2 - x_0 \leq 2 \\ & x_1 - 5x_2 - x_0 \leq -4 \\ & x_0, x_1, x_2 \geq 0\end{array}$$



$$z = -x_0$$

$$\begin{array}{l}x_3 = 2 + x_0 - 2x_1 + x_2 \\ x_4 = -4 + x_0 - x_1 + 5x_2\end{array}$$



$$z = -x_0$$

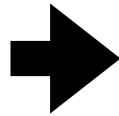
$$\begin{array}{l}x_2 = \frac{4}{5} - \frac{1}{5}x_0 + \frac{1}{5}x_1 + \frac{1}{5}x_4 \\ x_3 = \frac{14}{5} + \frac{4}{5}x_0 - \frac{9}{5}x_1 + \frac{1}{5}x_4\end{array}$$



The constraints
are equivalent
when $x_0 = 0$

Original LP:

$$\begin{array}{ll}\text{maximize} & 2x_1 - x_2 \\ \text{s.t.} & \\ & 2x_1 - x_2 \leq 2 \\ & x_1 - 5x_2 \leq -4 \\ & x_1, x_2 \geq 0\end{array}$$

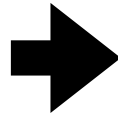


$$z = 2x_1 - x_2$$

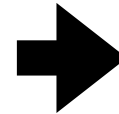
$$\begin{array}{l}x_3 = 2 - 2x_1 + x_2 \\ x_4 = -4 - x_1 + 5x_2\end{array}$$

Auxiliary LP:

$$\begin{array}{ll}
 \text{maximize} & -x_0 \\
 \text{s.t.} & \\
 & 2x_1 - x_2 - x_0 \leq 2 \\
 & x_1 - 5x_2 - x_0 \leq -4 \\
 & x_0, x_1, x_2 \geq 0
 \end{array}$$



$$\begin{array}{l}
 z = -x_0 \\
 x_3 = 2 + x_0 - 2x_1 + x_2 \\
 x_4 = -4 + x_0 - x_1 + 5x_2
 \end{array}$$



$$\begin{array}{l}
 z = -x_0 \\
 x_2 = \frac{4}{5} - \frac{1}{5}x_0 + \frac{1}{5}x_1 + \frac{1}{5}x_4 \\
 x_3 = \frac{14}{5} + \frac{4}{5}x_0 - \frac{9}{5}x_1 + \frac{1}{5}x_4
 \end{array}$$



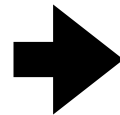
The constraints
are equivalent
when $x_0 = 0$



The constraints
are equivalent
when $x_0 = 0$

Original LP:

$$\begin{array}{ll}
 \text{maximize} & 2x_1 - x_2 \\
 \text{s.t.} & \\
 & 2x_1 - x_2 \leq 2 \\
 & x_1 - 5x_2 \leq -4 \\
 & x_1, x_2 \geq 0
 \end{array}$$



$$z = 2x_1 - x_2$$

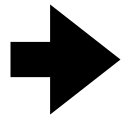
$$\begin{array}{l}
 x_3 = 2 - 2x_1 + x_2 \\
 x_4 = -4 - x_1 + 5x_2
 \end{array}$$



$$\begin{array}{l}
 x_2 = \frac{4}{5} + \frac{1}{5}x_1 + \frac{1}{5}x_4 \\
 x_3 = \frac{14}{5} - \frac{9}{5}x_1 + \frac{1}{5}x_4
 \end{array}$$

Original LP:

$$\begin{array}{ll}\text{maximize} & 2x_1 - x_2 \\ \text{s.t.} & \\ & 2x_1 - x_2 \leq 2 \\ & x_1 - 5x_2 \leq -4 \\ & x_1, x_2 \geq 0\end{array}$$



$$z = 2x_1 - x_2$$

$$\begin{array}{ll}x_3 = 2 & - 2x_1 + x_2 \\ x_4 = -4 & - x_1 + 5x_2\end{array}$$

