

## P3

- *Name & purpose:*
  - it computes  $n^k$  efficiently by reducing the number of multiplications using exponentiation by squaring
- *Inputs:*
  - **n**: integer (the base)
  - **k**: integer (the exponent)
- *Outputs:*
  - **p**: integer (the result of  $n^k$ )
- *Preconditions:*
  - $k \geq 0$
  - both **n** and **k** are integers
- *Postconditions:*
  - returns the exact value  $p = n^k$
- *High-level idea:*
  - if  $k$  is odd → multiply the result by  $n$  and subtract 1 from  $k$
  - if  $k$  is even → square  $n$  and halve  $k$
  - repeat until  $k$  becomes 0
  - it reduces the number of multiplications by halving  $k$  when possible
- *Pseudocode:*

```
# which is the best/ worst case of the following algorithm?
n = 2
k = 4
p = 1

while k > 0:
    if k % 2 == 1:
        p = p * n
        k = k - 1
    else:
        n = n * n
        k = k // 2
print(p)
```

- *Complexity:*
  - time:  $O(\log k)$
  - space:  $O(1)$
  - best case:  $k$  is a power of 2 (fewest multiplications)
  - worst case:  $k$  is odd (extra multiplications)
- *Correctness sketch:*
  - it works because each step keeps the result equal to  $n^k$ , and when  $k = 0$ ,  $p$  is the final value.

- *Edge cases:*
  - $k = 0 \rightarrow$  returns 1
  - $n = 0, k = 0 \rightarrow$  undefined, assumed 1
  - $n = 0, k > 0 \rightarrow$  returns 0
  - very large numbers may overflow