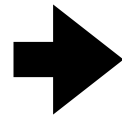


$$\begin{aligned}z &= 2x_1 - x_2 \\ &= 2x_1 - \left(\frac{4}{5} + \frac{1}{5}x_1 + \frac{1}{5}x_4\right) \\ &= -\frac{4}{5} + \frac{9}{5}x_1 - \frac{1}{5}x_4\end{aligned}$$

$$\begin{array}{ll}x_2 = \frac{4}{5} & + \frac{1}{5}x_1 + \frac{1}{5}x_4 \\ x_3 = \frac{14}{5} & - \frac{9}{5}x_1 + \frac{1}{5}x_4\end{array}$$

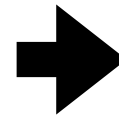
Original LP:

$$\begin{array}{ll}\text{maximize} & 2x_1 - x_2 \\ \text{s.t.} & \\ & 2x_1 - x_2 \leq 2 \\ & x_1 - 5x_2 \leq -4 \\ & x_1, x_2 \geq 0\end{array}$$



$$z = 2x_1 - x_2$$

$$\begin{array}{ll}x_3 = 2 & - 2x_1 + x_2 \\ x_4 = -4 & - x_1 + 5x_2\end{array}$$



$$z = -\frac{4}{5} + \frac{9}{5}x_1 - \frac{1}{5}x_4$$

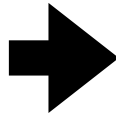
$$\begin{array}{ll}x_2 = \frac{4}{5} & + \frac{1}{5}x_1 + \frac{1}{5}x_4 \\ x_3 = \frac{14}{5} & - \frac{9}{5}x_1 + \frac{1}{5}x_4\end{array}$$

This slack-form LP
has a feasible basic solution.

Now we can solve it using
the SIMPLEX Algorithm.

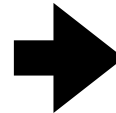
Auxiliary LP:

$$\begin{array}{ll}\text{maximize} & -x_0 \\ \text{s.t.} & \\ & 2x_1 - x_2 - x_0 \leq 2 \\ & x_1 - 5x_2 - x_0 \leq -4 \\ & x_0, x_1, x_2 \geq 0\end{array}$$



$$z = -x_0$$

$$\begin{array}{l}x_3 = 2 + x_0 - 2x_1 + x_2 \\ x_4 = -4 + x_0 - x_1 + 5x_2\end{array}$$



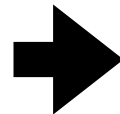
$$z = -x_0$$

$$\begin{array}{l}x_2 = \frac{4}{5} - \frac{1}{5}x_0 + \frac{1}{5}x_1 + \frac{1}{5}x_4 \\ x_3 = \frac{14}{5} + \frac{4}{5}x_0 - \frac{9}{5}x_1 + \frac{1}{5}x_4\end{array}$$

Last question: What if x_0 are a basic variable in these constraint equations?

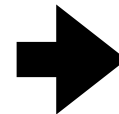
Original LP:

$$\begin{array}{ll}\text{maximize} & 2x_1 - x_2 \\ \text{s.t.} & \\ & 2x_1 - x_2 \leq 2 \\ & x_1 - 5x_2 \leq -4 \\ & x_1, x_2 \geq 0\end{array}$$



$$z = 2x_1 - x_2$$

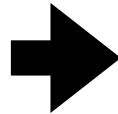
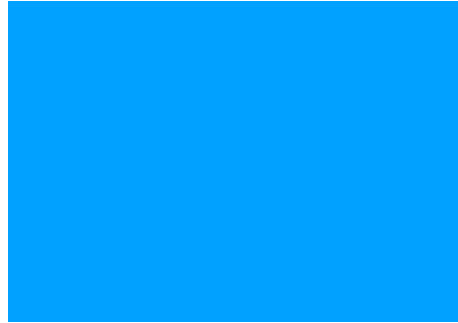
$$\begin{array}{l}x_3 = 2 - 2x_1 + x_2 \\ x_4 = -4 - x_1 + 5x_2\end{array}$$



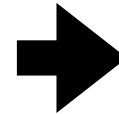
$$z = -\frac{4}{5} + \frac{9}{5}x_1 - \frac{1}{5}x_4$$

$$\begin{array}{l}x_2 = \frac{4}{5} + \frac{1}{5}x_1 + \frac{1}{5}x_4 \\ x_3 = \frac{14}{5} - \frac{9}{5}x_1 + \frac{1}{5}x_4\end{array}$$

Auxiliary LP:



$$z = -x_0$$



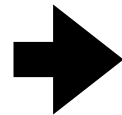
$$z = \dots$$

$$x_0 = a_0 + a_1x_1 + a_2x_2 + a_3x_3$$

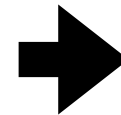
$$x_4 = b_0 + b_1x_1 + b_2x_2 + b_3x_3$$

For example ...

Original LP:



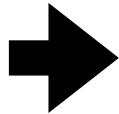
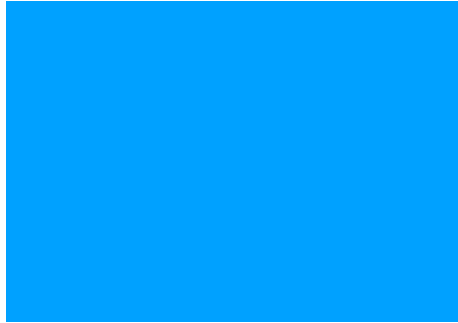
$$z = \dots$$



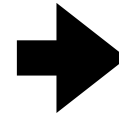
$$z = \dots$$



Auxiliary LP:



$$z = -x_0$$



$$z = \dots$$

$$x_0 = a_0 + a_1x_1 + a_2x_2 + a_3x_3$$

$$x_4 = b_0 + b_1x_1 + b_2x_2 + b_3x_3$$

Optimal solution:

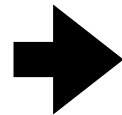
$$x_0 = a_0, x_1 = 0, x_2 = 0, x_3 = 0, x_4 = b_0$$

Optimal objective value:

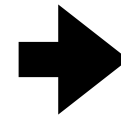
$$z = -x_0 = 0$$

For example ...

Original LP:



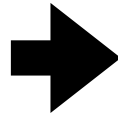
$$z = \dots$$



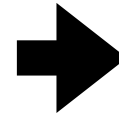
$$z = \dots$$



Auxiliary LP:



$$z = -x_0$$



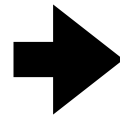
$$z = \dots$$

$$x_0 = a_0 + a_1x_1 + a_2x_2 + a_3x_3$$

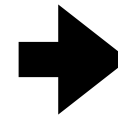
$$x_4 = b_0 + b_1x_1 + b_2x_2 + b_3x_3$$

$$a_0 = 0 \left\{ \begin{array}{l} \text{Optimal solution:} \\ x_0 = a_0, x_1 = 0, x_2 = 0, x_3 = 0, x_4 = b_0 \\ \text{Optimal objective value:} \\ z = -x_0 = 0 \end{array} \right.$$

Original LP:



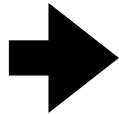
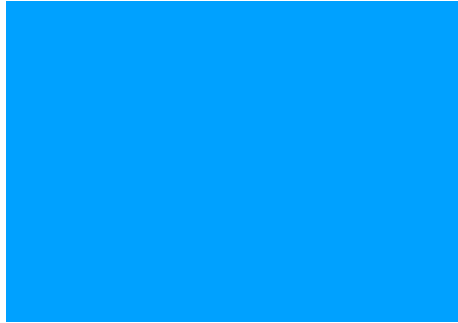
$$z = \dots$$



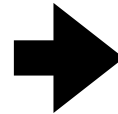
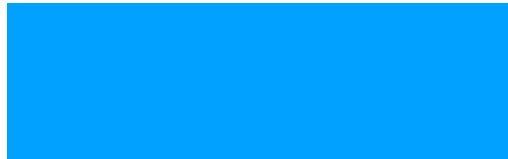
$$z = \dots$$



Auxiliary LP:



$$z = -x_0$$



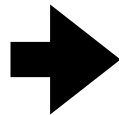
$$z = \dots$$

$$x_0 = a_1x_1 + a_2x_2 + a_3x_3$$

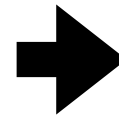
$$x_4 = b_0 + b_1x_1 + b_2x_2 + b_3x_3$$

$$a_0 = 0 \left\{ \begin{array}{l} \text{Optimal solution:} \\ x_0 = a_0, x_1 = 0, x_2 = 0, x_3 = 0, x_4 = b_0 \\ \text{Optimal objective value:} \\ z = -x_0 = 0 \end{array} \right.$$

Original LP:



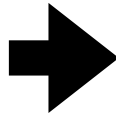
$$z = \dots$$



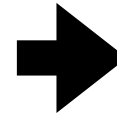
$$z = \dots$$



Auxiliary LP:



$$z = -x_0$$

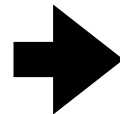


$$z = \dots$$

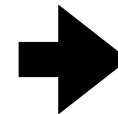
$$\boxed{\begin{array}{l} x_0 = a_1 x_1 + a_2 x_2 + a_3 x_3 \\ x_4 = b_0 + b_1 x_1 + b_2 x_2 + b_3 x_3 \end{array}}$$

$$a_0 = 0 \left\{ \begin{array}{l} \text{Optimal solution:} \\ x_0 = a_0, x_1 = 0, x_2 = 0, x_3 = 0, x_4 = b_0 \\ \text{Optimal objective value:} \\ z = -x_0 = 0 \end{array} \right.$$

Original LP:



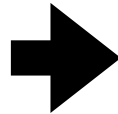
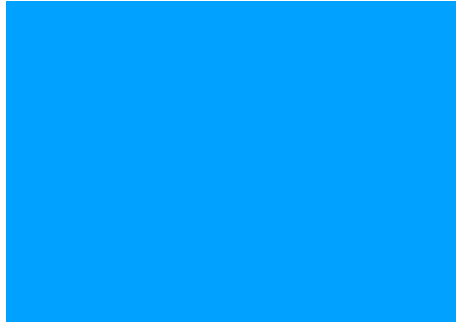
$$z = \dots$$



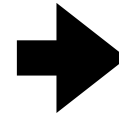
$$z = \dots$$



Auxiliary LP:



$$z = -x_0$$



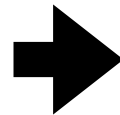
$$z = \dots$$

$$x_1 = \frac{1}{a_1}x_0 - \frac{a_2}{a_1}x_2 - \frac{a_3}{a_1}x_3$$

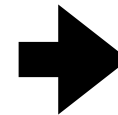
$$x_4 = b_0 + b_1x_1 + b_2x_2 + b_3x_3$$

$$a_0 = 0 \left\{ \begin{array}{l} \text{Optimal solution:} \\ x_0 = a_0, x_1 = 0, x_2 = 0, x_3 = 0, x_4 = b_0 \\ \text{Optimal objective value:} \\ z = -x_0 = 0 \end{array} \right.$$

Original LP:



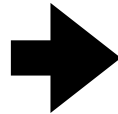
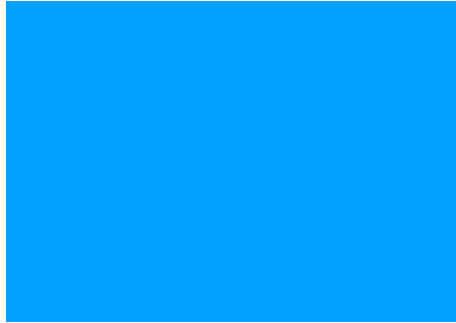
$$z = \dots$$



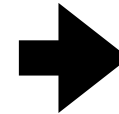
$$z = \dots$$



Auxiliary LP:



$$z = -x_0$$



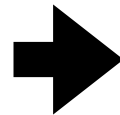
$$z = \dots$$

$$x_1 = \frac{1}{a_1}x_0 - \frac{a_2}{a_1}x_2 - \frac{a_3}{a_1}x_3$$

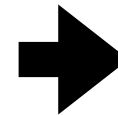
$$x_4 = b_0 + b_1x_1 + b_2x_2 + b_3x_3$$

$$a_0 = 0 \left\{ \begin{array}{l} \text{Optimal solution:} \\ x_0 = a_0, x_1 = 0, x_2 = 0, x_3 = 0, x_4 = b_0 \\ \text{Optimal objective value:} \\ z = -x_0 = 0 \end{array} \right.$$

Original LP:



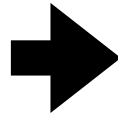
$$z = \dots$$



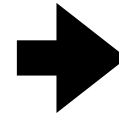
$$z = \dots$$



Auxiliary LP:



$$z = -x_0$$



The constants here are still 0 and $b_0 \geq 0$

$$z = \dots$$

$x_1 =$	$\frac{1}{a_1}x_0 - \frac{a_2}{a_1}x_2 - \frac{a_3}{a_1}x_3$
$x_4 =$	$b_0 + b_1x_1 + b_2x_2 + b_3x_3$

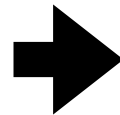
Optimal solution:

$a_0 = 0$ { $x_0 = a_0, x_1 = 0, x_2 = 0, x_3 = 0, x_4 = b_0$

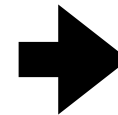
Optimal objective value:

$z = -x_0 = 0$

Original LP:



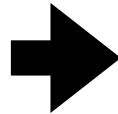
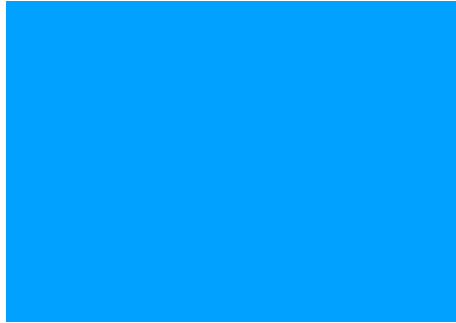
$$z = \dots$$



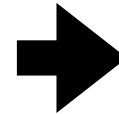
$$z = \dots$$



Auxiliary LP:



$$z = -x_0$$



$$z = \dots$$

$$x_1 = \frac{1}{a_1}x_0 - \frac{a_2}{a_1}x_2 - \frac{a_3}{a_1}x_3$$

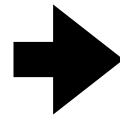
$$x_4 = b_0 + b_1x_1 + b_2x_2 + b_3x_3$$

The constants here are still 0 and $b_0 \geq 0$

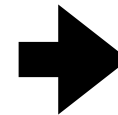


We now have a slack-form LP to start running SIMPLEX Algorithm with.

Original LP:



$$z = \dots$$



$$z = \dots$$

Let $x_0 = 0$

$$x_1 = \frac{1}{a_1}x_0 - \frac{a_2}{a_1}x_2 - \frac{a_3}{a_1}x_3$$

$$x_4 = b_0 + b_1x_1 + b_2x_2 + b_3x_3$$

Is the SIMPLEX Algorithm a polynomial-time algorithm? No

But in practice it is very efficient. Why?

Can LP be solved in polynomial time? Yes